# **DEPARTMENT OF CYBER SECURITY**



## **DESING AND ANALYSIS OF ALGORITHM ASSINGMENT**

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## a. Explanation for the Insertion Sort and Merge Sort Code

## i. <u>Insertion Sorting Code</u>

#### **Programming language:**

C++

```
#include<iostream>
using namespace std;
int main(){
    int A[10]={6,4,8,9,1,5,7,2,3,10},key,j;
    for (int i = 0; i<10; i++){
        key = A[i];
        j = i-1;
        while (j \ge 0 \&\& A[j]>key){
            A[j+1] = A[j];
            j=j-1;
    A[j+1] = key;
    for (int j = 0; j<10; j++){
        cout<<A[j]<<"\n";
```

The program is basically working on the user initialized array.

The working principal of the program is that it is working on the key and the element of an array and their comparison to sort the element.

For loop is being used to iterate over all the elements of the array. In each iteration, A[i] element is being taken as a key and the element behind that key is being taken as j.

While the j does not reach the 0 and we found out the greater element than the key, till then we will quit the loop else the loop will continue.

In the while loop condition, we are checking that if the element at the j = i-1 position is greater than the key, that means that it is already sorted and quit the loop.

In the body of the while loop, when we have found that element less than the key, than put this element in the sorted array. At the j+1 position.

## a. Merge Sorting code

The merge sorting code is working on the divide and conquer technique which is being done on the recursion.

The program has three parts:

- i. The main function
- ii. mergeSort function
- iii. merge function
- i. Main function

```
int main() {
  int arr[] = {6, 5, 12, 10, 9, 1};
  int size = sizeof(arr) / sizeof(arr[0]);

mergeSort(arr, 0, size - 1);
  cout << "Sorted array: \n";

for (int i = 0; i < size; i++)
  cout << arr[i] << " "<<endl;
}</pre>
```

The main function has the initialized array and it is calculating the size of the given array.

Then it is calling the mergeSort function will run the recursion and to divide the arrays to the pair to twos.

At-last the sorted array will be displayed to the user.

#### ii. MergeSort Function

```
void mergeSort(int arr[], int 1, int r) {
   if (1 < r) {

   // m is the point where the array is divided
     int m = 1 + (r - 1) / 2;

     mergeSort(arr, 1, m);
     mergeSort(arr, m + 1, r);

   // Merge the sorted subarrays
   merge(arr, 1, m, r);
}
</pre>
```

This function is taking the array and the index of the right and left sub-arrays.

Then mergeSort function is starting the recursion to break the sub-arrays to the sub-arrays and passing the indexes of the sub-arrays to the itself function.

First the left array is being sent to the mergeSort function, which it is dividing it to the more subparts of left and right arrays.

Then the mergeSort function is being run on the right sub-array to divide it to the sub-arrays.

When the both of them have been divided then the code will call the merge function to merge the left and right arrays.

#### iii. Merge function

```
□ void merge(int arr[], int p, int q, int r) {
    int n1 = q - p + 1, n2 = r - q;
    int L[n1], M[n2], i=0, j=0, k=p;
    for (int i = 0; i < n1; i++)
      L[i] = arr[p + i];
    for (int j = 0; j < n2; j++)
      M[j] = arr[q + 1 + j];
    while (i < n1 && j < n2) {
      if (L[i] <= M[j]) {</pre>
        arr[k] = L[i];
        i++;
       } else {
        arr[k] = M[j];
        j++;
      k++;
    while (i < n1) {
      arr[k] = L[i];
      i++;
      k++;
    while (j < n2) {
      arr[k] = M[j];
      j++;
      k++;
```

The function is making a new array for the left sub array and right sub arrays and filling the elements with the for loops.

Then while loop is being run to compare the element of the newly created arrays to sort them in the in the arr array.

The comparison is being done on the newly created left and right arrays L and M. Then the sorted element is being found and that is being put to the arr array. The arr array will be sorted array which will be merge and sorted.

The elements left out in any L or M array will instead to the arr array because the left out elements will be the one which were left out because the loop ended.

# b. Test cases used to check the code

In the test cases I have provided the input to the to the code by initializing the array with the int numbers.

i. <u>First test case for the insertion sorting</u>

**Input**: {6,4,8,9,1,5,7,2,3,10}

#### Output:

```
F:\University Work\5th semester\Analysis of Algorithm\Untitled1.exe

1 2 3 4 5 6 7 8 9 10

Process exited after 0.158 seconds with return value 0

Press any key to continue . . .
```

**Input:** int  $A[10] = \{9,101,88,1,34,11,99,1000,666,456\}$ 

# Output:

```
F:\University Work\5th semester\Analysis of Algorithm\Untitled1.exe

1 9 11 34 88 99 101 456 666 1000
------
Process exited after 0.2194 seconds with return value 0
Press any key to continue . . .
```

#### ii. Insertion Sorting

**Input:** int arr[] =  $\{6, 5, 12, 10, 9, 1, 7, 2, 13, 20\}$ 

#### **Output:**

F:\University Work\5th semester\Analysis of Algorithm\Untitled2.exe

```
Sorted array:
1 2 5 6 7 9 10 12 13 20
------
Process exited after 0.2325 seconds with return value 0
Press any key to continue . . . _
```

## Input:

int arr[] =  $\{200,300,100,400,900,502,444,302,101,999\}$ 

## **Output:**

```
F:\University Work\5th semester\Analysis of Algorithm\Untitled2.exe

Sorted array:
100 101 200 300 302 400 444 502 900 999

------

Process exited after 0.2237 seconds with return value 0

Press any key to continue . . . _
```

## c. Procedure for generation of inputs

#### The method for the input generation is the random input generator function

I am using the random generator in the code to generate the 10 elements and store them inside the array. Here the size of the input being generated is 10 and the numbers are less than the 100.

The code for the random input generator is:

The output of the merge sorting for the 10 random numbers less than 100 being sorted is

```
F:\University Work\5th semester\Analysis of Algorithm\Untitled2.exe

Sorted array:
0 24 34 41 58 62 64 67 69 78

------
Process exited after 0.1785 seconds with return value 0

Press any key to continue . . .
```

Now I am going to test the program on the 50 inputs which are less than 500, then the all the numbers where sorted perfectly

For the **insertion sorting** I am also doing the same and giving the **70 inputs** to the code which will be numbers between 0 to 1000. So the **merge sort result** before and after sorting can be seen is:

I have checked both the algorithms with the input size of 1000 and both the algorithms were performing very well as they were sorting the random inputs into the correct order.

```
----After Sorting------
2 3 3 3 6 7 7 8 8 8 8 9 10 10 11 11 15 15 17 18 18 19 20 21 21 21 22 22 23 24 25 27 28 28 28 30 30 31 31 35 35 35 35 36 37 37 38 38 39 40 40 40 41 41 41 43 44 44 49 50 5
53 53 53 55 58 58 60 60 60 61 64 67 67 70 71 71 72 72 72 72 73 75 75 75 77 80 80 82 82 84 84 85 86 87 87 88 88 90 93 93 97 98 99 101 102 103 105 106 107 108 109 10
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688 689 690 690 692 693 694 694 695 695 696 698 699 700 701 701 702 703 704 704 704 705 705 705 705 706 710 711 712 712 712 713 716 717 718 718 721 722 723 723 724 724 72
726 726 728 729 734 734 734 734 736 737 740 741 741 745 745 748 748 750 752 753 753 754 756 757 757 757 758 758 758 759 759 760 760 762 763 763 763 763 766 767 769 77
771 773 774 775 777 778 778 781 783 783 786 786 787 788 788 789 790 790 796 798 798 798 800 801 802 805 807 808 811 812 813 813 814 815 815 815 818 818 823 824 824 82
825 827 827 829 829 829 831 831 832 832 833 833 833 835 836 838 840 841 842 843 844 844 847 848 850 850 851 851 853 855 855 858 859 861 864 865 866 867 868 869 86
869 869 869 870 870 874 875 875 877 878 881 881 881 882 885 886 887 888 888 890 892 893 893 894 895 896 896 898 900 900 900 900 900 901 902 902 902 902 903 905 909 909 911 912 91
913 913 913 923 924 924 924 924 926 928 929 930 931 932 932 932 934 935 936 937 938 938 938 940 941 941 942 942 943 944 944 944 945 945 945 946 948 948 949 949 951 95
954 954 955 956 958 958 958 959 961 961 962 962 962 962 963 964 966 969 970 971 971 971 972 972 974 974 975 976 977 977 982 985 986 989 990 992 993 993 994 99
996 997 998 999
```

# d. Comparison of the time and memory complexity

## Time Complexity for Insertion Sorting:

**Best Case:** 

The best running time is O(n)

**Average Case:** 

The average running time is  $O(n^2)$ .

**Worst Case:** 

The Best case running time is  $O(n^2)$ .

Time complexity for the merge sorting

**Best Case:** 

The best running time is **O(nlogn)** 

**Average Case:** 

The average running time is **O(nlogn)**.

**Worst Case:** 

The Best case running time is O(nlogn).

## Memory Complexity for insertion sorting

Memory or space complexity for the insertion sorting is O(1). This is because the space required by the insertion sorting is constant and does not depend on the size of an input.

#### Memory Complexity for merge sorting

Memory or space complexity for the merge sorting is O(n). Merger sort being recursive takes up the additional space to store the elements in the memory.

# **Conclusion**

I compared the result time of the both the algorithm by giving both the input size of the 5000 ( five thousand ). The merge sort gave back the result in very short time but the insertion sorting took time. This is because of the reason that the merge sorting takes nlogn and the insertion sort takes n^2 time complexity. With the inputs growing to bigger and bigger, the time to sort is growing for the insertion sort while it does not effects merge sort pretty much. Merge sort takes memory but it does not affect the efficiency of the merge sort algorithm.