

Design and Analysis of Algorithms I

# Graph Algorithms Representing Graphs

#### Graphs

#### Two ingredients

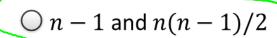
- <u>Vertices</u> aka nodes (V)
- Edges (E) = pairs of vertices
  - can be <u>undirected</u> [unordered pair] or <u>directed</u> [ordered pair] (aka <u>arcs</u>)



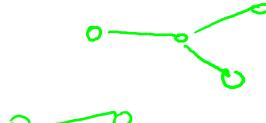


Examples: road networks, the Web, social networks, precedence constraints, etc.

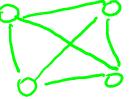
Consider an undirected graph that has n vertices, no parallel edges, and is connected (i.e., "in one piece"). What is the minimum and maximum number of edges that the graph could have, respectively?



- $\bigcirc n-1$  and  $n^2$
- $\bigcirc$  *n* and  $2^n$
- $\bigcirc$  *n* and  $n^n$



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#### Sparse vs. Dense Graphs

Let  $\underline{\mathbf{n}} = \#$  of vertices,  $\underline{\mathbf{m}} = \#$  of edges.

In most (but not all) applications, m is  $\Omega(n)$  and  $O(n^2)$ 

- in a "sparse" graph, m is or is close to O(n)
- in a "dense" graph, m is closer to  $\theta(n^2)$

# The Adjacency Matrix

Represent G by a n x n 0-1 matrix A where  $A_{ij} = 1 \Leftrightarrow G$  has an i-j edge  $\bigcirc$ 

#### **Variants**

- $A_{ij} = \#$  of i-j edges (if parallel edges)
- $A_{ij}$  = weight of i-j edge (if any)

• 
$$A_{ij} = \begin{bmatrix} +1 & \text{if } \bigcirc \bigcirc \bigcirc \bigcirc \\ -1 & \text{if } \bigcirc \bigcirc \bigcirc \bigcirc \end{bmatrix}$$

How much space does an adjacency matrix require, as a function of the number n of vertices and the number m of edges?

- $\bigcirc \theta(n)$
- $\bigcirc \theta(m)$
- $\bigcirc \theta(m+n)$
- $\bigcirc \theta(n^2)$

### **Adjacency Lists**

#### **Ingredients**

- array (or list) of vertices
- array (or list) of edges
- each edge points to its endpoints
- each vertex points to edges incident on it

How much space does an adjacency list representation require, as a function of the number n of vertices and the number m of edges?

- $\bigcirc \theta(n)$
- $\bigcirc \theta(m)$
- $\bigcirc \theta(m+n)$ 
  - $\bigcirc \theta(n^2)$

# **Adjacency Lists**

<u>Ingredients</u>	<u>Space</u>
<ul> <li>array (or list) of vertices</li> </ul>	$\theta(n)$
• array (or list) of edges one-to-one	heta(m)
• each edge points to its endpoints correspondence!	$\theta(m)$
<ul> <li>each vertex points to edges incident on it</li> </ul>	heta(m)
Question: which is better?  Answer: depends on graph density and operations	$\frac{\overline{\theta(m+n)}}{\theta(max\{m,n\})]}$

needed.

This course: focus on adjacency lists.

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