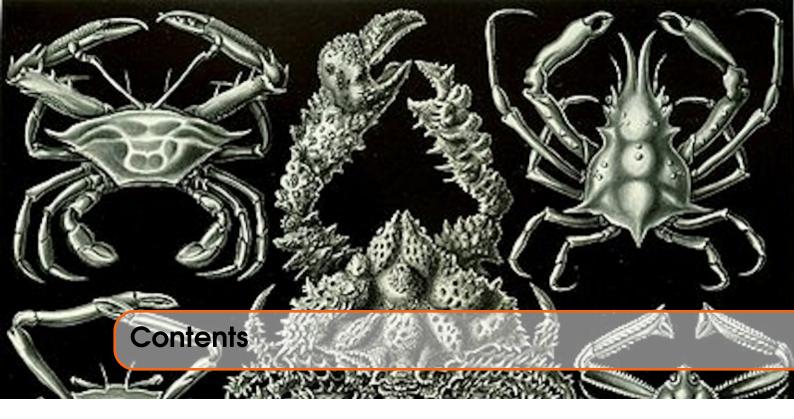


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1.1 Feedback

Please send all feedback to assembly@qdosmsq.dunbar-it.co.uk. You may also send articles to this address, however, please note that anything sent to this email address may be used in a future issue of the eMagazine. Please mark your email clearly if you do not wish this to happen.

This eMagazine is created in LATEX source format, aka plain text with a few formatting commands thrown in for good measure, so I can cope with almost any format you might want to send me. As long as I can get plain text out of it, I can convert it to a suitable source format with reasonable ease.

I use a Linux system to generate this eMagazine so I can read most, if not all, Word or MS Office documents, Quill, Plain text, email etc formats. Text87 might be a problem though!

1.2 Subscribing to The Mailing List

This eMagazine is available by subscribing to the mailing list. You do this by sending your favourite browser to http://qdosmsq.dunbar-it.co.uk/mailinglist and clicking on the link "Subscribe to our Newsletters".

On the next screen, you are invited to enter your email address *twice*, and your name. If you wish to receive emails from the mailing list in HTML format then tick the box that offers you that option. Click the Subscribe button.

An email will be sent to you with a link that you must click on to confirm your subscription. Once done, that is all you need to do. The rest is up to me!

1.3 Contacting The Mailing List

I'm rather hoping that this mailing list will not be a one-way affair, like QL Today appeared to be. I'm very open to suggestions, opinions, articles etc from my readers, otherwise how do I know what I'm doing is right or wrong?

I suspect George will continue to keep me correct on matters where I get stuff completely wrong, as before, and I know George did ask if the list would be contactable, so I've set up an email address for the list, so that you can make comments etc as you wish. The email address is:

assembly@qdosmsq.dunbar-it.co.uk

Any emails sent there will eventually find me. Please note, anything sent to that email address will be considered for publication, so I would appreciate your name at the very least if you intend to send something. If you do not wish your email to be considered for publication, please mark it clearly as such, thanks. I look forward to hearing from you all, from time to time.

If you do have an article to contribute, I'll happily accept it in almost any format - email, text, Word, Libre/Open Office odt, Quill, PC Quill, etc etc. Ideally, a LATEX source document is the best format, because I can simply include those directly, but I doubt I'll be getting many of those! But not to worry, if you have something, I'll hopefully manage to include it.



2.1 QDOS Companion

Derek Stewart has scanned Andy Pennell's *QDOS Companion* book and made it available as a PDF file. Andy gave his blessing to this, in case anyone is worried about copyright. There's a thread on the QL Forum, here¹, which you can follow for updates and corrections etc.

Derek has produced a PDF file containing the book, in as near to its original, format as possible, and the source code is also available as a plain Libre Office ODF file, in case you wish to add corrections and regenerate your own PDF file.

Files can be downloaded from the Sinclair QL account on GitHub, here². To do so, simply click on the file you want to download, then click the "download" button on the screen that appears next. If you are after the PDF, don;t wait for it to display, just click the button – it doesn't need to display to be downloaded.

A number of corrections have been made to the book already.

¹https://qlforum.co.uk/viewtopic.php?f=12&t=4137

²https://github.com/SinclairQL/QDOS-Companion



3.1 Introduction

In this issue of Beginners Corner, we are taking a look at Condition codes and, unavoidably, signed and unsigned numbers and tests. I say 'unavoidably' as I often find myself wondering which of the condition codes are signed and which are unsigned – and it all depends on what type of numbers you are dealing with.

3.2 The Numbers

You need to know about numbers first, before we dive into the condition codes.

Question: If a register holds the byte value 255_{Dec}/FF_{Hex}/111111111_{Bin} is it signed or unsigned?

Answer: Yes!

It can be either because it actually depends on what you, the programmer, wants it to be. In many SMSQ/E trap calls, there needs to be a timeout. This is usually a positive value but if infinite timeout is requested, the value -1 is used. 255 in Decimal is the same bit pattern as -1 in Decimal. So how does the system know the difference between -1 and 255?

The leftmost bit of the number, bit 31, 15 or 7, for long, word or byte values, represents the sign of the number.

If the sign bit is a 1, the number is negative, if it is a zero, the number is positive.

This representation of signed numbers is known as 2's *Compliment*. There is a lot of detail and explanation on Wikipedia¹ – if you are interested.

For unsigned numbers in *B* bits, the range of values is 0 through $2^B - 1$, while for a signed number, with the same number of bits, the range is $-\frac{2^B}{2}$ through $\frac{2^B}{2} - 1$.

¹https://en.wikipedia.org/wiki/Two%27s_complement

From this, we can work out that for 8, 16 and 32 bit numbers, byte, word and long, we get:

- 8 bit bytes, range from 0 through 255 unsigned, -128 through 127 signed;
- 16 bit words, range from 0 through 65,535 unsigned, -32,768 through 32,767 signed;
- 32 bit longs, range from 0 through 4,293,967,295 unsigned, -2,147,483,648 through 2,147,483,647 signed.

When you look at a number in binary, it's easy to determine the decimal value simply by adding up all the powers of two where there is a 1 bit present, for example:

10100101 is
$$2^7 + 2^5 + 2^2 + 2^0 = 165$$
.

If the number is signed, then it's slightly more complicated as there are at least three different methods.

Option 1: You take the leftmost power of two and subtract all the remaining powers of two where a 1 bit is present. For example:

• 10100101 is $2^7 - (2^5 + 2^2 + 2^0) = -91$.

Option 2: Flip all the bits, convert to an unsigned value, add 1 and change the sign to make the value negative, for example:

- 10100101 flipped is 01011010;
- Converted to an unsigned value gives $2^6 + 2^4 + 2^3 + 2^1 + 1 = 91$;
- Changing the sign gives -91 as before.

Option 3: Calculate the unsigned value, and subtract from 2^B where B is the number of bits, then change the sign to make the value negative. For example:

- 10100101 is still $2^7 + 2^5 + 2^2 + 2^0 = 165$;
- The number of bits is 8 and we know 2^8 is 256;
- Subtract 165 from 256 to get 91;
- Changing the sign gives -91, yet again!

Pick the version you prefer.

Right, that's enough about numbers, signed or otherwise, for now at least!

3.3 The Flags

So, moving slightly ahead, we now need to think about the Status Register, often known as SR, where the MC680xx CPUs hold their status flags. Some of these are available to the programmer for use in code to determine the outcome of a calculation, a register load, a test and so on.

If you need more information, grab hold of my eBook², based on many years of writing Assembly Language articles for the, sadly defunct, *QL Today* magazine. That will explain all.

In the meantime, all you need to know at the moment is that various flags in the SR are set or cleared according to the results of some operation. If the flags are set as the result of some arithmetic operation, or by loading a value into a register, for example, then some flags indicate whether the number is signed or unsigned – in a roundabout way!

Ok, I lied, slightly, sorry. What the flags indicate is whether the number is positive, negative – amongst others – regardless of whether you intended the number to be signed or unsigned. Table

²https://github.com/NormanDunbar/QLAssemblyLanguageBook/releases/latest

3.1 shows the full set of flags and Table 3.2 shows what they are set to for certain conditions which can be tested.

Flag	Flag Name	Description		
X	Extend	Usually the same as Carry, but used separately. (The manual is no		
		help!)		
N	Negative	Indicates if the top bit is set or clear.		
Z	Zero	Indicates the last operation resulted in a zero result.		
V	Overflow	The last operation caused arithmetic overflow.		
С	Carry	The last operation resulted in a carry.		

Table 3.1: MC680xx Status Register

Condition	Flags	Description	Signed
CC	C=0	Carry Clear	U
CS	C=1	Carry Set	U
EQ	Z=1	Equal/Zero	U and S
NE	Z=0	Not Equal/Not Zero	U and S
HI	C=0 and Z=0	Higher	U
LS	C=1 or Z=1	Lower or Same	U
PL	N=1	Plus/Positive	S
MI	N=1	Minus/Negative	S
GE	(N=1 and V=1) or	Greater Than or Equal	S
	(N=0 and V=0)		
GT	(N=1 and V=1 and Z=0) or	Greater Than	S
	(N=0 and V=0 and Z=0)		
LE	(Z=1) or	Less Than or Equal	S
	(N=1 and V=0) or		
	(N=0 and V=1)		
LT	(N=1 and V=0) or	Less Than	S
	(N=0 and V=1)		

Table 3.2: MC680xx Condition Codes

3.4 Condition Codes

Signed and unsigned condition codes.

Bcc, DBcc, CMP. Explain.

Signed: EQ, GE, GT, LS, LT, MI, NE, PL,

Unsigned: EQ, HI, LE, NE,

Others: F, T, VC, VS,

3.5 Assembling the Code

Assembling the code requires that you have an assembler installed. See Section 3.7, Tools and Manuals for links to allow you to obtain the assemblers that I frequently use for this eMagazine.

Once you have installed an assembler, assembling the code is simple and is described in the following sections. Pick the section for your own assembler.

3.5.1 With GWASS

- EXEC gwass60_bin to start the assembler;
- Select option 1 to start assembling;
- Type in the filename: ram1_????????_asm
- Wait.

3.5.2 With Qmac

To pass the commands directly via the command line:

```
• EX qmac; "ram1_????????asm -data 2048 -filetype 1 -nolink -bin ram1_????????_bin"
```

Note: The above command should be typed on one line - I've had to split it for the PDF page width.

Alternatively, you can type the command interactively:

- EX qmac
- Type the options: ram1_????????_asm -data 2048 -filetype 1 -nolink -bin ram1_????????_bin
- Wait

What you are doing here, in both cases, is telling the assembler to:

- Assemble the source file ram1_????????_asm;
- Create an executable file (-filetype 1), with 2,048 bytes of data space (-data 2048);
- Do not invoke the linker as it is not needed because everything is in the same source file (-nolink);
- Create the output file named ram1_????????_bin which will be in uppercase regardless of what you type here!

3.5.3 Executing the Code

After a successful assemble, and regardless of which assembler you used, ram1_????????pin will be the executable job. To run it:

```
• EX ram1_?????????_bin
```

On a successful execution,

3.6 Summary

That, hopefully, will help make writing structured Assembly Language programs more manageable. Let me know if there's anything you didn't follow or which may need more explanation, and I'll do my best to assist.

3.7 Tools and Manuals

Get hold of the SMSQ/E Reference Manual from Wolfgang Lenerz's web site³ for the official version. Alternatively, there are copies on Dilwyn's pages:

- Here⁴ for the PDF for version 4.5; or
- Here⁵ for the ODT (Libre Office) file for version 4.5.
- If you want to ensure that you have the most recent versions of those files, Wolfgang Lenerz's web site⁶ is the place to look.

Get hold of GWASS here⁷ for 68020 processors.

Download Qmac here⁸ for 68008 processors.

³https://www.wlenerz.com/qlstuff/#qdosms

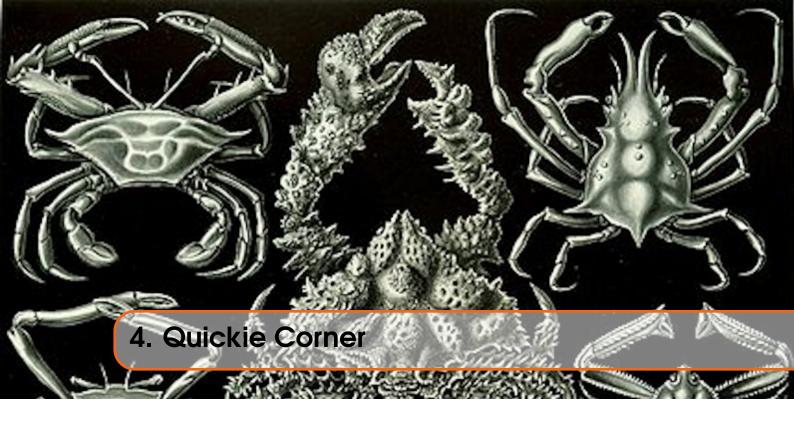
⁴http://www.dilwyn.me.uk/docs/manuals/QDOS_SMS%20Reference%20Guide%20v4.5.pdf

⁵http://www.dilwyn.me.uk/docs/manuals/QDOS_SMS%20Reference%20Guide%20v4.5.odt

⁶https://www.wlenerz.com/qlstuff/#qdosms

⁷http://gwiltprogs.info/gwassp22.zip

⁸http://www.dilwyn.me.uk/asm/gst/gstmacroquanta.zip



Following on from the resounding success¹ of last issue's Quickie Corner in which we found the fastest method of clearing a register to zero, this time we look at finding the quickest way to tell if a data register is zero, or not.

4.1 It's All About Nothing!

What's the fastest way to check if a data register is zero? The timings are taken from the Motorola Semiconductor's MC68000 Programmers Reference Guide, 4th Edition, Appendix E.

Table 4.1 shows a few instructions which can either test a data register for zero, or, cause the Z flag to be set if zero, and their timings.

- Dn is the data register in question.
- Dz is a data register containing zero.
- Dy is "just another data register"

It appears that TST.L Dn is the quickest at 4 clock cycles, closely followed by CMP.L Dn,Dz, NEG.L Dn, ADD.L Dz,Dn and SUB.L Dz,Dn all at 6 clocks.

¹By which I mean, nobody complained!

Instruction	Clock Cycles
CMPI.L #0,Dn	12
CMP.L Dn,Dz	6
TST.L Dn	4
ORI.L #0,Dn	14
OR.L Dn,Dn	22
ANDI.L Dn,Dn	14
AND.L Dn,Dn	22
MOVE.L Dn,Dy	4
NEG.L Dn	6
ADDQ.L #0,Dn	8
ADDI.L #0,Dn	14
ADD.L Dz,Dn	6
SUBI.L #0,Dn	14
SUB.L Dz,Dn	6
SUBQ.L #0,Dn	8

Table 4.1: Checking a data register for zero



Years ago I wrote an article for *QL Today* about the Maths Stack in QDOS and how the value in A1, on entry to a function or procedure, was *not* pointing to the top of the maths stack, as documented in the various QDOS manuals, and books around at that time. In brief, I discovered that on entry, we have three possibilities:

• If A1 is *negative*, it means that the function/procedure in question has been called as part of an expression, and A1 shows how much space has been used on the stack for the expression, so far. The code would resemble this:

```
PRINT 1234 * myFunction(...)
```

- If A1 is *positive*, it means that the maths stack has some free space available for use without any requirement to allocate more.
- If A1 is *zero*, it means that the function/procedure in question has been called on it's own, or at the beginning of an expression. A1 shows how much space has been used on the stack for the expression, so far. The code would resemble this: The code would resemble this:

```
PRINT myFunction(...)
```

or

```
PRINT myFunction(...) * 1234
```

or any expression where myFunction() is at the beginning.

This is covered in my eBook which you can download from GitHub¹ – section 7.8 is the location you will need.

¹https://github.com/NormanDunbar/QLAssemblyLanguageBook/releases/latest

There was a recent exercise by Derek Stewart to scan Andy Pennell's *QDOS Companion* book and make it available as a PDF. Andy did give his blessing to this, in case anyone is worried about copyright. There's a thread on the QL Forum, here², which shows the process and updates so far. See this issues News chapter for download details.

As part of the ensuing discussion, it was mentioned that Andy's book didn't correctly document A1 when entering a procedure or function. I responded with the above information and was brought before the select committee³ to explain:

- You don't say whether your findings apply to: JM, JS, Minerva, SMSQ/E etc. It could, in other words just be random.
- I'm no authority, but I would advise to not rely on this unless it were documented to be valid across the board, and to be there for a particular reason, ie as a result of the way things by necessity are done, or as a particular service to keyword writers. I wont go into reasons as I hope they are self evident.

Which are excellent points of order.

Let's dive into the maelstrom then!

5.1 The Process Outline

In order to test I will be following these few steps:

- Create a function with no parameters, and another with one parameters, as the test code. A TRAP #15 instruction will be placed at the very start of the function code. On normal running, this has no effect, but when I do some fiddling in the commands of QMON2, it will immediately jump into QMON2 and I can debug from there.
- When the debugger has control, examine the registers paying close attention to A1.
- The code will be tested on SMSQ/E under QPC2. It has already been tested in QDOS under the JS ROM, back in the day. I don't have Minerva or other ROMS, but I do have emulators which do!
- The functions will be tested stand-alone, and as part of an expression where it occurs at the beginning and another where it appears after the beginning.
- The functions will be tested in integer, float and string (coercion) expressions.

5.2 The Test Functions

Listing 5.1 is the code I'll be using for the test function FN_0 and Listing 5.2 is the code for function FN_1. The function names match the number of parameters they take. I have not shown the code that links the functions into S*BASIC, but this is present in the code download for this issue.

There's nothing special about the two functions, other than the TRAP #15 instructions which appear at the start of each function. This causes execution to jump into QMON2, if that has been loaded and set up correctly. FN_0 returns a word integer of zero as it's result, while FN_1 returns its parameter, incremented by 1, as another word integer.

```
    39 ;
    40 ; FN_0 takes no parameters and returns zero as an integer word
    41 ; of two bytes.
```

²https://qlforum.co.uk/viewtopic.php?f=12&t=4137

³Per Witte

```
fn0

trap #15 ; Jump into QMON2

moveq #0,d7 ; Result in D7

moveq #2,d6 ; How much space do I need?

bra.s fnResult ; Just return a result
```

Listing 5.1: Test Function: FN_0

```
50
   ; FN_1 takes one integer word parameter and returns it +1 as an
     integer word of two bytes. No validation here before fetching
51
   ; the parameter(s), but we do check for fetching only one.
53
54
   fn1
55
        trap #15
                                      ; Jump into QMON2
56
       move.w ca_gtint, a2
                                      ; Fetch word integers only
57
        jsr (a2)
                                      : Do it
58
        tst.1 d0
                                      : Ok?
59
                                      ; Yes, carry on
        beq.s fn1Check
60
                                      ; No, bale out
        rts
61
   fn1Check
62
63
        cmpi.w #1,d3
                                     ; How many? We need 1
        beq.s fn1Got1
                                     ; Yes, carry on
64
                                     ; Bad parameter
65
        moveq #err.ipar,d0
66
                                      ; Error out
        rts
67
68
   fn1Got1
69
        move.w 0(a6, a1.1), d7
                                      ; Fetch the parameter
70
        addq.w #1,d7
                                      ; Increment for return
71
        moveq #0,d6
                                      ; No space required
```

Listing 5.2: Test Function: FN_1

Obviously, as FN_0 takes no parameters, we need to allocate space on the maths stack before we can return a result. We do this by fetching the current maths stack pointer from SV_ARTHP(A6) into A1, and requesting D6.W bytes of space for the result. In this case, D6.W is 2 as we are returning a word integer. After the space has been allocated, it is possible that the maths stack, plus its contents, has been moved, so we need to refresh A1 from SV_ARTHP(A6).

FN_1, on the other hand, takes a word integer parameter and as such, has already got enough space on the maths stack to return its result. For FN_1 then, D6.W is zero to indicate that we don't need any space so we simply store the value in D7.W into the word pointed to by (A6,A1.L).

Listing 5.3 is the code that handles the returning of results and the allocation of maths stack space, as required, in order to do so.

```
; This code returns the function results. For FN_0 it has to
; allocate two bytes but for FN_1, there's already space as we
; used two bytes for the parameter. The result is in D7 and D6
; holds the space we need to allocate on the stack for the
; result.
;
fnResult
tst.w d6; Do I need space allocated?
```

```
82
                                       ; No, use existing space
        beq.s rtnFn1
83
        move. 1 sv_arthp(a6), a1
                                       ; Yes, get the stack pointer
84
        move.w d6, d1
                                       ; Space needed for result
85
        move.w qa.resri, a2
                                       ; Allocation vector
86
                                       ; Allocate - will not error out
        jsr (a2)
                                       ; Possible new maths stack
87
        move. 1 sv_arthp(a6), a1
88
        subq.1 #2, a1
                                       ; Make room for result
89
90
   rtnFn1
91
        move.w d7,0(a6,a1.1)
                                      ; Stack the result
92
        moveq #3, d4
                                       ; Signal word integer result
93
        move. 1 a1, sv_arthp(a6)
                                       ; Save top of stack
94
                                       ; No errors
        moveq #0,d0
95
        r t s
                                       ; Back to S*BASIC
96
                                        ; Uncomment for QMC assembler
97
         end
```

Listing 5.3: Returning results

5.3 Debugging with QMON2

After loading the QMON2 binary, we need to tell it to execute when a TRAP #15 instruction is executed. to do this we simply execute QMON2 in S*BASIC channel #1 – so that when it executed, it comes up in the same channel – and then tell it to intercept those TRAP #15 instructions. This is done as follows:

```
QMON #1

Qmon > TL 14

Qmon > g
```

That's all there is to do. Now whenever a TRAP #15 instruction is executed, QMON2 will intercept it, and break execution, giving us full control over the system at the instruction immediately after the TRAP #15 instruction.

The TL command in QMON2 will execute QMON2 every time a TRAP *higher* than that indicated, is executed.

After setting up QMON2, testing was done by calling PRINT with various expressions as its parameters, each using the appropriate test function in one of three places:

- As the only term in the expression.
- As the first term in the expression.
- As the final term in the expression.

My debugging session looked something like this:

```
X=1234
X% = 1234
x$="1234"

PRINT FN_0
PRINT FN_0 * X
```

5.4 Results 23

```
PRINT FN_O * X%

PRINT X * FN_O

PRINT X% * FN_O

PRINT X$ * FN_O

PRINT FN_1

PRINT FN_1 * X

PRINT FN_1 * X%

PRINT FN_1 * X%

PRINT FN_1 * X$

PRINT X * FN_1

PRINT X * FN_1

PRINT X * FN_1

PRINT X * FN_1

PRINT X * FN_1
```

5.4 Results

On SMSQ/E the results are slightly different from that in QDOS with the JS ROM. There was no difference when the function called was FN_0 or FN_1, they both showed the same values in A1.L when QMON2 intercepted execution. Table 5.1 shows the results of testing on SMSQ/E.

Test	A1 (Hex)	A1 (Decimal)
Function only	0	0
Function * integer	0	0
Function * float	0	0
Function * string	0	0
Integer * function	FFFFFFC	-4
Float * function	FFFFFF8	-8
String * function	FFFFFFA	-6

Table 5.1: Test function results

I repeated the tests with other arithmetic operators, the results are the same as Table 5.1 regardless of the operator in question.

Looking at the figures, it appears that on SMSQ/E, the value in A1 at the start of a procedure or function is 2 bytes greater than that used by QDOS under the same conditions.

- When A1 is zero, the SMSQ/E result is the same as the QDOS: JS result, no space yet used on the maths stack.
- When A1 is negative, the SMSQ/E result appears to show an extra 2 bytes over the size of the space used on the maths stack so far. So for a word sized integer, 2 bytes, the value in A1 is -4, for a 6 byte float it's -8 and for a string, which was never tested on QDOS: JS, it's always -6 regardless of the length of the string.
- I was unable to remember⁴ how I managed to get a positive value in A1 back in the QDOS: JS days, and I was unable to get one in SMSQ/E.

⁴I'm getting old. When I learn new stuff these days, something else has to be forgotten to make room for it!

5.5 Summary

Even without testing Minerva ROMs under QDOS, it is obvious that there are indeed differences – at least between QDOS and SMSQ/E. For this reason, it's best to ignore what I've said in the past, and simply treat A1 on entry to a procedure or function as *not being a suitable value for the top of the maths stack* so it must always be loaded from BV_RIP(A6) on QDOS or from SV_ARTHP(A6) on SMSQ/E.

- If your function takes no parameters, A1 *must* be loaded before attempting to request maths stack space to return the result.
- If your function takes any parameters, then until such time as you call the appropriate vector to fetch the parameters, A1 is not usable as the maths stack top. After fetching the parameters, it will be correctly set.



The front cover image on this ePeriodical is taken from the book *Kunstformen der Natur* by German biologist Ernst Haeckel. The book was published between 1899 and 1904. The image used is of various *Decapods*. The Decapoda or decapods (literally "ten-footed") are an order of crustaceans within the class Malacostraca, including many familiar groups, such as crabs, lobsters, crayfish, shrimp and prawns. Most decapods are scavengers. The order is estimated to contain nearly 15,000 species in around 2,700 genera, with around 3,300 fossil species.

I have also cropped the cover image for use on each chapter heading page.

You can read about Decapods on Wikipedia and there is a brief overview of the above book, also on Wikipedia, which shows a number of other images taken from the book. (Some of which I considered before choosing the current one!)

Decapods have absolutely nothing to do with the QL or computing in general - in fact, I suspect many of them died out before electricity was invented, and the rest probably don't care about electricity or computers! However, I liked the image, and decided that it would make a good cover for the book and a decent enough chapter heading image too.

Not that I am suggesting, in any way whatsoever, that we QL fans are 10 legged crustaceans.