

# Dividimos y No Conquistamos (D&!C)

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## 1 Template

### 1.1 C++ Template

```

1 #include <bits/stdc++.h>
2 using namespace std;
3 #define L(i, j, n) for (int i = (j); i < (int)n; i++)
4 #define SZ(x) int((x).size())
5 #define ALL(x) begin(x),end(x)
6 #define vec vector
7 #define pb push_back
8 #define eb emplace_back
9 using ll = long long;
10 using ld = long double;
11 void solve(){}
12 int main(){
13     ios::sync_with_stdio(0);cin.tie(0);
14     int TT = 1;
15     //cin >> TT;
16     while (TT--) {solve();}
17 }
18 // IF NEEDED FOR FILE READ
19 // freopen("in.txt", "r", stdin);
20 // freopen("out.txt", "w", stdout);

```

### 1.2 Bash CMD

```

1 co(){g++ $1.cpp -o $1 --std=c++20 -Wall -Wshadow -Wextra}
2 run(){for f in `ls *.txt`;do echo $f ;./$1 < $f; done}
3 #Build, template.cpp must exist!
4 for x in {A..Z}; do mkdir $x; cp template.cpp $x/$x.cpp;done

```

### 1.3 Python Template

```

1 import os, sys, io
2 finput = io.BytesIO(os.read(0, os.fstat(0).st_size)).readline
3 fprintf = sys.stdout.write

```

### 1.4 Java Template

```

1 import java.io.*;
2 import java.util.*;
3 import java.math.BigInteger;
4
5 public class Main {

```

```

6   static BufferedReader br;
7   static PrintWriter pw;
8   static StringTokenizer st;
9
10  public static void main(String[] args) throws IOException {
11      br = new BufferedReader(new FileReader("datos.txt"));
12      pw = new PrintWriter("salida.txt");
13      solve();
14      pw.close();
15  }
16
17  static void solve() throws IOException {
18      // Your code here
19      BigInteger a = nextBigInteger();
20      BigInteger b = nextBigInteger();
21      pw.println(a.add(b));
22  }
23
24  static String next() throws IOException {
25      while (st == null || !st.hasMoreTokens())
26          st = new StringTokenizer(br.readLine());
27      return st.nextToken();
28  }
29
30  static BigInteger nextBigInteger() throws IOException {
31      return new BigInteger(next());
32  }
33
34  static int nextInt() throws IOException {
35      return Integer.parseInt(next());
36  }
37
38  static long nextLong() throws IOException {
39      return Long.parseLong(next());
40  }
41
42  static double nextDouble() throws IOException {
43      return Double.parseDouble(next());
44  }
45
46  static String nextLine() throws IOException {
47      return br.readLine();
48  }
49

```

## 2 Search

### 2.1 Ternary

```

1 // Minimo de 'f' en '(l,r)'.
2 template<class Fun>ll ternary(Fun f, ll l, ll r) {
3     for (ll d = r-l; d > 2; d = r-l) {
4         ll a = l + d/3, b = r - d/3;
5         if (f(a) > f(b)) l = a; else r = b;
6     }
7     return l + 1;
8 }
9 // para error < EPS, usar iters=log((r-l)/EPS)/log(1.618)
10 template<class Fun>double golden(Fun f, double l, double r, int iters){
11     double const ratio = (3-sqrt(5))/2;
12     double x1=l+(r-l)*ratio, x2=r-(r-l)*ratio, f1=f(x1), f2=f(x2);
13     while (iters--) {
14         if (f1 > f2) l=x1, x1=x2, f1=f2, x2=r-(r-l)*ratio, f2=f(x2);
15         else           r=x2, x2=x1, f2=f1, x1=l+(r-l)*ratio, f1=f(x1);
16     }
17     return (l+r)/2;
18 }

```

### 2.2 Simulated Annealing

```

1 using my_clock = chrono::steady_clock;
2 struct Random {
3     mt19937_64 engine;
4     Random(): engine(my_clock::now().time_since_epoch().count()) {}
5     template<class Int>Int integer(Int n) {return integer<Int>(0, n);} // '[0,n)'
6     template<class Int>Int integer(Int l, Int r)
7         {return uniform_int_distribution{l, r-1}(engine);} // '[l,r)'
8     double real() {return uniform_real_distribution{}(engine);} // '[0,1)'
9 } rng;
10 struct Timer {
11     using time = my_clock::time_point;
12     time start = my_clock::now();
13     double elapsed() { // Segundos desde el inicio.
14         time now = my_clock::now();
15         return chrono::duration<double>(now - start).count();
16     }
17 } timer;

```

```

18 template<class See, class Upd>struct Annealing {
19     using energy = invoke_result_t<See>;
20     energy curr, low;
21     See see;
22     Upd upd;
23     Annealing(See _see, Upd _upd): see{_see}, upd{_upd}
24     {curr = low = see(), upd();}
25     void simulate(double s, double mult=1) { // Simula por 's' segundos.
26         double t0 = timer.elapsed();
27         for (double t = t0; t-t0 < s; t = timer.elapsed()) {
28             energy near = see();
29             auto delta = double(curr - near);
30             if (delta >= 0) upd(), curr = near, low = min(low, curr);
31             else {
32                 double temp = mult * (1 - (t-t0)/s);
33                 if (exp(delta/temp) > rng.real()) upd(), curr = near;
34             }
35         }
36     };
37     auto see = [&] -> double {
38         l = rng.integer(gsz); r = rng.integer(gsz);
39         swap(groups[l], groups[r]);
40         int ans = 0, rem = 0;
41         L(i, 0, gsz){
42             if (groups[i] > rem) {
43                 rem = x;
44                 ans++;
45             }
46             rem -= groups[i];
47         }
48         swap(groups[l], groups[r]);
49         return ans;
50     };
51     auto upd = [&] {swap(groups[l], groups[r]);};

```

### 3 Data structures

#### 3.1 Fenwick

```

1 #define LS0(S) (S & -S) //Least significant One
2 struct FT { // 1-Index
3     vec<int> ft; int n;

```

```

4     FT(vec<int> &v): ft(SZ(v)+1), n(SZ(v)+1) { // O(n)
5         L(i, 1, n){
6             ft[i] += v[i-1];
7             if (i + LS0(i) <= n) ft[i + LS0(i)]+=ft[i];
8         }
9     }
10    void update(int pos, int x){ for (int it=pos;it<=n;it+=LS0(it))ft[it]
11        ]+=x; }
12    int sum(int pos){
13        int res = 0;
14        for (int it=pos;it>0;it-=LS0(it))res+=ft[it];
15        return res;
16    }
17    int getSum(int l, int r){return sum(r) - sum(l - 1);}
18 };

```

#### 3.2 Fenwick - 2D

```

1 #define LS0(S) (S & -S)
2 struct BIT { // 1-Index
3     vec<vec<int>> B;
4     int n; // BUILD: N * N * log(N) * log(N)
5     BIT(int n_ = 1): B(n_+1,vec<int>(n_+1)), sz(n_) {}
6     void add(int i, int j, int delta){ // log(N) * log(N)
7         for (int x = i; x <= n; x += LS0(x))
8             for (int y = j; y <= n; y += LS0(y))
9                 B[x][y] += delta;
10    }
11    int sum(int i, int j){ // log(N) * log(N)
12        int tot = 0;
13        for (int x = i; x > 0; x -= LS0(x))
14            for(int y = j; y > 0; y -= LS0(y))
15                tot += B[x][y];
16        return tot;
17    }
18    int getSum(int x1, int y1, int x2, int y2) {return sum(x2, y2) - sum
19        (x2, y1) - sum(x1, y2) + sum(x1-1,y1-1);}
};


```

#### 3.3 DSU

```

1 struct DSU {
2     vec<int> par, sz; int n;
3     DSU(int n = 1): par(n), sz(n, 1), n(n) { iota(ALL(par), 0); }

```

```

4     int find(int a){return a == par[a] ? a : par[a] = find(par[a]);}
5     void join(int a, int b){
6         a=find(a);b=find(b);
7         if (a == b) return;
8         if (sz[b] > sz[a]) swap(a,b);
9         par[b] = a; sz[a] += sz[b];
10    }
11 };

```

### 3.4 Index Compression

```

1 template<class T>
2 struct Index{ // If only 1 use Don't need to copy T type
3     vec<T> d; int sz;
4     Index(const vec<T> &a): d(ALL(a)){
5         sort(ALL(d)); // Sort
6         d.erase(unique(ALL(d)), end(d)); // Erase continuous duplicates
7         sz = SZ(d); }
8     inline int of(T e) const{return lower_bound(ALL(d), e) - begin(d);}
9     // get index
10    inline T at(int i) const{return d[i];} // get value of index
11 };

```

### 3.5 Sparse Table

```

1 struct SPT {
2     vec<vec<int>> st;
3     SPT(vec<int> &a) {
4         int n = SZ(a), K = 0; while((1<<K)<=n) K++;
5         st = vec<vec<int>>(K, vec<int>(n));
6         L(i,0,n) st[0][i] = a[i];
7         L(i,1,K) for (int j = 0; j + (1 << i) <= n; j++)
8             st[i][j] = min(st[i-1][j], st[i - 1][j + (1 << (i - 1))]);
9             // change op
10    }
11    int get(int l, int r) {
12        int bit = log2(r - l + 1);
13        return min(st[bit][l], st[bit][r - (1<<bit) + 1]); // change op
14    }

```

### 3.6 Segment tree

```
#define LC(v) ((v<<1)
```

```

2 #define RC(v) ((v<<1)|1)
3 #define MD(L, R) (L+((R-L)>>1))
4 struct node { ll mx;ll cant; };
5 struct ST {
6     vec<node> st; vec<ll> lz; int n;
7     ST(int n = 1): st(4 * n + 10, {oo, oo}), lz(4 * n + 10, 0), n(n) {
8         build(1, 0, n - 1);}
9     node merge(node a, node b){
10         if (a.mx == oo) return b; if (b.mx == oo) return a;
11         if (a.mx == b.mx) return {a.mx, a.cant + b.cant};
12         return {max(a.mx, b.mx), a.mx > b.mx ? a.cant : b.cant};
13     }
14     void build(int v, int L, int R){
15         if (L == R){ st[v] = {0, 1}; return ;}
16         int m = MD(L, R);
17         build(LC(v), L, m); build(RC(v), m + 1, R);
18         st[v] = merge(st[LC(v)], st[RC(v)]);
19     }
20     void push(int v, int L, int R){
21         if (lz[v]){
22             if (L != R){
23                 st[LC(v)].mx += lz[v]; // Apply to left
24                 st[RC(v)].mx += lz[v]; // And right
25                 lz[LC(v)] += lz[v];
26                 lz[RC(v)] += lz[v];
27             }
28             lz[v] = 0;
29         }
30     }
31     void update(int v, int L, int R, int ql, int qr, ll w){
32         if (ql > R || qr < L) return;
33         push(v, L, R);
34         if (ql == L && qr == R){
35             st[v].mx += w; // Update actual node
36             lz[v] += w; // Add lazy
37             push(v, L, R); // Initial spread
38             return;
39         }
40         int m = MD(L, R);
41         update(LC(v), L, m, ql, min(qr, m), w);
42         update(RC(v), m + 1, R, max(m + 1, ql), qr, w);
43         st[v] = merge(st[LC(v)], st[RC(v)]);
44     }

```

```

44     node query(int v, int L, int R, int ql, int qr){
45         if (ql > R || qr < L) return {oo, oo};
46         push(v, L, R);
47         if (ql == L && qr == R) return st[v];
48         int m = MD(L, R);
49         return merge(query(LC(v), L, m, ql, min(m, qr)), query(RC(v), m
50             + 1, R, max(m + 1, ql), qr));
51     }
51     node query(int l, int r){return query(1, 0, n - 1, l, r);}
52     void update(int l, int r, ll w){update(1, 0, n - 1, l, r, w);}
53 };

```

### 3.7 Segment Tree Iterativo

```

1 struct STI {
2     vec<ll> st; int n, K;
3     STI(vec<ll> &a): n(SZ(a)), K(1) {
4         while(K<=n) K<<=1;
5         st.assign(2*K, 0); // 0 default
6         L(i,0,n) st[K+i] = a[i];
7         for (int i = K - 1; i > 0; i --) st[i] = st[i*2] + st[i*2+1];
8     void upd(int pos, ll w) {
9         pos += K; st[pos] += w;
10        while((pos>>=1) > 0) st[pos] = st[pos * 2] + st[pos * 2 + 1];
11    ll query(int l, int r) { // [l, r]
12        ll res = 0;
13        for (l += K, r += K + 1; l < r; l>>=1, r>>=1){
14            if (l & 1) res += st[l++];
15            if (r & 1) res += st[--r];
16        }
17        return res;
18    }
19 };

```

### 3.8 Segment Tree Persistente

```

1 struct Vertex{Vertex * l, *r;int sum;};
2 const int MVertex = 6000000; // ~= N * logN * 2
3 Vertex pool[MVertex]; // the idea is to keep versions on vec<Vertex*>
4     roots; roots.pb(build(ST_L, ST_R));
4 int p_num = 0;      //
5 Vertex * init_leaf(int x) {
6     pool[p_num].sum = x;
7     pool[p_num].l = pool[p_num].r = NULL;

```

```

8     return &pool[p_num++];
9 }
10 Vertex * init_node(Vertex * l, Vertex * r) {
11     int sum = 0;
12     if (l) sum += l->sum;
13     if (r) sum += r->sum;
14     pool[p_num].sum = sum; pool[p_num].l = l; pool[p_num].r = r;
15     return &pool[p_num++];
16 }
17 Vertex * build(int L, int R){
18     if (L == R){return init_leaf(0);}
19     int m = MD(L, R); return init_node(build(L, m), build(m + 1, R));
20 }
21 Vertex * update(Vertex * v, int L, int R, int pos, int w){
22     if (L == R) return init_leaf(v->sum + w);
23     int m = MD(L, R);
24     if (pos <= m) return init_node(update(v->l, L, m, pos, w), v->r);
25     return init_node(v->l, update(v->r, m + 1, R, pos, w));}
26 int query(Vertex * vl, Vertex * vr, int L, int R, int ql, int qr) {
27     if (!vl || !vr) return 0;
28     if (ql > R || qr < L) return 0;
29     if (ql == L && qr == R) {return vr->sum - vl->sum;}
30     int m = MD(L, R);
31     return query(vl->l, vr->l, L, m, ql, min(m, qr)) +
32         query(vl->r, vr->r, m + 1, R, max(m + 1, ql), qr);}

```

### 3.9 Policy Based

```

1 #include <ext/pb_ds/assoc_container.hpp>
2 using namespace __gnu_pbds;
3 template<typename Key, typename Val=null_type>
4 using indexed_set = tree<Key, Val, less<Key>, rb_tree_tag,
5                             tree_order_statistics_node_update>;
6 // indexed_set<char> s;
7 // char val = *s.find_by_order(0); // acceso por indice
8 // int idx = s.order_of_key('a'); // busca indice del valor
9 template<class Key, class Val=null_type>using htable=gp_hash_table<Key,
10                                     Val>;
// como unordered_map (o unordered_set si Val es vacio), pero sin metodo
// count

```

### 3.10 SQRT Decomposition

```

1 struct SQRTDecomp {
2     vec<int> B, Bs, Bid; int n;// DEFINE BLOCK_SIZE ~= sqrt N

```

```

3     SQRTDecomp(int n_): B(n_), Bid(n_), n(n_), Bs((n_ + BLOCK_SIZE - 1)/
4         BLOCK_SIZE) {
5         L(i,1,n) Bid[i] = Bid[i - 1] + (i % BLOCK_SIZE == 0);
6     } // useful if many updates not many queries, may be better than st
7     void upd(int ix, int w) { B[ix] += w; Bs[Bid[ix]] += w;} // O(1)
8     int query(int l, int r){ // O(BLOCK_SIZE)
9         int ans = 0;
10        for (int i = l; i <= r;) { // [l, r]
11            if (i + BLOCK_SIZE > r || (i % BLOCK_SIZE) != 0) ans += B[i
12                ++];
13            else { ans += Bs[Bid[i]]; i += BLOCK_SIZE;}
14        }
15        return ans;
16    }
17 };

```

### 3.11 Chull Trick

```

1     typedef ll tc;
2     const tc is_query=-(1LL<<62); // special value for query
3     struct Line {
4         tc m,b;
5         mutable multiset<Line>::iterator it,end;
6         const Line* succ(multiset<Line>::iterator it) const {
7             return (++it==end? NULL : &*it);}
8         bool operator<(const Line& rhs) const {
9             if(rhs.b!=is_query) return m<rhs.m;
10            const Line *s=succ(it);
11            if(!s) return 0;
12            return b-s->b<(s->m-m)*rhs.m;
13        }
14    };
15    struct HullDynamic : public multiset<Line> { // for maximum
16        bool bad(iterator y){
17            iterator z=next(y);
18            if(y==begin()){
19                if(z==end())return false;
20                return y->m==z->m&&y->b<=z->b;
21            }
22            iterator x=prev(y);
23            if(z==end())return y->m==x->m&&y->b<=x->b;
24            return 1.0*(x->b-y->b)*(z->m-y->m)>=1.0*(y->b-z->b)*(y->m-x->m);
25        }//Take care of overflow!

```

```

26     iterator next(iterator y){return ++y;}
27     iterator prev(iterator y){return --y;}
28     void add(tc m, tc b){
29         iterator y=insert((Line){m,b});
30         y->it=y;y->end=end();
31         if(bad(y)) erase(y);return;
32         while(next(y)!=end()&&bad(next(y)))erase(next(y));
33         while(y!=begin()&&bad(prev(y)))erase(prev(y));
34     }
35     tc eval(tc x){
36         Line l=*lower_bound((Line){x,is_query});
37         return l.m*x+l.b;
38     }
39 };

```

## 4 Graph

### 4.1 Bellman Ford

```

1     struct Edge {int a, b, cost;};
2     vector<Edge> edges;
3     int solve(int s) // Source
4     {
5         vector<int> d(n, INF);
6         d[s] = 0;
7         for (int i = 0; i < n - 1; ++i)
8             for (Edge e : edges)
9                 if (d[e.a] < INF)
10                     d[e.b] = min(d[e.b], d[e.a] + e.cost);
11     }

```

### 4.2 SCC

```

1     vec<int> dfs_num(N, -1), dfs_low(N, -1), in_stack(N);
2     int dfs_count = 0;
3     int numSCC = 0;
4     stack<int> st;
5     void dfs(int u){
6         dfs_low[u]=dfs_num[u]=dfs_count++;
7         st.push(u);
8         in_stack[u] = 1;
9         for(int v: G[u]) {
10             if (dfs_num[v] == -1) dfs(v);

```

```

11     if (in_stack[v]) dfs_low[u] = min(dfs_low[u], dfs_low[v]);
12 }
13 if (dfs_num[u] == dfs_low[u]){
14     numSCC++;
15     while(1){
16         int v = st.top(); st.pop();
17         in_stack[v] = 0;
18         if (u == v) break;
19     }
20 }
21 }
```

### 4.3 Bipartite Matching Hopcroft-Karp - With Konig

```

1 mt19937 rng((int) chrono::steady_clock::now().time_since_epoch().count())
2
3 struct hopcroft_karp {
4     int n, m; // n is Left Partition Size, m is Right Partition Size
5     vec<vec<int>> g;
6     vec<int> dist, nxt, ma, mb;
7     hopcroft_karp(int n_, int m_) : n(n_), m(m_), g(n),
8         dist(n), nxt(n), ma(n, -1), mb(m, -1) {}
9     void add(int a, int b) { g[a].pb(b); }
10    bool dfs(int i) {
11        for (int &id = nxt[i]; id < g[i].size(); id++) {
12            int j = g[i][id];
13            if (mb[j] == -1 or (dist[mb[j]] == dist[i]+1 and dfs(mb[j]))) {
14                ma[i] = j, mb[j] = i;
15                return true;
16            }
17        }
18        return false;
19    }
20    bool bfs() {
21        for (int i = 0; i < n; i++) dist[i] = n;
22        queue<int> q;
23        for (int i = 0; i < n; i++) if (ma[i] == -1) {
24            dist[i] = 0;
25            q.push(i);
26        }
27        bool rep = 0;
28        while (q.size()) {
29            int i = q.front(); q.pop();
30            for (int j : g[i]) {
31                if (mb[j] == -1) rep = 1;
32                else if (dist[mb[j]] > dist[i] + 1) {
33                    dist[mb[j]] = dist[i] + 1;
34                    q.push(mb[j]);
35                }
36            }
37        }
38        return rep;
39    }
40    int matching() {
41        int ret = 0;
42        for (auto& i : g) shuffle(ALL(i), rng);
43        while (bfs()) {
44            for (int i = 0; i < n; i++) nxt[i] = 0;
45            for (int i = 0; i < n; i++)
46                if (ma[i] == -1 and dfs(i)) ret++;
47        }
48        return ret;
49    }
50    vec<int> cover[2]; // if cover[i][j] = 1 -> node i, j is part of cover
51    int konig() {
52        cover[0].assign(n, 1); // n left size
53        cover[1].assign(m, 0); // m right size
54        auto go = [&](auto&& me, int u) -> void {
55            cover[0][u] = false;
56            for (auto v : g[u]) if (!cover[1][v]) {
57                cover[1][v] = true;
58                me(me, mb[v]);
59            }
60        };
61        L(u, 0, n) if (ma[u] < 0) go(go, u);
62        return size;
63    };
64 }
```

```

29     for (int j : g[i]) {
30         if (mb[j] == -1) rep = 1;
31         else if (dist[mb[j]] > dist[i] + 1) {
32             dist[mb[j]] = dist[i] + 1;
33             q.push(mb[j]);
34         }
35     }
36 }
37
38     return rep;
39 }
40 int matching() {
41     int ret = 0;
42     for (auto& i : g) shuffle(ALL(i), rng);
43     while (bfs()) {
44         for (int i = 0; i < n; i++) nxt[i] = 0;
45         for (int i = 0; i < n; i++)
46             if (ma[i] == -1 and dfs(i)) ret++;
47     }
48     return ret;
49 }
50 vec<int> cover[2]; // if cover[i][j] = 1 -> node i, j is part of cover
51 int konig() {
52     cover[0].assign(n, 1); // n left size
53     cover[1].assign(m, 0); // m right size
54     auto go = [&](auto&& me, int u) -> void {
55         cover[0][u] = false;
56         for (auto v : g[u]) if (!cover[1][v]) {
57             cover[1][v] = true;
58             me(me, mb[v]);
59         }
60     };
61     L(u, 0, n) if (ma[u] < 0) go(go, u);
62     return size;
63 }
```

### 4.4 Hungarian

```

1 using vi = vec<int>;
2 using vd = vec<ld>;
3 const ld INF = 1e100; // Para max asignacion, INF = 0, y negar costos
4 bool zero(ld x) {return fabs(x) < 1e-9;} // Para int/ll: return x==0;
5 vec<pii> ans; // Guarda las aristas usadas en el matching: [0..n)x[0..m]
```

```

6 struct Hungarian{
7     int n; vec<vd> cs; vi vL, vR;
8     Hungarian(int N, int M) : n(max(N,M)), cs(n, vd(n)), vL(n), vR(n){
9         L(x, 0, N) L(y, 0, M) cs[x][y] = INF;
10    }
11    void set(int x, int y, ld c) { cs[x][y] = c; }
12    ld assign(){
13        int mat = 0; vd ds(n), u(n), v(n); vi dad(n), sn(n);
14        L(i, 0, n) u[i] = *min_element(ALL(cs[i]));
15        L(j, 0, n){
16            v[j] = cs[0][j]-u[0];
17            L(i, 1, n) v[j] = min(v[j], cs[i][j] - u[i]);
18        }
19        vL = vR = vi(n, -1);
20        L(i, 0, n) L(j, 0, n) if(vR[j] == -1 and zero(cs[i][j] - u[i] - v[j]))
21            ){
22                vL[i] = j; vR[j] = i; mat++; break;
23            }
24        for(; mat < n; mat ++){
25            int s = 0, j = 0, i;
26            while(vL[s] != -1) s++;
27            fill(ALL(dad), -1); fill(ALL(sn), 0);
28            L(k, 0, n) ds[k] = cs[s][k]-u[s]-v[k];
29            while(true){
30                j = -1;
31                L(k, 0, n) if(!sn[k] and (j == -1 or ds[k] < ds[j])) j = k;
32                sn[j] = 1; i = vR[j];
33                if(i == -1) break;
34                L(k, 0, n) if(!sn[k]){
35                    auto new_ds = ds[j] + cs[i][k] - u[i]-v[k];
36                    if(ds[k] > new_ds) ds[k]=new_ds, dad[k]=j;
37                }
38                L(k, 0, n) if(k!=j and sn[k]){
39                    auto w = ds[k]-ds[j]; v[k] += w, u[vR[k]] -= w;
40                }
41                u[s] += ds[j];
42                while(dad[j] >= 0){ int d = dad[j]; vR[j] = vR[d]; vL[vR[j]] = j;
43                    j = d; }
44                vR[j] = s; vL[s] = j;
45            }
46            ld value = 0; L(i, 0, n) value += cs[i][vL[i]], ans.pb({i, vL[i]});
47            return value;
48        }
49    }
50}

```

```

47    }
48 }

```

## 4.5 Flow - Dinics

```

1 const int oo = (int)1e9;
2 struct Dinic {
3     bool scaling = false; // com scaling -> O(nm log(MAXCAP)),
4                                         // com constante alta
5     int lim;
6     struct edge {
7         int to, cap, rev, flow;
8         bool res;
9         edge(int to_, int cap_, int rev_, bool res_)
10            : to(to_), cap(cap_), rev(rev_), flow(0), res(res_) {}
11    };
12    vec<vec<edge>> g;
13    vec<int> lev, beg;
14    ll F;
15    Dinic(int n) : g(n), F(0) {}
16    void add(int a, int b, int c) {
17        g[a].emplace_back(b, c, g[b].size(), false);
18        g[b].emplace_back(a, 0, g[a].size()-1, true);
19    }
20    bool bfs(int s, int t) {
21        lev = vector<int>(g.size(), -1); lev[s] = 0;
22        beg = vector<int>(g.size(), 0);
23        queue<int> q; q.push(s);
24        while (q.size()) {
25            int u = q.front(); q.pop();
26            for (auto& i : g[u]) {
27                if (lev[i.to] != -1 or (i.flow == i.cap)) continue;
28                if (scaling and i.cap - i.flow < lim) continue;
29                lev[i.to] = lev[u] + 1;
30                q.push(i.to);
31            }
32        }
33        return lev[t] != -1;
34    }
35    int dfs(int v, int s, int f = oo) {
36        if (!f or v == s) return f;
37        for (int& i = beg[v]; i < g[v].size(); i++) {
38            auto& e = g[v][i];
39            if (lev[e.to] != lev[v] + 1) continue;
40            int d = min(f, e.cap - e.flow);
41            f -= d;
42            e.flow += d;
43            g[e.to][e.rev].flow -= d;
44        }
45    }
46}

```

```

39     int foi = dfs(e.to, s, min(f, e.cap - e.flow));
40     if (!foi) continue;
41     e.flow += foi, g[e.to][e.rev].flow -= foi;
42     return foi;
43   }
44   return 0;
45 }
46 ll max_flow(int s, int t) {
47   for (lim = scaling ? (1<<30) : 1; lim; lim /= 2)
48     while (bfs(s, t)) while (int ff = dfs(s, t)) F += ff;
49   return F;
50 }
51 };
52 vec<pair<int, int>> get_cut(Dinic& g, int s, int t) {
53   g.max_flow(s, t);
54   vec<pair<int, int>> cut;
55   vec<int> vis(g.g.size(), 0), st = {s};
56   vis[s] = 1;
57   while (st.size()) {
58     int u = st.back(); st.pop_back();
59     for (auto e : g.g[u]) if (!vis[e.to] and e.flow < e.cap)
60       vis[e.to] = 1, st.push_back(e.to);
61   }
62   for (int i = 0; i < g.g.size(); i++) for (auto e : g.g[i])
63     if (vis[i] and !vis[e.to] and !e.res) cut.emplace_back(i, e.to);
64   return cut;
65 }

```

## 4.6 Flow - MinCostMaxFlow

```

1 // O(nm + f * m log n)
2 // const ll oo = (ll)1e18;
3 template<typename T> struct mcmf {
4   struct edge {
5     int to, rev, flow, cap; // para, id da reversa, fluxo, capacidade
6     bool res; // se eh reversa
7     T cost; // custo da unidade de fluxo
8     edge() : to(0), rev(0), flow(0), cap(0), cost(0), res(false) {}
9     edge(int to_, int rev_, int flow_, int cap_, T cost_, bool res_)
10      : to(to_), rev(rev_), flow(flow_), cap(cap_), res(res_), cost(
11        cost_) {}
12   };
13   vec<vec<edge>> g;

```

```

13   vec<int> par_idx, par;
14   T inf;
15   vec<T> dist;
16   mcmf(int n) : g(n), par_idx(n), par(n), inf(numeric_limits<T>::max())
17     /3) {}
18   void add(int u, int v, int w, T cost) { // de u pra v com cap w e
19     edge a = edge(v, g[v].size(), 0, w, cost, false);
20     edge b = edge(u, g[u].size(), 0, 0, -cost, true);
21     g[u].push_back(a);
22     g[v].push_back(b);
23   }
24   vec<T> spfa(int s) { // nao precisa se nao tiver custo negativo
25     deque<int> q;
26     vec<bool> is_inside(g.size(), 0);
27     dist = vec<T>(g.size(), inf);
28     dist[s] = 0;
29     q.push_back(s);
30     is_inside[s] = true;
31     while (!q.empty()) {
32       int v = q.front();
33       q.pop_front();
34       is_inside[v] = false;
35       for (int i = 0; i < g[v].size(); i++) {
36         auto [to, rev, flow, cap, res, cost] = g[v][i];
37         if (flow < cap and dist[v] + cost < dist[to]) {
38           dist[to] = dist[v] + cost;
39           if (is_inside[to]) continue;
40           if (!q.empty() and dist[to] > dist[q.front()]) q.push_back(to)
41             ;
42           else q.push_front(to);
43           is_inside[to] = true;
44         }
45       }
46     }
47   }
48   bool dijkstra(int s, int t, vec<T>& pot) {
49     priority_queue<pair<T, int>, vec<pair<T, int>>, greater<>> q;
50     dist = vec<T>(g.size(), inf);
51     dist[s] = 0;
52     q.emplace(0, s);

```

```

53     while (q.size()) {
54         auto [d, v] = q.top();
55         q.pop();
56         if (dist[v] < d) continue;
57         for (int i = 0; i < g[v].size(); i++) {
58             auto [to, rev, flow, cap, res, cost] = g[v][i];
59             cost += pot[v] - pot[to];
60             if (flow < cap and dist[v] + cost < dist[to]) {
61                 dist[to] = dist[v] + cost;
62                 q.emplace(dist[to], to);
63                 par_idx[to] = i, par[to] = v;
64             }
65         }
66     }
67     return dist[t] < inf;
68 }
69 pair<int, T> min_cost_flow(int s, int t, int flow = (int)1e9) {
70     vec<T> pot(g.size(), 0);
71     pot = spfa(s); // mudar algoritmo de caminho minimo aqui
72     int f = 0;
73     T ret = 0;
74     while (f < flow and dijkstra(s, t, pot)) {
75         for (int i = 0; i < g.size(); i++)
76             if (dist[i] < inf) pot[i] += dist[i];
77         int mn_flow = flow - f, u = t;
78         while (u != s){
79             mn_flow = min(mn_flow,
80                           g[par[u]][par_idx[u]].cap - g[par[u]][par_idx[u]].flow);
81             u = par[u];
82         }
83         ret += pot[t] * mn_flow;
84         u = t;
85         while (u != s) {
86             g[par[u]][par_idx[u]].flow += mn_flow;
87             g[u][g[par[u]][par_idx[u]].rev].flow -= mn_flow;
88             u = par[u];
89         }
90         f += mn_flow;
91     }
92     return make_pair(f, ret);
93 }
94 // Opcional: retorna as arestas originais por onde passa flow = cap
95 vec<pair<int,int>> recover() {

```

```

96     vec<pair<int,int>> used;
97     for (int i = 0; i < g.size(); i++) for (edge e : g[i])
98         if(e.flow == e.cap && !e.res) used.push_back({i, e.to});
99     return used;
100 }
101 };

```

## 4.7 2 Sat

```

1 struct TwoSat {
2     int n, v_n;
3     vec<bool> vis, assign;
4     vec<int> order, comp;
5     vec<vec<int>> g, g_t;
6     TwoSat(int n_): n(n_), v_n(2 * n_), vis(v_n), assign(n_), comp(v_n
7         , -1), g(v_n), g_t(v_n) {
8         order.reserve(v_n);
9     }
10    void add_disj(int a, bool na, int b, bool nb) { // negated_a,
11        negated_b
12        a = 2 * a ^ na;
13        b = 2 * b ^ nb;
14        int neg_a = a ^ 1;
15        int neg_b = b ^ 1;
16        g[neg_a].pb(b);
17        g[neg_b].pb(a);
18        g_t[a].pb(neg_b);
19        g_t[b].pb(neg_a);
20    }
21    void dfs1(int u){
22        vis[u] = 1;
23        for (int v: g[u]) if (!vis[v]) dfs1(v);
24        order.pb(u);
25    }
26    void dfs2(int u, int cc) {
27        comp[u] = cc;
28        for (int v: g_t[u]) if (comp[v] == -1) dfs2(v, cc);
29    }
30    bool solve() {
31        order.clear();
32        vis.assign(v_n, 0);
33        L(i,0, v_n) if (!vis[i]) dfs1(i);
34        comp.assign(v_n, -1);
35    }
36 };

```

```

33     int cc = 0;
34     L(i, 0, v_n) {
35         int v = order[v_n - 1 - i];
36         if (comp[v] == -1) dfs2(v, cc++);
37     }
38     assign.assign(n, false);
39     for (int i = 0; i < v_n; i += 2) {
40         if (comp[i] == comp[i+1]) return false;
41         assign[i / 2] = comp[i] > comp[i + 1];
42     }
43     return true;
44 }
45 };

```

## 4.8 Euler Tour

```

1 // Directed version (uncomment commented code for undirected)
2 struct edge {
3     int y;
4     // list<edge>::iterator rev;
5     edge(int y):y(y){}
6 };
7 list<edge> g[N];
8 void add_edge(int a, int b){
9     g[a].push_front(edge(b)); //auto ia=g[a].begin();
10    // g[b].push_front(edge(a));auto ib=g[b].begin();
11    // ia->rev=ib;ib->rev=ia;
12 }
13 vec<int> p;
14 void go(int x){
15     while(g[x].size()){
16         int y=g[x].front().y;
17         //g[y].erase(g[x].front().rev);
18         g[x].pop_front();
19         go(y);
20     }
21     p.push_back(x);
22 }
23 vec<int> get_path(int x){ // get a path that begins in x
24 // check that a path exists from x before calling to get_path!
25     p.clear();go(x);reverse(p.begin(),p.end());
26     return p;
27 }

```

## 5 Trees

### 5.1 Heavy Light Decomposition

```

1 int ans[N], par[N], depth[N], head[N], pos[N];
2 vec<int> heavy(N, -1);
3 int t = 0;
4 vec<int> g[N];
5 int dfs(int u) {
6     int size = 1;
7     int max_size = 0;
8     for (int v: g[u]) if (v != par[u]) {
9         par[v] = u;
10        depth[v] = depth[u] + 1;
11        int cur_size = dfs(v);
12        size += cur_size;
13        if (cur_size > max_size) {
14            max_size = cur_size;
15            heavy[u] = v;
16        }
17    }
18    return size;
19 }
20 void decompose(int u, int h){
21     head[u] = h;
22     pos[u] = t++;
23     if (heavy[u] != -1){ decompose(heavy[u], h); }
24     for (int v: g[u]) if (v != par[u] && v != heavy[u]) {
25         decompose(v, v);
26     }
27 }
28 int query(int a, int b) {
29     int resp = -1;
30     for (; head[a] != head[b]; b = par[head[b]]){ // Subi todo el heavy
31         path y a su padre // Next
32         if (depth[head[a]] > depth[head[b]]) swap(a, b);
33         resp = max(resp, st.query(pos[head[b]], pos[b])); // pos[head[b]]
34         if (depth[a] > depth[b]) swap(a, b); // Una vez misma path(head)
35         entones es una query [a,b]
36         resp = max(resp, st.query(pos[a], pos[b]));
37     }
38 }

```

```

37 }
38 dfs(root);
39 decompose(root, root);

```

## 5.2 Centroid

```

1 int sz[N];
2 bool removed[N];
3 int getSize(int u, int p){
4     sz[u] = 1;
5     for(int v: G[u]) if (v != p && !removed[v]){
6         sz[u] += getSize(v, u);
7     }
8     return sz[u];
9 }
10 int centroid(int u, int p, int tz){
11     for (int v: g[u])
12         if (v != p && !removed[v] && sz[v] * 2 > tz) return centroid(v,
13             u, tz);
14     return u;
15 }
16 int build(int u){
17     int c = centroid(u, -1, getSize(u, -1));
18     removed[c] = 1;
19     for (int v: G[c]) if (!removed[v]) { build(v); }
20     return c;
}

```

## 5.3 LCA - Binary exponentiation

```

1 vec<int> g[N];
2 int K; // K should be (1<<K) > n
3 int jump[20][N];
4 int depth[N];
5
6 void dfs(int u, int p){
7     for (int v: g[u]) if (v != p) {
8         jump[0][v] = u;
9         L(i, 1, K + 1) {
10             jump[i][v] = -1;
11             if (jump[i - 1][v] != -1) {
12                 jump[i][v] = jump[i - 1][jump[i - 1][v]];
13             }
14         }
}

```

```

15     depth[v] = depth[u] + 1;
16     dfs(v, u);
17 }
18
19
20 int LCA(int u, int v){
21     if (depth[u] < depth[v]) swap(u, v); // Make u the deepest
22     for (int i = K; i >= 0; i --){ // make them same depth
23         if (jump[i][u] != -1 && depth[jump[i][u]] >= depth[v]){
24             u = jump[i][u];
25         }
26     }
27     if (u == v) return u; // u is parent of v
28     for (int i = K; i >= 0; i --){
29         if (jump[i][u] != jump[i][v] && jump[i][u] != -1 && jump[i][v]
30             != -1){
31             u = jump[i][u];
32             v = jump[i][v];
33         }
34     }
35     return jump[0][u];
}

```

## 5.4 LCA - Const Time

```

1 struct LCA {
2     vec<int> depth, in, euler;
3     vec<vec<int>> g, st;
4     int K, n;
5     inline int Min(int i, int j) {return depth[i] <= depth[j] ? i : j;}
6     void dfs(int u, int p) {
7         in[u] = SZ(euler);
8         euler.pb(u);
9         for (int v: g[u]) if (v != p){
10             depth[v] = depth[u] + 1;
11             dfs(v, u);
12             euler.pb(u);
13         }
14     }
15     LCA(int n_): depth(n_), g(vec<vec<int>>(n_)), K(0), n(n_), in(n_) {
16         euler.reserve(2 * n); }
17     void add_edge(int u, int v) {g[u].pb(v);}
    void build(int root){}
}

```

```

18     dfs(root, -1);
19     int ln = SZ(euler);
20     while((1<<K)<=ln)K++;
21     st = vec<vec<int>>(K, vec<int>(ln));
22     L(i,0,ln) st[0][i] = euler[i];
23     for (int i = 1; (1 << i) <= ln; i++) {
24         for (int j = 0; j + (1<<i) <= ln; j++) {
25             st[i][j] = Min(st[i-1][j], st[i-1][j + (1<<(i-1))]);
26         }
27     }
28 }
29 int get(int u, int v) {
30     int su = in[u];
31     int sv = in[v];
32     if (sv < su) swap(sv, su);
33     int bit = log2(sv - su + 1);
34     return Min(st[bit][su], st[bit][sv - (1<<bit) + 1]);
35 }
36 };

```

## 6 Dynamic Programming

### 6.1 Knapsack

```

1 int knapsack(vector<int>& values, vector<int>& weights, int W) {
2     int n = values.size();
3     vector<vector<int>> dp(n + 1, vector<int>(W + 1, 0));
4
5     for(int i = 1; i <= n; i++) {
6         for(int w = 0; w <= W; w++) {
7             if(weights[i-1] <= w) {
8                 dp[i][w] = max(dp[i-1][w],
9                                 dp[i-1][w-weights[i-1]] + values[i-1]);
10            } else {
11                dp[i][w] = dp[i-1][w];
12            }
13        }
14    }
15    return dp[n][W];
16 }

```

### 6.2 LIS

```

1 vector<int> getLIS(vector<int>& arr) {
2     int n = arr.size();
3     vector<int> dp(n + 1, INT_MAX); // dp[i] = smallest value that ends
4                                         an LIS of length i
5     vector<int> len(n);           // Length of LIS ending at each
6                                         position
7     dp[0] = INT_MIN;
8     for(int i = 0; i < n; i++) {
9         int j = upper_bound(dp.begin(), dp.end(), arr[i]) - dp.begin();
10        dp[j] = arr[i];
11        len[i] = j;
12    }
13 // Find maxLen and reconstruct sequence
14 int maxLen = 0;
15 for(int i = n-1; i >= 0; i--) maxLen = max(maxLen, len[i]);
16 vector<int> lis;
17 for(int i = n-1, currLen = maxLen; i >= 0; i--) {
18     if(len[i] == currLen) {
19         lis.push_back(arr[i]);
20         currLen--;
21     }
22 }
23 reverse(lis.begin(), lis.end());
24 return lis;
25 }

```

### 6.3 Edit Distance

```

1 int editDistance(string& s1, string& s2) {
2     int n = s1.length(), m = s2.length();
3     vector<vector<int>> dp(n + 1, vector<int>(m + 1));
4     for(int i = 0; i <= n; i++) dp[i][0] = i;
5     for(int j = 0; j <= m; j++) dp[0][j] = j;
6     for(int i = 1; i <= n; i++) {
7         for(int j = 1; j <= m; j++) {
8             if(s1[i-1] == s2[j-1]) {
9                 dp[i][j] = dp[i-1][j-1];
10            } else {
11                dp[i][j] = 1 + min({dp[i-1][j], // deletion
12                                     dp[i][j-1], // insertion
13                                     dp[i-1][j-1]}); // replacement
14            }
15        }
16    }
17 }

```

```

16     }
17     return dp[n][m];
18 }
```

## 6.4 Kadane

```

1 pair<int, pair<int,int>> kadane(vector<int>& arr) {
2     int maxSoFar = arr[0], maxEndingHere = arr[0];
3     int start = 0, end = 0, s = 0;
4
5     for(int i = 1; i < arr.size(); i++) {
6         if(maxEndingHere + arr[i] < arr[i]) {
7             maxEndingHere = arr[i];
8             s = i;
9         } else {
10            maxEndingHere += arr[i];
11        }
12
13        if(maxEndingHere > maxSoFar) {
14            maxSoFar = maxEndingHere;
15            start = s;
16            end = i;
17        }
18    }
19    return {maxSoFar, {start, end}}; // max, l, r
20 }
```

# 7 Strings

## 7.1 Hashing

```

1 static constexpr ll ms[] = {1'000'000'007, 1'000'000'403};
2 static constexpr ll b = 500'000'000;
3 struct StrHash { // Hash polinomial con exponentes decrecientes.
4     vec<ll> hs[2], bs[2];
5     StrHash(string const& s) {
6         int n = SZ(s);
7         L(k, 0, 2) {
8             hs[k].resize(n+1), bs[k].resize(n+1, 1);
9             L(i, 0, n) {
10                hs[k][i+1] = (hs[k][i] * b + s[i]) % ms[k];
11                bs[k][i+1] = bs[k][i] * b % ms[k];
12            }
13        }
14    }
15 }
```

```

13     }
14 }
15 ll get(int idx, int len) const { // Hashes en 's[idx, idx+len)'.
16     ll h[2];
17     L(k, 0, 2) {
18         h[k] = hs[k][idx+len] - hs[k][idx] * bs[k][len] % ms[k];
19         if (h[k] < 0) h[k] += ms[k];
20     }
21     return (h[0] << 32) | h[1];
22 }
23 }
24
25 pll union_hash(vec<pll> hs, vec<ll> lens){ //use arrays makes it slower
26     ll len = 0;
27     for(int i = hs.size()-1; i > 0; i--) {
28         len += lens[i];
29         pll& [l1, l2] = hs[i];
30         pll& [r1, r2] = hs[i-1];
31         l1 = ((l1 * binpow(b, len, ms[0])) % ms[0] + r1) % ms[0];
32         l2 = ((l2 * binpow(b, len, ms[1])) % ms[1] + r2) % ms[1];
33     }
34
35     return hs[0];
36 }
```

## 7.2 Trie

```

1 struct Trie {
2     map<char, int> ch;
3     bool eee;
4     Trie(): eee(0) {}
5 };
6 vec<Trie> t;
7 void initTrie(){t.clear(); t.pb(Trie());}
8 void insert(string &word) {
9     int v = 0;
10    for(char c : word) {
11        if(!t[v].ch[c]) {
12            t[v].ch[c] = SZ(t);
13            t.pb(Trie());
14        }
15        v = t[v].ch[c];
16    }
17 }
```

```

17     t[v].eee = 1;
18 }
```

### 7.3 KMP

```

1 vec<int> KMP(const string &s){
2     int n = SZ(s); vec<int> pi(n);
3     L(i,1,n){
4         int j = pi[i - 1];
5         while(j>0&&s[i]!=s[j]) j = pi[j-1];
6         if (s[i]==s[j])j++;
7         pi[i]=j;
8     }
9     return pi;
10 }
```

### 7.4 LPS

```

1 vec<int> getLps(string pat){ //geek4geeks implementatio with some
2     changes
3     vec<int> lps(pat.length(), 0);
4     int len = 0;
5     int i = 1;
6     lps[0] = 0;
7     while(i < pat.length()){
8         if(pat[i] == pat[len]){
9             len++;
10            lps[i] = len;
11            i++;
12        }
13        else //pat[i] != pat[len]
14        {
15            lps[i] = 0;
16            i++;
17        }
18    }
19    return lps;
}
```

### 7.5 Z-FUNCTION

```

1 template<class Char=char>vec<int> zfun(const basic_string<Char>& w) {
2     int n = SZ(w), l = 0, r = 0; vec<int> z(n);
3     z[0] = w.length();
```

```

4     L(i, 1, n) {
5         if (i <= r) {z[i] = min(r - i + 1, z[i - 1]);}
6         while (i + z[i] < n && w[z[i]] == w[i + z[i]]) {++z[i];}
7         if (i + z[i] - 1 > r) {l = i, r = i + z[i] - 1;}
8     }
9     return z;
10 }
```

### 7.6 Manacher

```

1 struct Manacher {
2     vec<int> p;
3     Manacher(string const& s) {
4         int n = SZ(s), m = 2*n+1, l = -1, r = 1;
5         vec<char> t(m); L(i, 0, n) t[2*i+1] = s[i];
6         p.resize(m); L(i, 1, m) {
7             if (i < r) p[i] = min(r-i, p[l+r-i]);
8             while (p[i] <= i && i < m-p[i] && t[i-p[i]] == t[i+p[i]]) ++p[i];
9             if (i+p[i] > r) l = i-p[i], r = i+p[i];
10        }
11    } // Retorna palindromos de la forma {comienzo, largo}.
12    pii at(int i) const {int k = p[i]-1; return pair{i/2-k/2, k};}
13    pii odd(int i) const {return at(2*i+1);} // Mayor centrado en s[i].
14    pii even(int i) const {return at(2*i);} // Mayor centrado en s[i-1,i].
15 }
```

### 7.7 Aho-Corasick

```

1 bool vis[N], r[N];
2 struct ACvertex {
3     map<char,int> next,go;
4     int p,link;
5     char pch;
6     vec<int> leaf;
7     ACACvertex(int p=-1, char pch=-1):p(p),pch(pch),link(-1){}
8 };
9 vec<ACvertex> t;
10 void aho_init(){ //do not forget!!
11     t.clear();t.pb(ACvertex());
12 }
13 void add_string(string &s, int id){
14     int v=0;
15     for(char c:s){
16         if(!t[v].next.count(c)){
```

```

17     t[v].next[c]=t.size();
18     t.pb(ACvertex(v,c));
19 }
20 v=t[v].next[c];
21 }
22 t[v].leaf.pb(id);
23 }
24 int go(int v, char c);
25 int get_link(int v){ // Failure link
26     if(t[v].link<0)
27         if(!v||!t[v].p)t[v].link=0;
28     else t[v].link=go(get_link(t[v].p),t[v].pch);
29     return t[v].link;
30 }
31 int go(int v, char c){ // state = go(state, ch) this state is ACvertex
32     id
33     if(!t[v].go.count(c))
34         if(t[v].next.count(c))t[v].go[c]=t[v].next[c];
35     else t[v].go[c]=v==0?0:go(get_link(v),c);
36     return t[v].go[c];
37 }
38 void proc(int x){
39     if (x == -1|| vis[x]) return;
40     vis[x] = 1;
41     L(i,0,SZ(t[x].leaf)) r[t[x].leaf[i]] = 1;
42     proc(get_link(x));
}

```

## 7.8 Suffix-Array

```

1 #define RB(x) ((x) < n ? r[x] : 0)
2 void csort(vec<int>& sa, vec<int>& r, int k) {
3     int n = SZ(sa);
4     vec<int> f(max(255, n)), t(n);
5     L(i,0, n) ++f[RB(i+k)];
6     int sum = 0;
7     L(i,0, max(255, n)) f[i] = (sum += f[i]) - f[i];
8     L(i,0, n) t[f[RB(sa[i]+k)]++] = sa[i];
9     sa = t;
10 }
11 vec<int> compute_sa(string& s){ // O(n*log2(n))
12     int n = SZ(s) + 1, rank;
13     vec<int> sa(n), r(n), t(n);

```

```

14     iota(all(sa), 0);
15     L(i,0, n) r[i] = s[i];
16     for (int k = 1; k < n; k *= 2) {
17         csort(sa, r, k), csort(sa, r, 0);
18         t[sa[0]] = rank = 0;
19         L(i, 1, n) {
20             if(r[sa[i]] != r[sa[i-1]] || RB(sa[i]+k) != RB(sa[i-1]+k)) ++rank;
21             t[sa[i]] = rank;
22         }
23         r = t;
24         if (r[sa[n-1]] == n-1) break;
25     }
26     return sa; // sa[i] = i-th suffix of s in lexicographical order
27 }
28 vec<int> compute_lcp(string& s, vec<int>& sa){
29     int n = SZ(s) + 1, K = 0;
30     vec<int> lcp(n), plcp(n), phi(n);
31     phi[sa[0]] = -1;
32     L(i, 1, n) phi[sa[i]] = sa[i-1];
33     L(i,0,n) {
34         if (phi[i] < 0) { plcp[i] = 0; continue; }
35         while(s[i+K] == s[phi[i]+K]) ++K;
36         plcp[i] = K;
37         K = max(K - 1, 0);
38     }
39     L(i,0, n) lcp[i] = plcp[sa[i]];
40     return lcp; // lcp[i] = longest common prefix between sa[i-1] and sa[i]
41 }

```

## 8 Math

### 8.1 Euclidean Extended

```

1 ll extendedGCD(ll a, ll b, ll &x, ll &y) {
2     if (b == 0) {
3         x = 1;
4         y = 0;
5         return a;
6     }
7     ll x1, y1;
8     ll gcd = extendedGCD(b, a % b, x1, y1);
9     x = y1;

```

```

10     y = x1 - (a / b) * y1;
11     return gcd;
12 }
13
14 bool findSolutionWithConstraints(ll a, ll b, ll c, ll x_min, ll y_min,
15     ll &x, ll &y) {
16     ll g = extendedGCD(a, b, x, y);
17
18     if (c % g != 0) return false;
19
20     x *= c / g;
21     y *= c / g;
22
23     // Ajustamos las variables a/g y b/g para mover las soluciones
24     a /= g;
25     b /= g;
26
27     if (x < x_min) {
28         ll k = (x_min - x + b - 1) / b; // Redondeo hacia arriba
29         x += k * b;
30         y -= k * a;
31     } else if (x > x_min) {
32         ll k = (x - x_min) / b;
33         x -= k * b;
34         y += k * a;
35     }
36
37     if (y < y_min) {
38         ll k = (y_min - y + a - 1) / a; // Redondeo hacia arriba
39         x += k * b;
40         y -= k * a;
41     } else if (y > y_min) {
42         ll k = (y - y_min) / a;
43         x -= k * b;
44         y += k * a;
45     }
46
47     return x >= x_min && y >= y_min;
}

```

## 8.2 Euler Totient

```

2 using namespace std;
3 typedef long long ll;
4
5
6 vector<ll> compute_totients(ll n) {
7     vector<ll> phi(n + 1);
8     for (ll i = 0; i <= n; i++) {
9         phi[i] = i;
10    }
11
12    for (ll i = 2; i <= n; i++) {
13        if (phi[i] == i) { // i es primo
14            for (ll j = i; j <= n; j += i) {
15                phi[j] = phi[j] * (i - 1) / i;
16            }
17        }
18    }
19
20    return phi;
21 }

```

## 8.3 Josephus

```

1 #include <iostream>
2 using namespace std;
3
4 typedef long long ll;
5
6 ll josephus_iterative(ll n, ll k) {
7     ll result = 0;
8     for (ll i = 2; i <= n; ++i) {
9         result = (result + k) % i;
10    }
11    return result;
12 }
13
14
15 ll josephus_recursive(ll n, ll k) {
16
17     if (n == 1)
18         return 0;
19
20     return (josephus_recursive(n - 1, k) + k) % n;
}

```

```

21 }
22
23
24 ll josephus_power_of_2(ll n) {
25
26     ll power = 1;
27     while (power <= n) {
28         power <= 1;
29     }
30     power >= 1;
31
32     return 2 * (n - power);
33 }
34

```

## 8.4 Mobius

```

26 ll mobius(ll x) {
27     ll count = 0;
28     for (ll i = 2; i * i <= x; i++) {
29         if (x % (i * i) == 0)
30             return 0;
31         if (x % i == 0) {
32             count++;
33             x /= i;
34         }
35     }
36
37     if (x > 1) count++;
38
39     return (count % 2 == 0) ? 1 : -1;
40 }

```

## 8.5 NTT

```

1 #include <bits/stdc++.h>
2 using namespace std;
3 typedef long long ll;
4
5
6 vector<ll> compute_mobius(ll n) {
7     vector<ll> mu(n + 1, 1);
8     vector<bool> is_prime(n + 1, true);
9
10    for (ll i = 2; i <= n; i++) {
11        if (is_prime[i]) { // i es un primo
12            for (ll j = i; j <= n; j += i) {
13                mu[j] *= -1; // Multiplicamos por -1 para cada primo
14                is_prime[j] = false;
15            }
16            for (ll j = i * i; j <= n; j += i * i) {
17                mu[j] = 0; // Si tiene un cuadrado de un primo, se pone
18                en 0
19            }
20        }
21    }
22
23    return mu;
24 }
25

```

```

1 #include <bits/stdc++.h>
2 using namespace std;
3 using cd = complex<double>;
4 typedef long long ll;
5 const ll mod = 998244353;
6 const ll root = 31;
7 const ll root_1 = inverse(root, mod);
8 const ll root_pw = 1 << 23;
9
10 ll inverse(ll a, ll m) {
11     ll res = 1, exp = m - 2;
12     while (exp) {
13         if (exp % 2 == 1) res = (1LL * res * a) % m;
14         a = (1LL * a * a) % m;
15         exp /= 2;
16     }
17     return res;
18 }
19
20 void ntt(vector<ll> & a, bool invert) {
21     int n = a.size();
22
23     for (int i = 1, j = 0; i < n; i++) {
24         int bit = n >> 1;
25         for (; j & bit; bit >>= 1)

```

```

26     j ^= bit;
27     j ^= bit;
28
29     if (i < j)
30         swap(a[i], a[j]);
31 }
32
33 for (int len = 2; len <= n; len <= 1) {
34     int wlen = invert ? root_1 : root;
35     for (int i = len; i < root_pw; i <= 1)
36         wlen = (int)(1LL * wlen * wlen % mod);
37
38     for (int i = 0; i < n; i += len) {
39         int w = 1;
40         for (int j = 0; j < len / 2; j++) {
41             int u = a[i+j], v = (int)(1LL * a[i+j+len/2] * w % mod);
42             a[i+j] = u + v < mod ? u + v : u + v - mod;
43             a[i+j+len/2] = u - v >= 0 ? u - v : u - v + mod;
44             w = (int)(1LL * w * wlen % mod);
45         }
46     }
47
48     if (invert) {
49         int n_1 = inverse(n, mod);
50         for (auto & x : a)
51             x = (int)(1LL * x * n_1 % mod);
52     }
53 }
54
55 vector<ll> multiply(vector<ll> const &a, vector<ll> const &b) {
56     vector<ll> fa(a.begin(), a.end()), fb(b.begin(), b.end());
57     ll n = 1;
58     while (n < a.size() + b.size())
59         n <= 1;
60     fa.resize(n);
61     fb.resize(n);
62
63     ntt(fa, false);
64     ntt(fb, false);
65     for (ll i = 0; i < n; i++)
66         fa[i] = (fa[i] * fb[i]) % mod;
67     ntt(fa, true);
68 }
```

```

69
70     vector<ll> result(n);
71     for (ll i = 0; i < n; i++)
72         result[i] = fa[i];
73     return result;
74 }
```

## 8.6 FFT

```

1  typedef long long ll;
2  typedef complex<double> C;
3  typedef vector<double> vd;
4  typedef vector<ll> vll;
5  const double PI = acos(-1);
6
7  void fft(vector<C>& a) {
8      int n = a.size(), L = 31 - __builtin_clz(n);
9      static vector<C> R(2, 1);
10     static vector<C> rt(2, 1);
11     for (static int k = 2; k < n; k *= 2) {
12         R.resize(n); rt.resize(n);
13         auto x = polar(1.0, PI / k);
14         for (int i = k; i < 2 * k; i++)
15             rt[i] = R[i] = i & 1 ? R[i / 2] * x : R[i / 2];
16     }
17     vector<int> rev(n);
18     for (int i = 0; i < n; i++) rev[i] = (rev[i / 2] | (i & 1) << L) /
19         2;
20     for (int i = 0; i < n; i++) if (i < rev[i]) swap(a[i], a[rev[i]]);
21     for (int k = 1; k < n; k *= 2)
22         for (int i = 0; i < n; i += 2 * k) for (int j = 0; j < k; j++) {
23             auto x = (double*)&rt[j + k], y = (double*)&a[i + j + k];
24             C z(x[0] * y[0] - x[1] * y[1], x[0] * y[1] + x[1] * y[0]);
25             a[i + j + k] = a[i + j] - z;
26             a[i + j] += z;
27         }
28
29     vll multiply(const vll& a, const vll& b) {
30         if (a.empty() || b.empty()) return {};
31         vd fa(a.begin(), a.end()), fb(b.begin(), b.end());
32         int L = 32 - __builtin_clz(fa.size() + fb.size() - 1), n = 1 << L;
33         vector<C> in(n), out(n);
34 }
```

```

34
35     for (int i = 0; i < a.size(); i++) in[i] = C(fa[i], 0);
36     for (int i = 0; i < b.size(); i++) in[i].imag(fb[i]);
37
38     fft(in);
39     for (C& x : in) x *= x;
40     for (int i = 0; i < n; i++) out[i] = in[-i & (n - 1)] - conj(in[i]);
41         // Corregido aqui
42     fft(out);
43
44     vll res(a.size() + b.size() - 1);
45     for (int i = 0; i < res.size(); i++) {
46         res[i] = llround(imag(out[i]) / (4 * n));
47     }
48     return res;
}

```

## 8.7 Rho

```

1 ll mul(ll a, ll b, ll mod) {
2     return (__int128)a * b % mod;
3 }
4
5 ll power(ll a, ll b, ll mod) {
6     ll res = 1;
7     while (b) {
8         if (b & 1) res = mul(res, a, mod);
9         a = mul(a, a, mod);
10        b >= 1;
11    }
12    return res;
13 }
14
15 bool isPrime(ll n) {
16     if (n < 2) return false;
17     for (ll p : {2, 3, 5, 7, 11, 13, 17, 19, 23}) {
18         if (n % p == 0) return n == p;
19     }
20     ll d = n - 1, s = 0;
21     while ((d & 1) == 0) d >>= 1, ++s;
22     for (ll a : {2, 325, 9375, 28178, 450775, 9780504, 1795265022}) {
23         if (a % n == 0) continue;
24         ll x = power(a, d, n);

```

```

25         if (x == 1 || x == n - 1) continue;
26         bool ok = false;
27         for (int i = 0; i < s; ++i) {
28             x = mul(x, x, n);
29             if (x == n - 1) { ok = true; break; }
30         }
31         if (!ok) return false;
32     }
33     return true;
34 }
35
36 ll rho(ll n) {
37     if (n % 2 == 0) return 2;
38     while (true) {
39         ll c = rand() % (n - 1) + 1;
40         ll x = 2, y = 2, d = 1;
41         while (d == 1) {
42             x = (mul(x, x, n) + c) % n;
43             y = (mul(y, y, n) + c) % n;
44             y = (mul(y, y, n) + c) % n;
45             d = std::gcd((x > y ? x - y : y - x), n);
46         }
47         if (d != n) return d;
48     }
49 }
50
51 void fact(ll n, std::map<ll, int>& f) {
52     if (n == 1) return;
53     if (isPrime(n)) { f[n]++; return; }
54     ll d = rho(n);
55     if (d == n) {
56         f[n]++;
57         return;
58     }
59     fact(d, f);
60     fact(n / d, f);
61 }

```

## 8.8 Get Divisors

```

1 vector<ll> getDivisors(const map<ll, int>& f) {
2     vector<ll> divisors = { 1 };
3     for (auto [p, e] : f) {

```

```

4     vector<ll> next;
5     ll pe = 1;
6     for (int i = 0; i <= e; i++) {
7         for (ll d : divisors)
8             next.push_back(d * pe);
9         pe *= p;
10    }
11    divisors.swap(next);
12 }
13 sort(divisors.begin(), divisors.end());
14 return divisors;
15 }
```

## 8.9 Simpson

```

1 ld simpsonRule(function<ld(ld)> f, ld a, ld b, int n) {
2     // Asegurarse de que n sea par
3     if (n % 2 != 0) {
4         n++;
5     }
6     ld h = (b - a) / n;
7     ld s = f(a) + f(b);
8
9     // Suma de terminos interiores con los factores apropiados
10    for (int i = 1; i < n; i++) {
11        ld x = a + i * h;
12        s += (i % 2 == 1 ? 4.0L : 2.0L) * f(x);
13    }
14    // Multiplica por h/3
15    return (h / 3.0L) * s;
16 }
17 // Ejemplo: integrar la funcion x^2 entre 0 y 3
18 auto f = [&](ld x){ return x * x; };
19 ld a = 0.0L, b = 3.0L;
20 int n = 1000; // numero de subintervalos
21 ld resultado = simpsonRule(f, a, b, n);
```

## 8.10 Simplex

```

1 pair<ld, vec<ld>> simplex(vec<vec<ld>> A, vec<ld> b, vec<ld> c) {
2     const ld EPS = (ld)1e-9;
3     int n = SZ(b), m = SZ(c);
4
5     vec<int> X(m), Y(n);
```

```

6     L(j, 0, m) X[j] = j;
7     L(i, 0, n) Y[i] = m + i;
8
9     ld z = 0;
10
11    auto pivot = [&](int x, int y) {
12        swap(X[y], Y[x]);
13
14        ld inv = (ld)1 / A[x][y];
15        b[x] *= inv;
16        L(j, 0, m) if (j != y) A[x][j] *= inv;
17        A[x][y] = inv;
18
19        L(i, 0, n) if (i != x && fabs1(A[i][y]) > EPS) {
20            ld coef = A[i][y];
21            b[i] -= coef * b[x];
22            L(j, 0, m) if (j != y) A[i][j] -= coef * A[x][j];
23            A[i][y] = -coef * A[x][y];
24        }
25
26        z += c[y] * b[x];
27        L(j, 0, m) if (j != y) c[j] -= c[y] * A[x][j];
28        c[y] = -c[y] * A[x][y];
29    };
30
31    while (true) {
32        int x = -1, y = -1;
33        ld mn = -EPS;
34        L(i, 0, n) if (b[i] < mn) { mn = b[i]; x = i; }
35        if (x < 0) break;
36        L(j, 0, m) if (A[x][j] < -EPS) { y = j; break; }
37        if (y < 0) {
38            return { numeric_limits<ld>::quiet_NaN(), {} };
39        }
40        pivot(x, y);
41    }
42
43    while (true) {
44        int y = -1, x = -1;
45        ld mx = EPS;
46        L(j, 0, m) if (c[j] > mx) { mx = c[j]; y = j; }
47        if (y < 0) break;
```

```

49     ld best = numeric_limits<ld>::infinity();
50     L(i, 0, n) if (A[i][y] > EPS) {
51         ld val = b[i] / A[i][y];
52         if (val < best) { best = val; x = i; }
53     }
54     if (x < 0) {
55         return { numeric_limits<ld>::infinity(), {} };
56     }
57     pivot(x, y);
58 }

59 vec<ld> sol(m, 0);
60 L(i, 0, n) if (Y[i] < m) sol[Y[i]] = b[i];
61 return { z, sol };
62 }
```

## 9 Geometry

### 9.1 Convex Hull

```

1  typedef pair<ll, ll> Point;
2  ll cross_product(Point O, Point A, Point B) {
3      return (A.first - O.first) * (B.second - O.second) - (A.second - O.
4          second) * (B.first - O.first);
5 }
6 vector<Point> convex_hull(vector<Point>& points) {
7     sort(points.begin(), points.end());
8     points.erase(unique(points.begin(), points.end()), points.end());
9     vector<Point> hull;
10    // Parte inferior
11    for (const auto& p : points) {
12        while (hull.size() >= 2 && cross_product(hull[hull.size() - 2],
13            hull[hull.size() - 1], p) < 0)
14            hull.pop_back();
15        if (hull.empty() || hull.back() != p) {
16            hull.push_back(p);
17        }
18    }
19    // Parte superior
20    int t = hull.size() + 1;
21    for (int i = points.size() - 1; i >= 0; --i) {
22        while (hull.size() >= t && cross_product(hull[hull.size() - 2],
23            hull[hull.size() - 1], points[i]) < 0)
```

```

21         hull.pop_back();
22         if (hull.empty() || hull.back() != points[i]) {
23             hull.push_back(points[i]);
24         }
25     }
26     hull.pop_back();
27     return hull;
28 }
```

### 9.2 Operations

```

1  ll cross_product(pair<ll, ll> P1, pair<ll, ll> P2, pair<ll, ll> P3) {
2      ll x1 = P2.first - P1.first;
3      ll y1 = P2.second - P1.second;
4      ll x2 = P3.first - P1.first;
5      ll y2 = P3.second - P1.second;
6      return x1 * y2 - y1 * x2;
7 }
8 double distancia(pair<ll, ll> P1, pair<ll, ll> P2) {
9     return sqrt((P2.first - P1.first) * (P2.first - P1.first) +
10                 (P2.second - P1.second) * (P2.second - P1.second));
11 }
12 ll dot_product(pair<ll, ll> P1, pair<ll, ll> P2, pair<ll, ll> P3) {
13     ll x1 = P2.first - P1.first;
14     ll y1 = P2.second - P1.second;
15     ll x2 = P3.first - P1.first;
16     ll y2 = P3.second - P1.second;
17     return x1 * x2 + y1 * y2;
18 }
```

### 9.3 Polygon Area

```

1  typedef pair<ll, ll> Point;
2  double polygon_area(const vector<Point>& polygon) {
3      ll area = 0;
4      int n = polygon.size();
5      for (int i = 0; i < n; ++i) {
6          ll j = (i + 1) % n;
7          area += (polygon[i].first * polygon[j].second - polygon[i].
8              second * polygon[j].first);
9      }
10     return abs(area) / 2.0;
11 }
```

## 9.4 Ray Casting

```

1 | typedef pair<ll, ll> Point;
2 | bool is_point_in_polygon(const vector<Point>& polygon, Point p) {
3 |     bool inside = false;
4 |     int n = polygon.size();
5 |     for (int i = 0, j = n - 1; i < n; j = i++) {
6 |         if ((polygon[i].second > p.second) != (polygon[j].second > p.
7 |             second) &&
8 |                 p.first < (polygon[j].first - polygon[i].first) * (p.second
9 |                     - polygon[i].second) /
10 |                         (polygon[j].second - polygon[i].second) + polygon[
11 |                             i].first) {
12 |             inside = !inside;
13 |         }
14 |     }
15 |     return inside;
16 | }
```

## 10 Other

### 10.1 Mo's algorithm

```

1 | const int BLOCK_SIZE = 450; using U64 = uint64_t;
2 | struct query {int l, r, id;U64 order;};
3 | U64 hilbertorder(U64 x, U64 y) {
4 |     const U64 logn = __lg(max(x, y) * 2 + 1) + 1;
5 |     const U64 maxn = (1ull << logn) - 1;
6 |     U64 res = 0;
7 |     for (U64 s = 1ull << (logn - 1); s, s >>= 1) {
8 |         bool rx = x & s, ry = y & s;
9 |         res = (res << 2) | (rx ? ry ? 2 : 1 : ry ? 3 : 0);
10 |        if (!rx) {
11 |            if (ry) x ^= maxn, y ^= maxn;
12 |            swap(x, y);
13 |        }
14 |    }
15 |    return res;
16 | } // sort by this order
17 | auto add = [&](int ix) { /* Add A[ix] to state*/};
18 | auto rem = [&](int ix) { /* Remove A[ix] from state*/};
19 | int c_l = 0, c_r = -1; // Cursors [0,-1] so r add 0 on first q
20 | for(const auto &qr: qs){
```

```

21 |     while(c_l > qr.l) add(--c_l);
22 |     while(c_r < qr.r) add(++c_r);
23 |     while(c_l < qr.l) rem(c_l++);
24 |     while (c_r > qr.r) rem(c_r--);
25 |     ans[qr.id] = /*State.Answer()*/;
26 | }
```

## 11 Ecuations

### 11.1 Combinatorics

$$\binom{n}{k} = \binom{n}{n-k}$$

$$\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}, \quad (1 \leq k \leq n)$$

$$\sum_{k=0}^n \binom{n}{k} = 2^n$$

$$\sum_{m=0}^n \binom{m}{k} = \binom{n+1}{k+1}, \quad (n \geq k \geq 0)$$

$$\sum_{k=0}^m \binom{n+k}{k} = \binom{n+m+1}{m}$$

$$\binom{n}{0}^2 + \binom{n}{1}^2 + \cdots + \binom{n}{n}^2 = \binom{2n}{n}$$

$$\sum_{k=1}^n k \binom{n}{k} = n2^{n-1}$$

$$F_n = \sum_{k=0}^{\lfloor \frac{n-1}{2} \rfloor} \binom{n-k-1}{k}$$

$$F_{n+1} = \sum_{k=0}^{\lfloor \frac{n}{2} \rfloor} \binom{n-k}{k}$$

## 11.2 Discreta

Vandermonde convolution:

$$\sum_{k=0}^n \binom{r}{k} \binom{s}{n-k} = \binom{r+s}{n}$$

Multinomial theorem:

$$(x_1 + \cdots + x_m)^n = \sum_{\substack{a_1+\cdots+a_m=n \\ a_i \geq 0}} \frac{n!}{a_1! \cdots a_m!} x_1^{a_1} \cdots x_m^{a_m}$$

Binomial inversion (sequence form):

$$g(n) = \sum_{k=0}^n \binom{n}{k} f(k) \iff f(n) = \sum_{k=0}^n (-1)^{n-k} \binom{n}{k} g(k)$$

Stars and bars (nonnegative):

$$x_1 + \cdots + x_k = n, x_i \geq 0 \Rightarrow \# = \binom{n+k-1}{k-1}$$

Positive parts:

$$x_1 + \cdots + x_k = n, x_i \geq 1 \Rightarrow \# = \binom{n-1}{k-1}$$

Compositions of n:

$$\#\{\text{ordered positive sum of } n \text{ into } k \text{ parts}\} = \binom{n-1}{k-1}, \quad \#\{\text{all compositions}\} = 2^{n-1}$$

Upper bounds via inclusion-exclusion:

$$x_1 + \cdots + x_k = n, 0 \leq x_i \leq u_i \Rightarrow \# = \sum_{S \subseteq \{1, \dots, k\}} (-1)^{|S|} \binom{n - \sum_{i \in S} (u_i + 1) + k - 1}{k-1}$$

(toma  $\binom{t}{k-1} = 0$  si  $t < k-1$ )

Multiset combinations:

$$\#\{k\text{-multicombinations from } n \text{ types}\} = \binom{n+k-1}{k}$$

Multiset permutations:

$$\#\{\text{perm of multiset with counts } m_1, \dots, m_r\} = \frac{(m_1 + \cdots + m_r)!}{m_1! \cdots m_r!}$$

Circular permutations:

$$\#\{\text{distinct cyclic orders of } n \text{ items}\} = (n-1)!$$

Surjections count (onto functions):

$$\#\{f : [m] \rightarrow [n] \text{ onto}\} = \sum_{j=0}^n (-1)^j \binom{n}{j} (n-j)^m = n! S(m, n)$$

Derangements:

$$!n = n! \sum_{i=0}^n \frac{(-1)^i}{i!} \quad \text{and} \quad !n \approx \frac{n!}{e}$$

Stirling numbers (second kind):

$$S(n, k) = k S(n-1, k) + S(n-1, k-1), \quad S(0, 0) = 1$$

**Stirling numbers (first kind, unsigned):**

$$c(n, k) = c(n-1, k-1) + (n-1)c(n-1, k), \quad c(0, 0) = 1$$

**Expansions with falling powers:**

$$x^n = \sum_{k=0}^n s(n, k) x^k, \quad x^n = \sum_{k=0}^n S(n, k) x^k$$

(here  $x^k = x(x-1)\cdots(x-k+1)$ ,  $s(n, k) = (-1)^{n-k} c(n, k)$ )

**Bell numbers:**

$$B_n = \sum_{k=0}^n S(n, k), \quad \sum_{n \geq 0} B_n \frac{x^n}{n!} = \exp(e^x - 1)$$

**Cayley trees:**

$$\#\{\text{labeled trees on } n \text{ vertices}\} = n^{n-2}$$

**Perfect matchings in complete graph:**

$$\#\{\text{perfect matchings in } K_{2n}\} = (2n-1)!! = \frac{(2n)!}{2^n n!}$$

**Grid shortest paths:**

$$\#\{\text{monotone paths from } (0, 0) \text{ to } (a, b)\} = \binom{a+b}{a}$$

**Ballot (Bertrand special case):**

$$p > q \Rightarrow \#\{\text{prefix-wise leading sequences}\} = \frac{p-q}{p+q} \binom{p+q}{q}$$

**Alternating binomial sums:**

$$\sum_{k=0}^n (-1)^k \binom{n}{k} = 0 \quad (n \geq 1), \quad \sum_{k=0}^n (-1)^k \binom{n}{k} k^m = 0 \quad (0 \leq m < n)$$

**Lucas theorem (mod prime p):**

$$n = \sum n_i p^i, \quad k = \sum k_i p^i \Rightarrow \binom{n}{k} \equiv \prod_i \binom{n_i}{k_i} \pmod{p}$$

## 11.3 Trigonometry

$$\begin{aligned} \sin(-x) &= -\sin x, & \cos(-x) &= \cos x, & \tan(-x) &= -\tan x \\ \sin^2 x + \cos^2 x &= 1, & 1 + \tan^2 x &= \sec^2 x, & 1 + \cot^2 x &= \csc^2 x \end{aligned}$$

$$\begin{aligned} \sin(\alpha \pm \beta) &= \sin \alpha \cos \beta \pm \cos \alpha \sin \beta \\ \cos(\alpha \pm \beta) &= \cos \alpha \cos \beta \mp \sin \alpha \sin \beta \\ \tan(\alpha \pm \beta) &= \frac{\tan \alpha \pm \tan \beta}{1 \mp \tan \alpha \tan \beta} \end{aligned}$$

$$\begin{aligned} \sin(2x) &= 2 \sin x \cos x \\ \cos(2x) &= \cos^2 x - \sin^2 x = 1 - 2 \sin^2 x = 2 \cos^2 x - 1 \end{aligned}$$

$$\begin{aligned} \tan(2x) &= \frac{2 \tan x}{1 - \tan^2 x} \\ \sin(3x) &= 3 \sin x - 4 \sin^3 x, & \cos(3x) &= 4 \cos^3 x - 3 \cos x \\ \sin^2 \frac{x}{2} &= \frac{1 - \cos x}{2}, & \cos^2 \frac{x}{2} &= \frac{1 + \cos x}{2}, & \tan \frac{x}{2} &= \frac{\sin x}{1 + \cos x} = \frac{1 - \cos x}{\sin x} \end{aligned}$$

$$\begin{aligned} \sin \alpha \sin \beta &= \frac{1}{2} [\cos(\alpha - \beta) - \cos(\alpha + \beta)] \\ \cos \alpha \cos \beta &= \frac{1}{2} [\cos(\alpha - \beta) + \cos(\alpha + \beta)] \\ \sin \alpha \cos \beta &= \frac{1}{2} [\sin(\alpha + \beta) + \sin(\alpha - \beta)] \end{aligned}$$

$$\begin{aligned} \sin \alpha \pm \sin \beta &= 2 \sin \frac{\alpha \pm \beta}{2} \cos \frac{\alpha \mp \beta}{2} \\ \cos \alpha + \cos \beta &= 2 \cos \frac{\alpha + \beta}{2} \cos \frac{\alpha - \beta}{2} \\ \cos \alpha - \cos \beta &= -2 \sin \frac{\alpha + \beta}{2} \sin \frac{\alpha - \beta}{2} \end{aligned}$$

$$\sum_{k=0}^{n-1} \cos(a + kd) = \frac{\sin(\frac{nd}{2})}{\sin(\frac{d}{2})} \cos\left(a + \frac{(n-1)d}{2}\right)$$

$$\sum_{k=0}^{n-1} \sin(a + kd) = \frac{\sin(\frac{nd}{2})}{\sin(\frac{d}{2})} \sin\left(a + \frac{(n-1)d}{2}\right)$$

$$\text{Law of sines: } \frac{a}{\sin A} = \frac{b}{\sin B} = \frac{c}{\sin C} = 2R$$

$$\text{Law of cosines: } c^2 = a^2 + b^2 - 2ab \cos C \quad (\text{and cyclic})$$

$$\text{Area: } \Delta = \frac{1}{2} ab \sin C = \frac{1}{2} bc \sin A = \frac{1}{2} ca \sin B$$

$$(x', y') = (x \cos \theta - y \sin \theta, x \sin \theta + y \cos \theta)$$

$$e^{ix} = \cos x + i \sin x, \quad \cos x = \frac{e^{ix} + e^{-ix}}{2}, \quad \sin x = \frac{e^{ix} - e^{-ix}}{2i}$$

$$\pi \text{ rad} = 180^\circ, \quad 1^\circ = \frac{\pi}{180} \text{ rad}$$

## 11.4 Catalan Numbers

Recursive definition:

$$C_0 = C_1 = 1$$

$$C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}, \quad n \geq 2$$

Closed form:

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

Combinatorial equivalent:

$$C_n = \binom{2n}{n} - \binom{2n}{n-1} = \frac{1}{n+1} \binom{2n}{n}, \quad n \geq 0$$

Combinatorial meaning:

Number of ways to: (i) arrange  $n$  balanced parenthesis pairs; (ii) full binary trees with  $n+1$  leaves;  
 (iii) Dyck paths of length  $2n$  that never cross the diagonal.

Generalized form ( $k$ ):

$$C_n^{(k)} = \frac{k+1}{n+k+1} \binom{2n+k}{n}$$

Extended recurrence:

$$C_n^{(k)} = \sum_{a_1+\dots+a_k=n} C_{a_1} C_{a_2} \cdots C_{a_k}, \quad C_0 = 1$$

Efficient recurrence (for computation):

$$C_n = \frac{2(2n-1)}{n+1} C_{n-1}, \quad n \geq 1$$

Generating function:

$$C(x) = \sum_{n=0}^{\infty} C_n x^n = \frac{1 - \sqrt{1 - 4x}}{2x}$$

Asymptotic behavior:

$$C_n \sim \frac{4^n}{n^{3/2} \sqrt{\pi}}$$

Examples:

$$C_0 = 1, C_1 = 1, C_2 = 2, C_3 = 5, C_4 = 14, C_5 = 42$$

## 11.5 Geometry

Rectangle:

$$A = b h$$

Area with base  $b$  and height  $h$ .

Triangle:

$$A = \frac{1}{2} b h$$

$$A = \sqrt{s(s-a)(s-b)(s-c)}, \quad s = \frac{a+b+c}{2} \quad (\text{Heron})$$

Base-height or Heron using side lengths  $a, b, c$ .

Parallelogram & rhombus:

$$A_{\text{parallelogram}} = b h, \quad A_{\text{rhombus}} = \frac{D d}{2}$$

$D, d$  are diagonals of a rhombus.

Trapezoid:

$$A = \frac{(B+b)}{2} h$$

$B$  and  $b$  are the parallel sides (bases).

### Regular $n$ -gon:

$$A = \frac{1}{2} P a = \frac{n l a}{2}$$

$P$  perimeter,  $l$  side,  $a$  apothem.

### Circle:

$$A = \pi r^2, \quad C = 2\pi r$$

### Circular sector (angle in radians):

$$A = \frac{1}{2} r^2 \theta, \quad \text{arc length } L = r\theta$$

### Circular segment (height $h$ ):

$$A = r^2 \arccos\left(\frac{r-h}{r}\right) - (r-h)\sqrt{2rh - h^2}$$

Region cut by a chord;  $0 < h < 2r$ .

### Annulus (circular crown):

$$A = \pi(R^2 - r^2)$$

Difference of two concentric disks ( $R > r$ ).

### Ellipse:

$$A = \pi ab$$

$a, b$  are semi-axes.

### Polygon by coordinates (shoelace):

$$A = \frac{1}{2} \left| \sum_{i=1}^n (x_i y_{i+1} - x_{i+1} y_i) \right|, \quad (x_{n+1}, y_{n+1}) = (x_1, y_1)$$

Works for any simple polygon in the plane.

### Triangle by coordinates:

$$A = \frac{1}{2} |x_1(y_2 - y_3) + x_2(y_3 - y_1) + x_3(y_1 - y_2)|$$

### Triangle from sides and circumradius/inradius:

$$A = \frac{abc}{4R}, \quad A = rs$$

$R$  circumradius,  $r$  inradius,  $s$  semiperimeter.

### Lune (difference of two circular sectors):

$$A_{\text{lune}} = \frac{1}{2}r_1^2\theta_1 - \frac{1}{2}r_2^2\theta_2$$

Two sectors overlapping with angles  $\theta_1, \theta_2$  matching the same chord.

### Lens (two equal circles radius $r$ , center distance $d$ ):

$$A = 2r^2 \arccos\left(\frac{d}{2r}\right) - \frac{d}{2} \sqrt{4r^2 - d^2}, \quad 0 < d < 2r$$

Intersection of two equal disks.

### Spherical cap (radius $R$ , height $h$ ):

$$A_{\text{cap}} = 2\pi Rh$$

Surface area of the cap on a sphere. (Volume:  $V = \frac{\pi h^2}{3}(3R - h)$ )

## 11.6 Useful math

### Arithmetic series:

$$\sum_{k=1}^n k = \frac{n(n+1)}{2}$$

### Squares & cubes:

$$\sum_{k=1}^n k^2 = \frac{n(n+1)(2n+1)}{6}, \quad \sum_{k=1}^n k^3 = \left(\frac{n(n+1)}{2}\right)^2$$

### Geometric series ( $r \neq 1$ ):

$$\sum_{i=0}^{n-1} r^i = \frac{r^n - 1}{r - 1}$$

For mod prime  $p$ : multiply by  $(r-1)^{-1} \equiv (r-1)^{p-2} \pmod{p}$ .

**Power sum of base  $a$ :**

$$1 + a + \dots + a^n = \frac{a^{n+1} - 1}{a - 1} \quad (a \neq 1), \quad = n + 1 \quad (a = 1)$$

**Harmonic numbers:**

$$H_n = \sum_{k=1}^n \frac{1}{k} \approx \ln n + \gamma + \frac{1}{2n}$$

$\gamma \approx 0.57721$  (Euler–Mascheroni). Useful for estimates.

**Basic mod rules:**

$$(a \pm b) \pmod{m} = ((a \pmod{m}) \pm (b \pmod{m})) \pmod{m}$$

$$(a \cdot b) \pmod{m} = ((a \pmod{m}) \cdot (b \pmod{m})) \pmod{m}$$

**Fermat little theorem (prime  $p$ ):**

$$a^{p-1} \equiv 1 \pmod{p} \quad \text{if } p \nmid a, \quad a^{-1} \equiv a^{p-2} \pmod{p}$$

**Euler theorem:**

$$a^{\varphi(m)} \equiv 1 \pmod{m} \quad \text{if } \gcd(a, m) = 1$$

**Chinese remainder (pairwise coprime):**

$$x \equiv a_i \pmod{m_i} \Rightarrow x \equiv \sum_i a_i M_i y_i \pmod{M}$$

$$M = \prod m_i, \quad M_i = M/m_i, \quad y_i \equiv M_i^{-1} \pmod{m_i}.$$

**gcd/lcm relation:**

$$\operatorname{lcm}(a, b) = \frac{|ab|}{\operatorname{gcd}(a, b)}$$

**Binomial theorem:**

$$(x + y)^n = \sum_{k=0}^n \binom{n}{k} x^{n-k} y^k$$

**Stars and bars (non-neg.):**

$$x_1 + \dots + x_k = n, \quad x_i \geq 0 \Rightarrow \# = \binom{n+k-1}{k-1}$$

**Permutations & combinations:**

$$P(n, k) = \frac{n!}{(n-k)!}, \quad \binom{n}{k} = \frac{n!}{k!(n-k)!}$$

**Derangements (approx):**

$$!n = \left\lfloor \frac{n!}{e} + \frac{1}{2} \right\rfloor, \quad !n = n! \sum_{i=0}^n \frac{(-1)^i}{i!}$$

**Stirling approximation:**

$$n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$$

**Inclusion-Exclusion (finite):**

$$\left| \bigcup_{i=1}^m A_i \right| = \sum_i |A_i| - \sum_{i < j} |A_i \cap A_j| + \dots + (-1)^{m+1} |A_1 \cap \dots \cap A_m|$$

**Dot and cross (2D):**

$$u \cdot v = u_x v_x + u_y v_y = |u||v| \cos \theta, \quad u \times v = u_x v_y - u_y v_x$$

$$|u \times v| = 2 \times \text{triangle area}(u, v). \quad \text{Orientation by sign of } u \times v.$$

**Distance point to line  $AB$ :**

$$\operatorname{dist}(P, AB) = \frac{|(B-A) \times (P-A)|}{|B-A|}$$

**Projection length on  $AB$ :**

$$\operatorname{proj}_{AB}(P) = \frac{(P-A) \cdot (B-A)}{|B-A|}$$

**AM-GM (non-neg.):**

$$\frac{x_1 + \cdots + x_n}{n} \geq (x_1 \cdots x_n)^{1/n}$$

**Cauchy-Schwarz:**

$$\left( \sum_i a_i b_i \right)^2 \leq \left( \sum_i a_i^2 \right) \left( \sum_i b_i^2 \right)$$

**Log rules:**

$$\log_a b = \frac{\ln b}{\ln a}, \quad \log(ab) = \log a + \log b$$

**Fast exponent splits:**

$$a^{x+y} = a^x a^y, \quad a^{2^k} = (\underbrace{\cdots (a^2)^2 \cdots}_k \text{ times})^2$$

**Divisor count/sum (multiplicative):**

$$n = \prod p_i^{e_i} \Rightarrow \tau(n) = \prod (e_i + 1), \quad \sigma(n) = \prod \frac{p_i^{e_i+1} - 1}{p_i - 1}$$

 $\tau(n)$  = number of divisors,  $\sigma(n)$  = sum of divisors.**Linearity of expectation:**

$$\mathbb{E}\left[\sum_i X_i\right] = \sum_i \mathbb{E}[X_i] \quad (\text{no independence needed})$$

**Binomial distribution:**

$$X \sim \text{Bin}(n, p) \Rightarrow \mathbb{E}[X] = np, \quad \text{Var}(X) = np(1-p)$$

## 11.7 Mobius

**Mobius function mu:**

$$\mu(1) = 1$$

$$\mu(n) = 0 \text{ if } \exists p^2 \mid n, \quad \mu(n) = (-1)^k \text{ if } n \text{ is square-free with } k \text{ distinct primes.}$$

**Basic convolutions:**

$$\sum_{d|n} \mu(d) = \begin{cases} 1, & n = 1, \\ 0, & n > 1, \end{cases} \quad (\mu * \mathbf{1})(n) = \varepsilon(n).$$

**Mobius inversion (divisor-sum):**  
If

$$g(n) = \sum_{d|n} f(d),$$

then

$$f(n) = \sum_{d|n} \mu(d) g\left(\frac{n}{d}\right) = \sum_{d|n} \mu\left(\frac{n}{d}\right) g(d).$$

**Inversion over multiples:**  
If

$$G(n) = \sum_{k \geq 1} f(kn),$$

then

$$f(n) = \sum_{k \geq 1} \mu(k) G(kn).$$

**Euler totient via mu:**

$$\varphi(n) = \sum_{d|n} \mu(d) \frac{n}{d}.$$

**Counting coprimes up to x:**

$$\#\{1 \leq m \leq x : \gcd(m, n) = 1\} = \sum_{d|n} \mu(d) \left\lfloor \frac{x}{d} \right\rfloor.$$

**Square-free count up to x:**

$$Q(x) = \#\{n \leq x : n \text{ square-free}\} = \sum_{k \leq \sqrt{x}} \mu(k) \left\lfloor \frac{x}{k^2} \right\rfloor.$$

**GCD=1 k-tuples:**

$$\#\{1 \leq x_1, \dots, x_k \leq n : \gcd(x_1, \dots, x_k) = 1\} = \sum_{d=1}^n \mu(d) \left\lfloor \frac{n}{d} \right\rfloor^k.$$

**Dirichlet series:**

$$\sum_{n=1}^{\infty} \frac{\mu(n)}{n^s} = \frac{1}{\zeta(s)}, \quad \Re(s) > 1.$$

**Mertens function:**

$$M(x) = \sum_{n \leq x} \mu(n).$$

## 11.8 Burnside

Necklaces under cyclic  $C_n$  (gcd form):

$$N_{\text{rot}} = \frac{1}{n} \sum_{k=0}^{n-1} m^{\gcd(n,k)}.$$

**Equivalent divisor form (for  $C_n$ ):**

$$N_{\text{rot}} = \frac{1}{n} \sum_{d|n} \varphi(d) m^{n/d}.$$

Necklaces/bracelets under dihedral  $D_n$  (gcd form):

$$N = \begin{cases} \frac{1}{2n} \left( \sum_{k=0}^{n-1} m^{\gcd(n,k)} + n m^{(n+1)/2} \right), & n \text{ odd}, \\ \frac{1}{2n} \left( \sum_{k=0}^{n-1} m^{\gcd(n,k)} + \frac{n}{2} m^{n/2} + \frac{n}{2} m^{n/2+1} \right), & n \text{ even}. \end{cases}$$

Here gcd is the greatest common divisor. A rotation by  $k$  positions has  $\gcd(n, k)$  cycles.