

### **Ex 1)**

- a. First and last bit decides row of  $S_1$ , while the rest decides column. The function is linear for input  $(x, y)$  if  $S(x) + S(y) = S(x+y)$ .

$$(x, y) = (000000, 000001): \quad row_x = 00 = 0, col_x = 0000 = 0$$

$$row_y = 01 = 1, col_y = 0000 = 0 \quad S(x) = 14 \quad S(y) = 0$$

$$S(x+y) = S(000001) = S(y) = 0 \quad \rightarrow \text{Not linear} \quad (x, y) = (111111, 100000)$$

$$row_x = 11 = 3, col_x = 1111 = 15 \quad row_y = 10 = 2, col_y = 0000 = 0 \quad S(x) = 13$$

$$S(y) = 4 \quad S(x+y) = S(011111) = 8 \quad \rightarrow \text{Not linear}$$

$$(x, y) = (101010, 010101) \quad row_x = 10 = 2, col_x = 0101 = 5$$

$$row_y = 01 = 1, col_y = 1010 = 10 \quad S(x) = 6 \quad S(y) = 12$$

$$S(x+y) = S(101010 + 010101) = S(111111) = 13 \quad \rightarrow \text{Not linear}$$

- b. See python code  
c. See python code

## **Ex 2)**

Verify  $IP(IP^{-1}(x))=x$  for  $(x_1, x_2, x_3, x_4, x_5)$  :

$$\begin{array}{cccccc} IP(x_1) \rightarrow x_{58} & IP^{-1}(x_{58}) \rightarrow x_1 & IP(x_2) \rightarrow x_{50} & IP^{-1}(x_{50}) \rightarrow x_2 & IP(x_3) \rightarrow x_{42} & \\ IP^{-1}(x_{42}) \rightarrow x_3 & IP(x_4) \rightarrow x_{34} & IP^{-1}(x_{34}) \rightarrow x_4 & IP(x_5) \rightarrow x_{26} & IP^{-1}(x_{26}) \rightarrow x_5 & \end{array}$$

**Ex 3)**

3.1)

$$P = (0, 0, 0, 0, \dots, 0) - 64 \text{ zeros} \quad K = (0, 0, 0, 0, \dots, 0) - 64 \text{ zeros} \quad \text{Initial permutation:}$$

$IP(P) = P$      $\square$  All zero input

First round:

$$L_i = (0, 0, 0, \dots, 0) - 32 \text{ zeros} \quad R_i = (0, 0, 0, \dots, 0) - 32 \text{ zeros} \quad K_i = (0, 0, 0, \dots, 0) - 48 \text{ zeros}$$
$$L_{i+1}=R_i \quad R_{i+1}=F(R_i, K_i) XOR L_i$$

Feistel:

- Expansion
  - o  $R_i = E(R_i) = (0, 0, 0, \dots, 0) - 48 \text{ zeros}$
- XOR with  $K_i$ :
  - o  $R_i = R_i \text{ XOR } K_i = (0, 0, 0, \dots, 0) - 48 \text{ zeros}$
- S\_boxes:
  - o Assuming every s\_box =  $S_1$  :
  - o Every s\_box outputs 14 = 1110
  - o  $R_i = (1, 1, 1, 0, 1, 1, 1, 0, \dots) - 1110 \times 8$
- Permutation:
  - o  $P(R_i) = (0, 1, 0, 1, 1, 0, 0, 1, 1, 1, 1, 1, 1, 1, 1, 0, 0, 1, 0, 1, 1, 1, 1, 1, 1, 1, 0, 1)$
- $R_i \text{ XOR } L_i$  :
  - o  $R_i \text{ XOR } L_i = R_i (L_i \text{ is all zeros})$

$$Output = \begin{pmatrix} L_{i+1}, R_{i+1} \end{pmatrix} \quad \begin{matrix} 0,0 \\ 0,1,0,1,1,0,0,1,1,1,1,1,1,1,0,0,1,0,1,1,1,1,1,1,0,1 \end{matrix}$$

### 3.2)

Since only one bit gets flipped, the first round is almost the same. Here I am just listing the differences.

$$P(x_{57})=1 \rightarrow IP(P) \rightarrow x_{63}=1$$

First round:

$$R_i \rightarrow x_{31}=1$$

Feistel:

$$\text{Expansion } E(R_i) \rightarrow x_{46}=1 \quad \text{XOR with } K_i - \text{No change}$$

S boxes: Last six bits of  $R_i = (0,0,0,1,0,0) \rightarrow S(R_{\text{last six bits}}) = 13 = 1101$   $R_i$  after the S boxes is the same as in the last, except the last four bits are 1101 instead of 1110.

$$\text{Permutation: } P(R_i) = (0,1,0,1,1,0,0,1,1,1,1,1,1,0,1,1,0,0,1,1,1,1,1,1,1,1,1,1,0,1)$$

$$R_i \text{ XOR } L_i = R_i \quad (L_i \text{ is still all zeros})$$

$$\text{Output} = (L_{i+1}, R_{i+1}) \quad 0,$$

$$0,1,0,1,1,0,0,1,1,1,1,1,1,0,1,1,0,0,1,1,1,1,1,1,1,1,1,1,0,1,1$$

As we can see, only a few bits flipped differently when we changed only one bit in the plaintext.

### 3.3)

Nothing will be different from 3.1 until we XOR with the key in the Feistel network.

Round 1 key generation:

$K = (1, 0, 0, 0, \dots, 0) - \text{length } 64$      $PC-1: K \rightarrow \text{Remove 8 parity bits, all of them 0, which gives}$

$K = (1, 0, 0, 0, \dots, 0) - \text{length } 56$      $K = (C_i, D_i)$      $C_i = (1, 0, 0, 0, \dots, 0) - \text{length } 28$

$D_i = \text{all zeros} - \text{length } 28$     *Since it's round 1,  $C_i \wedge D_i$  are 1 by 1, which gives:*

$C_i = (0, 0, 0, \dots, 0, 1), C(x_{28}) = 1$  (same length)     $D_i - \text{same as before}$

Permuting  $K_i$  following  $PC-2$ , we get:     $PC-2(K_i) \rightarrow K(x_{37}) = 1$

The position of the 1 gets permuted from  $x_{28}$  to  $x_{37}$  of the key.

Feistel:

$R_i \text{ XOR } K_i = K_i$ , since  $R_i$  is only zeros

S box: since  $R_i(x_{37}, x_{38}, x_{39}, x_{40}, x_{41}, x_{42}) = (1, 0, 0, 0, 0, 0)$  the 7th S\_box will be different from 3.1.

The seventh S\_box outputs 4 = 0100, which slightly affects the permutation.

$R_i = S(R_i) = (11101110111011101110111001001110)$

$R_i = P(R_i) = (0, 1, 0, 1, 1, 0, 0, 1, 1, 1, 1, 1, 1, 1, 1, 1, 0, 0, 1, 0, 1, 1, 1, 1, 1, 1, 1, 1, 0, 0)$

Output =  $(L_{i+1}, R_{i+1})$     0,

0, 1, 0, 1, 1, 0, 0, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 0, 0, 1, 0, 1, 1, 1, 1, 1, 1, 1, 0, 0

As we can see again, only one bit difference in the key does not change much from 3.1 after only one round.

#### **Ex 4)**

See python file