

# PROJECT: THE BLAHUT-ARIMOTO ALGORITHM

## I. THE ALGORITHM

Let  $p(y|x)$  be a discrete memoryless channel with input alphabet  $\mathcal{X}$  and output alphabet  $\mathcal{Y}$ . Recall that the capacity of  $p(y|x)$  is given by

$$C = \max_{p(x)} I(X; Y). \quad (1)$$

For any distribution  $p(x)$  on  $\mathcal{X}$  and any conditional distribution  $q(x|y)$  on  $\mathcal{X} \times \mathcal{Y}$ , define

$$f(p(x), q(x|y)) = \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} p(x) p(y|x) \log \frac{q(x|y)}{p(x)}.$$

Prove the following facts:

- 1) The conditional distribution  $q(x|y)$  that maximizes  $f(p(x), q(x|y))$  is given by

$$q^*(x|y) = \frac{p(x)p(y|x)}{\sum_{x' \in \mathcal{X}} p(x')p(y|x')}.$$

- 2) The distribution  $p(x)$  that maximizes  $f(p(x), q(x|y))$  is given by

$$p^*(x) = \frac{\prod_{y \in \mathcal{Y}} q(x|y)^{p(y|x)}}{\sum_{x' \in \mathcal{X}} \prod_{y \in \mathcal{Y}} q(x'|y)^{p(y|x')}}.$$

- 3) The capacity of  $p(y|x)$  can be expressed as

$$C = \max_{p(x)} \max_{q(x|y)} f(p(x), q(x|y)).$$

The above facts naturally suggest the following iterative algorithm (known as the Blahut-Arimoto algorithm) for computing the channel capacity as well as the capacity-achieving input distribution. Let  $p^{(0)}(x) = \frac{1}{|\mathcal{X}|}$  for all  $x \in \mathcal{X}$ . For  $k \geq 0$ , let

$$\begin{aligned} q^{(k)}(x|y) &= \frac{p^{(k)}(x)p(y|x)}{\sum_{x' \in \mathcal{X}} p^{(k)}(x')p(y|x')}, \\ p^{(k+1)}(x) &= \frac{\prod_{y \in \mathcal{Y}} q^{(k)}(x|y)^{p(y|x)}}{\sum_{x' \in \mathcal{X}} \prod_{y \in \mathcal{Y}} q^{(k)}(x'|y)^{p(y|x')}}. \end{aligned}$$

It can be shown that  $p^{(k)}(x)$  converges to a capacity-achieving distribution  $p^\dagger(x)$  (i.e.,  $p^\dagger(x)$  is an optimal solution to the maximization problem in (1)) as  $k \rightarrow \infty$ .

Remark: You are encouraged to consult the relevant literature on the Blahut-Arimoto algorithm.

## II. ADDITIVE GAUSSIAN NOISE CHANNEL WITH PEAK POWER CONSTRAINT

Consider an additive Gaussian noise channel  $Y = X + Z$ , where  $Z \sim \mathcal{N}(0, 1)$  is independent of  $X$ . Use the Blahut-Arimoto algorithm to compute the capacity-achieving distribution of this channel subject to the peak power constraint  $|X| \leq A$  for  $A = 0.1$ ,  $A = 1$ , and  $A = 10$ . Record and explain your findings.

Remark: You may need to discretize the input alphabet and the output alphabet in order to apply the Blahut-Arimoto algorithm.