

INTRODUCTION TO AMD GPU PROGRAMMING WITH HIP

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Agenda

- Introduction (5 minutes)
- AMD GPU Hardware (10 minutes)
- GPU Programming Concepts (45 minutes)
- GPU Programming Software (15 minutes)
- Porting existing CUDA codes to HIP (15 mins)



What we won't cover today (but is still important)

- Profiling:
 - rocprofiler / roctracer: libraries for collecting GPU hardware counters and application traces
 - Install: sudo apt install rocprofiler-dev roctracer-dev
- Debugging:
 - rocr debug agent: print state of wavefronts on memory violation / signals
 - HIP debugging tips
 - In kernel printf



What we won't cover today (but is still important)

- AOMP (AMD OpenMP Compiler):
 - OpenMP 4.5+ support, "target" pragmas, device offloading
- GPU Libraries:
 - hipBLAS: BLAS functionality on GPUs
 - rocFFT: FFTs on GPUs
 - rocRAND: random number generation
 - rocPRIM / hipCUB: high performance GPU primitives
 - <u>Tensile</u>: GEMMs, tensor contractions
 - <u>hipSPARSE</u>: BLAS for sparse matrices / vectors
 - rocALUTION: iterative sparse solvers
 - MIOpen, TensorFlow, PyTorch: machine learning



Comments

- Slides will be provided
- Ask questions in the google doc; we will be monitoring it
- Focus is on single node / device:
 - Little discussion of MPI, or multi-node

Please share any feedback or ask questions in the Google Doc





AMD GCN GPU Hardware

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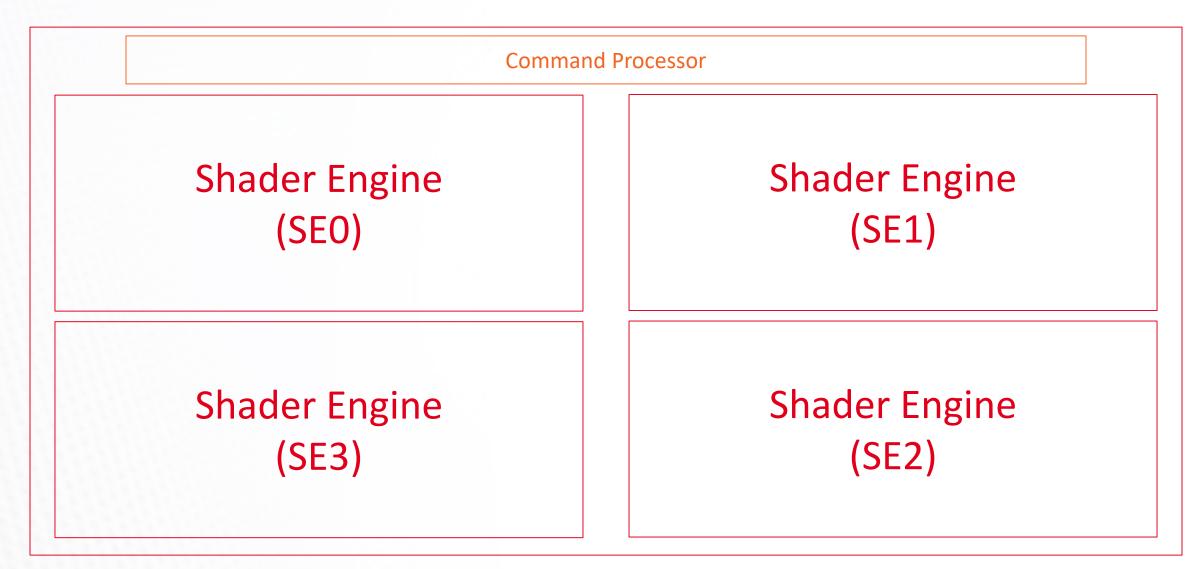
AMD GCN GPU HARDWARE

AGENDA

| Introduction | GCN Hardware Overview | |
|----------------------|-----------------------------|--|
| | AMD GPU Compute Terminology | |
| AMD GPU Architecture | GPU Memory and I/O System | |
| | GCN Compute Unit Internals | |



AMD GCN GPU Hardware Layout



AMD GCN GPU Hardware Layout

Command Queue

Queues reside in user-visible DRAM

Command Queue

Command Processor

| CU | CU |
|----|----|
| CU | CU |
| | |

Workload Manager

Workload

Manager

| CU | CU |
|----|----|
| CU | CU |

Workload Manager

| | CU | CU |
|---|----|----|
| | CU | CU |
| _ | | |

Workload Manager

| CU | CU |
|----|----|
| CU | CU |
| | |

Hardware Configuration Parameters on Modern AMD GPUs

| GPU SKU | Shader Engines | CUs / SE |
|--|----------------|----------|
| AMD Radeon Instinct™ MI60 | 4 | 16 |
| AMD Radeon Instinct™ MI50 | 4 | 15 |
| AMD Radeon™ VII | 4 | 15 |
| AMD Radeon Instinct™ MI25 AMD Radeon™ Vega 64 | 4 | 16 |
| AMD Radeon™ Vega 56 | 4 | 14 |
| AMD Radeon Instinct™ MI6 | 4 | 9 |
| AMD Ryzen™ 5 2400G | 1 | 11 |



AMD

AMD GPU Compute Terminology



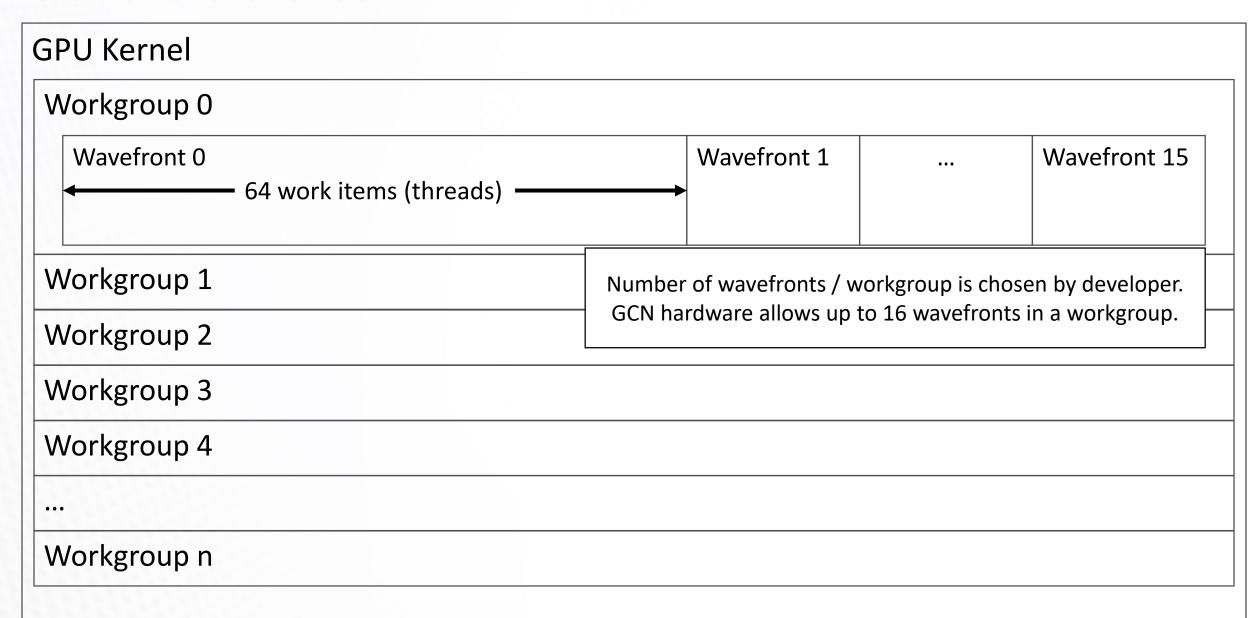
GPU Kernel

Functions launched to the GPU that are executed by multiple parallel workers

Examples: GEMM, triangular solve, vector copy, scan, convolution

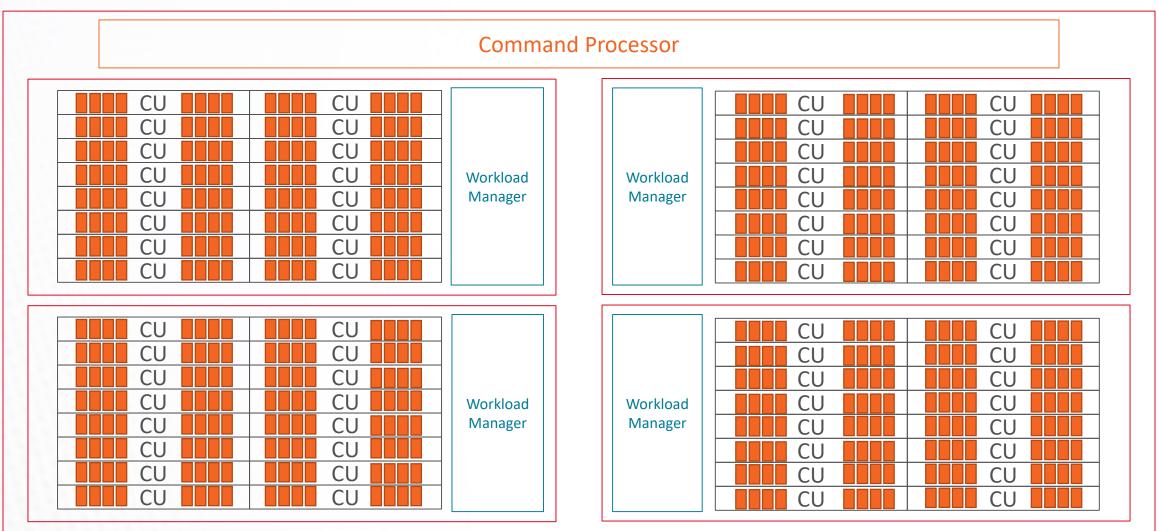
GPU Kernel Workgroup 0 Group of threads that are on the GPU at the same time. **CUDA Terminology** Also on the same compute unit. Thread Block Can synchronize together and communicate through memory in the CU. Workgroup 1 Workgroup 2 Workgroup 3 Workgroup 4 Workgroup n Programmer controls the number of workgroups – it's usually a function of problem size.

GPU Kernel Workgroup 0 Wavefront Collection of resources that execute in lockstep, run the same instructions, **CUDA Terminology** and follow the same control-flow path. Individual lanes can be masked off. Warp Can think of this as a vectorized thread. Workgroup 1 Workgroup 2 Workgroup 3 Workgroup 4 Workgroup n

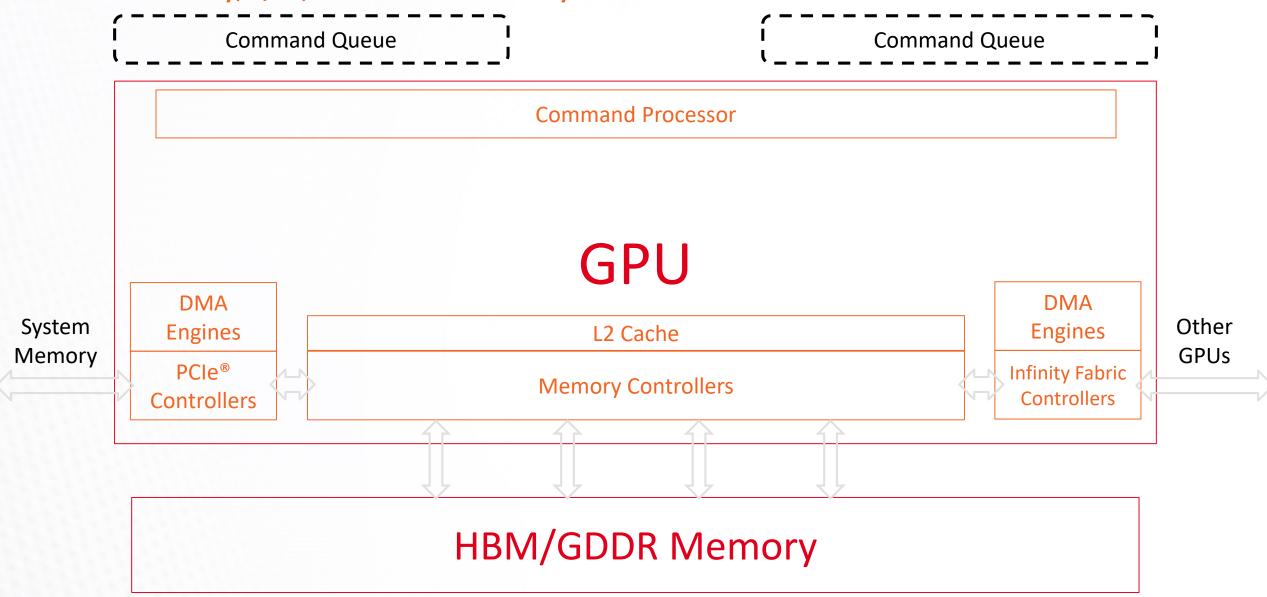


Scheduling work to a GPU

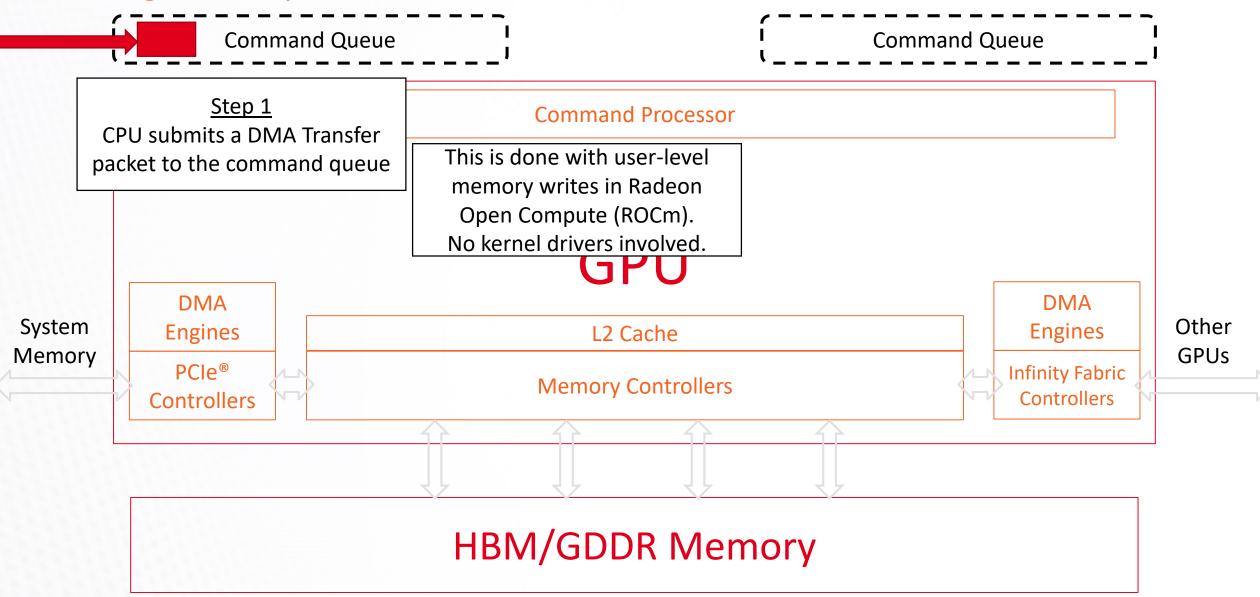




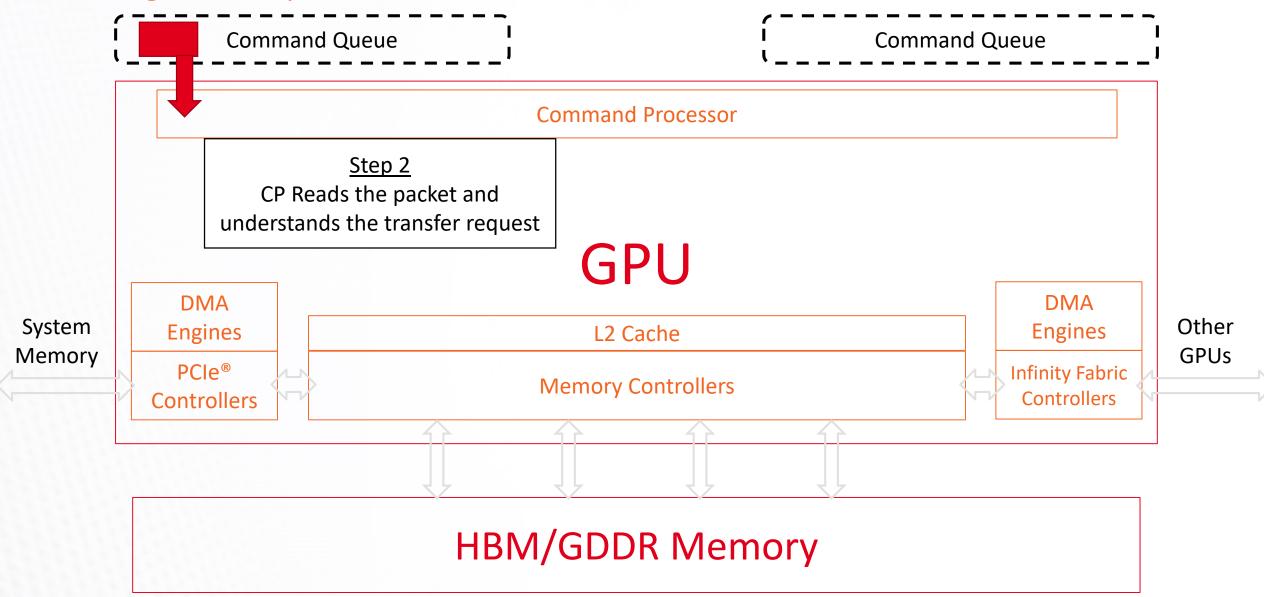
GPU Memory, I/O, and Connectivity



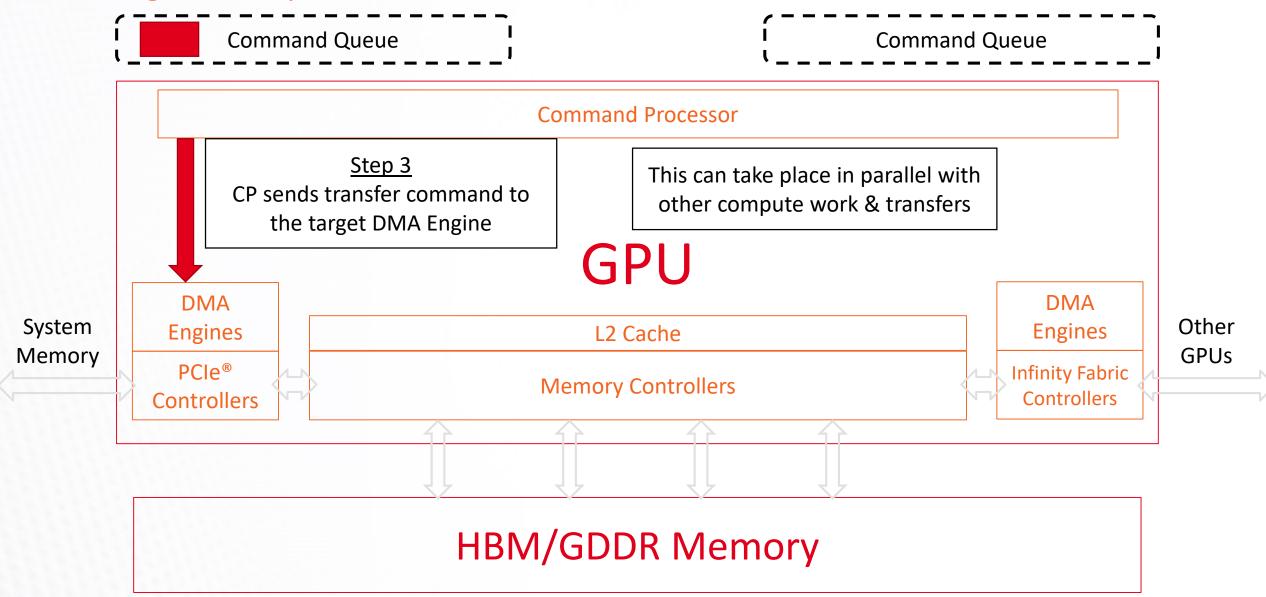




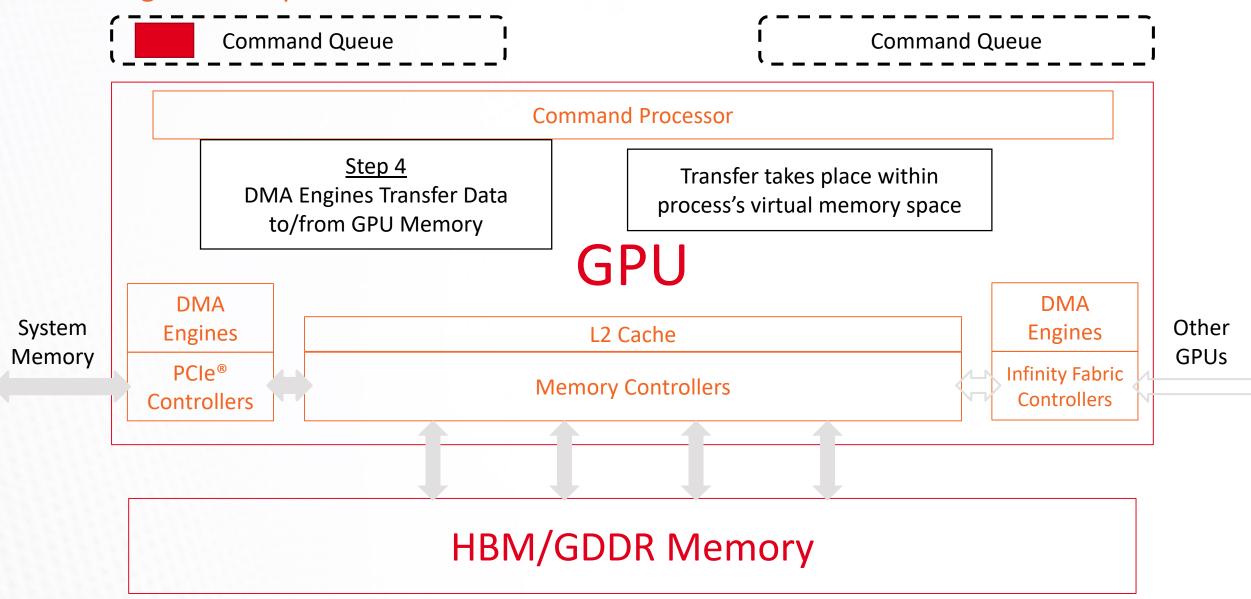














Compute Unit (CU)

- The workload manager sends work packages (i.e. blocks of threads) to the Compute Units (CUs)
 - Blocks are executed in wavefronts (groups of 64 threads in SIMD)
 - All wavefronts in a block are guaranteed to reside in the same CU
 - The CU's scheduler can hold wavefronts from many blocks
 - At most 40 wavefronts total per CU

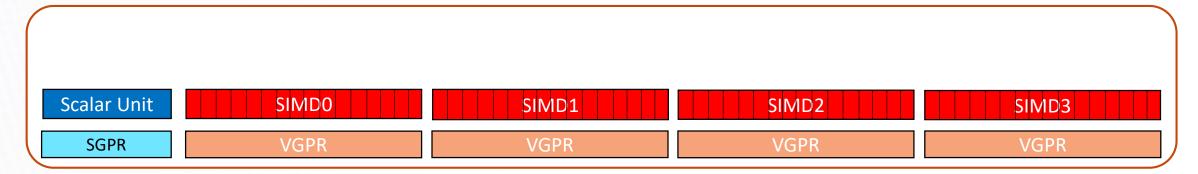


Scalar Unit

SGPR

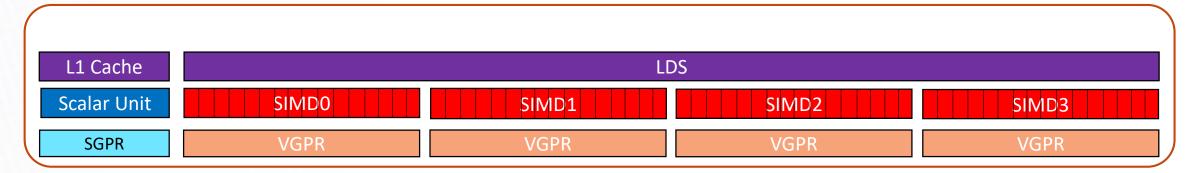
- The Scalar Unit (SU)
 - Shared by all threads in each wavefront, accessed on a per-wavefront level
 - Threads in a wavefront performing the exact same operation can offload this instruction to the SU
 - Used for control flow, pointer arithmetic, loading a common value, etc.
 - Has its own pool of Scalar General-Purpose Register (SGPR) file, 12.5KiB per CU





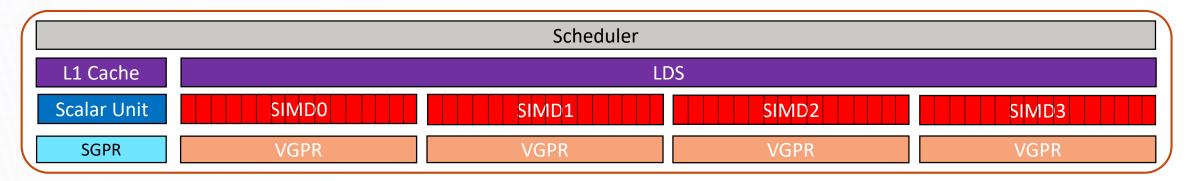
- SIMD Units
 - 4x SIMD vector units (each 16 lanes wide)
 - 4x 64KiB (256KiB total) Vector General-Purpose Register (VGPR) file
 - A maximum of 256 total registers per SIMD lane each register is 64x 4-byte entries
 - Instruction buffer for 10 wavefronts on each SIMD unit
 - Each wavefront is local to a single SIMD unit, not spread among the 4 (more on this in a moment)





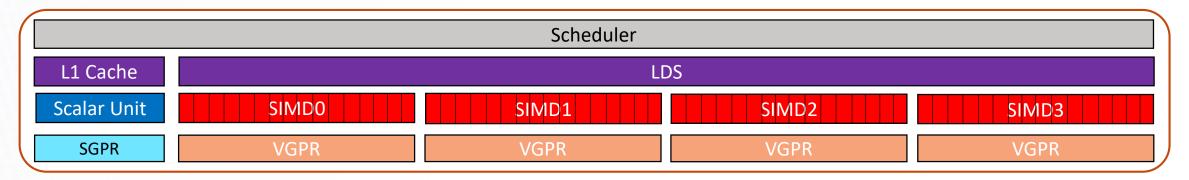
- 64KiB Local Data Share (LDS, or shared memory)
 - 32 banks with conflict resolution
 - Can share data between all threads in a block
- 16 KiB Read/Write L1 vector data cache
 - Write-through; L2 cache is the coherence point shared by all CUs





- Scheduler
 - Buffer for 40 wavefronts 2560 threads
 - Separate decode/issue for
 - VALU, VGPR load/store
 - SALU, SGPR load/store
 - LDS load/store
 - Global mem load/store
 - Special instructions (NoOps, barriers, branch instructions)





Scheduler

- At each clock, waves on 1 SIMD unit are considered for execution (Round Robin scheduling among SIMDs)
- At most 1 instruction per wavefront may be issued
- At most 1 instruction from each category may be issued (S/V ALU, S/V GPR, LDS, global, branch, etc)
- Maximum of 5 instructions issued to wavefronts on a single SIMD, per cycle per CU
- Some instructions take 4 or more cycles to retire (e.g. FP32VALU instruction on 1 wavefront using 16-wide SIMD)
 - Round robin scheduling of SIMDs hides execution latency
 - Programmer can still 'pretend' CU operates in 64-wide SIMD



Software Terminology

| Nvidia/CUDA Terminology | AMD Terminology | Description |
|-----------------------------|--------------------|--|
| Streaming Multiprocessor | Compute Unit (CU) | One of many parallel vector processors in a GPU that contain parallel ALUs. All waves in a workgroups are assigned to the same CU. |
| Kernel | Kernel | Functions launched to the GPU that are executed by multiple parallel workers on the GPU. Kernels can work in parallel with CPU. |
| Warp | Wavefront | Collection of operations that execute in lockstep, run the same instructions, and follow the same control-flow path. Individual lanes can be masked off. Think of this as a vector thread. A 64-wide wavefront is a 64-wide vector op. |
| Thread Block | Workgroup | Group of wavefronts that are on the GPU at the same time. Can synchronize together and communicate through local memory. |
| Thread | Work Item / Thread | Individual lane in a wavefront. On AMD GPUs, must run in lockstep with other work items in the wavefront. Lanes can be individually masked off. |
| | | GPU programming models can treat this as a separate thread of execution, though you do not necessarily get forward sub-wavefront progress. |



Software Terminology

| Nvidia/CUDA Terminology | AMD Terminology | Description |
|----------------------------|-----------------|---|
| Global Memory | Global Memory | DRAM memory accessible by the GPU that goes through some layers cache |
| Shared Memory | Local Memory | Scratchpad that allows communication between wavefronts in a workgroup. |
| Local Memory | Private Memory | Per-thread private memory, often mapped to registers. |





AMD GPU Programming Concepts

Programming with HIP: Kernels, blocks, threads, and more

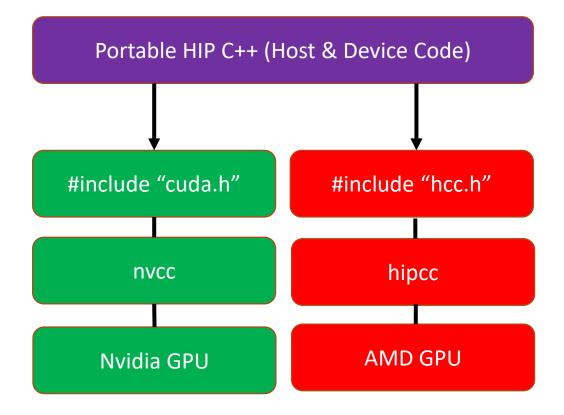
Damon McDougall <damon.mcdougall@amd.com>

What is HIP?

AMD's Heterogeneous-compute Interface for Portability, or **HIP**, is a C++ runtime API and kernel language that allows developers to create portable applications that can run on AMD's accelerators as well as CUDA devices.

HIP:

- Is open-source.
- Provides an API for an application to leverage GPU acceleration for both AMD and CUDA devices.
- Syntactically similar to CUDA. Most CUDA API calls can be converted in place: cuda -> hip
- Supports a strong subset of CUDA runtime functionality.



A Tale of Host and Device

Source code in HIP has two flavors: Host code and Device code

- The Host is the CPU
- Host code runs here
- Usual C++ syntax and features
- Entry point is the 'main' function
- HIP API can be used to create device buffers, move between host and device, and launch device code.

- The Device is the GPU
- Device code runs here
- C-like syntax
- Device codes are launched via "kernels"
- Instructions from the Host are enqueued into "streams"





HIP API

- Device Management:
 - hipSetDevice(), hipGetDevice(), hipGetDeviceProperties()
- Memory Management
 - hipMalloc(), hipMemcpy(), hipMemcpyAsync(), hipFree()
- Streams
 - hipStreamCreate(), hipSynchronize(), hipStreamSynchronize(), hipStreamFree()
- Events
 - hipEventCreate(), hipEventRecord(), hipStreamWaitEvent(), hipEventElapsedTime()
- Device Kernels
 - __global___, __device___, hipLaunchKernelGGL()
- Device code
 - threadIdx, blockIdx, blockDim, shared
 - 200+ math functions covering entire CUDA math library.
- Error handling
 - hipGetLastError(), hipGetErrorString()



AMD

Kernels, memory, and structure of host code



Device Kernels: The Grid

- In HIP, kernels are executed on a 3D "grid"
 - You might feel comfortable thinking in terms of a mesh of points, but it's not required
- The "grid" is what you can map your problem to
 - It's not a physical thing, but it can be useful to think that way
- AMD devices (GPUs) support 1D, 2D, and 3D grids, but most work maps well to 1D
- Each dimension of the grid partitioned into equal sized "blocks"
- Each block is made up of multiple "threads"
- The grid and its associated blocks are just organizational constructs
 - The threads are the things that do the work
- If you're familiar with CUDA already, the grid+block structure is very similar in HIP



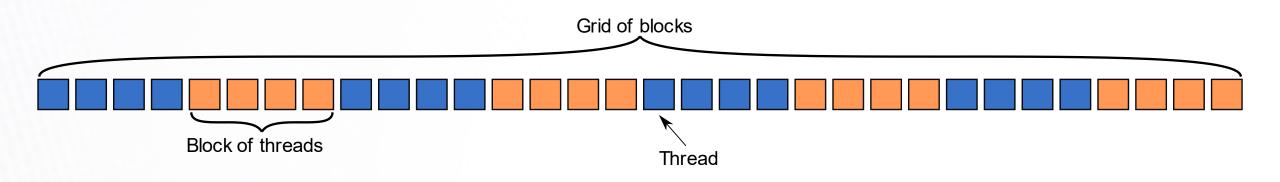
Device Kernels: The Grid

Some Terminology:

| CUDA | HIP | OpenCL™ |
|--------|--------------------|------------|
| grid | grid | NDRange |
| block | block | work group |
| thread | work item / thread | work item |
| warp | wavefront | sub-group |



The Grid: blocks of threads in 1D



Threads in grid have access to:

- Their respective block: blockIdx.x
- Their respective thread ID in a block: threadIdx.x
- Their block's dimension: blockDim.x
- The number of blocks in the grid: gridDim.x

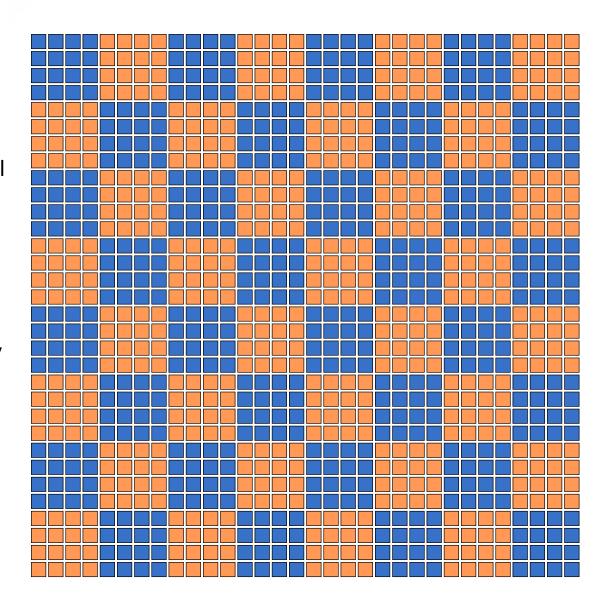


The Grid: blocks of threads in 2D

- Each color is a block of threads
- Each small square is a thread
- The concept is the same in 1D and 2D
- In 2D each block and thread now has a two-dimensional index

Threads in grid have access to:

- Their respective block IDs: blockIdx.x, blockIdx.y
- Their respective thread IDs in a block: threadIdx.x, threadIdx.y
- Etc.





Kernels

A simple embarrassingly parallel loop

```
for (int i=0;i<N;i++) {</pre>
  h a[i] *= 2.0;
```

Can be translated into a GPU kernel:

```
global__ void myKernel(int N, double *d_a) {
int i = threadIdx.x + blockIdx.x*blockDim.x;
if (i<N) {
  d_a[i] *= 2.0;
```

- A device function that will be launched from the host. program is called a kernel and is declared with the _global___ attribute
- Kernels should be declared void
- All pointers passed to kernels must point to memory on the device (more later)
- All threads execute the kernel's body "simultaneously"
- Each thread uses its unique thread and block IDs to compute a global ID
- There could be more than N threads in the grid (we'll see why in a minute)

Kernels

Kernels are launched from the host:

```
dim3 threads(256,1,1);
                                      //3D dimensions of a block of threads
dim3 blocks((N+256-1)/256,1,1);
                                     //3D dimensions the grid of blocks
hipLaunchKernelGGL(myKernel,
                                      //Kernel name ( global void function)
                                     //Grid dimensions
                   blocks,
                   threads,
                                     //Block dimensions
                   0,
                                     //Bytes of dynamic LDS space (see extra slides)
                   0,
                                     //Stream (0=NULL stream)
                   N, a);
                                     //Kernel arguments
```

Analogous to CUDA kernel launch syntax:

```
myKernel<<<blooks, threads, 0, 0>>>(N,a);
```



SIMD operations

Why blocks and threads?

Natural mapping of kernels to hardware:

- Blocks are dynamically scheduled onto CUs
- All threads in a block execute on the same CU
- Threads in a block share LDS memory and L1 cache
- Threads in a block are executed in **64-wide** chunks called "wavefronts"
- Wavefronts execute on SIMD units (Single Instruction Multiple Data)
- If a wavefront stalls (e.g. data dependency) CUs can quickly context switch to another wavefront

A good practice is to make the block size a multiple of 64 and have several wavefronts (e.g. 256 threads)



Device Memory

The host instructs the device to allocate memory in VRAM and records a pointer to device memory:

```
int main() {
  int N = 1000;
  size_t Nbytes = N*sizeof(double);
  double *h_a = (double*) malloc(Nbytes);
                                                       //Host memory
 double *d a = NULL;
 hipMalloc(&d_a, Nbytes);
                                                       //Allocate Nbytes on device
 free(h_a);
                                                       //free host memory
 hipFree(d_a);
                                                       //free device memory
```

Device Memory

The host queues memory transfers:

```
//copy data from host to device
hipMemcpy(d_a, h_a, Nbytes, hipMemcpyHostToDevice);
//copy data from device to host
hipMemcpy(h_a, d_a, Nbytes, hipMemcpyDeviceToHost);
//copy data from one device buffer to another
hipMemcpy(d_b, d_a, Nbytes, hipMemcpyDeviceToDevice);
```



Device Memory

Can copy strided sections of arrays:

```
hipMemcpy2D(d_a, //pointer to destination
          DLDAbytes, //pitch of destination array
          h_a, //pointer to source
          LDAbytes, //pitch of source array
          Nbytes, //number of bytes in each row
          Nrows, //number of rows to copy
          hipMemcpyHostToDevice);
```



Error Checking

Most HIP API functions return error codes of type hipError t

```
hipError t status1 = hipMalloc(...);
hipError t status2 = hipMemcpy(...);
```

- If API function was error-free, returns hipSuccess, otherwise returns an error code.
- Can also peek/get at last error returned with

```
hipError t status3 = hipGetLastError();
hipError t status4 = hipPeekLastError();
```

Can get a corresponding error string using hipGetErrorString(status). Helpful for debugging, e.g.

```
#define HIP CHECK(command) {
 hipError t status = command;
 if (status!=hipSuccess) {
    std::cerr << "Error: HIP reports " << hipGetErrorString(status) << std::endl; \</pre>
    std::abort(); } }
```

Putting it all together

```
#include "hip/hip runtime.h"
int main() {
  int N = 1000;
  size t Nbytes = N*sizeof(double);
  double *h a = (double*) malloc(Nbytes); //host memory
  double *d a = NULL;
  HIP CHECK(hipMalloc(&d a, Nbytes));
  HIP_CHECK(hipMemcpy(d_a, h_a, Nbytes, hipMemcpyHostToDevice)); //copy data to device
  hipLaunchKernelGGL(myKernel, dim3((N+256-1)/256,1,1), dim3(256,1,1), 0, 0, N, d a); //Launch kernel
  HIP CHECK(hipGetLastError());
 HIP CHECK(hipMemcpy(h a, d a, Nbytes, hipMemcpyDeviceToHost)
 free(h a);
                         //free host memory
 HIP_CHECK(hipFree(d_a)); //free device memory
```

```
global void myKernel(int N, double *d a) {
  int i = threadIdx.x + blockIdx.x*blockDim.x;
  if (i<N) {</pre>
    d a[i] *= 2.0;
#define HIP_CHECK(command) {
```

hipError t status = command;

std::cerr << "Error: HIP reports"</pre>

<< std::endl;

<< hipGetErrorString(status)</pre>

if (status!=hipSuccess) {

std::abort(); } }

AMDA

Device management and asynchronous computing



Device Management

Multiple GPUs in system? Multiple host threads/MPI ranks? What device are we running on?

Host can query number of devices visible to system:

```
int numDevices = 0;
hipGetDeviceCount(&numDevices);
```

Host tells the runtime to issue instructions to a particular device:

```
int deviceID = 0;
hipSetDevice(deviceID);
```

Host can guery what device is currently selected:

```
hipGetDevice(&deviceID);
```

- The host can manage several devices by swapping the currently selected device during runtime.
- MPI ranks can set different devices or over-subscribe (share) devices.



Device Properties

The host can also query a device's properties:

```
hipDeviceProp_t props;
hipGetDeviceProperties(&props, deviceID);
```

- hipDeviceProp_t is a struct that contains useful fields like the device's name, total VRAM, clock speed, and GCN architecture.
 - See "hip/hip_runtime_api.h" for full list of fields.

Blocking vs Nonblocking API functions

- The kernel launch function, hipLaunchKernelGGL, is non-blocking for the host.
 - After sending instructions/data, the host continues immediately while the device executes the kernel
 - If you know the kernel will take some time, this is a good area to do some work (i.e. MPI comms) on the host
- However, hipMemcpy is blocking.
 - The data pointed to in the arguments can be accessed/modified after the function returns.
- The non-blocking version is hipMemcpyAsync

```
hipMemcpyAsync(d_a, h_a, Nbytes, hipMemcpyHostToDevice, stream);
```

- Like hipLaunchKernelGGL, this function takes an argument of type hipStream_t
- It is not safe to access/modify the arguments of hipMemcpyAsync without some sort of synchronization.

Putting it all together

```
#include "hip/hip_runtime.h"
                                                                  global void myKernel(int N, double *d a) {
int main() {
                                                                  int i = threadIdx.x + blockIdx.x*blockDim.x;
  int N = 1000;
                                                                  if (i<N) {</pre>
  size t Nbytes = N*sizeof(double);
                                                                    d a[i] *= 2.0;
  double *h a = (double*) malloc(Nbytes); //host memory
  double *d a = NULL;
  HIP CHECK(hipMalloc(&d a, Nbytes));
  HIP CHECK(hipMemcpy(d a, h a, Nbytes, hipMemcpyHostToDevice)); //copy data to device
  hipLaunchKernelGGL(myKernel, dim3((N+256-1)/256,1,1), dim3(256,1,1), 0, 0, N, d_a); //Launch kernel
  HIP CHECK(hipGetLastError());
                                                           The host waits for the kernel to finish here
 HIP_CHECK(hipMemcpy(h_a, d_a, Nbytes, hipMemcpyDeviceToHost));
                                                                   //copy results back to host
 free(h_a);
                         //free host memory
 HIP_CHECK(hipFree(d_a)); //free device memory
```

- A stream in HIP is a queue of tasks (e.g. kernels, memcpys, events).
 - Tasks enqueued in a stream complete in order on that stream.
 - Tasks being executed in different streams are allowed to overlap and share device resources.
- Streams are created via:

```
hipStream t stream;
hipStreamCreate(&stream);
```

And destroyed via:

```
hipStreamDestroy(stream);
```

- Passing 0 or NULL as the hipStream t argument to a function instructs the function to execute on a stream called the 'NULL Stream':
 - No task on the NULL stream will begin until all previously enqueued tasks in all other streams have completed.
 - Blocking calls like hipMemcpy run on the NULL stream.



Suppose we have 4 small kernels to execute:

```
hipLaunchKernelGGL(myKernel1, dim3(1), dim3(256), 0, 0, 256, d_a1);
hipLaunchKernelGGL(myKernel2, dim3(1), dim3(256), 0, 0, 256, d a2);
hipLaunchKernelGGL(myKernel3, dim3(1), dim3(256), 0, 0, 256, d_a3);
hipLaunchKernelGGL(myKernel4, dim3(1), dim3(256), 0, 0, 256, d_a4);
```

Even though these kernels use only one block each, they'll execute in serial on the NULL stream:

NULL Stream myKernel1 myKernel2 myKernel3 myKernel4

With streams we can effectively share the GPU's compute resources:

```
hipLaunchKernelGGL(myKernel1, dim3(1), dim3(256), 0, stream1, 256, d_a1);
hipLaunchKernelGGL(myKernel2, dim3(1), dim3(256), 0, stream2, 256, d a2);
hipLaunchKernelGGL(myKernel3, dim3(1), dim3(256), 0, stream3, 256, d_a3);
hipLaunchKernelGGL(myKernel4, dim3(1), dim3(256), 0, stream4, 256, d_a4);
```

| NULL Stream | | |
|-------------|-----------|--|
| Stream1 | myKernel1 | |
| Stream2 | myKernel2 | |
| Stream3 | myKernel3 | |
| Stream4 | myKernel4 | |

Note 1: Check that the kernels modify different parts of memory to avoid data races.

Note 2: With large kernels, overlapping computations may not help performance.

- There is another use for streams besides concurrent kernels:
 - Overlapping kernels with data movement.
- AMD GPUs have separate engines for:
 - Host->Device memcpys
 - Device->Host memcpys
 - Compute kernels.
- These three different operations can overlap without dividing the GPU's resources.
 - The overlapping operations should be in separate, non-NULL, streams.
 - The host memory should be **pinned.**



Pinned Memory

Host data allocations are pageable by default. The GPU can directly access Host data if it is pinned instead.

• Allocating pinned host memory:

```
double *h a = NULL;
hipHostMalloc(&h_a, Nbytes);
```

Free pinned host memory:

```
hipHostFree(h a);
```

- Host<->Device memcpy bandwidth increases significantly when host memory is pinned.
 - It is good practice to allocate host memory that is frequently transferred to/from the device as pinned memory.

```
Suppose we have 3 kernels which require moving data to and from the device:
    hipMemcpy(d_a1, h_a1, Nbytes, hipMemcpyHostToDevice));
    hipMemcpy(d_a2, h_a2, Nbytes, hipMemcpyHostToDevice));
    hipMemcpy(d a3, h a3, Nbytes, hipMemcpyHostToDevice));
    hipLaunchKernelGGL(myKernel1, blocks, threads, 0, 0, N, d a1);
    hipLaunchKernelGGL(myKernel2, blocks, threads, 0, 0, N, d a2);
    hipLaunchKernelGGL(myKernel3, blocks, threads, 0, 0, N, d a3);
    hipMemcpy(h_a1, d_a1, Nbytes, hipMemcpyDeviceToHost);
    hipMemcpy(h a2, d a2, Nbytes, hipMemcpyDeviceToHost);
    hipMemcpy(h_a3, d_a3, Nbytes, hipMemcpyDeviceToHost);
```

| NULL Stream HToD1 HToD | HToD3 myKernel1 | myKernel2 myKernel3 [| DToH1 DToH2 DToH3 |
|------------------------|-----------------|-----------------------|-------------------|
|------------------------|-----------------|-----------------------|-------------------|

Changing to asynchronous memcpys and using streams:

```
hipMemcpyAsync(d_a1, h_a1, Nbytes, hipMemcpyHostToDevice, stream1);
hipMemcpyAsync(d_a2, h_a2, Nbytes, hipMemcpyHostToDevice, stream2);
hipMemcpyAsync(d_a3, h_a3, Nbytes, hipMemcpyHostToDevice, stream3);
hipLaunchKernelGGL(myKernel1, blocks, threads, 0, stream1, N, d_a1);
hipLaunchKernelGGL(myKernel2, blocks, threads, 0, stream2, N, d a2);
hipLaunchKernelGGL(myKernel3, blocks, threads, 0, stream3, N, d a3);
hipMemcpyAsync(h_a1, d_a1, Nbytes, hipMemcpyDeviceToHost, stream1);
hipMemcpyAsync(h a2, d a2, Nbytes, hipMemcpyDeviceToHost, stream2);
hipMemcpyAsync(h a3, d a3, Nbytes, hipMemcpyDeviceToHost, stream3);
```

| NULL Stream | | | | | | |
|-------------|-------|-----------|-----------|-----------|-------|--|
| Stream1 | HToD1 | myKernel1 | DToH1 | | | |
| Stream2 | | HToD2 | myKernel2 | DToH2 | | |
| Stream3 | | | HToD3 | myKernel3 | DToH3 | |

Synchronization

How do we coordinate execution on device streams with host execution? Need some synchronization points.

- hipDeviceSynchronize();
 - Heavy-duty sync point.
 - Blocks host until all work in all device streams has reported complete.
- hipStreamSynchronize(stream);
 - Blocks host until all work in stream has reported complete.

Can a stream synchronize with another stream? For that we need 'Events'.

Events

```
A hipEvent t object is created on a device via:
    hipEvent_t event;
    hipEventCreate(&event);
We queue an event into a stream:
    hipEventRecord(event, stream);
```

- The event records what work is currently enqueued in the stream.
- When the stream's execution reaches the event, the event is considered 'complete'.

At the end of the application, event objects should be destroyed:

```
hipEventDestroy(event);
```

Events

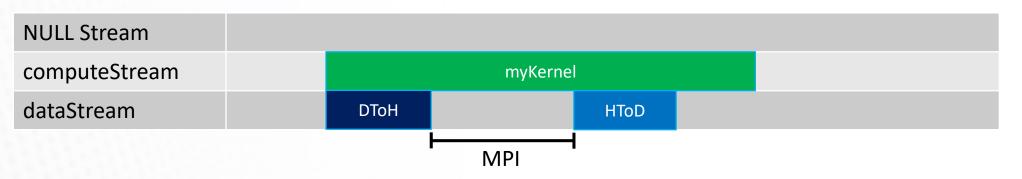
What can we do with queued events?

- hipEventSynchronize(event);
 - Block host until event reports complete.
 - Only a synchronization point with respect to the stream where event was enqueued.
- hipEventElapsedTime(&time, startEvent, endEvent);
 - Returns the time in ms between when two events, startEvent and endEvent, completed
 - Can be very useful for timing kernels/memcpys
- hipStreamWaitEvent(stream, event);
 - Non-blocking for host.
 - Instructs all future work submitted to stream to wait until event reports complete.
 - Primary way we enforce an 'ordering' between tasks in separate streams.



```
A common use-case for streams is MPI traffic:
```

```
//Queue local compute kernel
hipLaunchKernelGGL(myKernel, blocks, threads, 0, computeStream, N, d_a);
//Copy halo data to host
hipMemcpyAsync(h commBuffer, d commBuffer, Nbytes, hipMemcpyDeviceToHost, dataStream);
hipStreamSynchronize(dataStream); //Wait for data to arrive
//Exchange data with MPI
MPI_Data_Exchange(h_commBuffer);
//Send new data back to device
hipMemcpyAsync(d_commBuffer, h_commBuffer, Nbytes, hipMemcpyHostToDevice, dataStream);
```



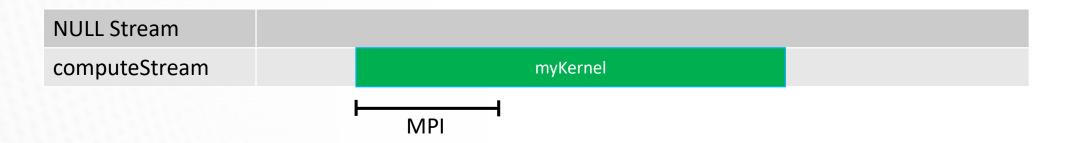
```
With a GPU-aware MPI stack, the Host<->Device traffic can be omitted:
```

```
//Some synchronization so that data on GPU and local compute are ready
hipDeviceSynchronize();

//Exchange data with MPI (with device pointer)
MPI_Data_Exchange(d_commBuffer, &request);

//Queue local compute kernel
hipLaunchKernelGGL(myKernel, blocks, threads, 0, computeStream, N, d_a);

//Wait for MPI request to complete
MPI_Wait(&request, &status);
```





AMDA

Device code, shared memory, and thread synchronization



Function Qualifiers

hipcc makes two compilation passes through source code. One to compile host code, and one to compile device code.

- global functions:
 - These are entry points to device code, called from the host
 - Code in these regions will execute on SIMD units
- device functions:
 - Can be called from global and other device functions.
 - Cannot be called from host code.
 - Not compiled into host code essentially ignored during host compilation pass
- host device functions:
 - Can be called from __global___, __device___, and host functions.
 - Will execute on SIMD units when called from device code!



SIMD Execution

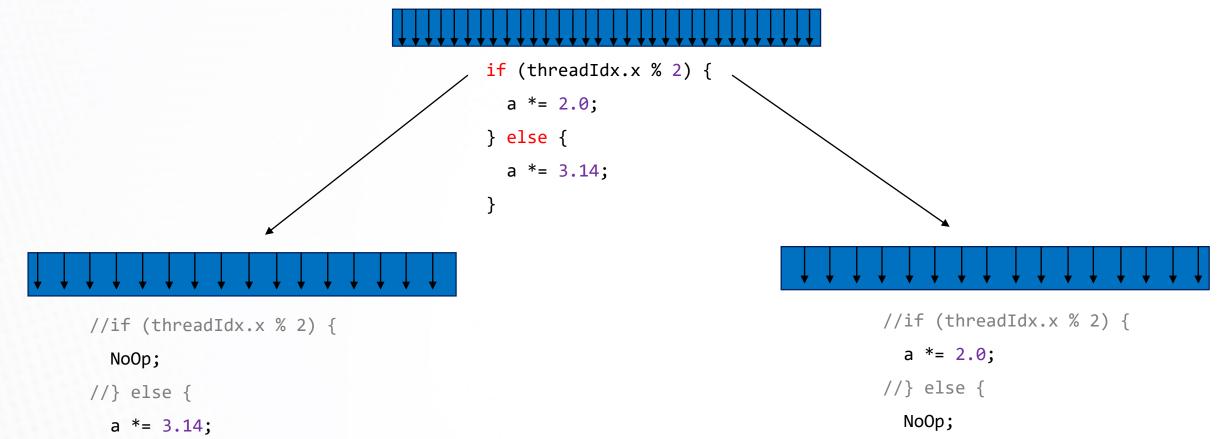
On SIMD units, be aware of divergence.

- Branching logic (if else) can be costly:
 - Wavefront encounters an if statement
 - Evaluates conditional
 - If true, continues to statement body
 - If false, also continues to statement body with all instructions replaced with NoOps.
 - Known as 'thread divergence'
- Generally, wavefronts diverging from each other is okay.
- Thread divergence within a wavefront can impact performance.



SIMD Execution

//}





//}

Memory declarations in Device Code

- Malloc/free not supported in device code.
- Variables/arrays can be declared on the stack.
- Stack variables declared in device code are allocated in registers and are private to each thread.
- Threads can all access common memory via device pointers, but otherwise do not share memory.
 - Important exception: shared memory
- Stack variables declared as shared :
 - Allocated once per block in LDS memory
 - Shared and accessible by all threads in the same block
 - Access is faster than device global memory (but slower than register)
 - Must have size known at compile time



```
Shared Memory
__global__ void reverse(double *d_a) {
 __shared__ double s_a[256]; //array of doubles, shared in this block
 int tid = threadIdx.x;
 s_a[tid] = d_a[tid]; //each thread fills one entry
 //all wavefronts must reach this point before any wavefront is allowed to continue.
                         //something is missing here...
syncthreads();
 d_a[tid] = s_a[255-tid]; //write out array in reverse order
int main() {
 hipLaunchKernelGGL(reverse, dim3(1), dim3(256), 0, 0, d_a); //Launch kernel
```

Thread Synchronization

- syncthreads():
 - Blocks a wavefront from continuing execution until all wavefronts have reached syncthreads()
 - Memory transactions made by a thread before syncthreads() are visible to all other threads in the block after __syncthreads()
 - Can have a noticeable overhead if called repeatedly

- **Best practice:** Avoid deadlocks by checking that **all** threads in a block execute **the same** syncthreads() instruction.
- Note 1: So long as at least one thread in the wavefront encounters __syncthreads(), the whole wavefront is considered to have encountered syncthreads().
- Note 2: Wavefronts can synchronize at different syncthreads() instructions, and if a wavefront exits a kernel completely, other wavefronts waiting at a __syncthreads() may be allowed to continue.

HIP API

- Device Management:
 - hipSetDevice(), hipGetDevice(), hipGetDeviceProperties()
- Memory Management
 - hipMalloc(), hipMemcpy(), hipMemcpyAsync(), hipFree()
- Streams
 - hipStreamCreate(), hipSynchronize(), hipStreamSynchronize(), hipStreamFree()
- Events
 - hipEventCreate(), hipEventRecord(), hipStreamWaitEvent(), hipEventElapsedTime()
- Device Kernels
 - __global___, __device___, hipLaunchKernelGGL()
- Device code
 - threadIdx, blockIdx, blockDim, shared
 - 200+ math functions covering entire CUDA math library.
- Error handling
 - hipGetLastError(), hipGetErrorString()



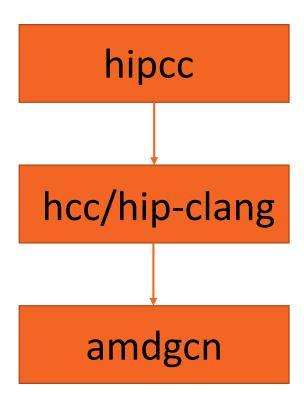


AMD GPU Software

René van Oostrum <rene.vanoostrum@amd.com>

AMD GPU Compilers

- AMD supports several compilers that emit AMDGCN assembly
 - hcc
 - A fork of clang, will eventually be deprecated.
 - hip-clang
 - Based on upstream version of clang, currently in development
 - AOMP
- hipcc
 - A perl script that will invoke nvcc or hcc (or hip-clang).
 - Will automatically set the command line options needed by each compiler.
- hcc
 - Compiles HIP applications
 - Fork of clang
 - Generates and links AMDGCN from HIP device code
 - All the x86 host code is handled as usual by clang
- AOMP (AMD OpenMP Compiler)
 - Compiles C/C++ code with OpenMP "target" pragmas
 - Links with libomptarget to produce a binary that can offload work to the GPU





AMD GPU Compilers: the FORTRAN story

- FORTRAN is a technology important to the US Department of Energy
- AMD has plans to support OpenMP 4.5+ target offload from FORTRAN with two open source options
 - F18 (based on llvm)
 - gfortran
- FORTRAN compiler work is an ongoing effort
- See the Frontier spec sheet for what is expected to be supported on Frontier
 - https://www.olcf.ornl.gov/wp-content/uploads/2019/05/frontier_specsheet.pdf



hipcc usage

Usage is straightforward. Accepts all/any flags that vanilla clang accepts, e.g.

```
hipcc dotprod.cpp -o dotprod
```

Set HIPCC_VERBOSE=7 to see a bunch of useful information

- Compile and link lines
- Various paths

```
$ HIPCC_VERBOSE=7 hipcc dotprod.cpp -o dotprod
HIP_PATH=/opt/rocm
HIP_PLATFORM=hcc
HSA_PATH=/opt/rocm/hsa
HCC_HOME=/opt/rocm/hcc
hipcc-args: dotprod.cpp -o dotprod
hipcc-cmd: /opt/rocm/hcc/bin/hcc -hc -D__HIPCC__ -isystem /opt/rocm/hcc/include
-isystem /opt/rocm/include/hip/hcc_detail/cuda -isystem /opt/rocm/hsa/include -Wno-deprecated-register
-isystem /opt/rocm/profiler/CXLActivityLogger/include -isystem /opt/rocm/include
-DHIP_VERSION_MAJOR=1 -DHIP_VERSION_MINOR=5 -DHIP_VERSION_PATCH=19284
-D__HIP_ARCH_GFX900__=1 -D__HIP_ARCH_GFX906__=1
dotprod.cpp -o dotprod
```

Inspecting the AMD GCN ISA

- You can inspect the AMD GCN ISA that was emitted by hcc (remember this is just clang)
- The command you need is extractkernel
- This is roughly equivalent to objdump (except it only dumps the GCN assembly, not x86)
- The GCN ISA is publicly available: https://developer.amd.com/wp-content/resources/Vega Shader ISA 28July2017.pdf

```
$ /opt/rocm/bin/extractkernel -i vectoradd
Generated GCN ISA for gfx900 at: vectoradd-000-gfx900.isa
$ grep v add vectoradd.000-gfx900.isa
        v add u32 e32 v1, s1, v1
                                                                    // 00000001138: 68020201
        v add3 u32 v0, s0, v0, v1
                                                                    // 00000001154: D1FF0000 04060000
                                                                    // 0000000119C: 32040002
        v add co u32 e32 v2, vcc, s2, v0
        v addc co u32 e32 v3, vcc, v3, v1, vcc
                                                                    // 000000011A0: 38060303
        v add co u32 e32 v4, vcc, s4, v0
                                                                   // 000000011A8: 32080004
        v addc co u32 e32 v5, vcc, v5, v1, vcc
                                                                   // 000000011AC: 380A0305
        v add co u32 e32 v0, vcc, s0, v0
                                                                   // 000000011c0: 3200000
        v addc co u32 e32 v1, vcc, v6, v1, vcc
                                                                   // 000000011C4: 38020306
        v add f32 e32 v2, v2, v3
                                                                    // 000000011CC: 02040702
```

Installing ROCm

- Requirements:
 - Linux[®]!
 - Ubuntu: ROCm can be installed using a Debian repo.
 - CentOS/RHEL: ROCm can be installed using a yum repo.
 - Other distros: you must build from source (support planned for SUSE based distributions).
 - To run on AMD hardware you need discrete GPUs in families GFX8 (Polaris) or GFX9 (Vega)
 - APUs aren't currently officially supported
- ROCm has been on a monthly release cycle
- ROCm is now compatible with AMD drivers in some upstream linux kernels.
- ROCm can be installed with:
 - ROCK kernel driver (from ROCm repos)
 - Only supported on Ubuntu, CentOS/RHEL
 - AMD drivers in some upstream kernels
 - Should work on more distributions
- Latest install instructions can be found here: https://github.com/RadeonOpenCompute/ROCm
- Also check out: https://github.com/RadeonOpenCompute/Experimental ROC



Querying System

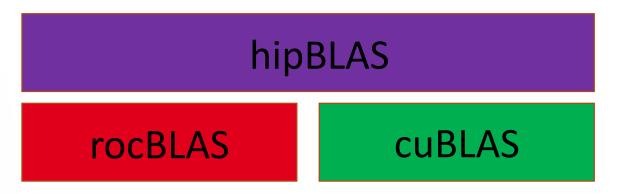
- rocminfo: Queries and displays information on the system's hardware
 - More info at: https://github.com/RadeonOpenCompute/rocminfo
- Querying ROCm version:
 - If you install ROCm in the standard location (/opt/rocm) version info is at: /opt/rocm/.info/version-dev
 - Can also run the command 'dkms status' and the ROCm version will be displayed
- rocm-smi: Queries and sets AMD GPU frequencies, power usage, and fan speeds
 - sudo privileges are needed to set frequencies and power limits
 - sudo privileges are not needed to query information
 - Get more info by running 'rocm-smi -h' or looking at: https://github.com/RadeonOpenCompute/ROC-smi

| ==== | | | =====ROCm | System | Manage | ement Int | erface== | | |
|------|-------|--------|-----------|--------|---------|-----------|----------|-------|------|
| GPU | Temp | AvgPwr | SCLK | MCLK | Fan | Perf | PwrCap | VRAM% | GPU% |
| 1 | 38.0c | 18.0W | 1440Mhz | 945Mhz | 0.0% | manual | 220.0W | 0% | 0% |



AMD GPU Libraries

- A note on naming conventions:
 - roc* -> AMGCN library usually written in HIP
 - cu* -> NVIDIA PTX libraries
 - hip* -> usually interface layer on top of roc*/cu* backends
- hip* libraries:
 - Can be compiled by hipcc and can generate a call for the device you have:
 - hipcc->hcc->AMD GCN ISA
 - hipcc->nvcc (inlined)->NVPTX
 - Just a thin wrapper that marshals calls off to a "backend" library:
 - corresponding roc* library backend containing optimized GCN
 - corresponding cu* library backend containing NVPTX for NVIDIA devices
 - E.g., hipBLAS is a marshalling library:





Decoder ring: Math library equivalents

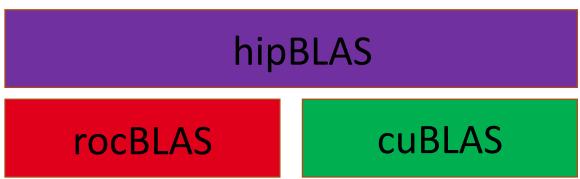
| CUBLAS | ROCBLAS | Basic Linear Algebra Subroutines |
|--------|-----------|--|
| CUFFT | ROCFFT | Fast Fourier Transforms |
| THRUST | ROCTHRUST | Deep Learning Library |
| CUB | ROCPRIM | Optimized Parallel Primitives |
| EIGEN | EIGEN | C++ Template Library for Linear Algebra |

MORE INFO AT: GITHUB.COM/ROCM-DEVELOPER-TOOLS/HIP → HIP_PORTING_GUIDE.MD



AMD GPU Libraries: BLAS

- rocBLAS `sudo apt install rocblas`
 - Source code: https://github.com/ROCmSoftwarePlatform/rocBLAS
 - Documentation: https://rocblas.readthedocs.io/en/latest/
 - Basic linear algebra functionality
 - axpy, gemv, trsm, etc
 - Use hipBLAS if you need portability between AMD and NVIDIA devices
- hipBLAS `sudo apt install hipblas`
 - Documentation: https://github.com/ROCmSoftwarePlatform/hipBLAS/wiki/Exported-functions
 - Use this if you need portability between AMD and NVIDIA
 - It is just a thin wrapper:
 - It can dispatch calls to rocBLAS for AMD devices
 - It can dispatch calls to cuBLAS for NVIDIA devices





AMD GPU Libraries: rocBLAS example

- rocBLAS
 - Documentation: https://rocblas.readthedocs.io/en/latest/
 - Level 1, 2, and 3 functionality
 - axpy, gemv, trsm, etc
 - Note: rocBLAS syntax matches BLAS closer than hipBLAS or cuBLAS
 - Use hipBLAS only if you need portability between AMD and NVIDIA devices
 - Link with: -lrochlas

```
#include <rochlas.h>
int main(int argc, char ** argv) {
  rocblas int N = 500000;
  // Allocate device memory
  double * dx, * dy;
  hipMalloc(&dx, sizeof(double) * N);
  hipMalloc(&dy, sizeof(double) * N);
  // Allocate host memory (and fill up the arrays) here
  std::vector<double> hx(N), hy(N);
  // Copy host arrays to device
  hipMemcpy(dx, hx.data(), sizeof(double) * N, hipMemcpyHostToDevice);
  hipMemcpy(dy, hy.data(), sizeof(double) * N, hipMemcpyHostToDevice);
  const double alpha = 1.0;
  rocblas handle handle;
  rocblas create handle(&handle);
  rocblas status status;
  status = rocblas daxpy(handle, N, &alpha, dx, 1, dy, 1);
  rocblas destroy handle(handle);
  // Copy result back to host
  hipMemcpy(hy.data(), dy, sizeof(double) * N, hipMemcpyDeviceToHost);
  hipFree(dx);
  hipFree(dy);
  return 0;
```

Some Links to Key Libraries

- BLAS
 - rocBLAS (https://github.com/ROCmSoftwarePlatform/rocBLAS)
 - hipBLAS (hipBLAS (hitps://github.com/ROCmSoftwarePlatform/hipBLAS (hitp
- FFTs
 - rocFFT (https://github.com/ROCmSoftwarePlatform/rocFFT)
- Random number generation
 - rocRAND (https://github.com/ROCmSoftwarePlatform/rocRAND)
 - hipRAND (https://github.com/ROCmSoftwarePlatform/hipRAND)
- Sparse linear algebra
 - rocSPARSE (https://github.com/ROCmSoftwarePlatform/rocSPARSE)
 - hipSPARSE (https://github.com/ROCmSoftwarePlatform/hipSPARSE)
- Iterative solvers
 - rocALUTION (https://github.com/ROCmSoftwarePlatform/rocALUTION)
- Parallel primitives
 - rocPRIM (https://github.com/ROCmSoftwarePlatform/rocPRIM)
 - hipCUB (https://github.com/ROCmSoftwarePlatform/hipCUB)



More links to key libraries

Machine Learning libraries and Frameworks

- Tensorflow: https://github.com/ROCmSoftwarePlatform/tensorflow-upstream
- Pytorch: https://github.com/ROCmSoftwarePlatform/pytorch
- MIOpen (similar to cuDNN): https://github.com/ROCmSoftwarePlatform/MIOpen
- Tensile: https://github.com/ROCmSoftwarePlatform/Tensile
- RCCL (ROCm analogue of NCCL): https://github.com/ROCmSoftwarePlatform/rccl





Porting CUDA Applications to HIP

Paul Bauman <paul.bauman@amd.com>

Objectives

This training:

- Demonstrates how to convert CUDA codes into HIP
- Explains the meaning of the term 'hipify'
- Provides a simple means to examine port quality
- Provides an idea of the common 'gotchas' of porting apps



Getting started with HIP

CUDA VECTOR ADD

```
global void add(int n,
                  double *x,
                  double *y){
int index = blockIdx.x * blockDim.x
             + threadIdx.x;
int stride = blockDim.x * gridDim.x;
for (int i = index; i < n; i += stride){</pre>
  y[i] = x[i] + y[i];
```

HIP VECTOR ADD

```
_global__ void add(int n,
                   double *x,
                   double *y){
int index = blockIdx.x * blockDim.x
             + threadIdx.x;
int stride = blockDim.x * gridDim.x;
for (int i = index; i < n; i += stride){</pre>
  y[i] = x[i] + y[i];
```

KERNELS ARE SYNTACTICALLY THE SAME

CUDA APIs vs HIP API

| CUDA | HIP |
|---|--|
| cudaMalloc(&d_x, N*sizeof(double)); | hipMalloc(&d_x, N*sizeof(double)); |
| <pre>cudaMemcpy(d_x, x, N*sizeof(double),</pre> | <pre>hipMemcpy(d_x, x, N*sizeof(double), hipMemcpyHostToDevice);</pre> |
| cuda Device Synchronize (); | hipDeviceSynchronize(); |

Launching a kernel

CUDA KERNEL LAUNCH SYNTAX

HIP KERNEL LAUNCH SYNTAX

```
some_kernel<<<gridsize, blocksize,</pre>
               shared_mem_size, stream>>>
               (arg0, arg1, ...);
```

```
hipLaunchKernelGGL(some_kernel,
                   gridsize, blocksize,
                   shared_mem_size, stream,
                   arg0, arg1, ...);
```

Difference between HIP and CUDA

Some things to be aware of writing HIP, or porting from CUDA:

- AMD GCN hardware 'warp' size = 64 (warps are referred to as 'wavefronts' in AMD documentation)
- Device and host pointers allocated by HIP API use flat addressing
 - Unified virtual addressing is enabled by default
 - Unified memory is available, but does not perform optimally currently
- Dynamic parallelism not currently supported
- CUDA 9+ thread independent scheduling not supported (e.g., no syncwarp)
- Some CUDA library functions do not have AMD equivalents
- Shared memory and registers per thread can differ between AMD and Nvidia hardware
- Inline PTX or AMD GCN assembly is not portable

Despite differences, majority of CUDA code in applications can be simply translated.



Enter Hipify

- AMD provides 'Hipify' tools to automatically convert most CUDA code
 - Hipify-perl
 - Hipify-clang
- Good resource to help with porting: https://github.com/ROCm-Developer- Tools/HIP/blob/master/docs/markdown/hip porting guide.md
- In practice, large portions of many HPC codes have been automatically Hipified:
 - ~90% of CUDA code in CORAL-2 HACC
 - ~80% of CUDA code in CORAL-2 PENNANT
 - ~80% of CUDA code in CORAL-2 QMCPack
 - ~95% of CUDA code in CORAL-2 Laghos

The remaining code requires programmer intervention



Hipify tools

- Hipify-perl:
 - Easy to use –point at a directory and it will attempt to hipify CUDA code
 - Very simple string replacement technique: may make incorrect translations
 - sed -e 's/cuda/hip/g', (e.g., cudaMemcpy becomes hipMemcpy)
 - Recommended for quick scans of projects
- Hipify-clang:
 - Requires clang compiler
 - More robust translation of the code. Uses clang to parse files and perform semantic translation
 - Can generate warnings and assistance for code for additional user analysis
 - High quality translation, particularly for cases where the user is familiar with the make system



Hipify-perl

- Sits in \$HIP/bin/ (export PATH=\$PATH:[MYHIP]/bin)
- Command line tool: hipify-perl foo.cu > new_foo.cpp
- Compile: hipcc new_foo.cpp
- How does this this work in practice?
 - Hipify source code
 - Check it in to your favorite version control
 - Try to build
 - Manually work on the rest



Hipify-clang

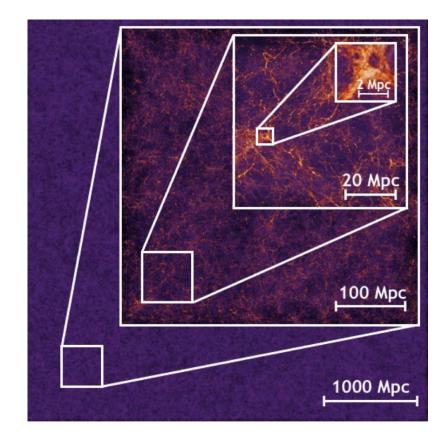
- Build from source
- hipify-clang has unit tests using LLVM lit/FileCheck (44 tests)
- Hipification requires same headers that would be needed to compile it with clang:
- ./hipify-clang foo.cu -I /usr/local/cuda-8.0/samples/common/inc

https://github.com/ROCm-Developer-Tools/HIP/tree/master/hipify-clang



Example: HACC

- Hardware Accelerated Cosmology Code
- Simulates time-evolution of universe
 - Mpc = Megaparsec = 3.09×10^{22} meters
- Our HIP success story:
 - Ported in an afternoon
- Profiling:
 - 10% of time is spent in the tree walk
 - >80% in the short force kernels
 - (GPU kernel)
 - 5% in the 3d Transposes / FFTs



$$f_{SR} = (s + \epsilon)^{-3/2} - f_{grid}(s)$$

where,

$$s = \mathbf{r} \cdot \mathbf{r}$$

and.

$$f_{grid}(s) = POLY_5(s)$$



HACC: What made it a success

What was easy?

- Simple GPU kernel
- Few library dependencies (FFTW, not in kernel)
- No advanced CUDA features

What was difficult?

- Inline PTX: required translation to AMD GCN
- Hand-written wave-32 code (for a reduction)



Porting HACC

CUDA

```
cudaMemcpyAsync(d_npos,h_npos,Nposbytes,
                 cudaMemcpyHostToDevice,stream);
cudaMemcpyAsync(d_mask,h_mask,NmaskBytes,
                 cudaMemcpyHostToDevice,stream);
calcHHCullenDehnen<<<br/>blocksPerGrid,
               threadsPerBlock, 0, stream>>>
               (cnt, SIZE, d_npos, d_mask, rsm);
cudaMemcpyAsync(h pos,
                 d npos+(SIZE-cnt), cntBytes,
                 cudaMemcpyDeviceToHost,stream);
cudaMemcpyAsync(h_mask,d_mask,NmaskBytes,
                 cudaMemcpyDeviceToHost,stream);
```

HIP

```
hipMemcpyAsync(d npos,h npos,Nposbytes,
               hipMemcpyHostToDevice, stream);
hipMemcpyAsync(d mask,h mask,NmaskBytes,
               hipMemcpyHostToDevice, stream);
hipLaunchKernelGGL((calcHHCullenDehnen),
      blocksPerGrid, threadsPerBlock, 0, stream,
      cnt, SIZE, d npos, d mask, rsm);
hipMemcpyAsync(h_pos,
               d npos+(SIZE-cnt), cntBytes,
               hipMemcpyDeviceToHost, stream);
hipMemcpyAsync(h mask,d mask,NmaskBytes,
               hipMemcpyDeviceToHost,stream);
```

HACC: Comparison to the CUDA version

| MEASUREMENT | TITAN-V W/ CUDA 9.0 | MI-25 W/ HIP (ROCM 1.9) |
|---------------------|---------------------|-------------------------|
| REGISTER USE (VGPR) | 32 | 36 |
| REGISTER SPILLING | 0 | 0 |
| DIVERGENCE | 92.5% | 84% |



Performance portability

| MEASUREMENT | TITAN-V W/ CUDA 9.0 | MI-25 W/ HIP (ROCM 1.9) | | |
|---|---------------------|-------------------------|--|--|
| REGISTER USE (VGPR) | 30, 74, 44, 58, 53 | 12, 73, 25, 38, 24 | | |
| REGISTER SPILLING | 0, 0, 0, 0 | 0, 0, 0, 0, 0 | | |
| DIVERGENCE | 97, 95, 70, 99, 78 | 97, 90, 71, 98, 82 | | |
| Pennant has 5 gpu kernels: Multiple entries denote each GPU kernel (GPUMain1, etc,) | | | | |
| Similar numbers observed in Quicksilver, HACC, etc | | | | |



HIPified (ported) codes

QUICKSILVER HACC SW4LITE HPL PENNANT LAGHOS

AND (ALREADY) MANY MORE...



Portability layers using HIP

Several portability layers are already supporting, or implementing, HIP

- RAJA
 - HIP kernel execution policies syntactically identical to CUDA
 - Official PRs under review
- Kokkos
 - HIP kernel execution policies syntactically identical to CUDA
 - Support is in Alpha and under development by Kokkos and AMD developers
- OCCA
 - OKL kernels can compile for HIP devices
 - Available in OCCA's master branch
- OpenMP 5.0
 - gcc and AMD's aomp compilers support target offload regions, interop with HIP



Tips and tricks for performance

What to look for when porting:

- Inline PTX assembly
- **CUDA** intrinsics
- Hard-coded dependencies on warp size, shared memory size
 - "grep 32"
 - Do not use hard coded dependencies on warp size!
- Code geared toward limited size of register file on NVIDIA hardware
- Functions implicitly inlined
- **Unified Memory**





QUESTIONS?

Dynamic Shared Memory

- Can actually use <u>shared</u> arrays when sizes aren't known at compile time
 - Called dynamic shared memory
 - Declare one array using HIP_DYNAMIC_SHARED macro, use for all dynamic LDS space
 - Use the hipLaunchKernelGGL argument we haven't discussed yet



Dynamic Shared Memory

```
__global__ void reverse(double *d_a, int N) {
 HIP_DYNAMIC_SHARED(double, s_a); //dynamic array of doubles, shared in this block
 int tid = threadIdx.x;
 s_a[tid] = d_a[tid]; //each thread fills one entry
 //all wavefronts should reach this point before any wavefront is allowed to continue.
  __syncthreads();
 d a[tid] = s a[N-1-tid]; //write out array in reverse order
int main() {
  size t NsharedBytes = N*sizeof(double);
 hipLaunchKernelGGL(reverse, dim3(1), dim3(N), NsharedBytes, 0, d_a, N); //Launch kernel
```

Atomic Operations

Atomic functions:

- Perform a read+write of a single 32 or 64-bit word in device global or LDS memory
- Can be called by multiple threads in device code
- Performed in a conflict-free manner
- AMD GPUs support atomic operations on 32-bit integers in hardware
 - Float /double atomics implemented as atomicCAS (Compare And Swap) loops, may have poor performance
- Can check at compile time if 32 or 64-bit atomic instructions are supported on target device
 - #ifdef __HIP_ARCH_HAS_GLOBAL_INT32_ATOMICS___
 - #ifdef __HIP_ARCH_HAS_GLOBAL_INT64_ATOMICS___

Atomic Operations

Supported atomic operations in HIP:

| Operation | Type, T | Notes |
|---------------------------------|-----------------------------------|--|
| T atomicAdd(T* address, T val) | int, long long int, float, double | Adds val to *address |
| T atomicExch(T* address, T val) | int, long long int, float | Replace *address with val and return old value |
| T atomicMin(T* address, T val) | int, long long int | Replaces *address if val is smaller |
| T atomicMax(T* address, T val) | int, long long int | Replaces *address if val is larger |
| T atomicAnd(T* address, T val) | int, long long int | Bitwise AND between *address and val |
| T atomicOr(T* address, T val) | int, long long int | Bitwise OR between *address and val |
| T atomicXor(T* address, T val) | int, long long int | Bitwise XOR between *address and val |



AMD GPU programming resources

- ROCm platform: https://github.com/RadeonOpenCompute/ROCm/
 - With instructions for installing from Debian/CentOS/RHEL binary repositories
 - Has links to source repositories for all components, including HIP
- HIP porting guide: https://github.com/ROCm-Developer-Tools/HIP/blob/master/docs/markdown/hip_porting_guide.md
- ROCm/HIP libraries: https://github.com/ROCmSoftwarePlatform
- ROC-profiler: https://github.com/ROCm-Developer-Tools/rocprofiler
 - Collects application traces and performance counters
 - Trace timeline can be visualized with chrome://tracing
- AMD GPU ISA docs and more: https://developer.amd.com/resources/developer-guides-manuals/



CUDA features not supported by HIP

- CUDA 5.0:
 - Dynamic Parallelism (not supported)
 - culpc functions (under development).
- CUDA 5.5:
 - CUPTI (not directly supported, AMD GPUPerfAPI an alternative in some cases)
- **CUDA 6.0**
- Managed memory (under development)



CUDA features not in HIP, cont.

- **CUDA 7.0**
 - Per-thread-streams (under development)
- **CUDA 8.0**
 - Page Migration including cudaMemAdvise, cudaMemPrefetch, other cudaMem* APIs (not supported)

https://github.com/ROCm-Developer-Tools/HIP/blob/master/docs/markdown/hip_faq.md#what-specific-version-ofcuda-does-hip-support



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