Numerical methods for eigenvalue problems

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Introduction

Power Method

Inverse Power Method
Inverse Power Method with Shift
Rayleigh Quotient Iteration

Simultaneous Iteration

QR Algorithm

QR Algorithm with Shift

Summary

Introduction

Linear Algebra:

Definition

$$A\in\mathbb{C}^{n imes n},\quad ext{eigenvector } ec{x}\in\mathbb{C}^n\setminus\{0\}\;,\quad ext{eigenvalue } \lambda\in\mathbb{C}$$
 $Aec{x}=\lambdaec{x}$

Theorem

$$\lambda$$
 eigenvalue of $A \iff P_A(\lambda) := det(\lambda I - A) = 0$

- ► P_A characteristic polynomial
- ▶ $deg(P_A) \ge 5$: no explicit formula \longrightarrow numerical algorithm required
- using $P_A(\lambda)$ numerical instable

Definition

Rayleigh quotient: $\rho_A(\vec{x}) = \frac{\vec{x}^H A \vec{x}}{\vec{x}^H \vec{x}}$

Theorem

 \vec{x} eigenvector of $A \Longrightarrow \rho_A(\vec{x})$ is eigenvalue of A

Definition

- ▶ numbering of eigenvalues: $|\lambda_1| \ge |\lambda_2| \ge ... \ge |\lambda_n|$
- $ightharpoonup \vec{x}_j$ eigenvector for λ_j

Power Method

Given: $A \in \mathbb{C}^{n \times n}, \ \vec{y}_0 \in \mathbb{C}^n \setminus \{0\}, \ \|\vec{y}_0\|_2 = 1$

Assume: $\vec{y}_0^H \vec{x}_1 \neq 0$

Idea:

$$\vec{y}_0 = \sum_{j=1}^n \beta_j \vec{x}_j$$

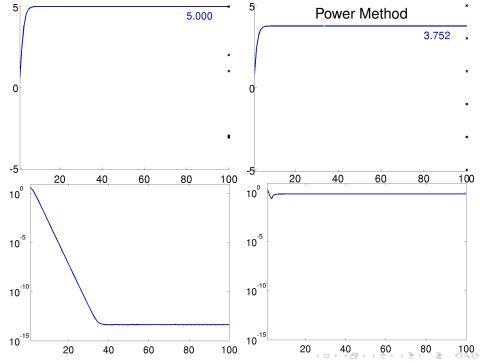
$$\Longrightarrow A\vec{y}_0 = \sum_{j=1}^n \beta_j A\vec{x}_j = \sum_{j=1}^n \beta_j \lambda_j \vec{x}_j$$

$$\Longrightarrow A^k \vec{y}_0 = \sum_{j=1}^n \beta_j \lambda_j^k \vec{x}_j = \beta_1 \lambda_1^k \vec{x}_1 + \mathcal{O}\left(\left|\frac{\lambda_2}{\lambda_1}\right|^k\right)$$

Power Method

Initial guess $\|\vec{y}_0\|_2 = 1$. for $k = 0, 1, \ldots$ do $\vec{z}_{k+1} = A\vec{y}_k$ // power $\rho_k = \vec{y}_k^H \vec{z}_{k+1}$ // Rayleigh quotient $\frac{\vec{y}_k^H A \vec{y}_k}{\vec{y}_k^H \vec{y}_k}$ $\vec{y}_{k+1} = \frac{\vec{z}_{k+1}}{\|\vec{z}_{k+1}\|_2}$ // avoid over and underflow end for

- ▶ Convergence if $\eta := \left| \frac{\lambda_2}{\lambda_1} \right| < 1$ and $\vec{x}_1^H \vec{y}_0 \neq 0$.
- ▶ Approximation $\rho(\vec{y}_k) = \lambda_1 + \mathcal{O}(\eta^k)$, slow if $\eta \approx 1$.
- Finds λ₁ only.



Interior Eigenvalues

Recall: $\lambda \neq 0$ eigenvalue of A $\iff \frac{1}{\lambda}$ eigenvalue of A^{-1} .

Idea: apply power method to $A^{-1} \rightarrow$ inverse power method.

Definition

LU decomposition: A = LU, L lower, U upper triangular matrix

$$A\vec{z} = \vec{y}$$

$$\iff L(U\vec{z}) = \vec{y}$$

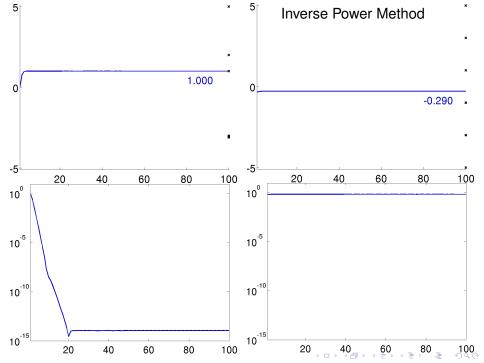
$$\iff L\vec{t} = \vec{y}, \ U\vec{z} = \vec{t}$$

two triangular linear systems

Inverse Power Method

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Initial guess \|\vec{y}_0\|_2 = 1. Calculate LU decomposition of A. for k = 0, 1, \ldots do Solve A\vec{z}_{k+1} = \vec{y}_k with LU decomposition. /\!/ \vec{z}_{k+1} = A^{-1}\vec{y}_k \vec{y}_{k+1} = \frac{\vec{z}_{k+1}}{\|\vec{z}_{k+1}\|_2} // avoid over and underflow \rho_k = \rho_A(\vec{y}_{k+1}) = \vec{y}_{k+1}^H A\vec{y}_{k+1} // Rayleigh quotient end for
```

- Finds eigenvalue closest to zero.
- ► Convergence if $\eta := \left| \frac{\lambda_n}{\lambda_{n-1}} \right| < 1$ and $\vec{x}_n^H \vec{y}_0 \neq 0$.
- ▶ Approximation $\rho(\vec{y}_k) = \lambda_n + \mathcal{O}(\eta^k)$, slow if $\eta \approx 1$.



What about the remaining eigenvalues?

Smallest eigenvalue of $A - \mu I$ is eigenvalue from A closest to μ .

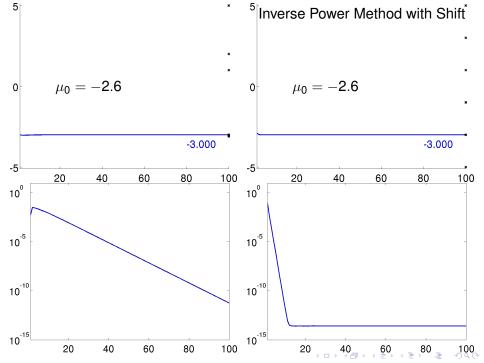
Definition

 μ is called Shift and $A - \mu I$ shifted matrix.

Inverse Power Method with Shift

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Initial guess \|\vec{y}_0\|_2 = 1. \mu_0 initial guess for desired eigenvalue. Calculate LU decomposition of (A-\mu_0I). for k=0,1,\ldots do Solve (A-\mu_0I)\vec{z}_{k+1}=\vec{y}_k with LU decomposition. \vec{y}_{k+1}=\frac{\vec{z}_{k+1}}{\|\vec{z}_{k+1}\|_2} // avoid over and underflow \rho_k=\rho_A(\vec{y}_{k+1})=\vec{y}_{k+1}^HA\vec{y}_{k+1} // Rayleigh quotient end for
```

- ▶ Finds eigenvalue λ_i closest to μ_0 .
- ▶ Convergence if $\eta := \max_{m \neq j} \left| \frac{\lambda_j \mu_0}{\lambda_m \mu_0} \right| < 1$ and $\vec{x}_j^H \vec{y}_0 \neq 0$.
- ▶ Approximation $\rho(\vec{y}_k) = \lambda_i + \mathcal{O}(\eta^k)$, slow if $\eta \approx 1$.



Rayleigh Quotient Iteration

Initial guess $\|\vec{y}_0\|_2 = 1$.

 μ_0 initial guess for desired eigenvalue.

for
$$k = 0, 1, ...$$
 do

Calculate LU decomposition of $(A - \mu_k I)$.

Solve $(A - \mu_k I)\vec{z}_{k+1} = \vec{y}_k$ with LU decomposition.

$$\vec{y}_{k+1} = \frac{\vec{z}_{k+1}}{\|\vec{z}_{k+1}\|_2}$$
 // avoid over and underflow

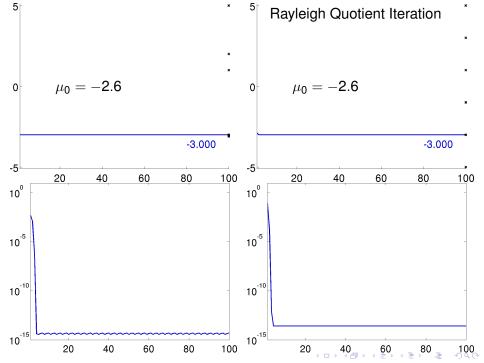
$$\rho_k = \rho_A(\vec{y}_{k+1}) = \vec{y}_{k+1}^H A \vec{y}_{k+1} // \text{Rayleigh quotient}$$

$$\mu_{\mathbf{k}} := \rho_{\mathbf{k}}$$

end for

- ▶ Finds eigenvalue λ_j closest to μ_0 .
- ▶ Convergence if $\eta_0 := \max_{m \neq j} \left| \frac{\lambda_j \mu_0}{\lambda_m \mu_0} \right| < 1$ and $\vec{x}_j^H \vec{y}_0 \neq 0$.
- ▶ Approximation $\rho(\vec{y}_k) = \lambda_j + \mathcal{O}(\eta_0^{2k})$.
- LU decompostion in each step expensive
- Convergence fast



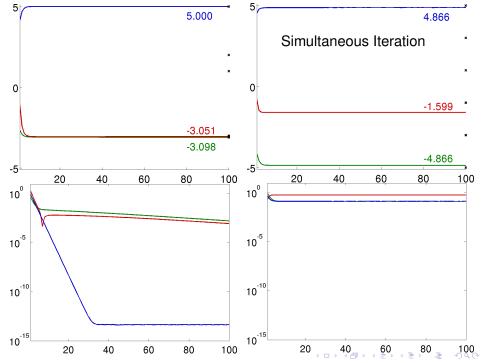


Simultaneous Iteration

Until now: calculating only one eigenvalue Next step: calculating many or all eigenvalues

$$m \leq n,\ U_0 \in \mathbb{C}^{n,m}$$
 unitary, $U_0^H U_0 = I_m,\ A_0 := A$ for $k=0,1,\ldots$ do
$$Y_{k+1} := A_k U_k \text{ // power of each column of } U_k$$
 Calculate QR decomposition $Y_{k+1} = U_{k+1} R_{k+1}.$ $A_{k+1} = U_{k+1}^H A U_{k+1} \text{ // similarity transformation}$ end for

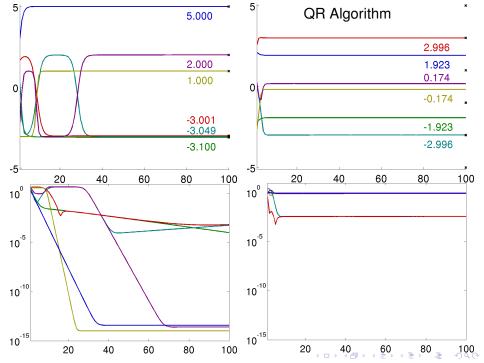
Orthogonalisation is necessary to avoid that all columns converge against x₁.



m = n: QR algorithm (equivalent to simultaneous iteration)

$$A_0:=A$$
 for $k=0,1,\ldots$ do Calculate QR Decomposition $A_k=Q_kR_k$. $A_{k+1}:=R_kQ_k$ end for

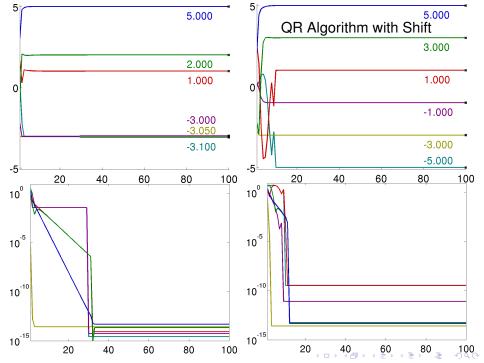
- ▶ $A_{k+1} = R_k Q_k = Q_k^H A_k Q_k = (Q_{k-1} Q_k)^H A_{k-1} (Q_{k-1} Q_k) = \widehat{Q}_k^H A_0 \widehat{Q}_k$, unitary similarity transformation.
- ▶ $A_k \longrightarrow$ upper triangular matrix, eigenvalues on diagonal (proof: schur decomposition).
- This way QR algorithm can be very slow.



QR Algorithm with Shift

```
A_0 := A for k = 0, 1, \ldots do Choose a shift \mu_k. Calculate QR decomposition A_k - \mu_k I = Q_k R_k. A_{k+1} := R_k Q_k + \mu_k I end for
```

this way convergence is fast in general



HENRICH HEINE UNIVERSITÄT QR Implementation

- ▶ complexity of one QR decomposition: $\frac{2}{3}n^3$ operations
- \triangleright complexity to calculate all eigenvalues: at least $\mathcal{O}(n^4)$
- reduce complexity for QR decomposition by $A \leftarrow H := U^H A U$ with H upper Hessenberg.

$$U^{H}AU = H = \begin{bmatrix} h_{1,1} & \dots & h_{1,n-1} & h_{1,n} \\ h_{2,1} & \ddots & \vdots & \vdots \\ & \ddots & h_{n-1,n-1} & h_{n-1,n} \\ 0 & & h_{n,n-1} & h_{n,n} \end{bmatrix}$$

U product of Givens-Rotations

$$G = \begin{bmatrix} \cos(\alpha) & \sin(\alpha) \\ -\sin(\alpha) & \cos(\alpha) \end{bmatrix}$$



- power method for biggest eigenvalue
- inverse power method for eigenvalue closest to zero.
- inverse power method with shift for desired eigenvalue
- Rayleigh quotient iteration, fast
- simultaneous iteration for many eigenvalues
- QR algorithm (with shift) for all eigenvalues (fast)

Thank you for your attention!