

# Evaluating QPU Performance with the Ising Model

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## Aim

To use state of the art classical computational solving techniques to evaluate Ising model instances, evaluate the performance, and compare against D-Wave's 2000Q.

## Claim

There exists at least one Ising model optimization problem which is evaluated with a smaller time-to-solution metric on D-Wave's 2000Q than with an optimized, classical simulated annealer.

# Optimization Problems

## Optimization

An optimization problem seeks to extremize a given function.

## Ising Problem

$$H = - \sum_{i=1}^N B_i s_i - \sum_i^N \sum_{j \in \text{nn}(i)} J_{ij} s_i s_j$$

Which vector  $\vec{S}$  (with each  $s_i \in \{-1, +1\}$ ) minimizes the Ising Hamiltonian  $H$ , given some input sets:

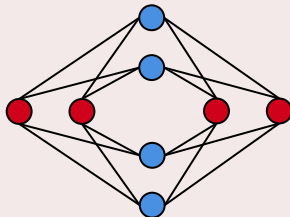
$$\vec{B} = \{B_1, B_2, \dots, B_N\}; \quad \vec{J} = \begin{pmatrix} J_{1,1} & J_{1,2} & \cdots & J_{1,N} \\ J_{2,1} & J_{2,2} & \cdots & J_{2,N} \\ \vdots & \vdots & \ddots & \vdots \\ J_{N,1} & J_{N,2} & \cdots & J_{N,N} \end{pmatrix}$$

# The Chimera Unit Cell

## Bipartite Cell

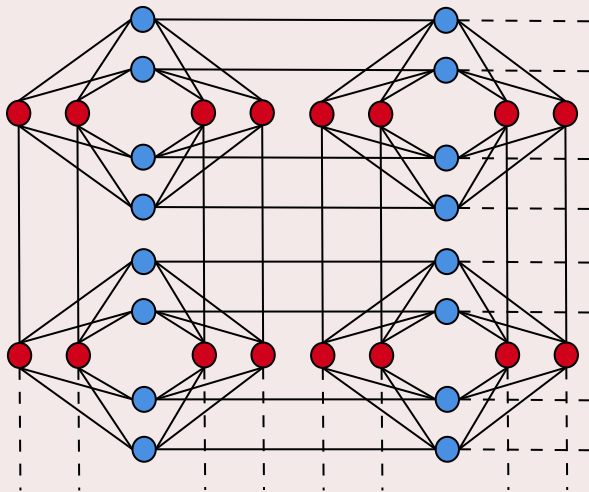
A graph with a set of vertices  $V$  and a set of non-zero edges  $E$  is **bipartite** if its vertices can be partitioned into two disjoint subsets  $A$  and  $B$  such that every edge  $e_i \in E$  connects one vertex from set  $A$  to one vertex in set  $B$ .

## One Chimera Unit Cell



# The Chimera Graph

## Several Interconnected Chimera Unit Cells



# From Ising Model to Chimera Graph

## Embedding Ising Problems onto the Chimera Graph

Features of the Ising problem need to be embedded onto the Chimera graph, solved, and then interpreted.

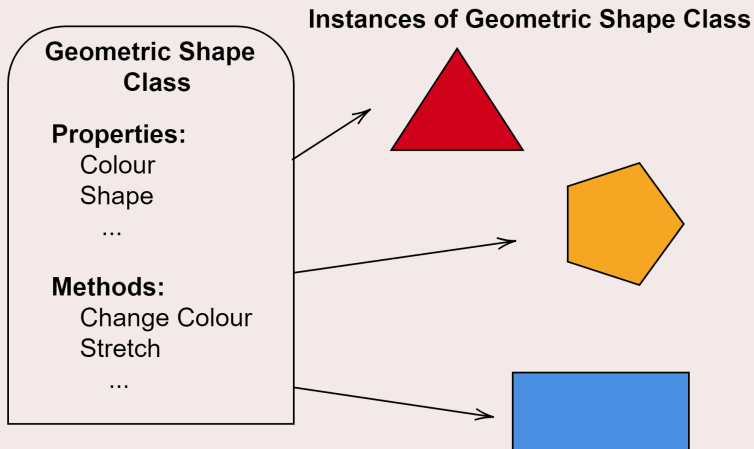
- The set of vertices (qubits)  $V$  are binary and therefore should correspond to the set of spins  $\vec{S}$ .
- The weight associated with each vertex should correspond to the magnetic field strength  $\vec{B}$  at the location each spin.
- The set of edges  $E$  correspond to the set of interspin coupling strengths  $J$ .

## Hardware Limitations

Only problems which can be embedded onto the Chimera graph represent feasible problems.

# Instance Definition

## Object-Oriented Programming Analogy



# Example Problem

## Embedding this Problem

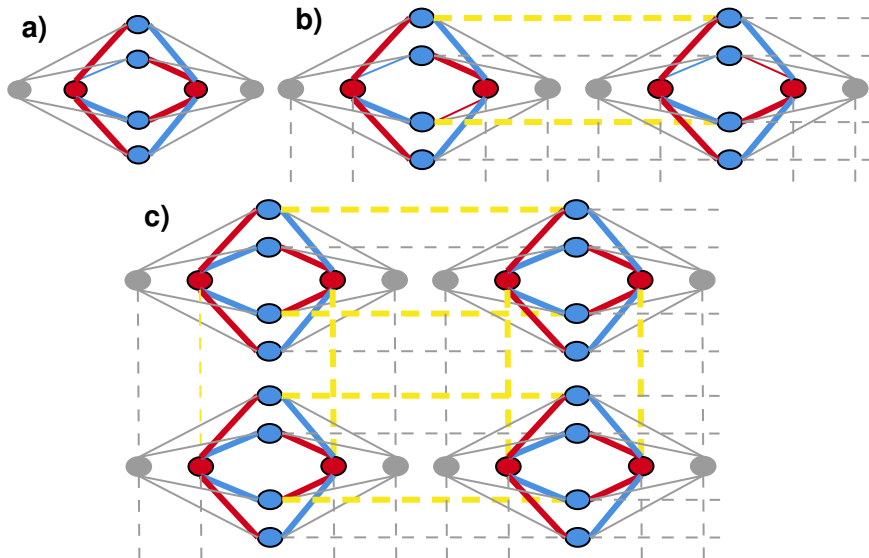
Consider the following problem:

$$\vec{B} = \{2, 2, 2, 2, 0, -2, -2, 0\};$$

$$\vec{J} = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & -2 & 2 & 0 \\ 0 & 0 & 0 & 0 & 0 & 2 & -2 & 0 \\ 0 & 0 & 0 & 0 & 0 & 2 & -2 & 0 \\ 0 & 0 & 0 & 0 & 0 & -2 & 2 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ -2 & 2 & 2 & -2 & 0 & 0 & 0 & 0 \\ 2 & -2 & -2 & 2 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

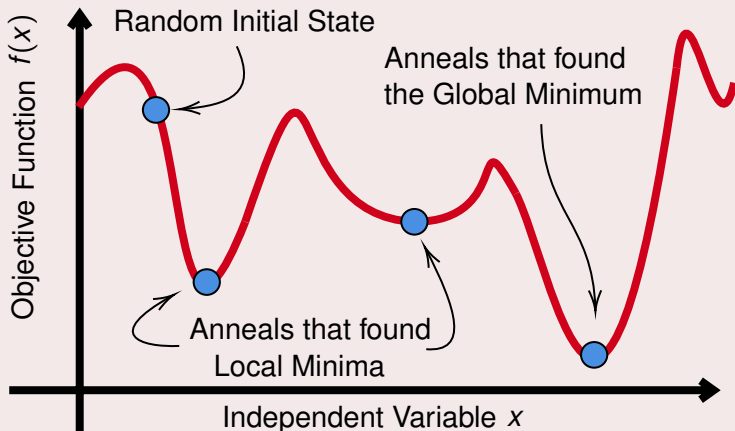


# Example Problem



# Simulated Annealing

## Annealers Don't Always find the Global Minimum



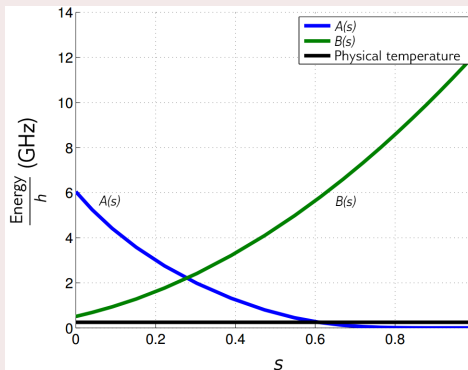
# D-Wave's Quantum Processing Unit (QPU)

## D-Wave's 2000Q

The QPU inside of the 2000Q is described by the Hamiltonian:

$$H_{\text{Ising}} = \underbrace{-\frac{A(s)}{2} \sum_i \hat{\sigma}_x^{(i)}}_{\text{initial Hamiltonian}} + \underbrace{\frac{B(s)}{2} \left( \sum_i h_i \hat{\sigma}_z^{(i)} + \sum_{i>j} J_{i,j} \hat{\sigma}_z^{(i)} \hat{\sigma}_z^{(j)} \right)}_{\text{final Hamiltonian}}$$

## Annealing Schedule



# Performance Metric

## Success probability $P_s$

The probability of finding the global minimum in a single annealing run of time  $t_A$ .

## Time-to-Solution

The time-to-solution (TTS) required to achieve a success probability of  $P_s$  at least once with a desired probability  $P$  is found using:

$$\text{TTS} = t_A \frac{\log(1 - P)}{\log(1 - P_s)}$$

# Problem Definition

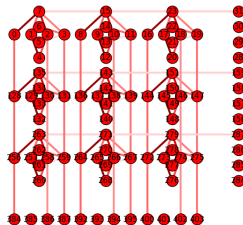
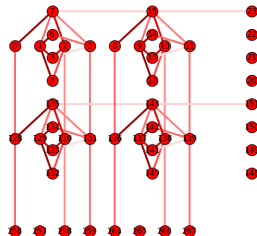
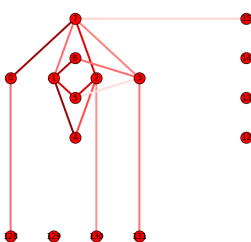
## Problem Definition

Consider the problem described by:

$$\vec{B} = \{0.5143, 0.9702, 0.7975, 0.5248, 0.5397, 0.1735, 0.6969, 0.6760\}$$

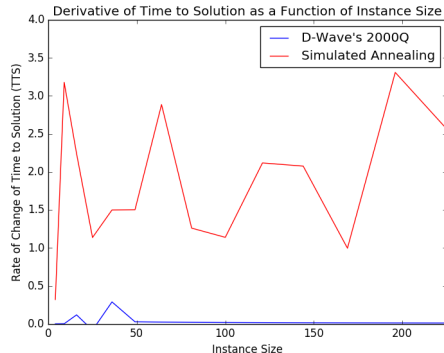
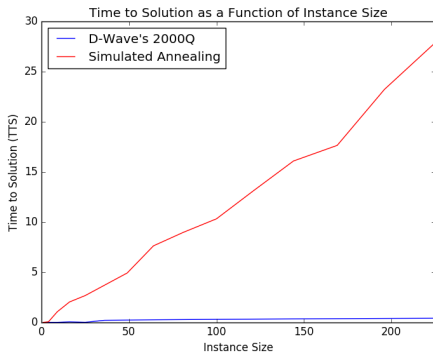
$$\vec{J} = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0.909 \\ 0 & 0 & 0 & 0 & 0.917 & 0.623 & 0.832 & 0.292 \\ 0 & 0 & 0 & 0 & 0.361 & 0.650 & 0 & 0.603 \\ 0 & 0 & 0 & 0 & 0 & 0.063 & 0.291 & 0.255 \\ 0 & 0.917 & 0.361 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0.623 & 0.650 & 0.063 & 0 & 0 & 0 & 0 \\ 0 & 0.832 & 0 & 0.291 & 0 & 0 & 0 & 0 \\ 0.909 & 0.292 & 0.603 & 0.255 & 0 & 0 & 0 & 0 \end{pmatrix}$$

# Chimera Graph Representation





# Solver Performance



# Summary

- Minimizing the Ising model's Hamiltonian constitutes an optimization problem.
- The Ising model Hamiltonian can be embedded onto the Chimera graph, which represents the 2000Q's hardware structure.
- The structure of the Chimera graph restricts the problems that can be solved by the 2000Q.
- Chimera problems can be instantiated with an integer multiple size of the smallest instance and grown by embedding the problem again in an adjacent unit of the Chimera graph.
- Both simulated annealing and quantum annealing can be used to solve optimization problems.
- There exists at least one problem which can be solved with a smaller time to solution metric on the 2000Q than with an optimized simulated annealer.