

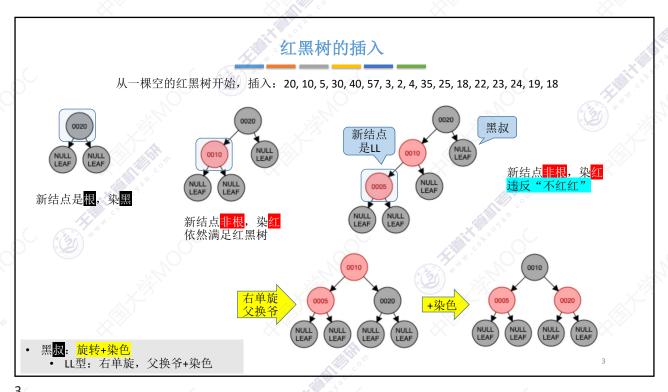
红黑树的插入

从一棵空的红黑树开始,插入: 20, 10, 5, 30, 40, 57, 3, 2, 4, 35, 25, 18, 22, 23, 24, 19, 18

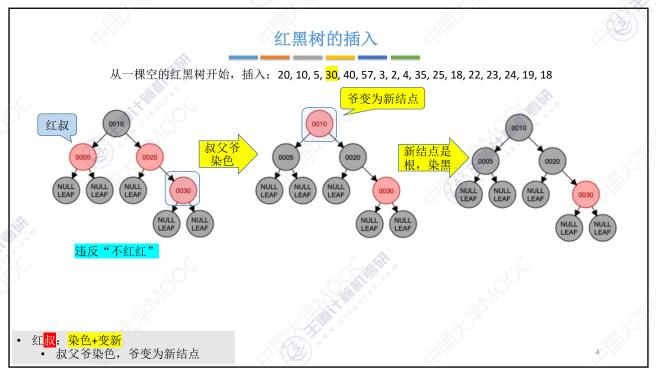
- 先查找,确定插入位置(原理同二叉排序树),插入新结点 <mark>新结点</mark>是<mark>根</mark>——染为<mark>黑色</mark>
- <mark>点非根</mark>——染为红色
 - 若插入新结点后依然满足红黑树定义,则插入结束
 - 若插入<u>新</u>结点后不满足红黑树定义,需要<mark>调整</mark>,使其重新满足红黑树定义
 - 黑叔:旋转+染色
 - LL型: 右单旋, 父换爷+染色
 - RR型: 左单旋, 父换爷+染色
 - LR型: 左、右双旋, 儿换爷+染色
 - RL型: 右、左双旋, 儿换爷+染色
 - 红叔: 染色+变新
 - 叔父爷染色,爷变为新结点

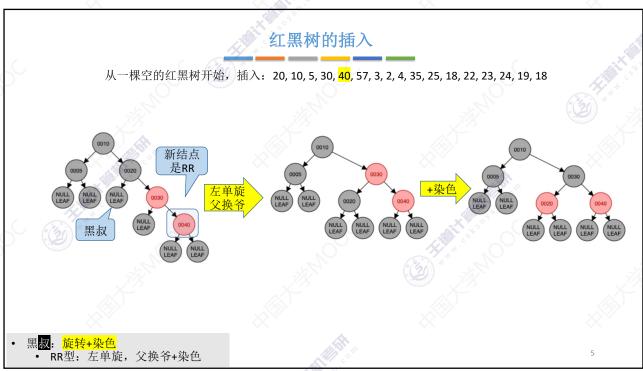
<mark>如何调整</mark>:看新结点<mark>叔叔</mark>的脸色

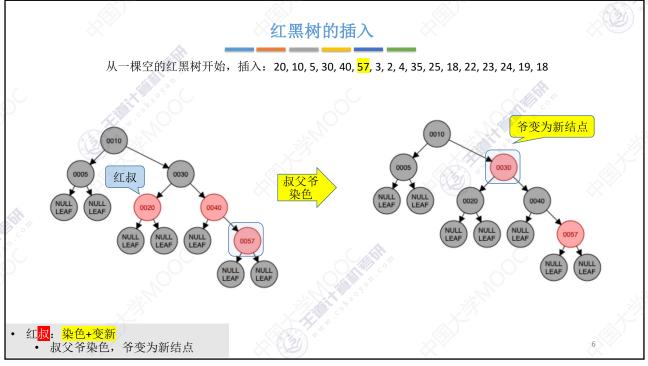


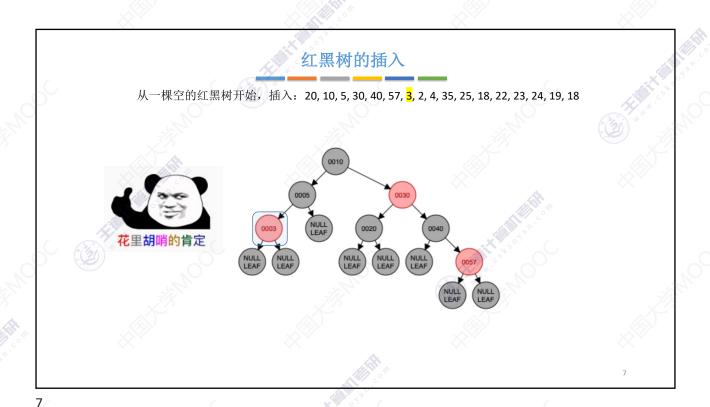


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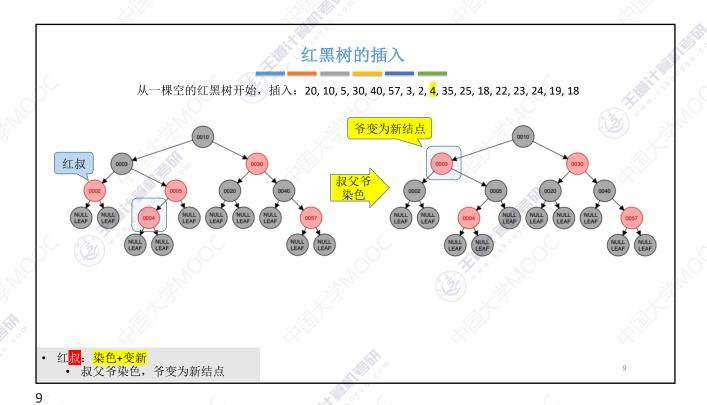




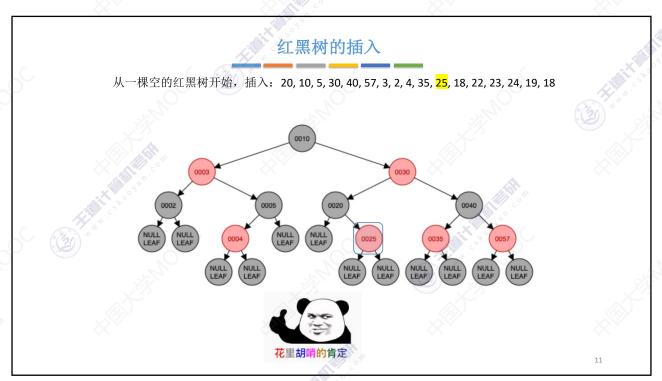


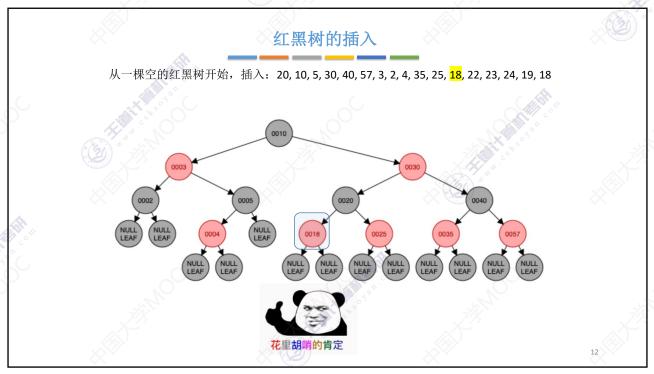


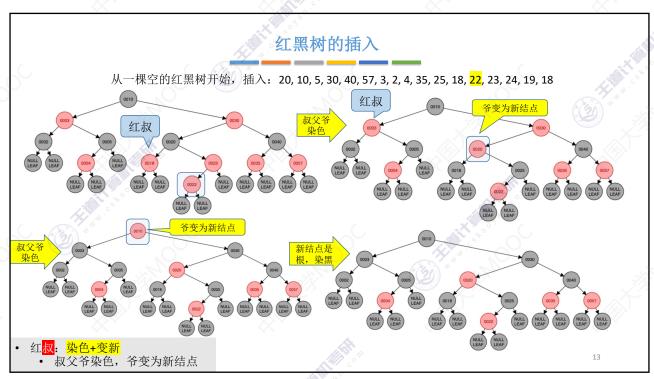
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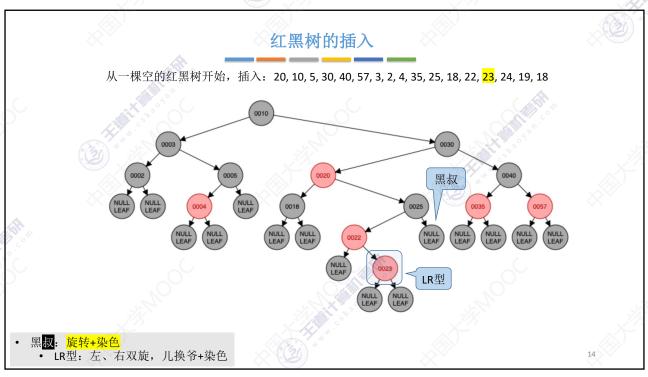


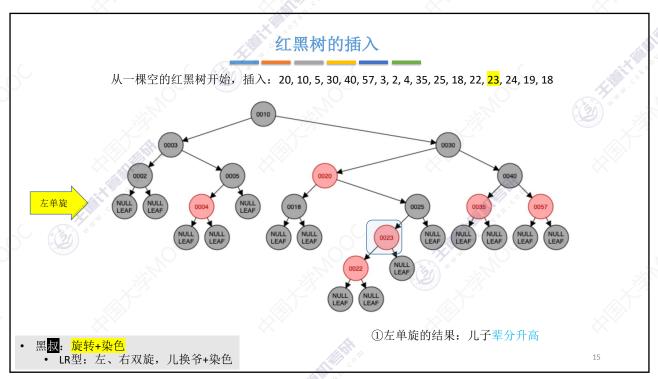
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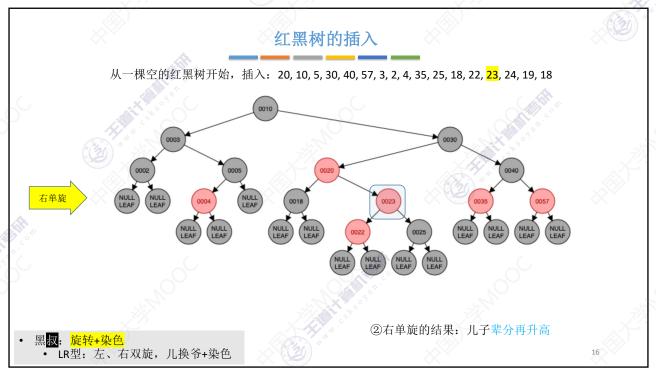


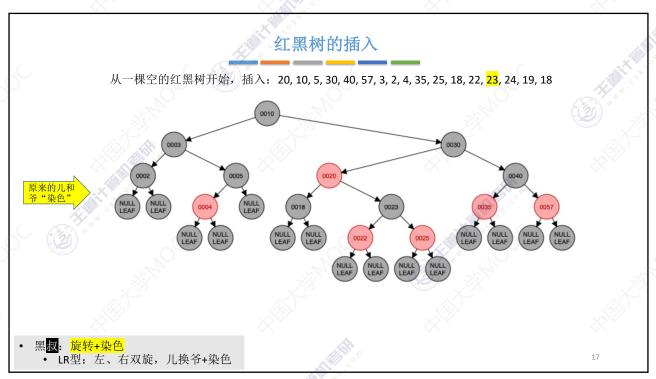


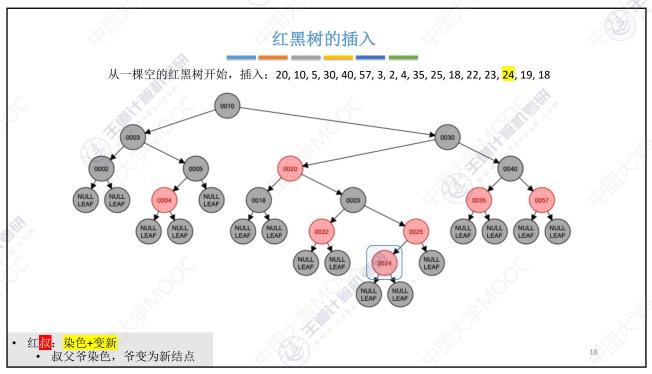


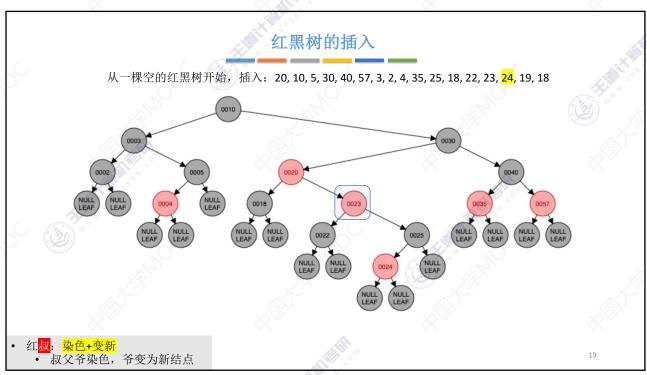


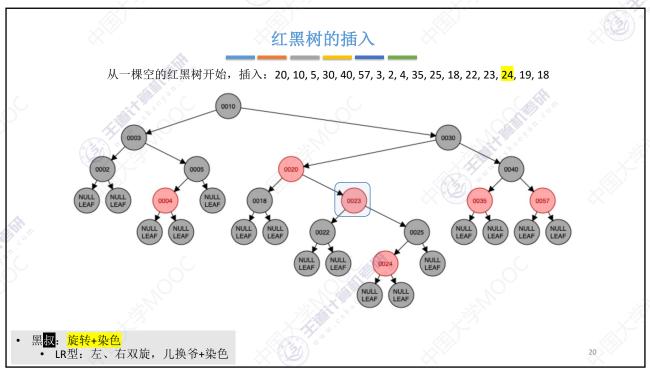


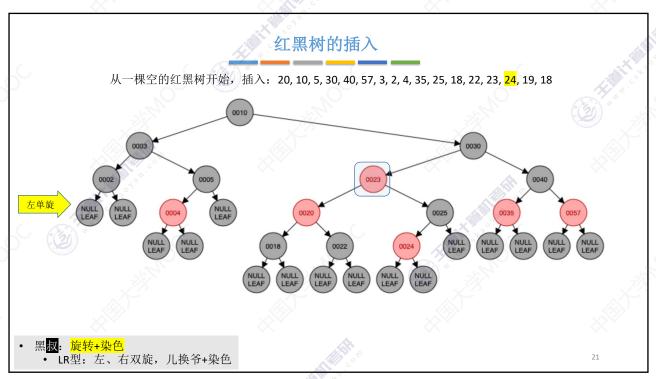


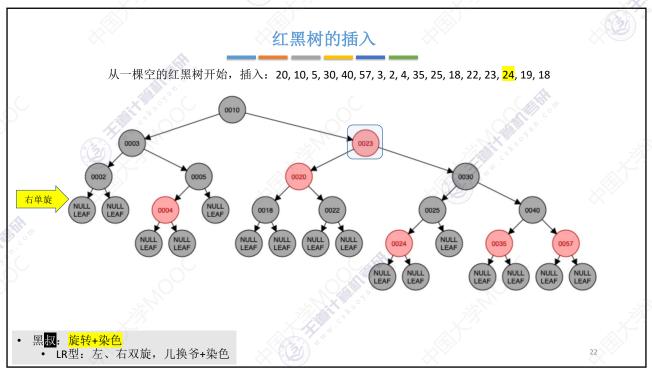


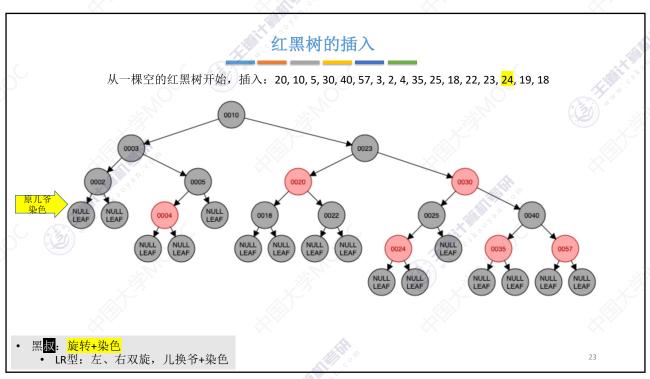


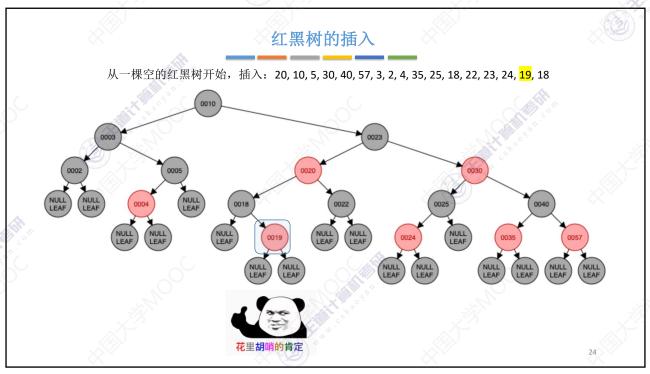


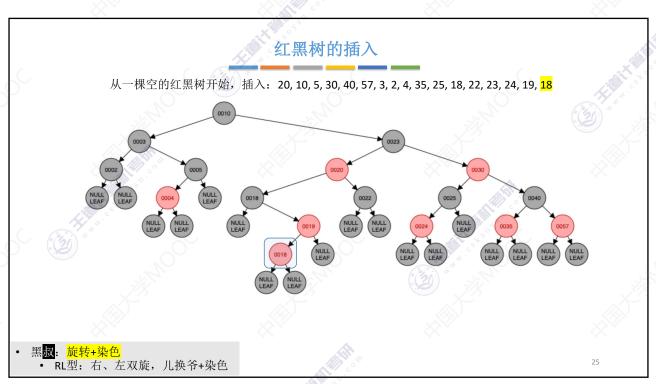


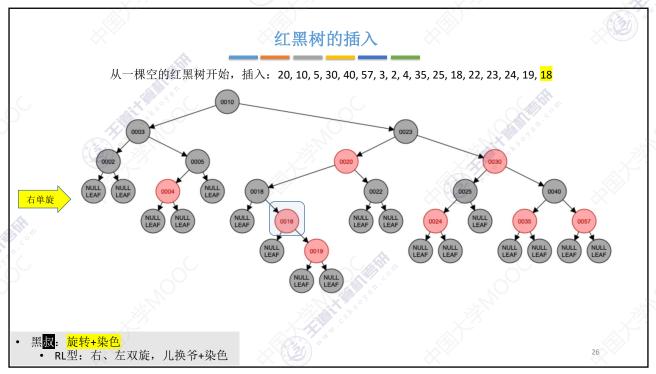


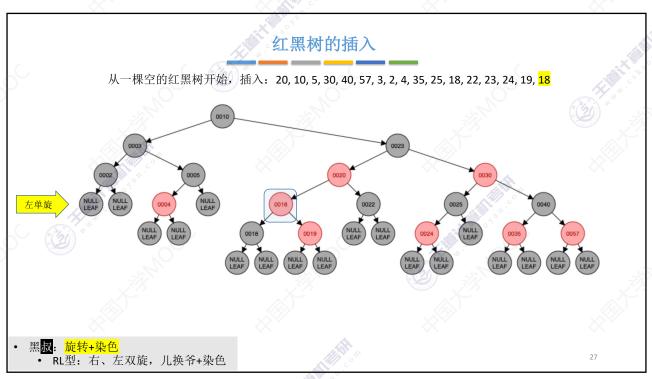


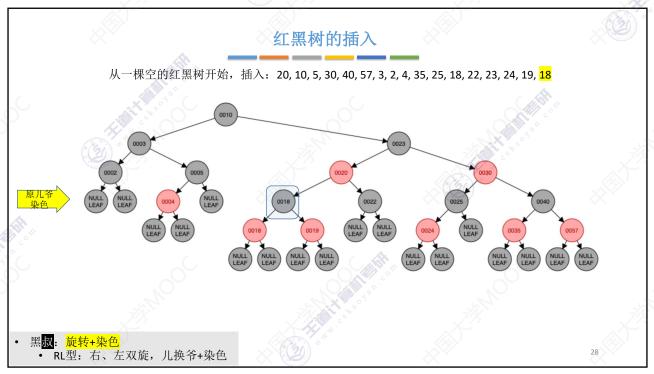


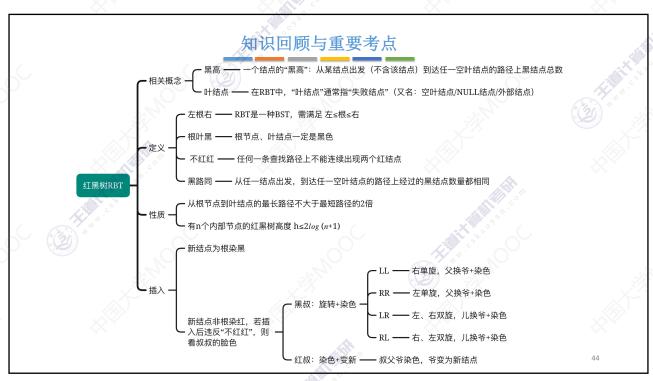


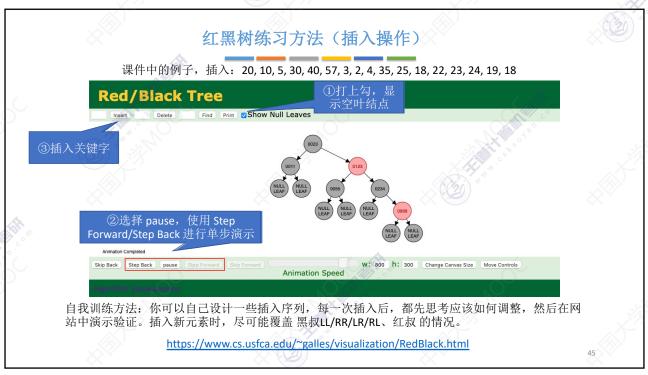


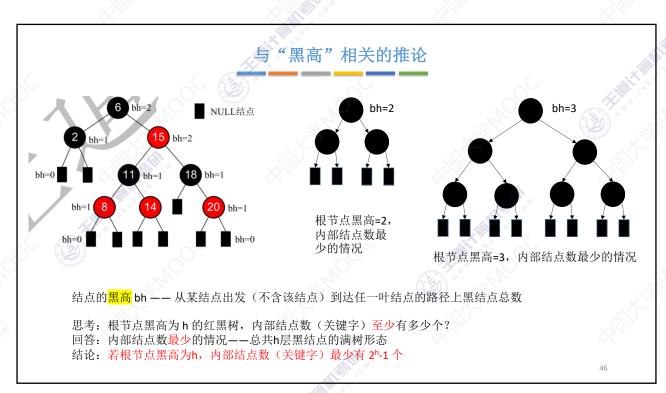












红黑树的定义→性质

红黑树是二叉排序树 ➡ 左子树结点值≤根结点值≤右子树结点值

与普通BST相比,有什么要求 ■

- → ①每个结点或是红色,或是黑色的
- ②根节点是黑色的
 - ③叶结点(外部结点、NULL结点、失败结点)均是黑色的
 - ④不存在两个相邻的红结点(即红结点的父节点和孩子结点均是黑色)
 - ⑤对每个结点,从该节点到任一叶结点的简单路径上,所含黑结点的数 目相同



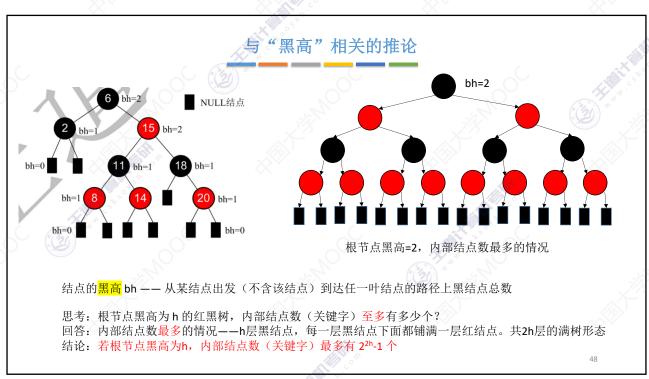
张口就是freestyle

性质1: 从根节点到叶结点的最长路径不大于最短路径的2倍性质2: 有n个内部节点的红黑树高度 $h \le 2log_2(n+1)$

→ 红黑树查找操作时间复杂度 = O(log₂n)

性质1证明:任何一条查找失败路径上黑结点数量都相同,而路径上不能连续出现两个红结点,即红结点只能穿插在各个黑结点中间

性质2证明: 若红黑树总高度=h,则根节点黑高 \geq h/2,因此内部结点数 $n \geq 2^{h/2}$ -1,由此推出 $h \leq 2log_2(n+1)$



/1Ω