

# Theory of Computation

G-4

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## ☰ Is it Regular? Context Free?...

$L = \{w \in \{0,1\}^* \mid w = \text{reverse}(w) \text{ and the length of } w \text{ is divisible by } 4\}$

- if you think its regular, try constructing a DFA/NFA..
- if you think its context free, try constructing its grammar/NPDA..

## Random exercises

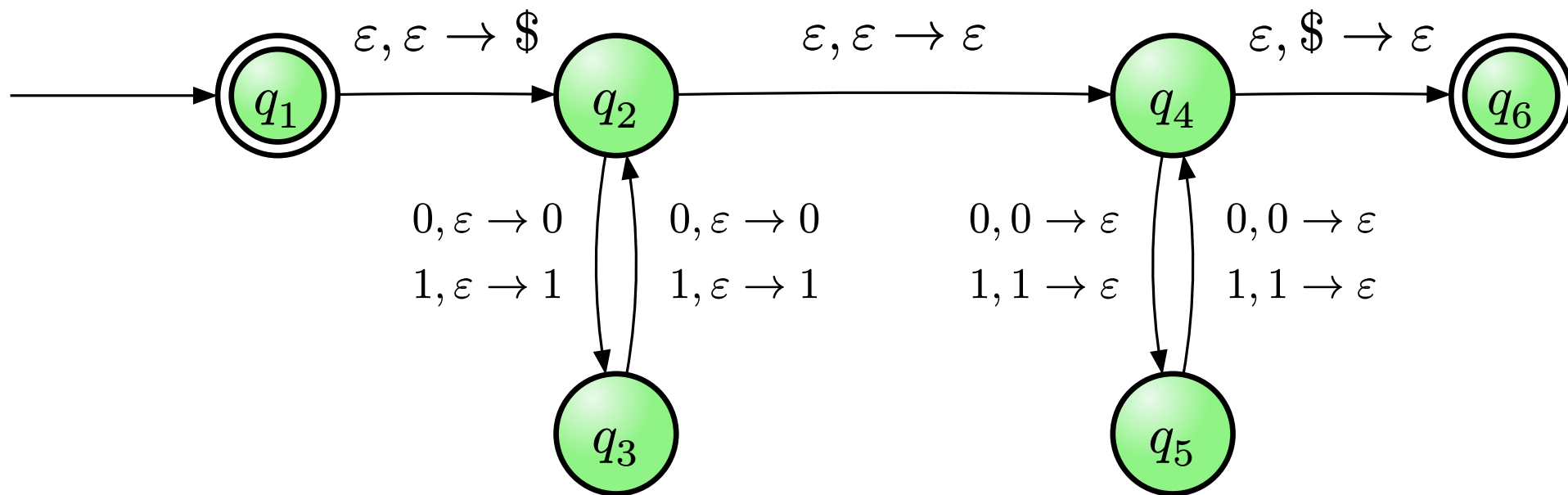
**Answer:**  $L$  is CFL,

$$G = (N, T, P, S), \quad N = \{S\}, \quad T = \{0, 1\}, \quad (1)$$

starting variable  $S$ , and rules:

$$P = \{S \rightarrow 00S00 \mid 01S10 \mid 10S01 \mid 11S11 \mid \varepsilon\} \quad (2)$$

# Random exercises



Is it regular?

# Random exercises

Lets prove that  $L$  is not regular by contradiction, suppose its regular, let  $p$  be the pumping length.. consider string  $s = 0^p 1^{2p} 0^p \in L$   $|s| = 4p > p$  so the conclusions of the pumping lemma must hold, thus we can split  $s = xyz$  (1)  $xy^i z \in L \forall i \geq 0$ , (2)  $|y| > 0$  and (3)  $|xy| \leq p$ . because all of the first  $p$  symbols of  $s$  are 0s, (3) implies that  $x$  and  $y$  must only consist of 0s. Also  $z$  must consist of the rest of the 0s at the beginning, followed by  $1^{2p} 0^p$ .

$x = 0^j, y = 0^k, z = 0^m 1^{2p} 0^p$  where  $j + k + m = p$ .. (2) implies that  $k > 0$  so by (1)  $xyyz$  must belong to  $L$ .

$$xyyz = 0^j 0^k 0^k 0^m 1^{2p} 0^p = 0^{p+k} 1^{2p} 0^p \notin L \quad \text{contradiction} \quad (3)$$

# More exercises

## Task 1

Can a transition function of a dfa be bijective?  
Explain..

## More exercises

$$\delta : Z \times A \rightarrow Z \quad (4)$$

For a bijective function between two sets, the domain and codomain must have the same size(cardinality)

$$|Z \times A| = |Z| \quad \rightarrow \quad |Z| \times |A| = |Z| \quad (5)$$

$$\implies |A| = 1. \quad (6)$$

So yes, transition function  $\delta$  of a dfa can be bijective, if alphabet contains just one symbol..

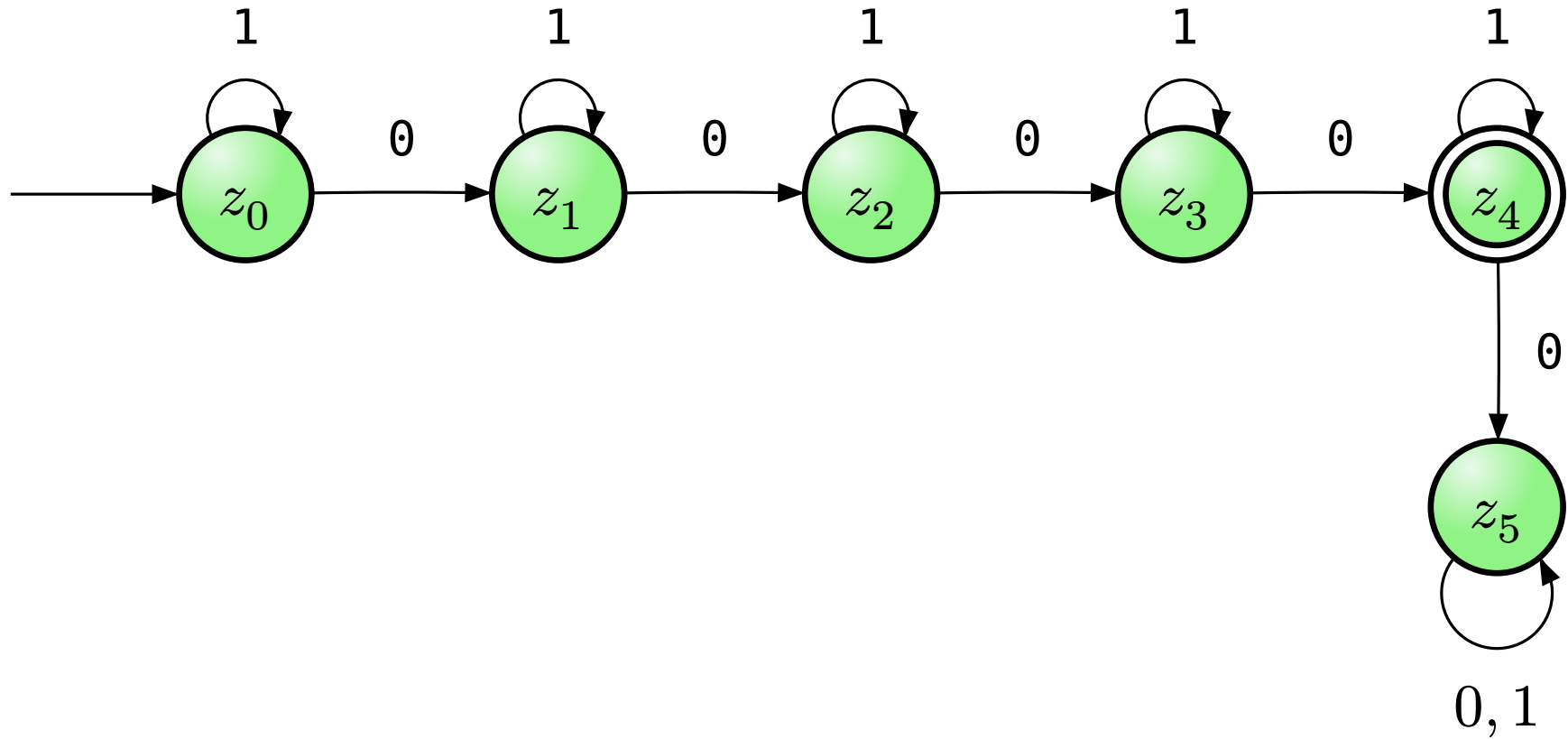


# More exercises

## ☰ Task 2

Sketch a dfa that accepts binary strings that have exactly 4 zeros

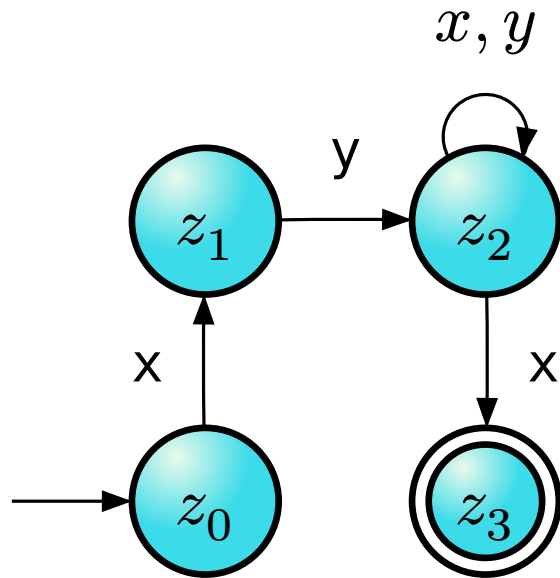
# More exercises



# More exercises

## ☰ Task 3

Give an example of an accepting computation (by writing down the sequence of configurations)



## More exercises

Accepting string can be for example  $xyx$

Configurations:

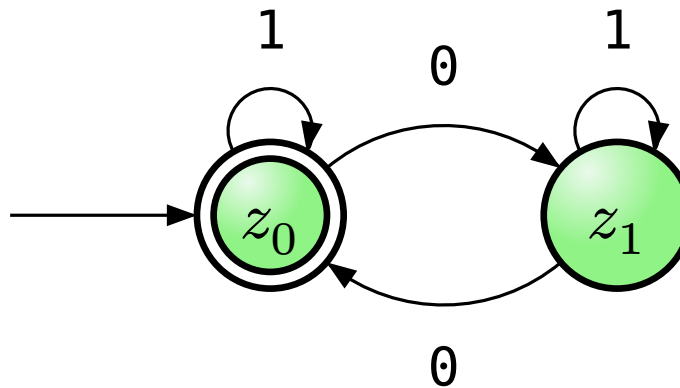
$$(z_0, xyx) \rightarrow (z_1, yx) \rightarrow (z_2, x) \rightarrow (z_3, \varepsilon)$$

thats it...

# More exercises

## ☰ Task 4

Describe the language accepted by DFA



# More exercises

its easy to see...

$$1^*(0(1^*)0)^*1^* \quad (7)$$

# More exercises

## ☰ Task 5

Say  $k = (z_4, 01100, s0a0)$  is the configuration of NPDA, and  $(z_1, \text{push } a) \in \delta(z_4, 0, \varepsilon)$ . Write down a possible successor configuration of  $k$ .

## More exercises

$$\begin{aligned} k &= (z_4, 01100, s0a0) \\ \delta(z_4, 0, \varepsilon) &\rightarrow (z_1, \text{push } a) \end{aligned} \tag{8}$$

means that if we are in the state  $z_4$  and read 0, we push  $a$  on top of the stack, no matter what is on top..  
successor configuration would look like this:

$$k_{\text{succ}} = (z_1, 1100, as0a0) \tag{9}$$



# More exercises

## ☰ Task 6

Let  $A, B$  be two alphabets. Does  $A^* \cup B^* = (A \cup B)^*$ ?

# More exercises

$$A^* \cup B^* = (A \cup B)^*?$$

Lets consider two alphabets:  $A = \{a\}$ ,  $B = \{b\}$

$$A^* = \{\varepsilon, a, aa, \dots\} \quad B^* = \{\varepsilon, b, bb, \dots\} \quad A^* \cup B^* = \{\varepsilon, a, b, aa, bb, \dots\}$$

Note that the string  $ab$  is in  $(A \cup B)^*$  but not in  $A^* \cup B^*$

$$A^* \cup B^* \neq (A \cup B)^*$$

$$A^* \cup B^* \subseteq (A \cup B)^*$$

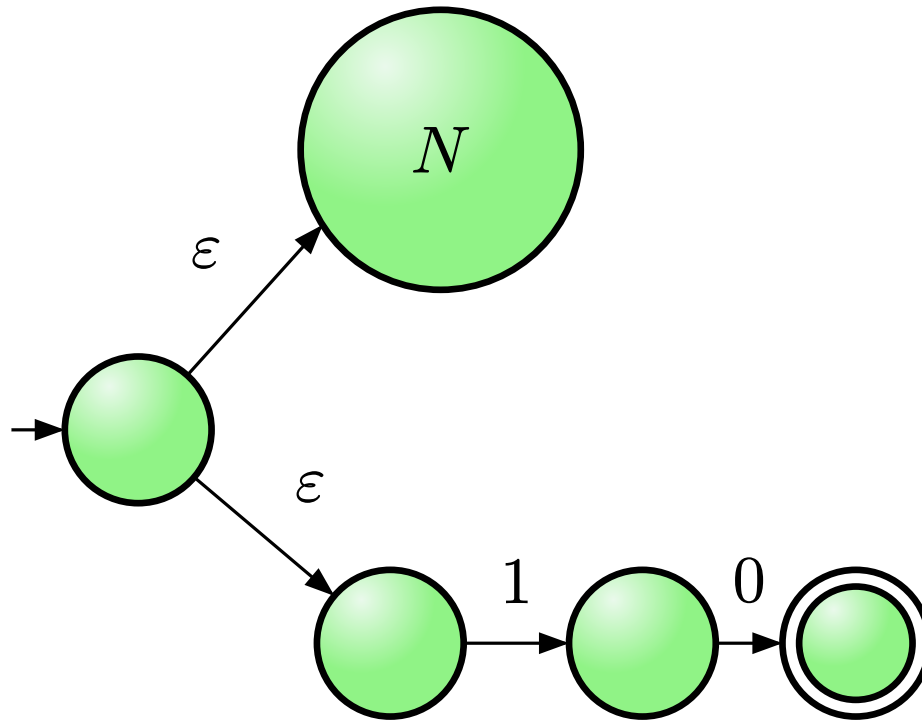
# More exercises

## ☰ Task 7

Say  $N$  is a given nfa that accepts the language  $L$  over the alphabet  $A = \{0, 1\}$ . Construct an nfa that accepts  $L \cup \{10\}$

# More exercises

$N$  is our given nfa, to construct union of two nfes we can just use  $\epsilon$  moves like that:



# More exercises

## ☰ Task 8

Prove that the set of integers

$$\mathbb{Z} = \{\dots - 2, -1, 0, 1, 2, 3\dots\} \quad (10)$$

is countable

# More exercises

Lets try to come up with a bijective function  $f: \mathbb{N} \rightarrow \mathbb{Z}$

$$f(n) = \begin{cases} \frac{n}{2} & \text{n is even} \\ -\frac{n+1}{2} & \text{n is odd} \end{cases} \quad (11)$$

$$0 \rightarrow 0, 1 \rightarrow -1, 2 \rightarrow 1, 3 \rightarrow -2, 4 \rightarrow 2 \text{ and so on ...} \quad (12)$$

# More exercises

## Task 9

Is every countably infinite language regular?  
(provide proof for your answer)

## More exercises

$$\{0^n 1^n \mid n \in \mathbb{N}_n\} \tag{13}$$

Is countable and not regular by pumping lemma.

The bijection is  $n \mapsto 0^n 1^n$ , thats it..



# More exercises

## ☰ Task 10

Find a regular expression which represents the set of strings over  $\{a,b\}$  which contain the two substrings **aa** and **bb**

## More exercises

$$(a \cup b)^* ((aa(a \cup b)^*bb) \cup (bb(a \cup b)^*aa))(a \cup b)^*$$

that might not be intuitive at first but...

(Which of the following L is regular)

$$\{a^n b^m \mid n \geq m \vee n \leq m\}$$

(Which of the following L is regular)

$$\{a^n b^m \mid n \geq m \vee n \leq m\}$$

is the same as  $a^*b^*$  so its regular..

(Which of the following L is regular)

$$\{a^n b^m \mid n > m \wedge n < m\}$$

(Which of the following L is regular)

$$\{a^n b^m \mid n > m \wedge n < m\}$$

is empty  $\implies$  regular

(Which of the following L is regular)

$$\{a^n b^m \mid n \geq m \wedge n \leq m\}$$

(Which of the following L is regular)

$$\{a^n b^m \mid n \geq m \wedge n \leq m\}$$

is the same as  $a^n b^n \Rightarrow$  not regular (already seen why)



(Which of the following L is regular)

$$\{a^n b^m \mid n > m \vee n < m\}$$

(Which of the following L is regular)

$$\{a^n b^m \mid n > m \vee n < m\}$$

It is the same language as  $\{a^n b^m \mid n \neq m\}$   
which is complement to  $\{a^n b^m \mid n = m\}$ .

Since regularity is **closed under complementation**, the considered language can not be regular.

(cuz it would imply that  $\{a^n b^m \mid n = m\}$  is regular)

# More exercises

## ☰ Task 999

Define a primitive recursive function  $f : \mathbb{N} \rightarrow \mathbb{N}$  that counts the number of occurrences of the digit 5 in a natural number

# More exercises

`</>` Code

```
count = 0
while n > 0:
    last_digit = n % 10
    if last_digit == 5:
        count += 1
    n //= 10
print(count)
```

# More exercises

At first we need some auxiliary primitive recursive functions

## More exercises

- quotient  $\left\lfloor \frac{x}{y} \right\rfloor : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$

$$quo(x, y) = \min\{m : m \leq x \wedge (m + 1)y > x\}$$

$$quo(5, 2) = \dots$$

for  $m = 0, 0 \leq 5, 1 \cdot 2 > 5$  ? no

for  $m = 1, 1 \leq 5, 2 \cdot 2 > 5$  ? no

for  $m = 2, 2 \leq 5, 3 \cdot 2 > 5$  ? yes

outputs  $quo(5, 2) = 2$

## More exercises

- remainder  $rem(x, y) : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$

$$rem(x, y) = x \dot{-} (y \cdot quo(x, y))$$

$$\begin{aligned} rem(17, 5) &= 17 \dot{-} (5 \cdot quo(17, 5)) \\ &= 17 \dot{-} 5 \cdot 3 \\ &= 2 \end{aligned}$$

## More exercises

- $\text{length}(\text{number of digits}) : \mathbb{N} \rightarrow \mathbb{N}$   
 $\text{length}(x) = \min\{j : j \leq x, 10^{j+1} > x\} + 1$   
 $\text{length}(12) = \dots$

$$j = 0, 0 \leq 12 \text{ but } 10^1 > 12?$$

$$j = 1, 1 \leq 12, \text{ and } 10^2 > 12$$

so  $\text{length}(12)$  outputs simply 2



## More exercises

$f : \mathbb{N} \rightarrow \mathbb{N}$  is then defined as:

$$f(m) = \sum_{i=1}^{length(m)} eq(5, rem(quo(m, 10^{i-1}), 10))$$

## More exercises

Example:

$$\begin{aligned} f(253) &= eq(5, rem(quo(253, 1), 10)) \\ &\quad + eq(5, rem(quo(253, 10), 10)) \\ &\quad + eq(5, rem(quo(253, 100), 10)) \\ &= 0 + 1 + 0 = 1 \end{aligned}$$

# More exercises

Useful book to understand  $\mu$  recursive functions

click it...

Good luck on midterm..

