

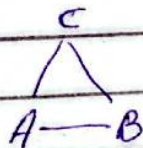
## Graph theory: Basic Definitions and Theorems

→ Def: A graph  $G = (V, E)$  consist of a set  $V$  of vertices (also called nodes) and a set  $E$  of edges.

(2) → If an edge connects to a vertex we say the edge is incident to the vertex and say the vertex is an endpoint edge.

- → incident describes the relationship b/w an edge and a vertex it touches.
- Think of it like a handshake the edge "shakes hands" (is incident) with both of its endpoints vertices.

ex:



ex:  $A-C$  is incident to  $A$  and  $C$ .

(3) If an edge has only one endpoint then it is called a loop edge.

ex  $A \rightarrow A$



(4) If two or more edges have the same endpoints then they are called multiple or parallel edges.



(5) Two vertices that are joined by an edge are called adjacent vertices.

ex:



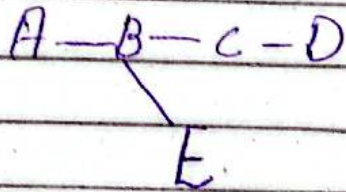
- A and B are adjacent (connected by an edge)

- B and C are adjacent (connected by an edge)

- A and C are not adjacent (no direct edge b/w them)

(6) A pendant vertex is a vertex that is connected to exactly one other vertex by a single edge.

ex:



- D is a pendant vertex (only connected to C)
- E is a pendant vertex (only connected to B)
- A is a pendant vertex (only connected to B)
- B and C are not pendant because they each have multiple connections.



⑦. A walk in a graph is a sequence of alternating vertices and edges  $v_1, v_2, e_2, \dots, v_n, v_{n+1}$  with  $n \geq 0$ . If  $v_i = v_{n+1}$  then the walk is closed. The length of  $P$  is the number of edges in the walk. A walk of length zero is a trivial walk.

-> simple:  
A walk is a seq of alternating vertices (nodes) and edges where you "travel" from one vertex to another by following edges.

ex:  $v_1 - (e_1) \rightarrow v_2 - (e_2) \rightarrow v_3 - (e_3) \rightarrow \dots - (e_n)$   
• start at  $v_1$ , move along edge  $e_1$  to  $v_2$ , then  $e_2$  to  $v_3$ , etc.

length = Number of edges ( $n$ )

• A trivial walk has zero edges (just a single vertex, no movement)



⑧ Trail  $\rightarrow$  No edge repeated  
Path  $\rightarrow$  No vertex repeated  
Circuit  $\rightarrow$  Trail that comes back to start  
Eulerian - covers every edge exactly once.

⑨ A cycle is a non-trivial circuit in which only repeated vertex is the first/last.  
exe

Graph:  $A-B-C-A$

Cycle:  $A \rightarrow B \rightarrow C \rightarrow A$

- Edges not repeated.

- only vertex A is repeated (start = end)

loop

$\rightarrow$  An edge that starts and ends at the same vertex.

exe - An edge from A back to A  
it has only one vertex and one edge.

cycle &

A closed Path with at least 3 vertices.

Ex -  $A \rightarrow B \rightarrow C \rightarrow A$

- only start/end vertex repeats, others are unique.

Simple & Diff

- Loop = 1 vertex, 1 edge (small self-connection)

- cycle = 3 or more vertices connected in a closed chain



→ (10) A simple graph is a graph with no loop edges. Edges in a simple graph may be specified by a set  $\{x, y\}$  of the two vertices that the edge makes adjacent.

A graph with more than one edge b/w a pair of vertices is called a multigraph while a graph with loop edges is called a pseudograph.

(11) A directed graph is a graph in which the edges may only be traversed in one direction.

(12) The degree of a vertex is the number of edges incident to the vertex and is denoted  $\deg(v)$ .

example:  $A-B, A-C, B-C, C-D$   
•  $\deg(A) = 2$  (edges  $A-B, A-C$ )  
•  $\deg(B) = 2$  (edges  $B-A, B-C$ )  
•  $\deg(C) = 3$  (edges  $C-A, C-B, C-D$ )  
•  $\deg(D) = 1$  (edge  $D-C$ )

(13) In a directed graph, the in-degree of a vertex is the number of edges



⇒ 14 A graph is connected if there is a walk  
b/w every pair of distinct vertices in the  
graph  
or

A graph is connected if you can travel (walk)  
from any vertex to any other vertex by following  
the edges.  
ex:

A-B-C-D

15 A subgraph is just a smaller graph  
taken from a bigger graph.

(16)