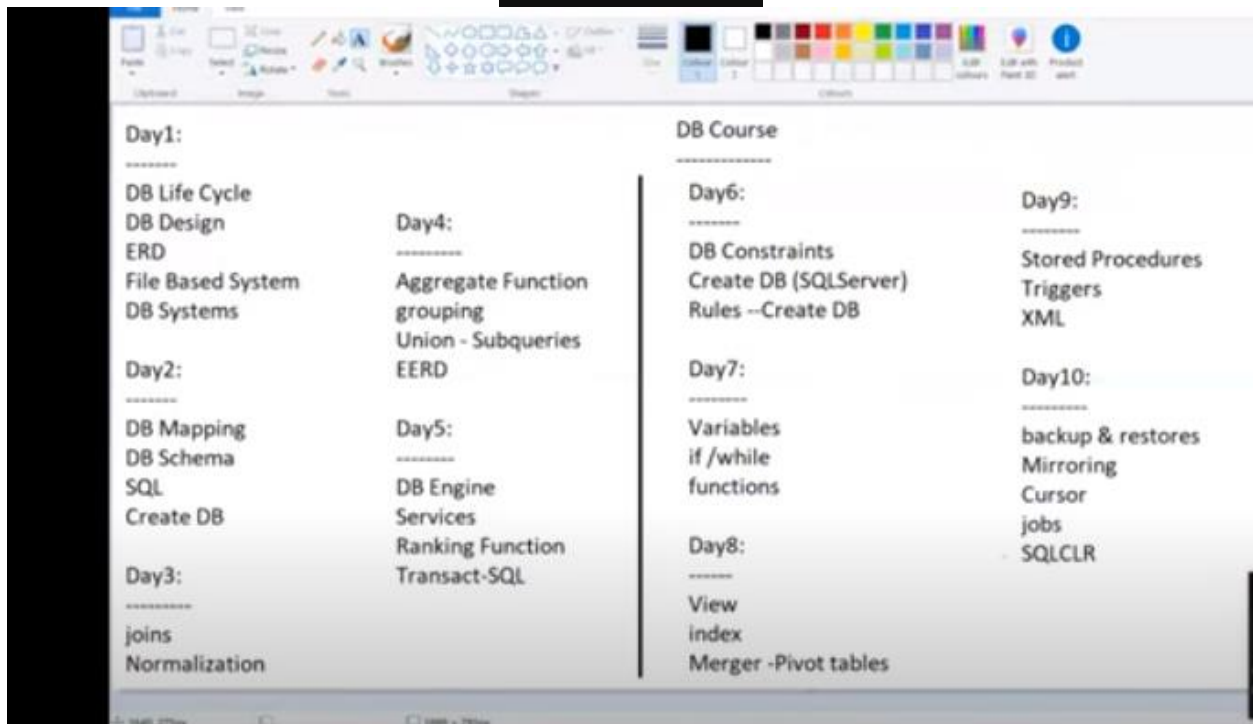


SQL



DB Course	
Day1: ----- DB Life Cycle DB Design ERD File Based System DB Systems	Day4: ----- Aggregate Function grouping Union - Subqueries EERD
Day2: ----- DB Mapping DB Schema SQL Create DB	Day5: ----- DB Engine Services Ranking Function Transact-SQL
Day3: ----- joins Normalization	Day6: ----- DB Constraints Create DB (SQLServer) Rules --Create DB
	Day7: ----- Variables if /while functions
	Day8: ----- View index Merger -Pivot tables
	Day9: ----- Stored Procedures Triggers XML
	Day10: ----- backup & restores Mirroring Cursor jobs SQLCLR

1-SQL introduction, DB Design, ERD, File Based System, DB system

1) هل كل applications لازم يكون ليها DB !؟

لا زي (html,css,js) static website , calculator , office

ولاكن 95% من Big applications معتمدة علي DB

Day 1

DB Life cycle ?!

مراحل DB

1) → analysis

→ "System analysis"

عمران تحصل على "require Document" ويكون فيه تفاصيل المشروع لتوفير ود الفهم بين developer client

2) → DB Design

→ "DB Designer"

تحويل الأفكار الى نموذج use case , ERD

- ERD → entity relationship Diagram

يقدم بنية action او methods ومن هو object

ولكن هو من Table ويكون فيه relationship بين entity

ERD هو يكون DB ؟

لا لازم يمر على DB mapping

3) → DB mapping

→ "DB Designer"

Set of Rules to ERD because

Actual schema "Table" هو من على الماثل الحقيقة

هل require Document هو على ERD واحد ؟

لا ممكن يطلع منه اكثر من ERD واحد كدة بترجع

system analysis هو تعرف المبح

هل ERD نفسه mapping هو مختلف ؟

لا فكل مرة mapping هو مختلف من roles مضافه

وظائف ثابتة

لا فكل مرة mapping هو مختلف من roles مضافه

DB Developer
DB implementation → up to running. (shared DB)

DB نفع
هذه المرحلة 1 - ما تحول DB من logical إلى Physical DB

يكون محتاجين Tool تطبق على الجهاز 1 -
RDBMS → relational DB management system
(MySQL, SQL server, Oracle)

* SQL → structure query language

4 - ما ينزل Tools سواء MySQL و SQL Server
الجهاز 1 يكون "DB server"

5 - application → GUI → interface →

هذه نفع ابن Desktop, mobile, website

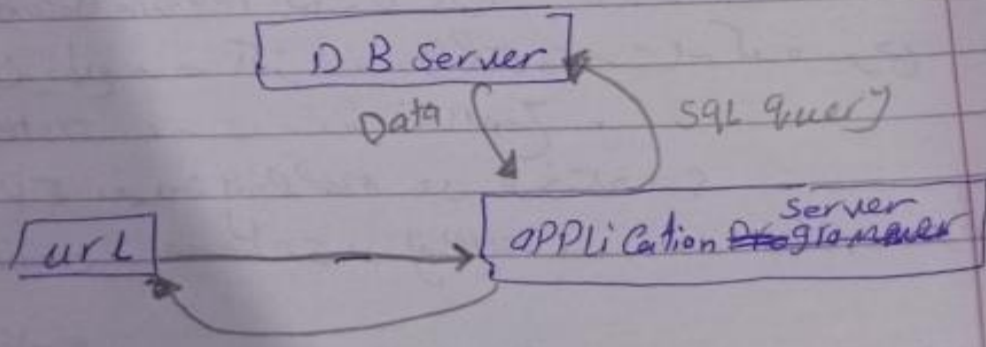
نفع على نفع DB من مع اختلاف اللغة البرمجة

هذه 1 DB Developer تحول إلى application programmer

→ (DB user)

و دفنوا الجهاز 1 يكون application server

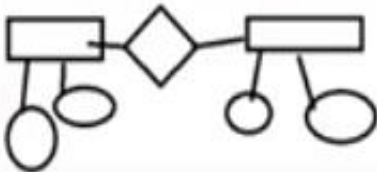
ويكون client يبعث connect



DB Life Cycle

1- Analysis --> System Analyst
Scope (Req Doc)
HR System

2- DB Design --> DB designer
ERD
Entity Relationship Diagram



3- DB Mapping --> DB Designer
Set of Rules
Actual Schema
Tables

4- DB Implementation (shared DB)

Physical DB
Tool
RDBMS
Relational DB Management System
(SQL Server - oracle - MySQL --Access)
SQL (Structure Query Lang)
DB Developer
DB Server

5-Application
GUI -- Interface
Web Site -- Mobile - Desktop
Application Programmer (DB User)
MVC C# java React Php
Application Server

6-Client --> EndUser --> browser -- URL

client → enduser → browser → url

هل client لديه علاقة بـ DB ؟ لا ، لا علاقة بـ browser
يكن url

نعم لأننا (oracle أو SQL server) أنظمة و غالباً ، إنجهوا
File Based system.

وهو عن طريق Text . File . البنية هي ، خيوط خيوط
و هي عبارة عن خيوط في الجوان .

و Format File يكون له Format



Delimited File

يقسم بين (و) ،
"Delimiter"

1. Nora , 25

2. Eman , 20

وهنا اننا نقسم من كل واحد بطور

Delimiter

Fixed width File

مقسمة بـ Bits معينة

1 Java 20 50

يعني بعض read , write

بعض Bits معينة

يعني id ← 20 Bits

Java ← 5 Bits

و 20 ← 3 Bits

وهكذا

لهذا كل ال هيكون موزونة ؟

① different search

صعوبة في البحث

② Low Performance

الكفاءة منخفضة

③ No relationship

مفاتيح علاقات

④ Separated Copies

نسخ منفصلة

⑤ No DB integrity

مفاتيح ترابط بين DB

- تكرار DB
- ⑥ DB Duplication
 - ⑦ long Development time.
لوعايزة اتقاضي خطأ او Duplication في الـ code كود
كبيرة عناء اتقاضي الخطأ
 - ⑧ Security 8 Permissions.
منه مأمونة
 - ⑨ Constrains 8 rules.
ممنوع في الـ "Tables" قواعد
 - ⑩ No Data quality
ممنوع منه تصريف يعرف بين "Data Type"
 - ⑪ manual Backup 8 restore.

- ⑫ No Standard.
ممنوع نظام معين في الـ
- ⑬ Different integrate.
ممنوع الـ "File Text" في الـ
لاخلاف الطريقة

* ونظراً لوجود كل الكثرة في الـ
→ "DB System"

وهنا تحويل ~~Tables relationships~~ إلى Text File
في تحويل Text File إلى Tables relationships
مميزات نظام "DB System"

- ① one standard.
ليجانف الشكل "Tables"
- ② meta Data + Data
الايمن موجود بين جوة DB
مفاتيح بين Data , meta Data ؟
- الـ Tables , علاقات
meta Data → attribute ,
Data → values "cells"
الوجود جوة "cells" values
ماهو الموجود في meta Data , er diagram , Data
"meta Data"

File Based System

Delimited File

1,ahmed,22,10 ok
2,ali,24,70
3,eman,25
3,eman,25
3,eman,25

Student.txt

Fixed width File

10 java 20 SD
30 admin 40 HR
50 IS

Dept.txt

1-Diff Search
2-low Performance
3-Separated Copies
4-No relationship
5-No DB Integrity
6-DB Duplication
7-long Development Time

8-Security & Permissions
9-Constraints & Rules
10-No Data Quality
11-Manual Backup & Restore
12-No Standard
13-Diff Integration

DB System

Tables & Relationship

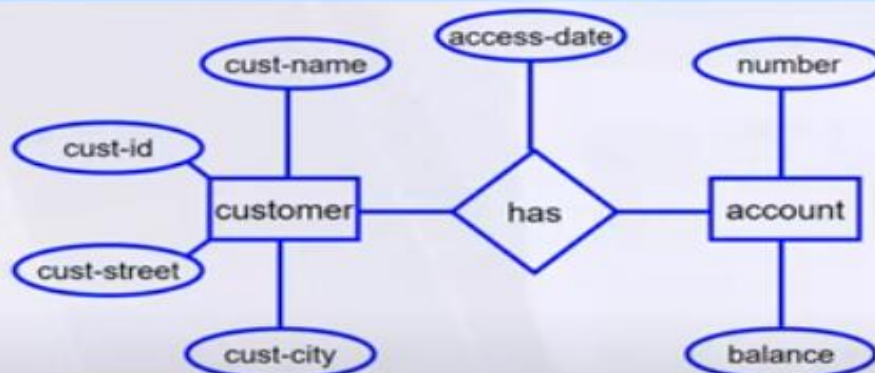
Sid	Sname	
1	ahmed	10
2	ali	70
3	eman	20
3	eman	30

did	dname
10	SD
20	IS
30	CS

foreign key

1- One Standard
2-MetaData + Data
3-Column --> DataType
4-Primary key (Unique -- not NULL)
5-Centralized DB

ER Diagram: Starting Example



- ▶ Rectangles: entity sets
- ▶ Diamonds: relationship sets
- ▶ Ellipses: attributes

5 Column → Data Type. كل Column سيكون له نوع

6 Primary key → (unique, not null)

7 centralized DB لأن يكون على Server واحد

8

* ما هو الفرق بين Database, Database management system و Database system?


11 Data base → create Table أي إنشاء قاعدة


12 Database management system → Tools
SQL Server أو Oracle أو MySQL

13 Database system → هو عبارة عن

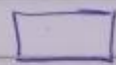
- Database + Tools + interface (GUI)


14 ERD Entity relationship Diagram * ما هو تكوين

15 Entity → object  rectangles

16 attributes → characteristic  ellipses

17 relationship → link → Primary key

* كيان entity 

* سمة attribute 

* علاقة relationship 

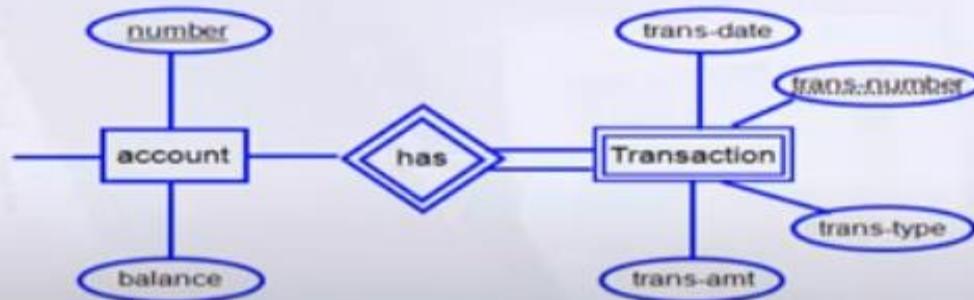
18 في ملف Requirement Document يكون اسماء الأفعال

الاسماء هي "entity"

الأفعال هي "relationship"

Strong Entity Vs Weak Entity

- A **Strong Entity**- An Entity set that has a primary key.
- A **Weak Entity**- An entity set that do not have sufficient attributes to form a primary key.



Partial key: A set of attributes that can be associated with P.K of an owner entity set to distinguish a weak entity.

⑥ هل بين entity Total entities التزم
- ايوة يقع ويكون لازم يكون بتعدده مختلف

* ماهر انواع

entity

⊛ weak entity

⑥ Strong entity

میں ملے اور میں نے کہا کہ -

سکانی

لوہنت قسم من Table

هتخلف فـ عـ عاد الموططين

Primary key ← ایکویٹو لیما

Strong is ~~the~~ entity

مفتی مہرود رسول تائی

بسم الله الرحمن الرحيم

در موقوفه نیز لازم نیست

۱۲۰۰ و ۱۲۰۱

Partial key. $\log_2 2^k$

لافتة ١.٩٩ من

ic attributes

④ ماهی انواع

1 Composite

② multi-value

③ Derived

④ Complex

⑤ Simple.

* الاختلاف بينهم ازاى ؟

① attributes میں سے

(Composite)

⑤ لومتي هيتيب في runtime

bēio (Derived) → dashed.

(٢) لومتي ببتكر لنفسك

double و (multi-value) يكون جوة e لكل يضاف

* ما هو نوع "multi-value 88 composite" ← Complex

mix بين الاثنين

(*) relationship :-

(1) - Degree of Relationship

(2) - Cardinality Constraint

(3) - Participation Constraint

أنواع

محدد الترابط

مشاركة

(*) (1) Degree of Relationship?

- unary → entity يتكون من كيان
- Binary → 2 entity يتكون من كيانين
- Ternary → 3 entity يتكون من ثلاثة كيانات

(*) (2) Cardinality Constraint?

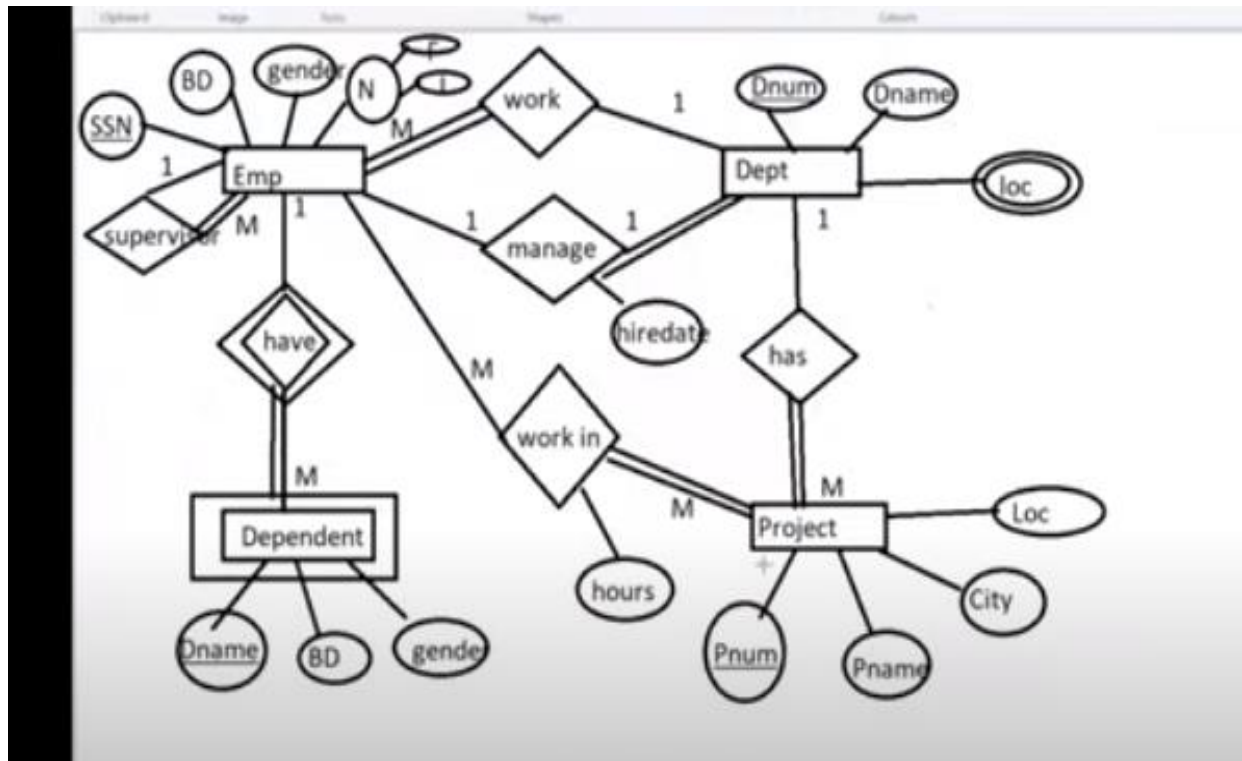
- one to one .
- one to many .
- ~~one~~ many to many .

(*) (3) Participation Constraint

- Total Participation must (=)
- Partial Participation may (—)
(zero or more)

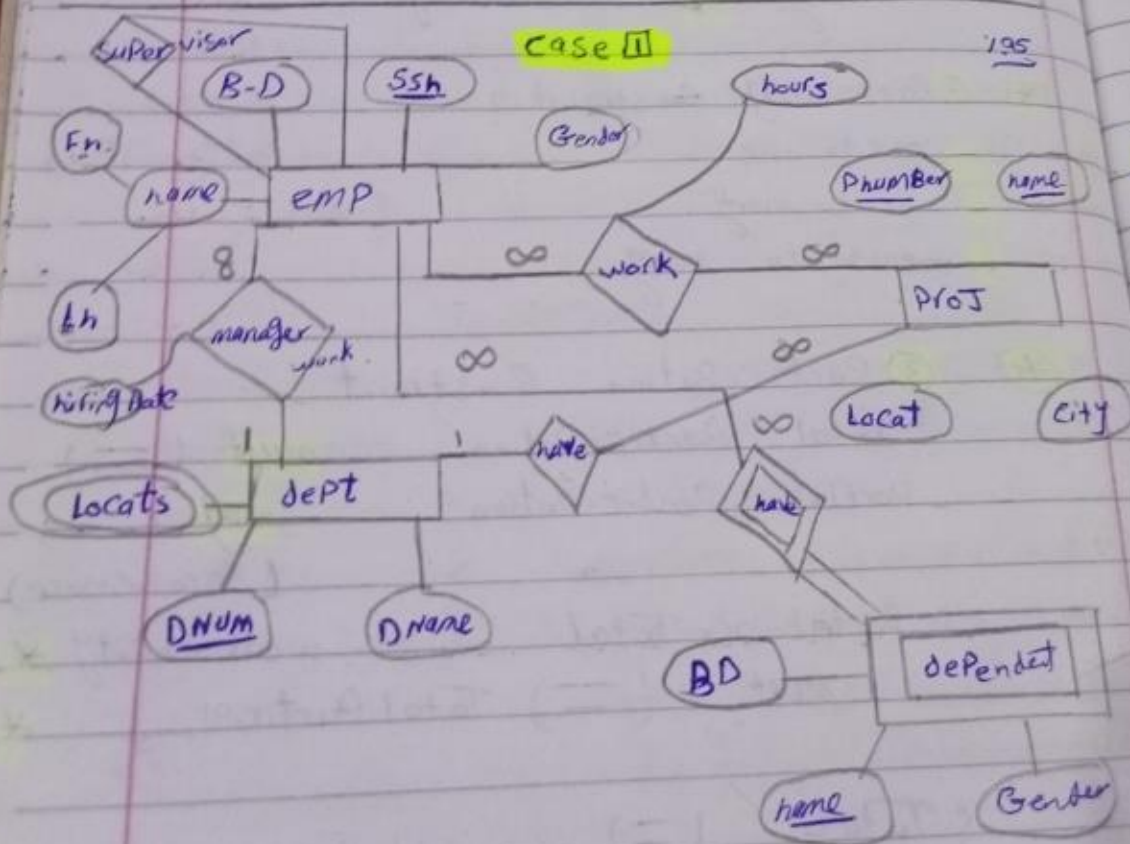
Partial و Total يتكون weak entity *
must (=) Total Participation يتكون *

Total . (=) ماضى اللى اللى اللى اللى *
(must (one or more (mandatory)
Partial (—) ماضى اللى اللى اللى اللى *
(may (zero or more (optional)



[2- SQL Mapping, DB Schema, SQL statements, Creating DB](#)

- * أنواع Key في DB
- 1 Candidate key →
 - 2 Primary key → under lined.
 - 3 Foreign key
 - 4 Composite key → attribute
 - 5 Partial key
 - 6 Alternate key
 - 7 Super key



Day 2

* DB mapping:

Foreign key : هو Primary Key في ملف ثاني وبتعبره
بين الجدول

* في DB لا يوجد غير علاقة
one → many
"Parent" → "Child"

* ممكن يكون Foreign key is null
على الجدول يكون Partial Participation لأنه اختيارى

(*) Server → DBs → Schema → object → Tables →
Columns or Rows → cells

* كل Column في DB ممكن له Domain يعني
وينعرفه من طريق "Data Type" ← "quality"

"Size" ←

* ممكن لو "Size" 1B ← Integer ← Type

"quality" - 128, 128 ←

(*) الـ PK الـ تابع من FK الـ PK بتكون عند PK

MAPPING

- ① mapping to strong entity
 - ② mapping to weak entity
 - ③ mapping of Binary 1:1 relation
 - ④ mapping of Binary 1:N relation
 - ⑤ mapping of Binary M:N relation
 - ⑥ mapping of N-ary relation
 - ⑦ mapping of unary relation
- Binary
Ternary
unary

Strong Entity

- الخطوات الأولى: -
لو في علاقة 1 → Total Partiat من الطرفين
احلها الأول .

① لو Composite attribute يعني (En. In name)
علاقة (En) في المودول و مـ فيكون في جاذة
name .

② multi-value attribute لو

و هيكون في جدول جديد مع
Composite key

③ لو Driven-key من هيكون في Table لان
بيتمسب في "runTime"

④ لو Complex ← هنعمل جدول جديد زي multi-value

❌ ولاكن ممكن Driven يكسب في Tables والحالات
هنا لو داخل في relationship و فيه عدد كبير من cell
و Performance هيكون قليل

الخطوات - weak entity

① علاقة Composite key وهو Partial key
Strong Entity. بتاع Primary key

٢ حالات " → Binary one to one

① one to one and 2 mandatory (Total)

له فيكون الى جدول واحد بين entity فيكون الى جدول واحد

② one to one x optional, y mandatory

له فيكون في جدولين ولكن FK يكون mandatory y

③ one to one x optional, y optional

له فيكون ٢ جداول و الجدول الثاني يكون الى

PK x < PK y

④ one to many x optional, y mandatory

له فيكون جدولين و FK في many mandatory حتى لو كانت mandatory

⑤ one to many 2 optional

له فيكون ٢ جداول و الجدول ٢ يكون في PK x < PK y

و Primary فيكون الى من ناحية many

⑥ many to many

من هنا في optional في غير الاعتبار و يكون في

٢ جداول PK x < PK y Composite

* لاحظ لو في attribute الى relationship يكون في ١٤

في الجدول الى جامع Primary key لجدولين

* Ternary

h>2

المتحول الى جدول على جدولين
"Cardinality: Pontical"

* unary relation ship

Fk يكون في نفس الجدول

"mapping"

Case II

4.1--

EMP

<u>SSN</u>	BD	Gender	NE	NL	Superid	work_Dept
------------	----	--------	----	----	---------	-----------

Dept

<u>Dum</u>	Dname	manf-id	hire Date-
------------	-------	---------	------------

Dept-loc	<u>Dum-id</u>	Locat
----------	---------------	-------

Project

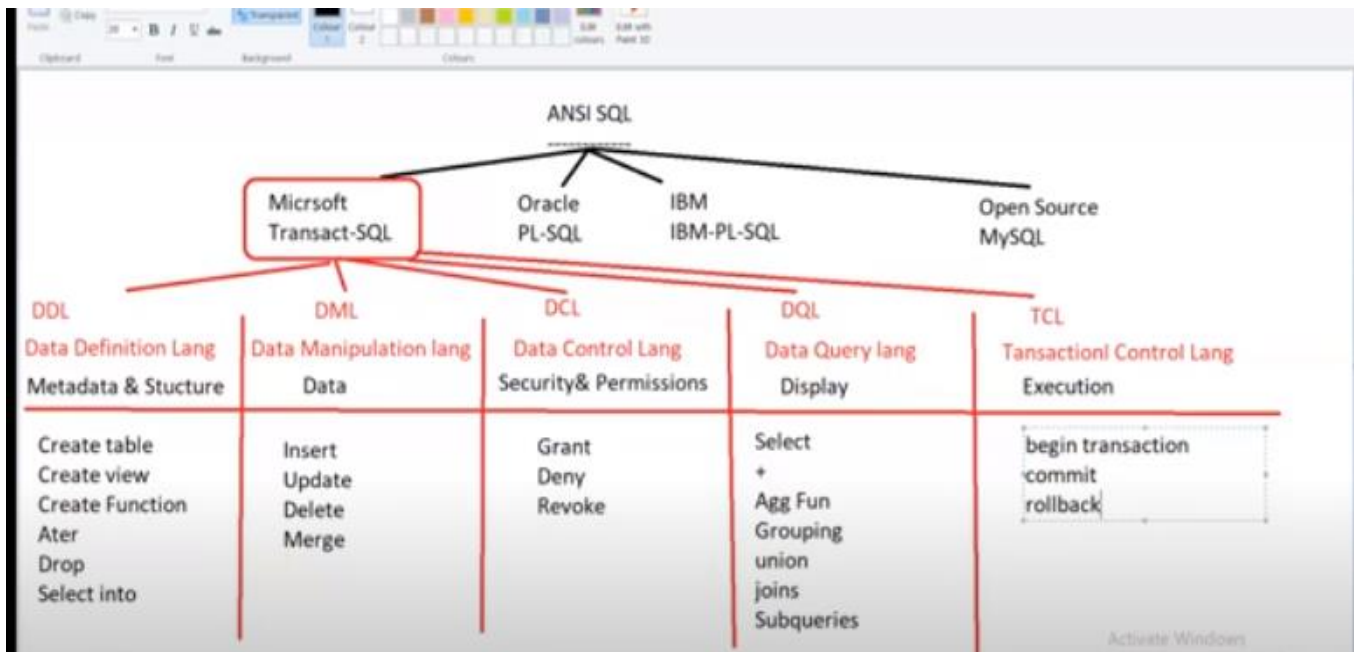
<u>Pnum</u>	Pname	city	loc	pebt-id
-------------	-------	------	-----	---------

work

<u>Pnum</u>	<u>EMP-id</u>	hours
-------------	---------------	-------

de Pent

<u>Dname</u>	<u>BD</u>	<u>Gender</u>	<u>EMP-id</u>
--------------	-----------	---------------	---------------



```

create table emp
(
  eid int Primary key,
  ename varchar(20) not null,
  eage int,
  eadd varchar(20) default 'cairo',
  hiredate date default getdate(),
  Dnum int
)

alter table emp add sal int

```

Create table

Alter table to add new attribute

* ans: SQL \rightarrow

American National standards institute
structured query language.

query, i Syntax $\vec{p} \in \mathbb{R}^n$ is a vector

المعلومة - صحة DB

*) النماذج مفتوحة المصدر open source أخذت mySQL

Transact-SQL ← microSoft

my son ✗ تفسیر

* DDL → Data Definition Language

(*) DML \rightarrow Data Manipulation lang.

(*) Dcl → Data Control lang

```
alter table emp add sal int
```

```
alter table emp alter column sal bigint
```

```
alter table emp drop column sal
```

```
drop table emp
```

--DMIL

```
--insert update delete
```

```
insert into emp
```

```
values(1, 'ali', NULL, 'alex', '1/1/2010', NULL)
```

```
insert into emp(ename,eid)
```

```
values('eman', 9)
```

```
--insert constructor
```

```
insert into emp(ename,eid)
```

```
values('eman',8), ('ali',12), ('nada',7)
```

* نوع عبارة `alter` `alter Table emp alter Column Gender Datatype int;`

* نوع عبارة `drop` `alter Table emp drop Column Gender;`

* نوع عبارة `drop` `Drop Table emp;`

DML: `insert` / `update` / `delete`

* `insert` `into` `emp` `values` (1, 21212020; female, morad);

أو

`insert into emp (ssn, Gender, h-f) values (2, male, Ali);`

* نوع عبارة `insert` `into` `emp` `values`

`(3, Female)` & `(4, male)` , `(5, male);`

* نوع عبارة `update` `emp` `set` `Gender = 'male' where ssn = 1;`

* نوع عبارة `delete` `from` `emp` `where ssn = 1;`

* نوع عبارة `select` `from` `emp`

DQL "select"

`Select * from emp;`

*) لو عبارة اعرف Data بترتيب
Select * From emp order by Date desc

تساربي

*) لو عبارة اجمع بين ~~attributes~~ 2
Select F-h, L-h as Fullname From emp,

*) هل يمنع اقرار ب Null عن طريق is
Select * From emp ~~where~~ where Gender is not null

*) لو عبارة اعرف Data بدون تكرار
distinct

Select distinct F-h From emp,

*) لو عبارة اعرف Condition اعرف city
Cairo او man Soura

*) Select * From emp where city in ('Cairo', 'man Soura')

*) لو عبارة age يكون من 30 الى 50
Select * From emp where age between
30 and 50,

3-SQL Joins, Normalization

Types of joins

- Cross join
 - Cartesian Product -- 16 rows
- Inner Join
 - Equi Join PK = FK
- Outer Join
 - Left Outer join -- 4 rows
 - Right Outer Join -- 5 rows
 - Full Outer Join -- 6 rows
- Self Join (Unary relationship)

Select Sname, Dname
From Student S left full outer join Dept D
On D.Did = S.Did

Student

Sid	Sname	did
1	ahmed	10
2	khalid	10
3	eman	20
3	omar	NULL

Dept

Did	Dname
10	SD
20	HR
30	IS
40	Admin

Result of Full Outer Join:

Sname	Dname
ahmed	SD
khalid	SD
eman	HR
omar	NULL
NULL	IS
NULL	Admin

Self join

Select X.Ename as EmpName , Y.Ename as SuperName
From Employee X , Employee Y
where Y.eid = X.superid

Employee

Eid	ename	Superid
1	ahmed	NULL
2	omar	1
3	eman	1
4	nada	2

Result of Self Join:

Empname	SuperName
omar	Ahmed
eman	Ahmed
nada	omar

Select X. Ename , Y.*
From employee X, employee Y
where y.eid=x.superid

Result of Self Join with all columns:

parent	Child
Super	employee

Example

SID	SName	Birthdate	City	Zip Code	Subject	Grade	Teacher
1	Ahmed	1/1/1980	Cairo	1010	DB	A	Hany
1	Ahmed	1/1/1980	Cairo	1010	Math	B	Eman
1	Ahmed	1/1/1980	Cairo	1010	WinXP	A	khalid
2	Ali	1/1/1983	Alex	1111	DB	B	Hany
2	Ali	1/1/1983	Alex	1111	SWE	B	Heba
3	Mohamed	1/1/1990	Cairo	1010	NC	C	Mona

Student(SID, Sname, Birthdate,)

SID	SName	Birthdate	ZipCode
1	Ahmed	1/1/1980	1010
2	Ali	1/1/1983	1111
3	Mohamed	1/1/1990	1010

Stud_City(City, Zip Code)

City	Zip Code
Cairo	1010
Alex	1111

Stud_Subject (SID, Subject, Grade)

SID	Subject	Grade
1	DB	A
1	Math	B
1	WinXP	A
2	DB	B
2	SWE	B
3	NC	C

Subject (Subject, Teacher)

Subject	Teacher
DB1	Hany
Math	Eman
WinXP	khalid
DB2	Hany
SWE	Heba
NC	Mona

ITI Example

ITI Students Sheet

Platform Name : SWE

Platform Description: Software Engineering

Graduate Manager: Dr.Baha

Appno	Name	F-code	Faculty	Address	Telno	Grade	Att. Hrs	Sdate
123	Ahmed	SC-phy	Science	Haram	3386842	A	600	14 Sep
124	Mona	Eng-cs	Engineering	Dokki	3389745, 3389744, 5123445	B	591	15 Sep
127	Ali	Com-ac	Commerce	Nasr City	2241593, 2222345	A	550	21 Sep
223	Karim	Med-bio	Medicine	Sheraton	2286845	C	600	14 Sep

3NF

- **Students:** appno, name , FCode, address
- **Fac_majors:** faculty , FCode

Unchanged Tables

- **Platform :** pfname , pfdesc , pfManager
- **Std_Tel:** appno, telno
- **Students_pf:** pfname, appno, grade, attd , start_date

4-SQL, Aggregate Function, Grouping, Union, Subqueries, EERD

Aggregate Function

Count, Max, Min, Avg, Sum

Select **Sum(Salary)**
From employee **62500**

Select **Min(Salary), Max(salary)**
From Employee **1000 9000**

Select **Count(eid)**
from employee **15**

Count(*) → 15
Count(eid) → 15
count(ename) → 14

Select **Avg(salary)**
from employee

Select **Min(salary), did**
from employee
group by did

1000	10
2000	20
1500	30

Select **Count(eid), Address**
from employee
group by Address

6	Cairo
5	Alex
4	Mansoura

Select **Count(eid), address**
from employee
where did in(10,30)
group by address

3	Cairo
3	Alex
4	Mansoura

Select **Min(salary), did**
from employee
where address like '_a%'
group by did

2000	10
2000	20
1500	30

Eid	Ename	Salary	Address	did
1	ahmed	3000	cairo	10
2	ali	5000	cairo	10
3	eman	2000	cairo	10
4	khalid	1000	alex	10
5	yousef	4000	alex	10
6	sameh	5000	alex	10
7	mohamed	6000	alex	20
8	alaa	7000	alex	20
9	ola	4000	cairo	20
10	reem	2000	cairo	20
11	nada	9000	cairo	20
12	sayed	8000	mansoura	30
13	reham	1500	mansoura	30
14	sally	2000	mansoura	30
15	on	3000	mansoura	30

NULL

Select **Sum(salary), did**
from employee
group by did
having Sum(salary) >= 22000

28000	20
-------	----

Select **Count(eid), address**
from employee
group by address
having Count(eid) >= 5

6	CAIRO
5	ALEX

Select **Sum(salary), did**
from employee
group by did
having Count(eid) > 5

20000	10
-------	----

Select **Sum(salary), did**
from employee
where address like '_a%'
group by did
having sum(salary) > 12000

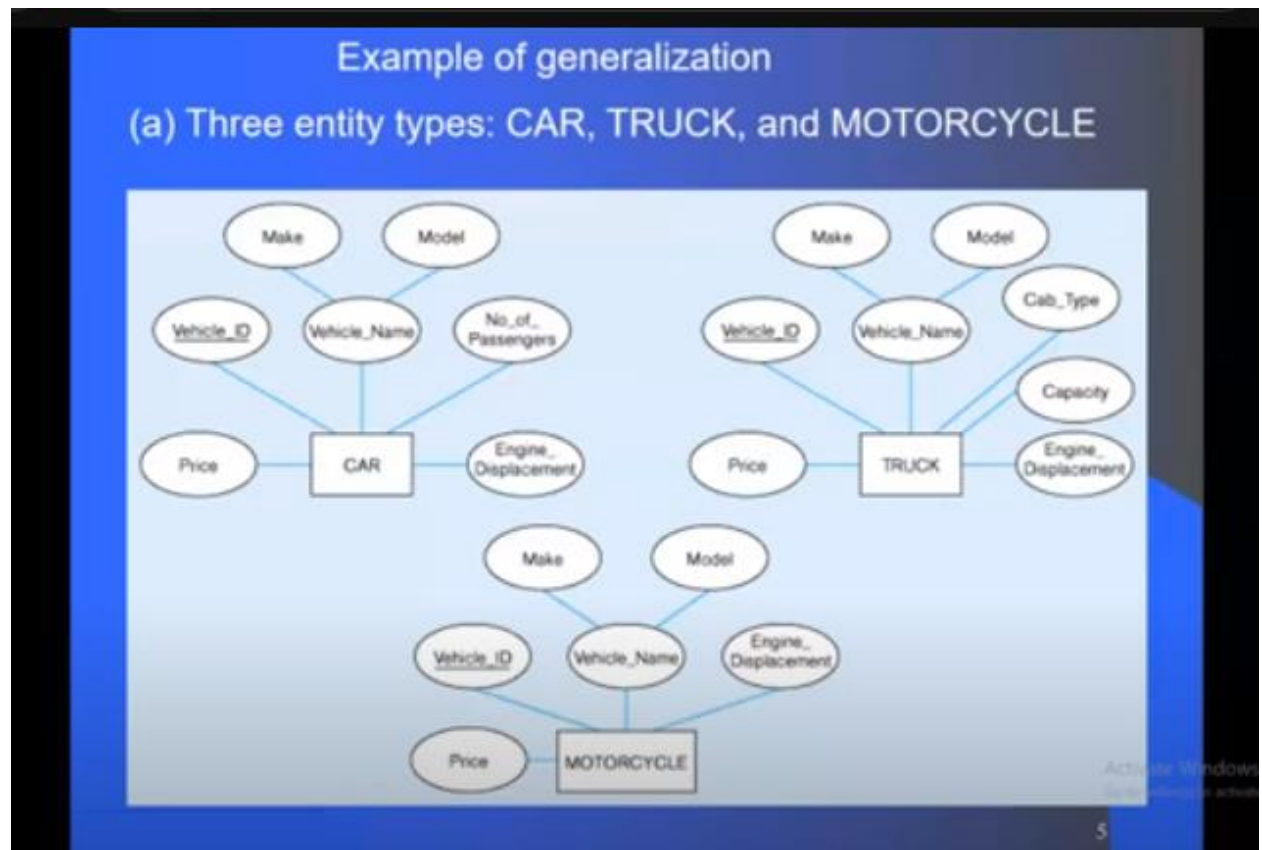
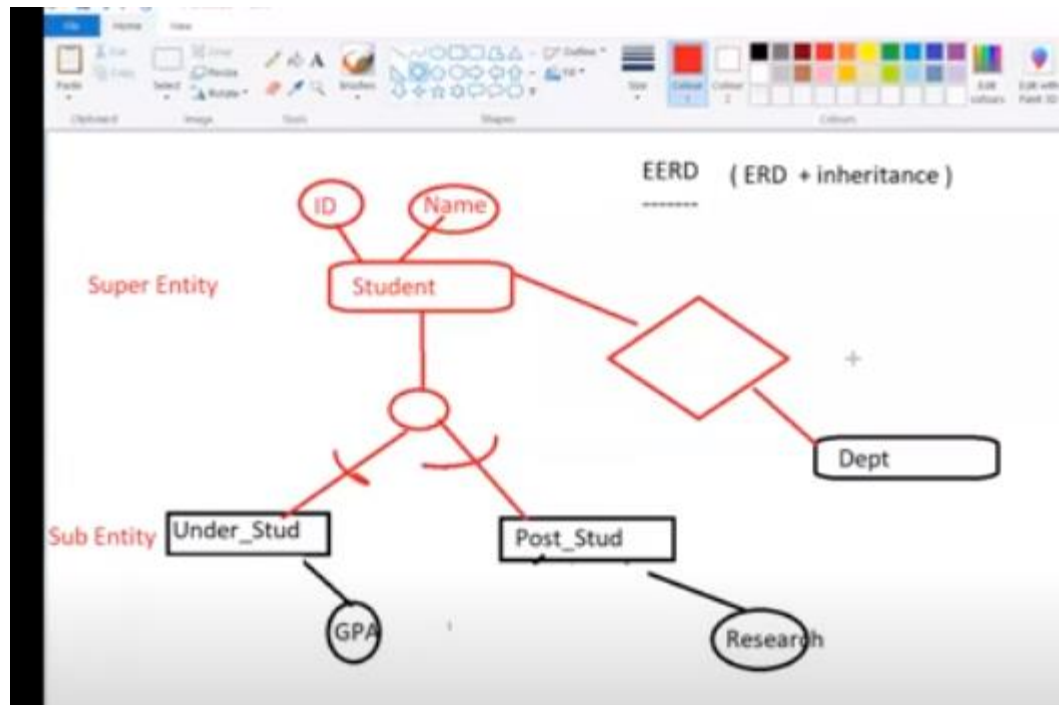
15000	20
14500	30

Select **Max(salary), address**
from employee
where did in(10,30)
group by address
having Count(eid) > 3

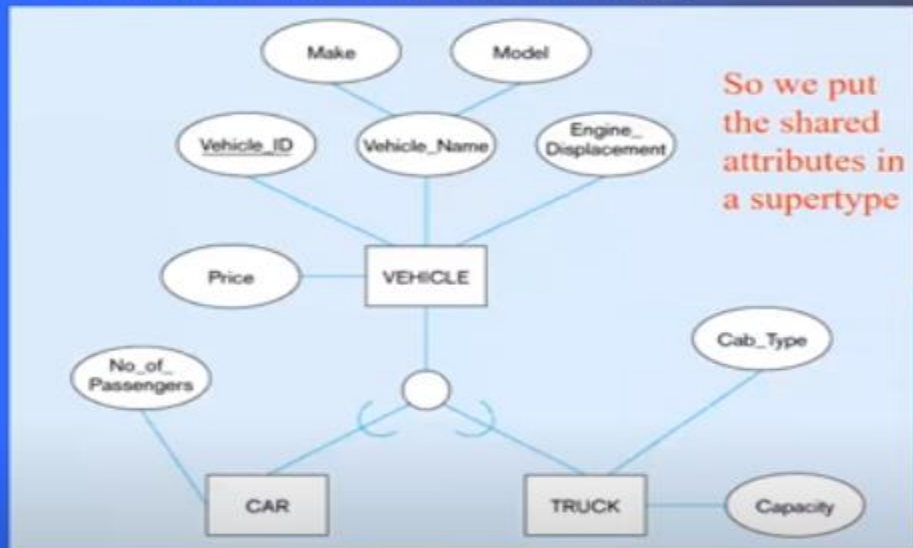
8000	Mansoura
------	----------

Eid	Ename	Salary	Address	did
1	ahmed	3000	cairo	10
2	ali	5000	cairo	10
3	eman	2000	cairo	10
4	khalid	1000	alex	10
5	yousef	4000	alex	10
6	sameh	5000	alex	10
7	mohamed	6000	alex	20
8	alaa	7000	alex	20
9	ola	4000	cairo	20
10	reem	2000	cairo	20
11	nada	9000	cairo	20
12	sayed	8000	mansoura	30
13	reham	1500	mansoura	30
14	sally	2000	mansoura	30
15	on	3000	mansoura	30

NULL



Generalization to VEHICLE supertype



Note: no subtype for motorcycle, since it has no unique attributes

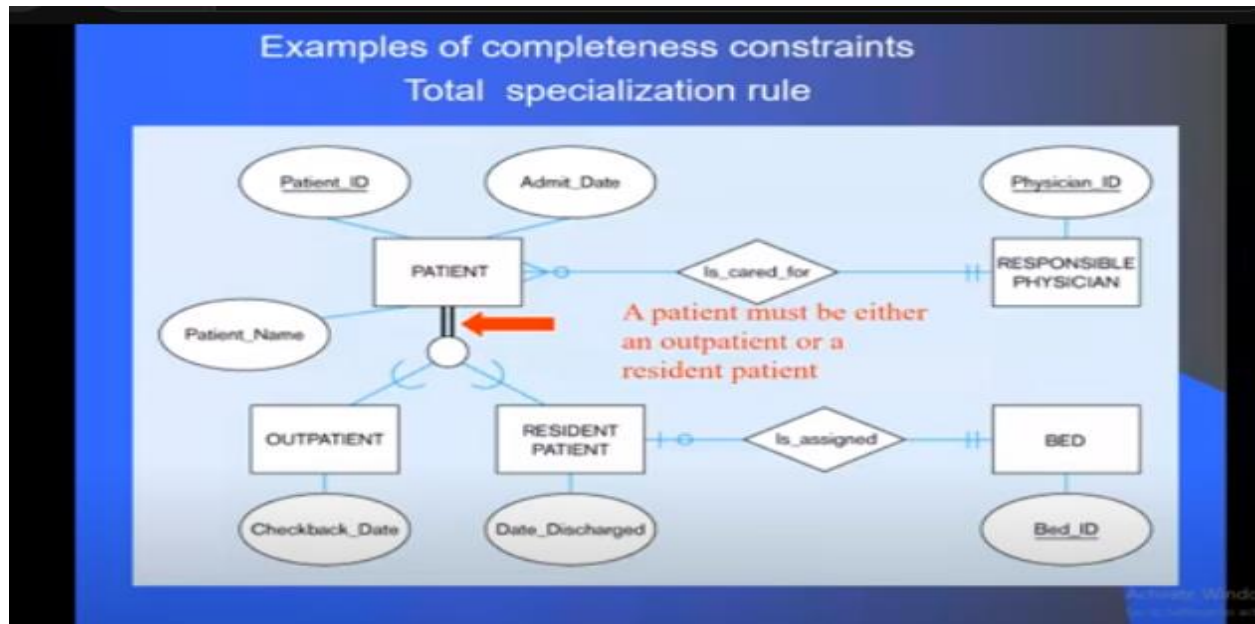
6

New example

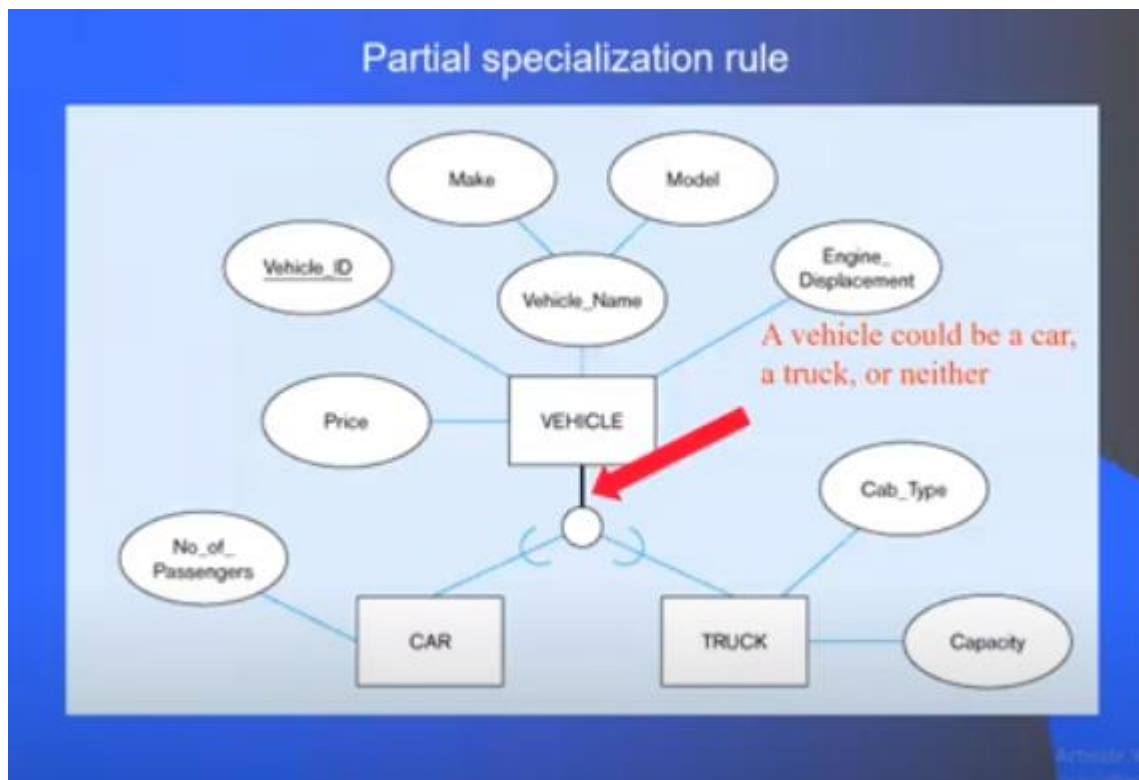
Employee supertype with three subtypes



Total

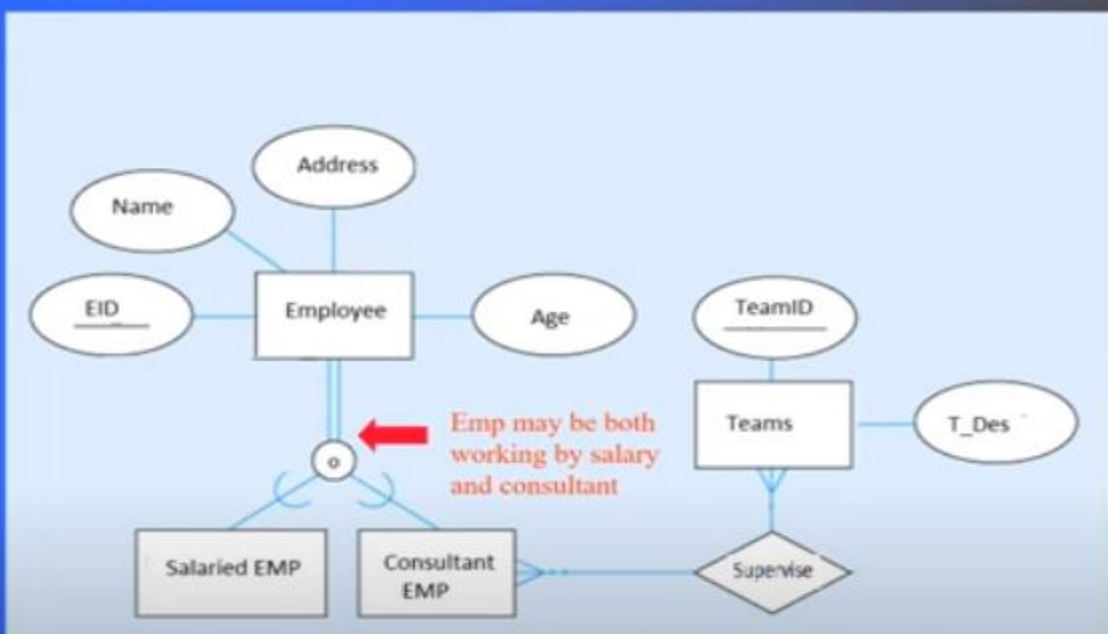
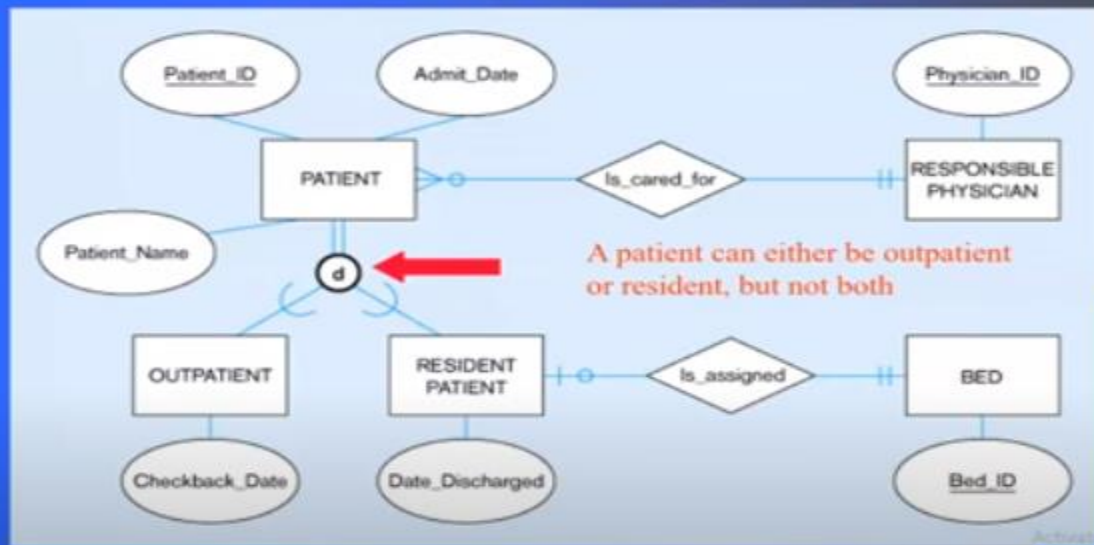


partial



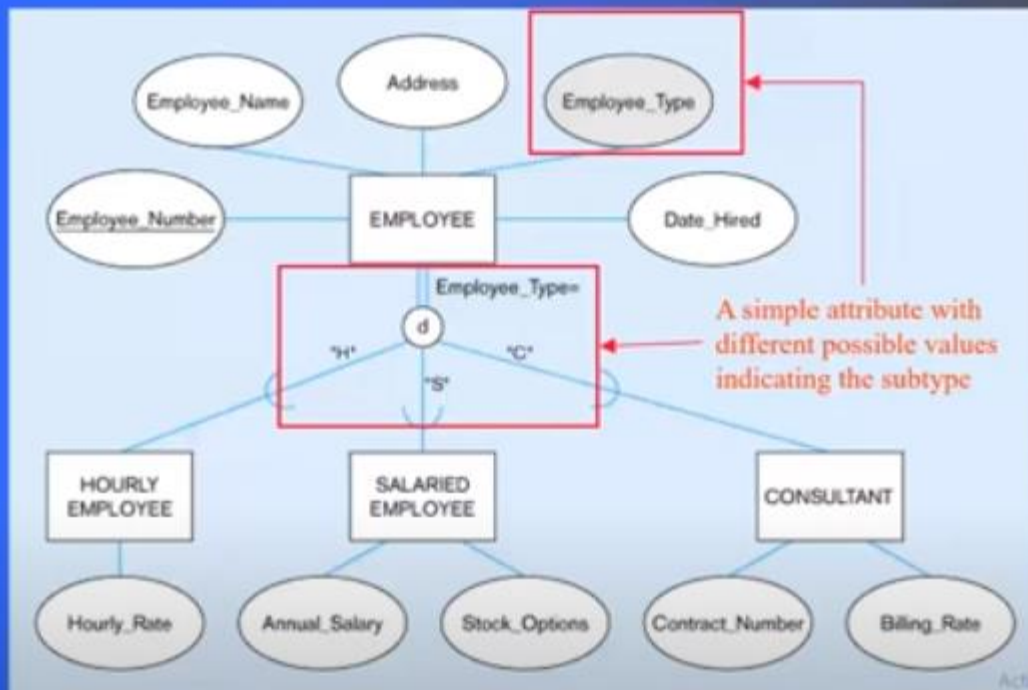
Examples of disjointness constraints

Disjoint rule

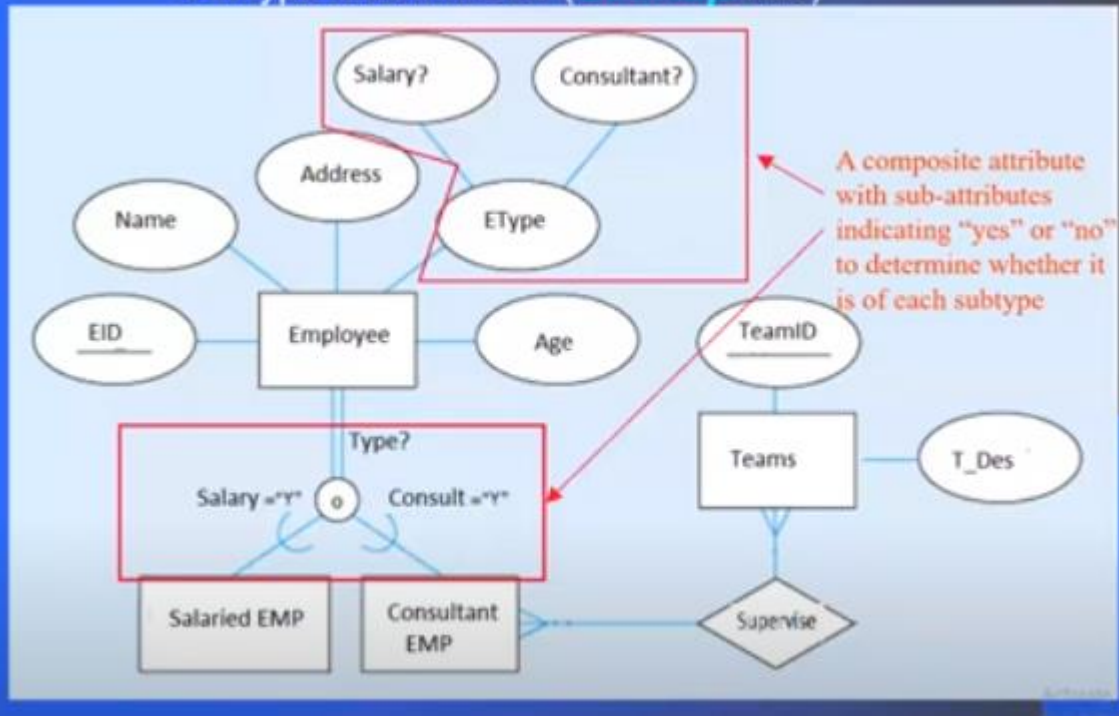


Overlap rule

Introducing a subtype discriminator (*disjoint* rule)



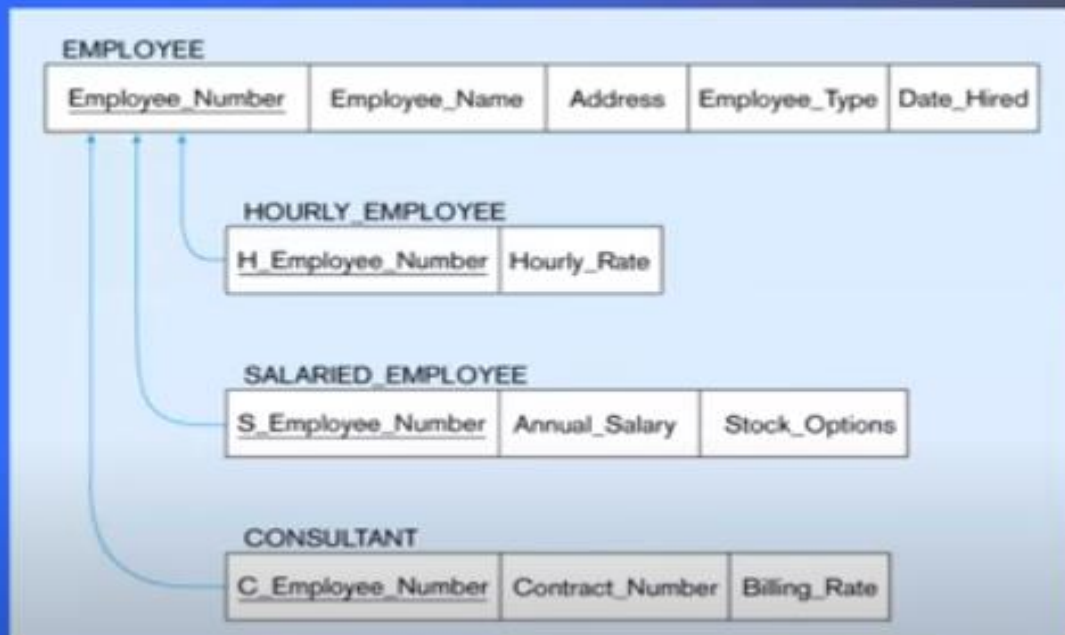
Subtype discriminator (*overlap* rule)

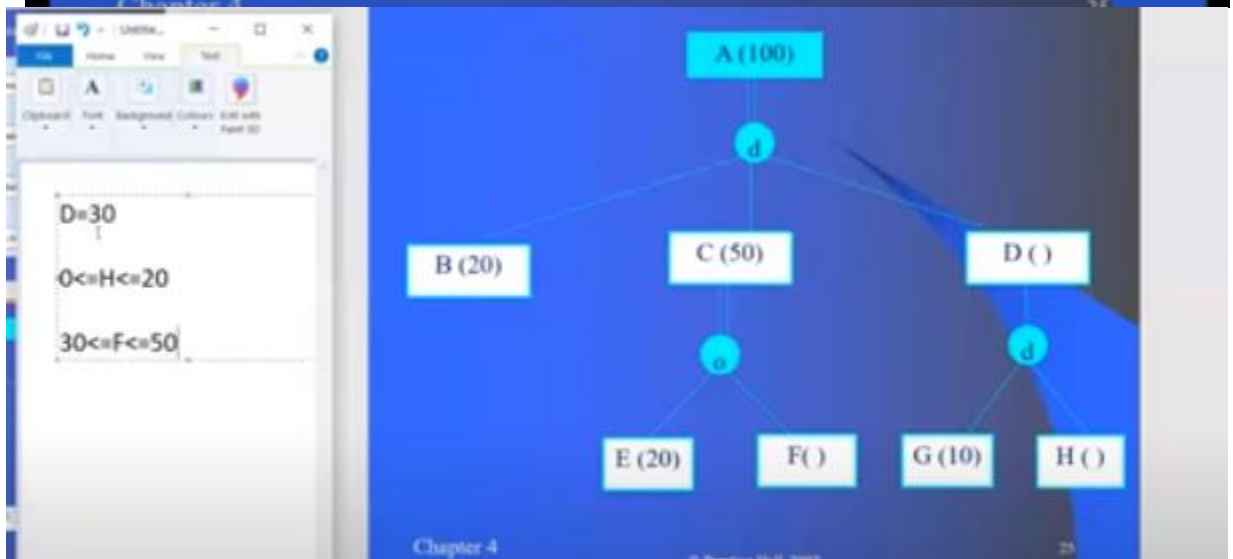
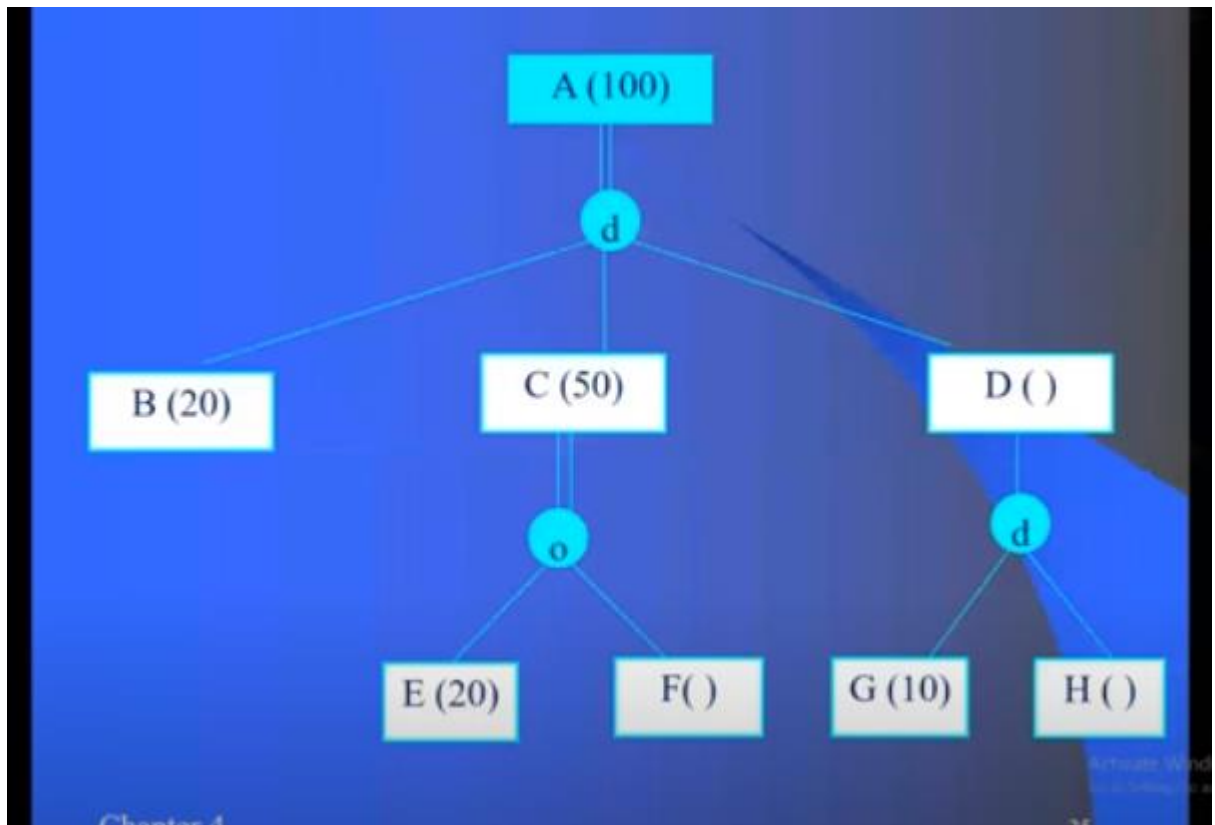


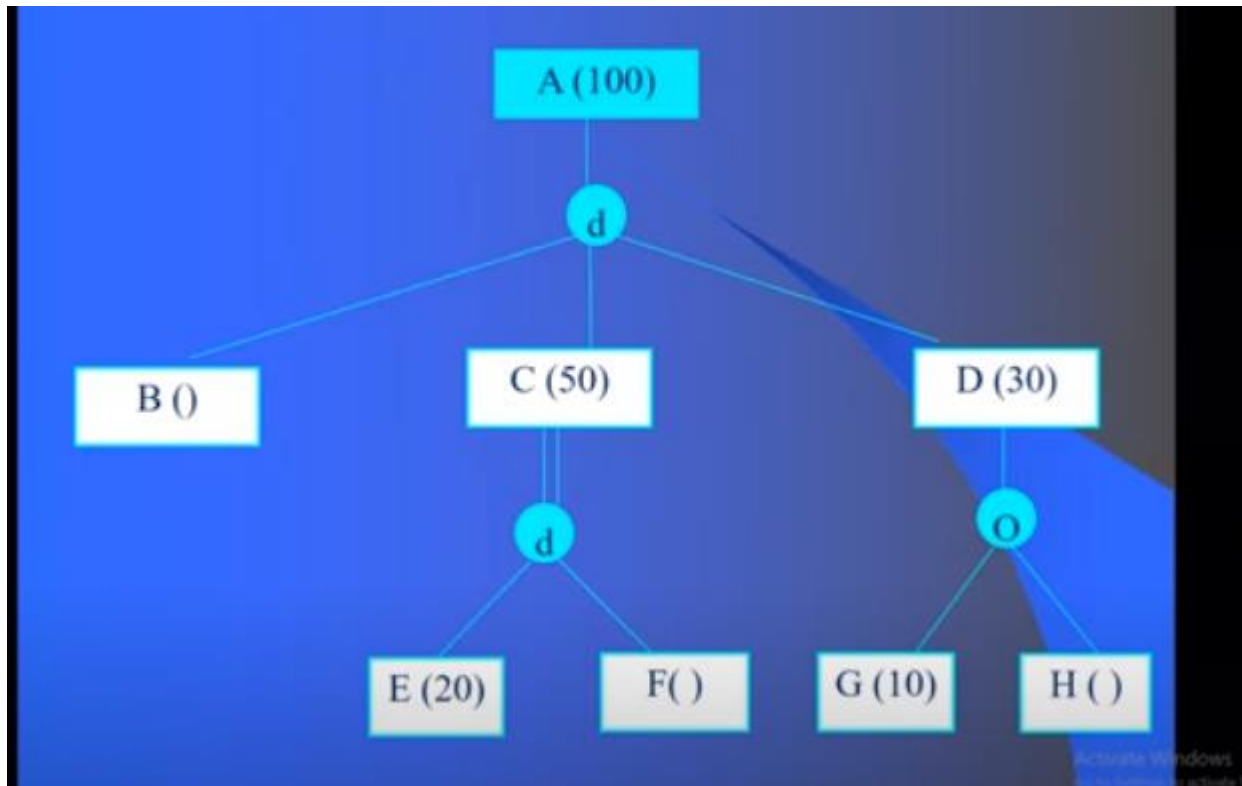
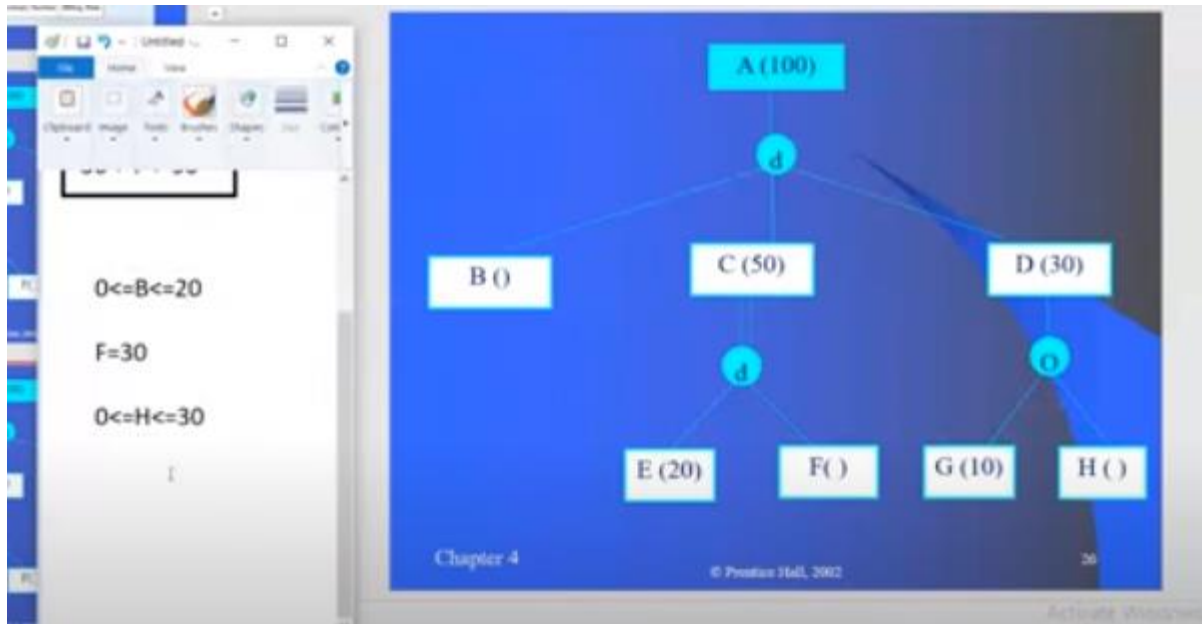
Transforming EER Diagrams into Relations



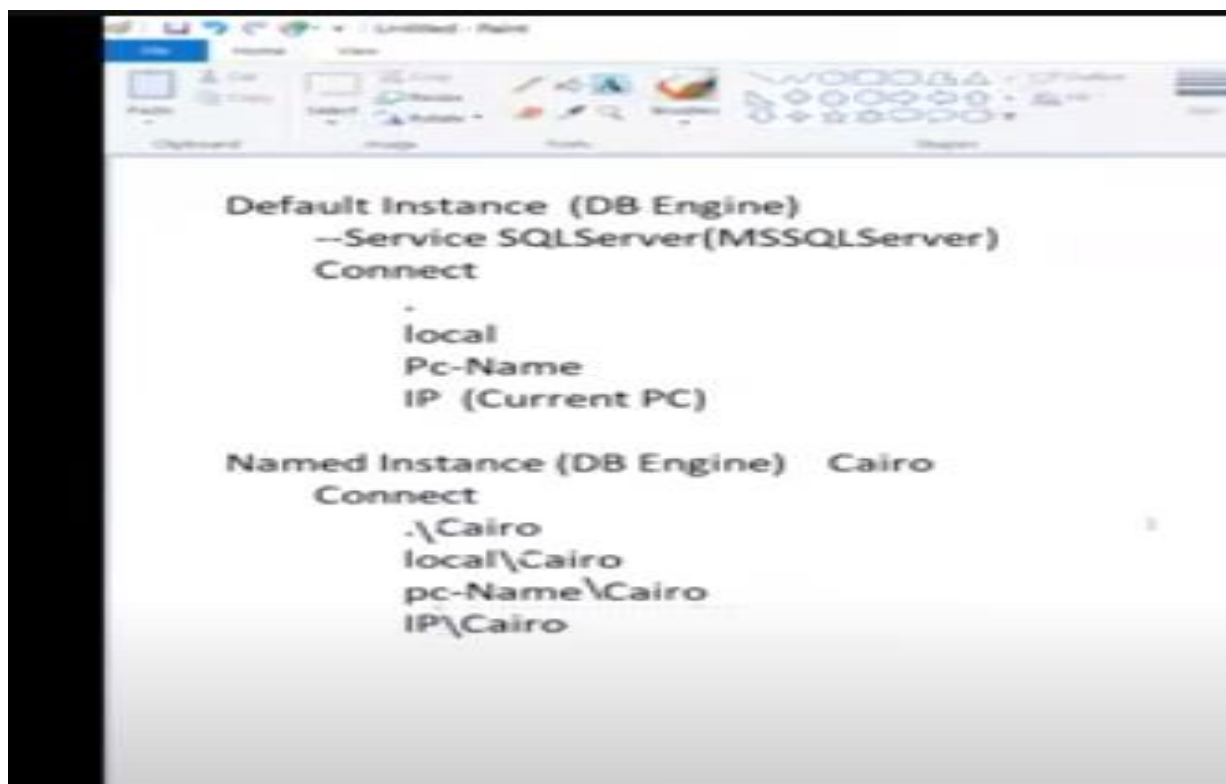
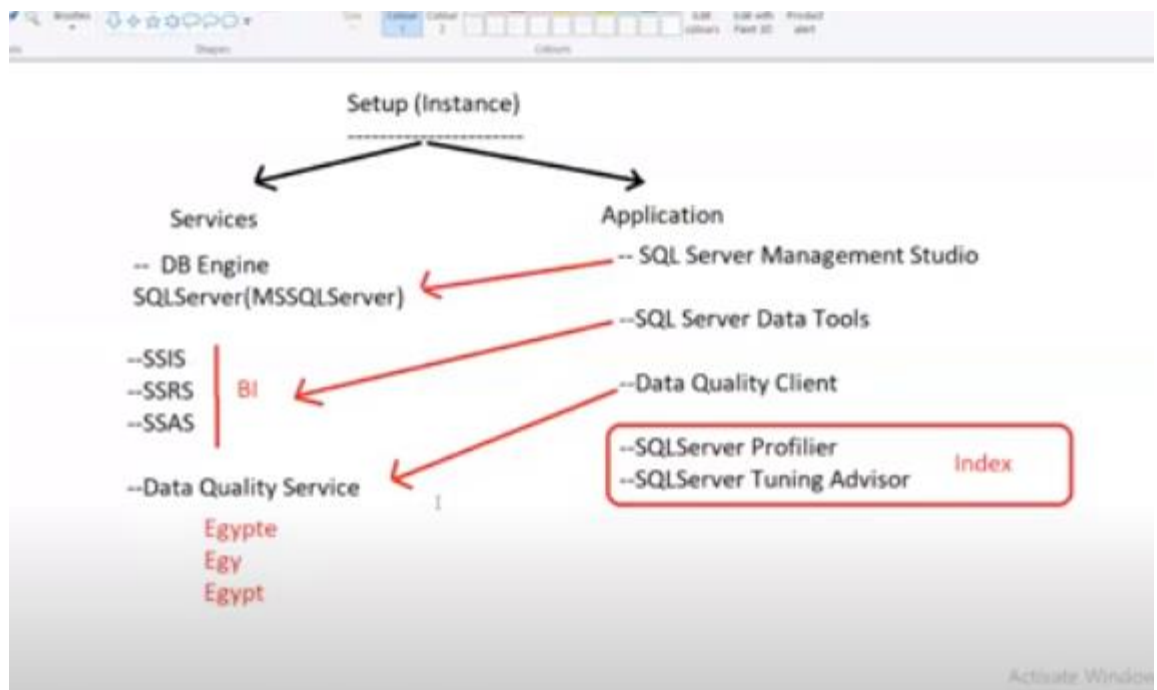
Mapping Supertype/subtype relationships to relations







5-SQL, DB Engine, SQL Services, Ranking Function, Transact SQL



Ranking Functions

- Row_Number()
- Dense_rank()
- NTiles(Group)
- Rank()**

```

Select *
From (
    Select *, Row_Number() over(order by esal desc) as RN
        , Dense_rank() over(order by esal desc) as DR
        , NTile(3) over(order by esal desc) as G
    From employee ) as Newtable

where RN=1          DR=1
       RN=3          DR<=2
       RN<=2         G=1
  
```

eid	ename	esal	did	RN	DR	G
15	ahmed	10000	10	1	1	1
14	ali	10000	10	2	1	1
12	eman	9000	10	3	2	1
1	nada	9000	10	4	2	1
2	reem	9000	10	5	2	1
3	khalid	8000	10	6	3	2
7	mohamed	7000	20	7	4	2
8	sayed	7000	20	8	4	2
6	hassan	6000	20	9	5	2
5	omar	6000	20	10	5	2
9	sally	5000	30	11	6	3
10	shimaa	4000	30	12	7	3
11	hana	4000	30	13	7	3
12	lama	3000	30	14	8	3

Ranking Functions

- Row_Number()
- Dense_rank()
- NTiles(Group)
- Rank()**

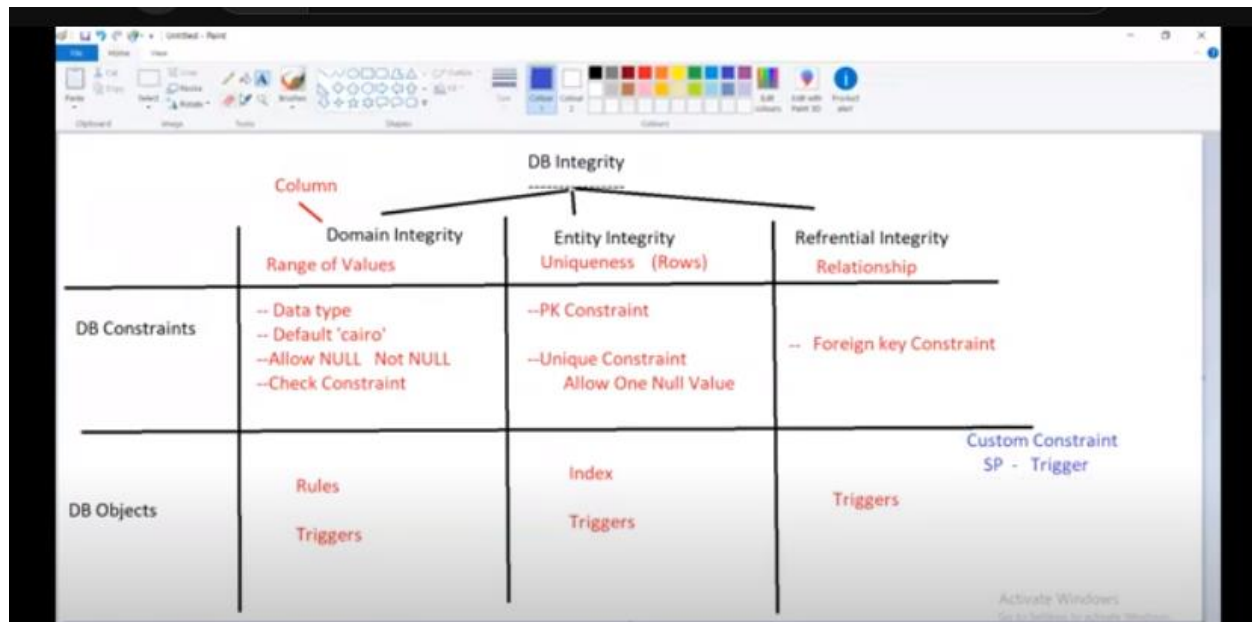
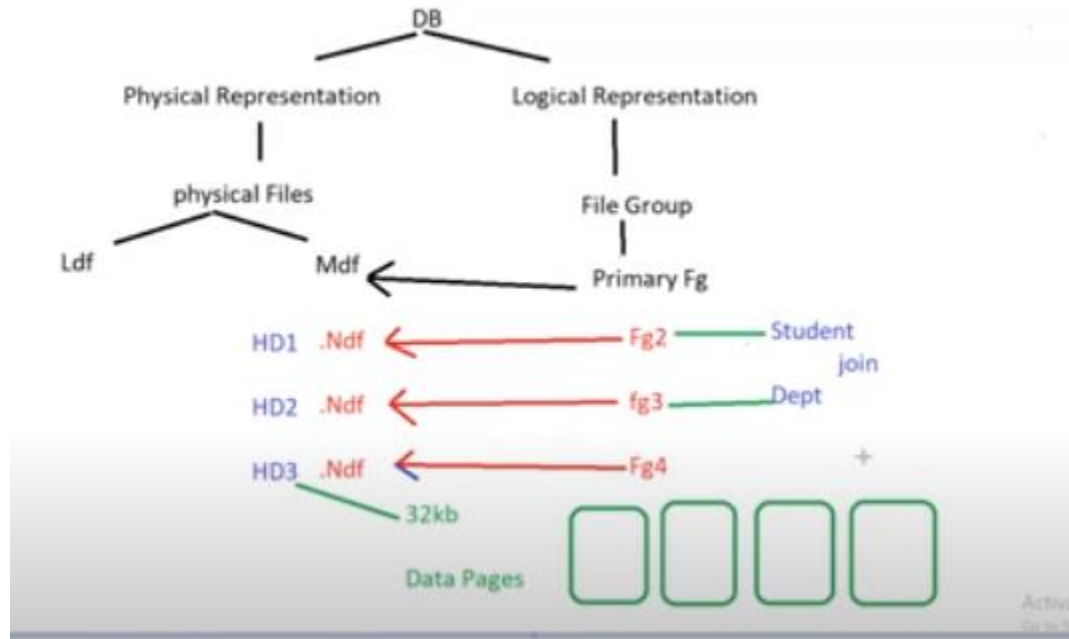
```

Select *
From (
    Select *, Row_Number() over(partition by did order by esal desc) as RN
        , Dense_Rank() over (Partition by did order by esal desc) as DR
    From employee ) as newtable

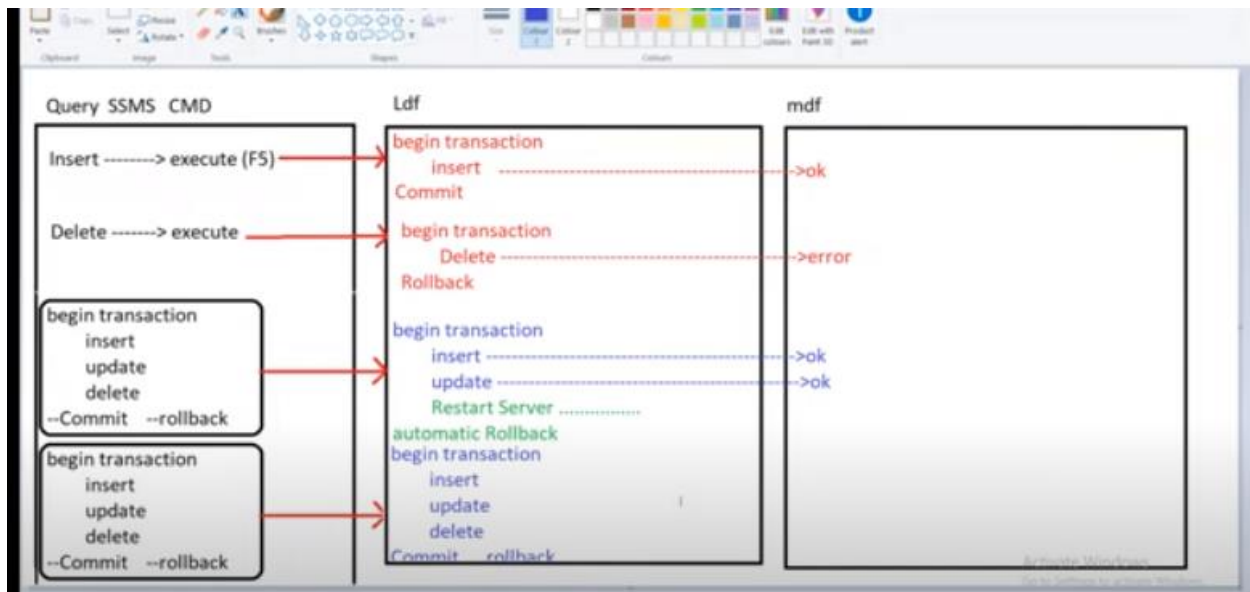
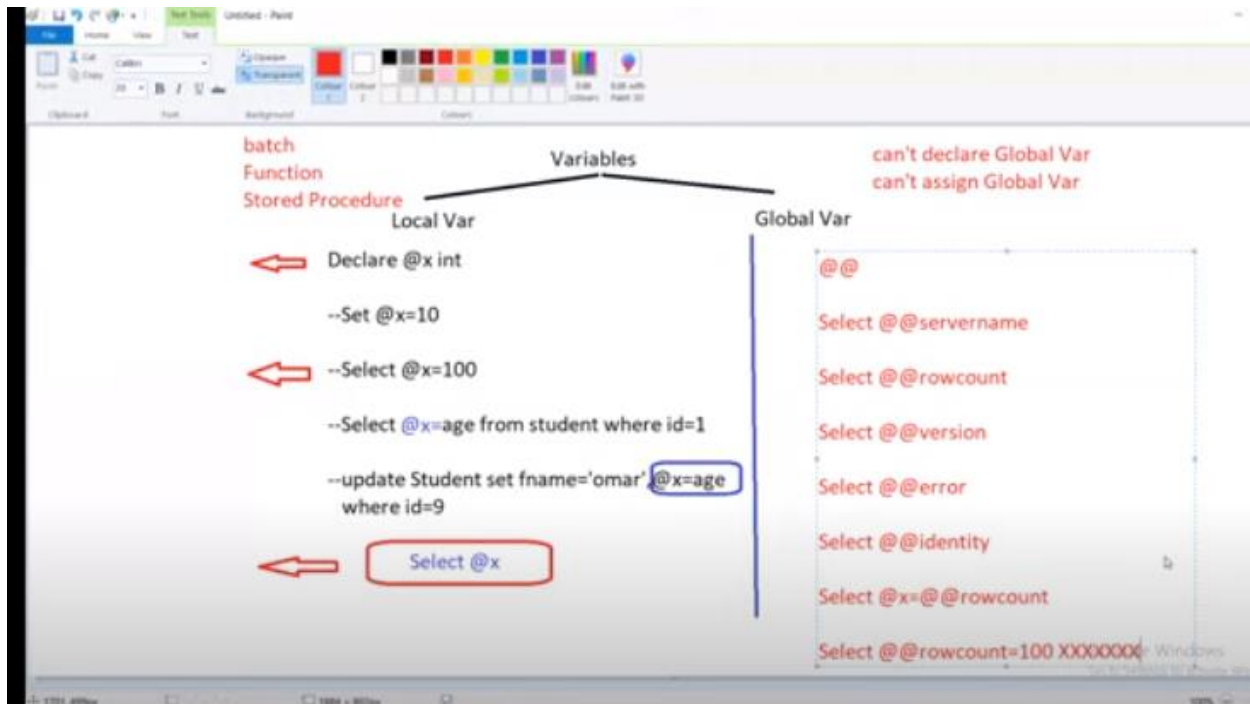
Where RN=1          DR=1
       RN=3          DR<=2
  
```

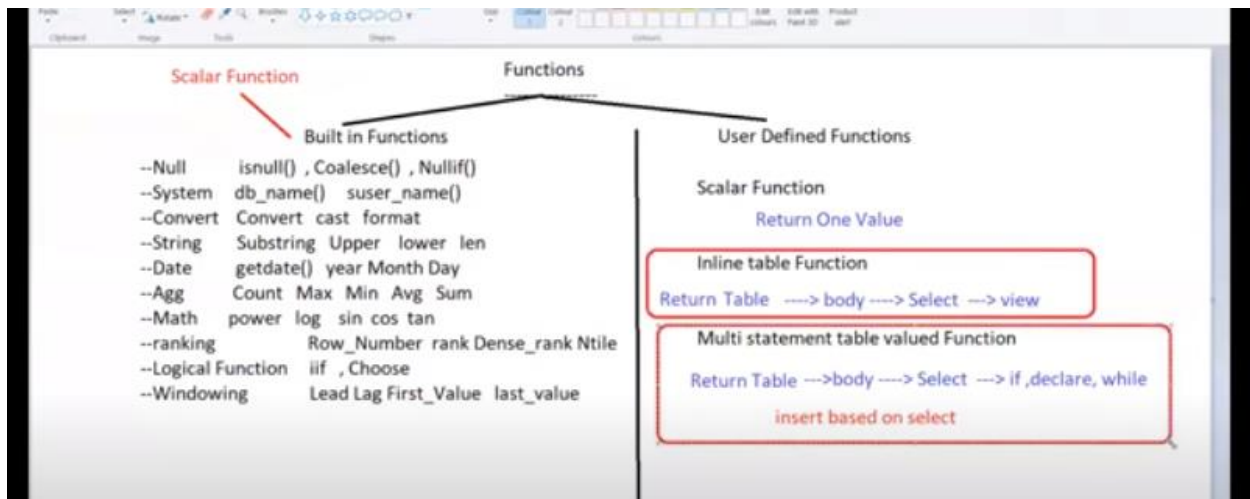
eid	ename	esal	did	RN	DR
15	ahmed	10000	10	1	1
14	ali	10000	10	2	1
12	eman	9000	10	3	2
1	nada	9000	10	4	2
2	reem	9000	10	5	2
3	khalid	8000	10	6	3
7	mohamed	7000	20	1	1
8	sayed	7000	20	2	1
6	hassan	6000	20	3	2
5	omar	6000	20	4	2
9	sally	5000	30	1	1
10	shimaa	4000	30	2	2
11	hana	4000	30	3	2
12	lama	3000	30	4	3

6-SQL, Database Constraints, Rules-- Create DB

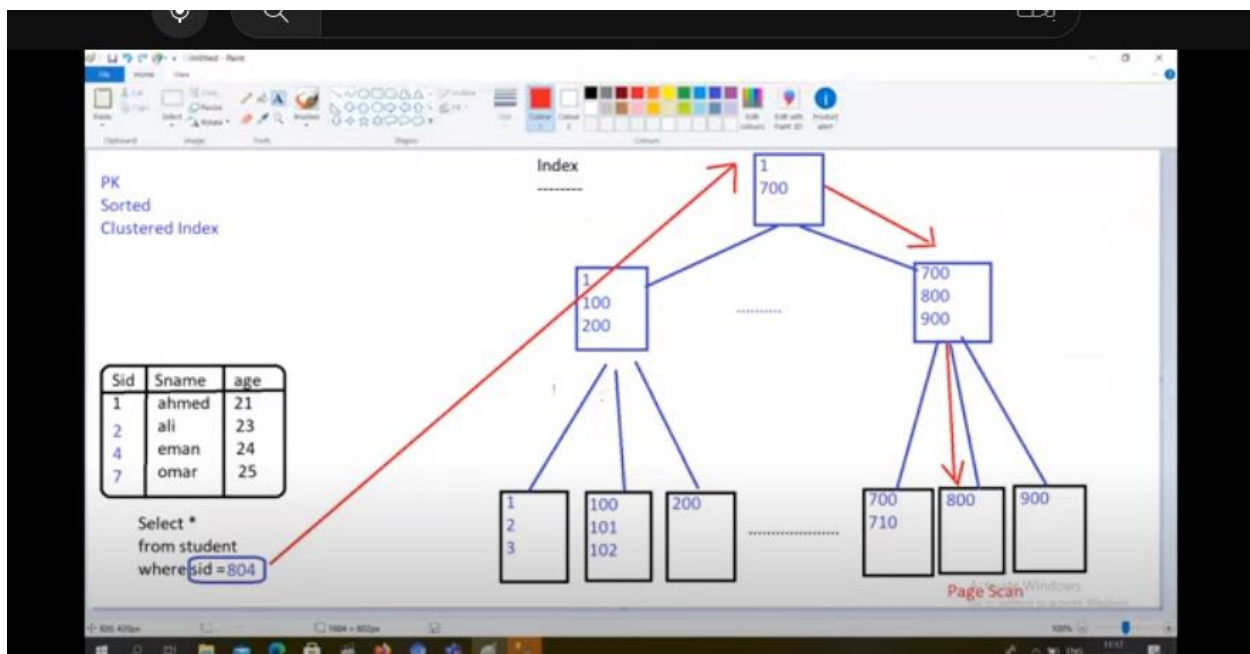


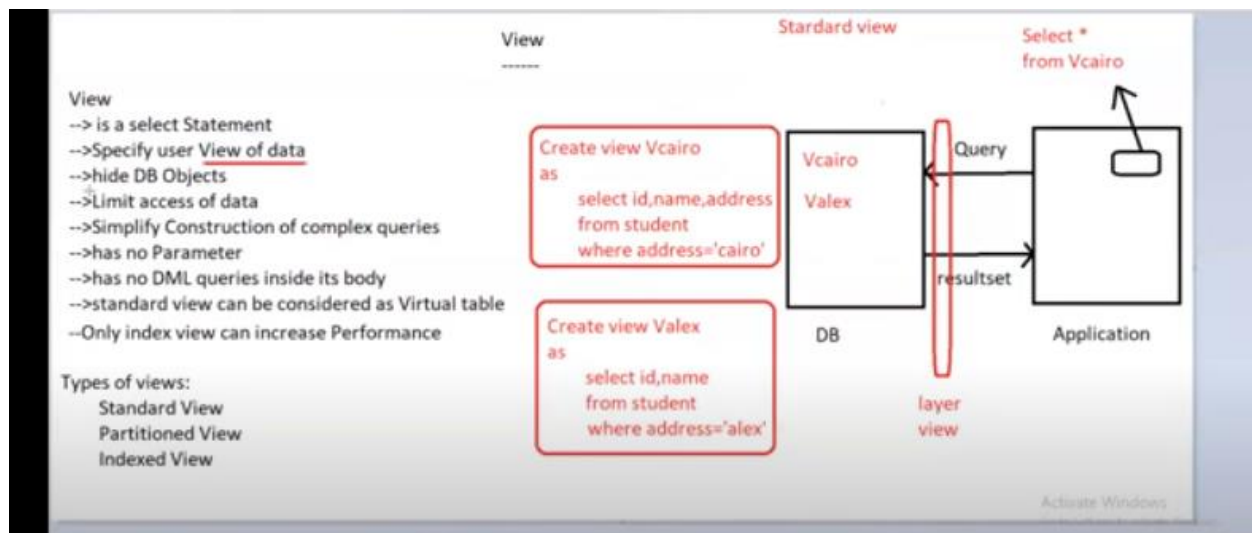
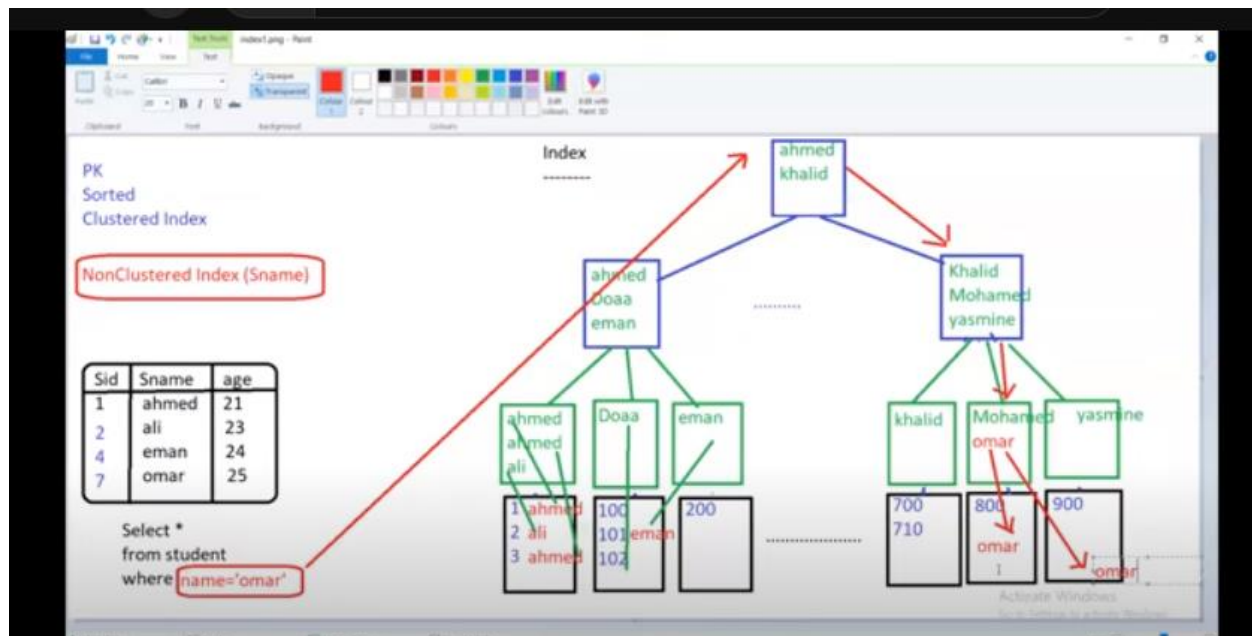
7-SQL, Variables, If, While, functions





8-SQL, View, Index, Merge Statements, Pivot tables.





View

- > is a select Statement
- > Specify user View of data
- > hide DB Objects
- > Limit access of data
- > Simplify Construction of complex queries
- > has no Parameter
- > has no DML queries inside its body
- > standard view can be considered as Virtual table
- > Only index view can increase Performance

Types of views:

- Standard View
- Partitioned View
- Indexed View

Create view VStuds

```
as
select *
from Mans_Server.iti.dbo.students
union all
Select *
From SohagServer.iti2.HR.studs
```

Diagram:

Standard view

Select * from Vcairo

Merge into LastTrasaction as T

using Daily Transaction as S

On T.id = S.did

when Matched then //1,2

update

Set T.myvalue=S.dval

When Not Matched then //10

insert

values(S.did,S.dname,s.dval) ;

LastTransaction

id	name	myvalue
1	ahmed	4000 9000
2	ali	2000 1000
3	omar	6000
4	eman	7000
10	nada	3000

DailyTransaction

did	dname	dval
1	ahmed	9000
2	ali	1000
10	nada	3000

Merge into LastTrasaction as T
 using ^{subquery} Daily Transaction as S
 On T.id = S.did

when Matched and S.dval > T.myvalues then //1
 update
 Set T.myvalue = S.dval

When Not Matched by target then //10
 insert
 values(S.did, S.dname, s.dval)

when not Matched by Source then //3,4
 Delete;

LastTransaction

id	name	myvalue
1	ahmed	9000
2	ali	2000

Taret

10	nada	3000
----	------	------

DailyTransaction

did	dname	dval
1	ahmed	9000
2	ali	1000
10	nada	3000

Source

9-SQL, Stored Procedure, Triggers, XML Tables

Stored Procedures & Triggers

Query → Parsing (syntax) → optimize (Metadata) → Query Tree → Execution Plan

From where Select → Memory

Create Procedure GetSt @id int
 as
 Select *
 From Student
 where id=@id 4

Create Proc InstSt @id int, @name varchar(20)
 as
 if not exists(select id from student where id=@id)
 insert into student(id,name)
 values(@id,@name)
 else
 select 'Duplicate ID'

First Call
 Network Traffic
 InstSt 7, 'ali'

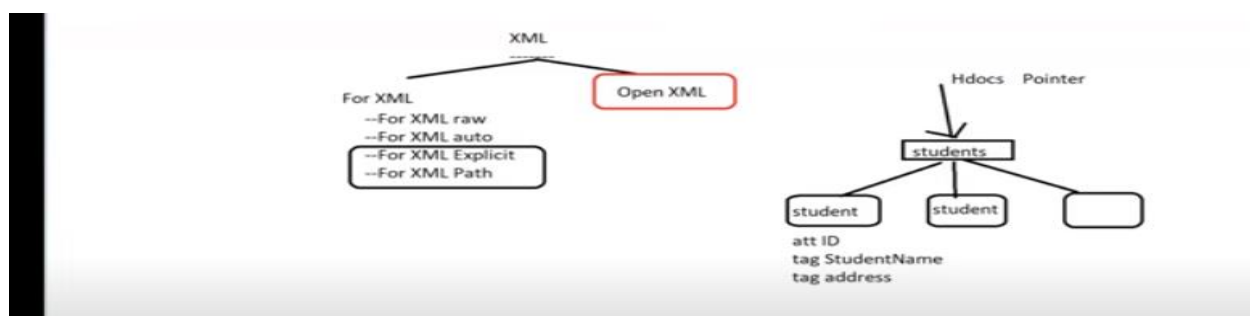
DB

Application

Getst
 InstSt

Query
 Resultset

id 7
 name ali
 btn



10-SQL, Cursor, Database Backups and Restore, SQL Jobs, snapshot, SQL CLR

