KUTAISI INTERNATIONAL UNIVERSITY

Exercises for TTF

Introduction to Theory of Computation Summer semester 2025

Exercises below are your homework; after submission, they will also be discussed during exercise classes.

Week 13

- 1. Give an example of a space bound Turing machine that does not halt.
- 2. Show that $\{0^n 1^n \mid n \in \mathbb{N}_0\} \in \text{TIME}_2(n) \cap \text{SPACE}_2(n)$.
- 3. Prove Lemma 1 and Lemma 2 from the lecture on Deterministic Complexity Classes; that is, prove

Lemma 1. If L is decided by a t(n)-time bounded k-tape Turing machine, then L is decided by state by a t(n) + O(1)-time bounded k-tape Turing machine.

Lemma 2. If L is decided by state by a t(n)-time bounded k-tape Turing machine, then L is decided by a O((t(n))-time bounded k-tape Turing machine.

4. Let $x \notin \mathbb{B}$ and

$$L = \{wx^n w \mid w \in \mathbb{B}^n, n \in \mathbb{N}_0\}.$$

Show that $L \in \text{TIME}_2(n)$ and $L \in \text{TIME}_1(n^2)$.

- 5. Sketch the proof of Lemma 12 from the lecture on Deterministic Complexity Classes; that is, sketch the argument for
- **Lemma 12.** There is a universal 2-tape Turing machine U which simulates t steps of 1-tape Turing machine M_u started with input v of length |v| = n in time $O(|u| \cdot (n+t))$. If M_u uses space s then U uses space $O(|u| \cdot s)$.
- 6. Read and understand the proof of Space Hierarchy Theorem (including the necessary Lemma).
- 7. Read and understand the construction of the decidable non-time constructable function f(n) > n.
- 8. Try: conclude the tighter time-hierarchy theorem then Lemma 17 from the lecture on Deterministic Complexity Classes. Recall, (1 p)

Lemma 17. Let t be time constructible. Then

$$TIME(t(n)) \subseteq TIME_2(t(n)\log(t(n)).$$