

IF120 Discrete Mathematics

09 Recurrence Relation

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REVIEW

- Discrete Probability
- Binomial Coefficient
- Combinatorial Identity

<u>OUTLINE</u>

- Recursive Algorithm
- Recurrence Relations
- Solving Recurrence Relations

Recursive Algorithm

- A recursive function is a function that invokes itself.
- A recursive algorithm is an algorithm that contains a recursive function.
- Recursion is a powerful, elegant, and natural way to solve a large class of problems.
- A problem in this class can be solved using a divide-and-conquer technique in which the problem is decomposed into problems of the same type as the original problem.

Recursive Algorithm

 \square A robot can take steps of I meter or 2 meters. We write an algorithm to calculate the number of ways the robot can walk n meters. As examples:

Distance	Sequence of Steps	Number of Ways to Walk
1	1	1
2	1,1 or 2	2
3	1, 1, 1 or 1, 2 or 2, 1	3
4	1, 1, 1, 1 or 1, 1, 2	5
	or 1, 2, 1 or 2, 1, 1 or 2, 2	

Let walk(n) denote the number of ways the robot can walk n meters. We have observed that

$$walk(1) = 1, walk(2) = 2.$$

Now suppose that n > 2. The robot can begin by taking a step of I meter or a step of 2 meters. If the robot begins by taking a I-meter step, a distance of n-1 meters remains; but, by definition, the remainder of the walk can be completed in walk(n-1) ways.

Similarly, if the robot begins by taking a 2-meter step, a distance of n-2 meters remains and, in this case, the remainder of the walk can be completed in walk(n-2) ways. Since the walk must begin with either a 1-meter or a 2-meter step, all of the ways to walk n meters are accounted for. We obtain the formula

Robot Walking

This algorithm computes the function defined by

$$walk(n) = \begin{cases} 1, & n = 1 \\ 2, & n = 2 \\ walk(n-1) + walk(n-2) & n > 2. \end{cases}$$

```
Input: n

Output: walk(n)

walk(n) {

if (n == 1 \lor n == 2)

return n

return walk(n - 1) + walk(n - 2)
}
```

```
walk(n)
= walk(n-1) + walk(n-2).
For example,
walk(4) = walk(3) + walk(2)
= 3 + 2 = 5
We can write a recursive algorithm to
compute walk(n) by translating the
equation
walk(n)
= walk(n-1) + walk(n-2)
directly into an algorithm.
The base cases are n = 1 and n = 2.
```

Recurrence Relations

- Definition 9.1:A **recurrence relation** for the sequence a_0, a_1, \ldots is an equation that relates a_n to certain of its predecessors $a_0, a_1, \ldots, a_{n-1}$.
- Initial conditions for the sequence a_0, a_1, \ldots are explicitly given values for a finite number of the terms of the sequence.
- ☐ The Fibonacci sequence is defined by the recurrence relation

$$f_n = f_{n-1} + f_{n-2}, \qquad n \ge 3,$$

and initial conditions

$$f_1 = 1$$
, $f_2 = 1$.

Let S_n denote the number of subsets of an n-element set. Since going from an (n-1)-element set to an n-element set doubles the number of subsets (remember the Set Theory), we obtain the recurrence relation

$$S_n = 2S_{n-1}$$
.

The initial condition is

$$S_0 = 1.$$

Recurrence Relations

 \square A person invests \$1000 at 12 percent interest compounded annually. If A_n represents the amount at the end of n years, find a recurrence relation and initial conditions that define the sequence $\{A_n\}$.

At the end of n-1 years, the amount is A_{n-1} . After one more year, we will have the amount A_{n-1} plus the interest. Thus

$$A_n = A_{n-1} + (0.12)A_{n-1} = (1.12)A_{n-1}, \qquad n \ge 1.$$

Computing Compound Interest

This recursive algorithm computes the amount of money at the end of n years assuming an initial amount of \$1000 and an interest rate of 12 percent compounded annually.

```
Input: n, the number of years

Output: The amount of money at the end of n years
```

```
    compound_interest(n) {
    if (n == 0)
    return 1000
    return 1.12 * compound_interest(n = 1)
    }
```

To apply this recurrence relation for n = 1, we need to know the value of A_0 .

Since A_0 is the beginning amount, we have the initial condition

$$A_0 = 1000.$$

Solving Recurrence Relations

- To solve a recurrence relation involving the sequence a_0, a_1, \ldots is to find an explicit formula for the general term a_n .
- In this section we discuss two methods of solving recurrence relations:
- I. iteration and
- 2. linear homogeneous recurrence relations with constant coefficients.
- To solve a recurrence relation involving the sequence a_0, a_1, \ldots by **iteration**, we use the recurrence relation to write the *n*th term a_n in terms of certain of its predecessors a_{n-1}, \ldots, a_0 .
- We then successively use the recurrence relation to replace each of a_{n-1} , ... by certain of their predecessors.
- We continue until an explicit formula is obtained.

Iteration Method

■We can solve the recurrence relation

$$a_n = a_{n-1} + 3$$
, (9.1)

subject to the initial condition

$$a_1 = 2$$
,

by iteration. Replacing n by n - 1 in (9.1), we obtain

$$a_{n-1} = a_{n-2} + 3.$$

If we substitute this expression for a_{n-1} into (9.1), we obtain

$$a_n = a_{n-2} + 3 + 3$$

= $a_{n-2} + 2 \cdot 3$. (9.2)

Replacing n by n-2 in (9.1), we obtain

$$a_{n-2} = a_{n-3} + 3.$$

If we substitute this expression for a_{n-2} into (9.2), we obtain

$$a_n = a_{n-3} + 3 + 2 \cdot 3$$

= $a_{n-3} + 3 \cdot 3$.

Iteration Method

In general, we have

$$a_n = a_{n-k} + k \cdot 3.$$

If we set k = n - 1 in this last expression, we have

$$a_n = a_1 + (n - 1) \cdot 3.$$

Since aI = 2, we obtain the explicit formula

$$a_n = 2 + 3(n - 1)$$

for the sequence a.

We can solve the recurrence relation

$$S_n = 2S_{n-1}$$

subject to the initial condition

$$S_0 = 1$$
,

by iteration:

$$S_n = 2S_{n-1} = 2(2S_{n-2}) = \dots = 2^n S_0 = 2^n$$
.

Iteration Method

Assume that the deer population of Rustic County is 1000 at time n = 0 and that the increase from time n - 1 to time n is 10 percent of the size at time n - 1. Write a recurrence relation and an initial condition that define the deer population at time n and then solve the recurrence relation.

Let d_n denote the deer population at time n. We have the initial condition $d_0 = 1000$.

The increase from time n-1 to time n is d_n-d_{n-1} . Since this increase is 10 percent of the size at time n-1, we obtain the recurrence relation

$$d_n - d_{n-1} = 0.1 \ d_{n-1},$$

which may be rewritten

$$d_n = 1.1d_{n-1}.$$

The recurrence relation may be solved by iteration:

$$d_n = 1.1d_{n-1} = 1.1(1.1d_{n-2}) = (1.1)^2 d_{n-2}$$

= $\cdots = (1.1)^n d_0 = (1.1)^n 1000$.

The assumptions imply exponential population growth.

 Definition 9.2: A linear homogeneous recurrence relation of order k with constant coefficients is a recurrence relation of the form

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + ... + c_k a_{n-k}, c_k \neq 0.$$

• Notice that a linear homogeneous recurrence relation of order k with constant coefficients, together with the k initial conditions

$$a_0 = C_0, a_1 = C_1, \dots, a_{k-1} = C_{k-1},$$

uniquely defines a sequence a_0, a_1, \dots

The recurrence relations

$$S_n = 2S_{n-1}$$

and

$$f_n = f_{n-1} + f_{n-2},$$

which defines the Fibonacci sequence, are both linear homogeneous recurrence relations with constant coefficients. The first recurrence relation is of **order 1** and the second is of **order 2**.

☐ The recurrence relation

$$a_n = 3a_{n-1}a_{n-2}$$

is **not** a linear homogeneous recurrence relation with constant coefficients. In a linear homogeneous recurrence relation with constant coefficients, each term is of the form ca_k . Terms such as $a_{n-1}a_{n-2}$ are not permitted. Recurrence relations such as the example above are said to be **nonlinear**.

The recurrence relation

$$a_n - a_{n-1} = 2n$$

is **not** a linear homogeneous recurrence relation with constant coefficients because the expression on the right side of the equation is not zero. (Such an equation is said to be **inhomogeneous**.)

The recurrence relation

$$a_n = 3na_{n-1}$$

is **not** a linear homogeneous recurrence relation with constant coefficients because the coefficient 3n is not constant. It is a linear homogeneous recurrence relation with nonconstant coefficients.

■ Theorem 9.1: Let

$$a_n = c_1 a_{n-1} + c_2 a_{n-2}$$
 (9.3)

be a second-order, linear homogeneous recurrence relation with constant coefficients. If S and T are solutions of (9.3), then U = bS + dT is also a solution of (9.3). If r is a root of

$$t^2 - c_1 t - c_2 = 0, (9.4)$$

then the sequence r^n , n = 0, 1, ..., is a solution of (9.3).

If a is the sequence defined by (9.3),

$$a_0 = C_0, \ a_1 = C_1,$$
 (9.5)

and r_1 and r_2 are roots of (9.4) with $r_1 \neq r_2$, then there exist constants b and d such that

$$a_n = br_1^n + dr_2^n, \qquad n = 0,1,....$$

Assume that the deer population of Rustic County is 200 at time n = 0 and 220 at time n = 1 and that the increase from time n = 1 to time n is twice the increase from time n = 2 to time n = 1. Write a recurrence relation and an initial condition that define the deer population at time n and then solve the recurrence relation.

Let d_n denote the deer population at time n. We have the initial conditions

$$d_0 = 200, d_1 = 220.$$

The increase from time n-1 to time n is d_n-d_{n-1} , and the increase from time n-2 to time n-1 is $d_{n-1}-d_{n-2}$. Thus we obtain the recurrence relation

$$d_n - d_{n-1} = 2 (d_{n-1} - d_{n-2}),$$

which may be rewritten

$$d_n = 3d_{n-1} - 2d_{n-2}.$$



To solve this recurrence relation, we first solve the quadratic equation

$$t^2 - 3t + 2 = 0$$

to obtain roots I and 2. The sequence d is of the form

$$d_n = b \cdot 1^n + c \cdot 2^n = b + c2^n$$
.

To meet the initial conditions, we must have

$$200 = d_0 = b + c$$
, $220 = d_1 = b + 2c$.

Solving for b and c, we find that b = 180 and c = 20. Thus d_n is given by $d_n = 180 + 20 \cdot 2^n$.

The growth is exponential.

☐ Find an explicit formula for the Fibonacci sequence.

The Fibonacci sequence is defined by the linear homogeneous, second-order recurrence relation

$$f_n - f_{n-1} - f_{n-2} = 0, \qquad n \ge 3,$$

and initial conditions

$$f_1 = 1$$
, $f_2 = 1$.

We begin by using the quadratic formula to solve

$$t^2 - t - 1 = 0.$$

The solutions are

$$t = \frac{1 \pm \sqrt{5}}{2}.$$

Thus the solution is of the form

$$f_n = b \left(\frac{1 + \sqrt{5}}{2} \right)^n + d \left(\frac{1 - \sqrt{5}}{2} \right)^n.$$



To satisfy the initial conditions, we must have

$$b\left(\frac{1+\sqrt{5}}{2}\right) + d\left(\frac{1-\sqrt{5}}{2}\right) = 1$$

$$b\left(\frac{1+\sqrt{5}}{2}\right)^2 + d\left(\frac{1-\sqrt{5}}{2}\right)^2 = 1.$$

Solving these equations for b and d, we obtain

$$b = \frac{1}{\sqrt{5}}, \qquad d = -\frac{1}{\sqrt{5}}.$$

Therefore, an explicit formula for the Fibonacci sequence is

$$f_n = \frac{1}{\sqrt{5}} \left(\frac{1 + \sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left(\frac{1 - \sqrt{5}}{2} \right)^n$$

Surprisingly, even though f_n is an integer, the preceding formula involves the irrational number $\sqrt{5}$.

■ Theorem 9.2: Let

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} (9.6)$$

be a second-order linear homogeneous recurrence relation with constant coefficients.

Let a be the sequence satisfying (9.6) and

$$a_0 = C_0$$
, $a_1 = C_1$

If both roots of

$$t^2 - c_1 t - c_2 = 0, (9.7)$$

are equal to r, then there exist constants b and d such that

$$a_n = br^n + dnr^n$$
, $n = 0,1,...$

□ Solve the recurrence relation

$$d_n = 4(d_{n-1} - d_{n-2}) (9.8)$$

subject to the initial conditions

$$d_0 = 1 = d_1$$
.

According to Theorem 9.1, $S_n = r^n$ is a solution of (9.8), where r is a solution of $t^2 - 4t + 4 = 0$. (9.9)

Thus we obtain the solution

$$S_n = 2^n$$

of (9.8). Since 2 is the only solution of (9.9), by Theorem 9.2, $T_n = n2^n$.

is also a solution of (9.8). Thus the general solution of (9.8) is of the form U = aS + bT.

We must have

$$U_0 = 1 = U_1$$
.

These last equations become

$$aS_0 + bT_0 = a + 0b = 1,$$
 $aS_1 + bT_1 = 2a + 2b = 1.$

Solving for a and b, we obtain

$$a = 1,$$
 $b = -\frac{1}{2}.$

Therefore, the solution of (9.8) is

$$d_n = 2^n - n2^{n-1}.$$

PRACTICE

PRACTICE I

I. Use the formulas

$$s_1 = 1, s_n = s_{n-1} + n$$
 for all $n \ge 2$, to write a recursive algorithm that computes $s_n = 1 + 2 + 3 + \cdots + n$.

2. Give a proof using mathematical induction that your algorithm for part (1) is correct.

PRACTICE 2

- In Exercises below, find a recurrence relation and initial conditions that generate a sequence that begins with the given terms.
- 1. 3, 7, 11, 15, ...
- 2. 3, 6, 9, 15, 24, 39, ...

PRACTICE 3

- In Exercises below, solve the given recurrence relation for the initial conditions given.
- 1. $a_n = 2na_{n-1}$; $a_0 = 1$
- 2. $n = 6a_{n-1} 8a_{n-2}$, $a_0 = 1 = a_1$
- 3. $a_n = 7a_{n-1} 10a_{n-2}$, $a_0 = 5$, $a_1 = 16$

NEXT WEEK'S OUTLINE

- Graph Terminologies
- Path and Cycle
- Euler Cycle and Hamiltonian Cycle

<u>REFERENCES</u>

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- Rosen, Kenneth H., 2005, Discrete Mathematics and Its Applications, 6th edition, McGraw-Hill.
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- Lipschutz, Seymour, Lipson, Marc Lars, Schaum's Outline of Theory and Problems of Discrete Mathematics, McGraw-Hill.
- Liu, C.L., 1995, Dasar-Dasar Matematika Diskret, Jakarta: Gramedia Pustaka Utama.
- Other offline and online resources.

Visi

Menjadi Program Studi Strata Satu Informatika **unggulan** yang menghasilkan lulusan **berwawasan internasional** yang **kompeten** di bidang Ilmu Komputer (*Computer Science*), **berjiwa wirausaha** dan **berbudi pekerti luhur**.



Misi

- . Menyelenggarakan pembelajaran dengan teknologi dan kurikulum terbaik serta didukung tenaga pengajar profesional.
- 2. Melaksanakan kegiatan penelitian di bidang Informatika untuk memajukan ilmu dan teknologi Informatika.
- 3. Melaksanakan kegiatan pengabdian kepada masyarakat berbasis ilmu dan teknologi Informatika dalam rangka mengamalkan ilmu dan teknologi Informatika.