Design and Analysis of Algorithm

Advanced Data Structure (Red Black Tree)

(Properties, Rotation and Insertion)

LECTURE 28 -31

Overview

- A variation of binary search trees.
- Balanced: height is O(lg n), where n is the number of nodes.
- Operations will take O(lg n) time in the worst case.

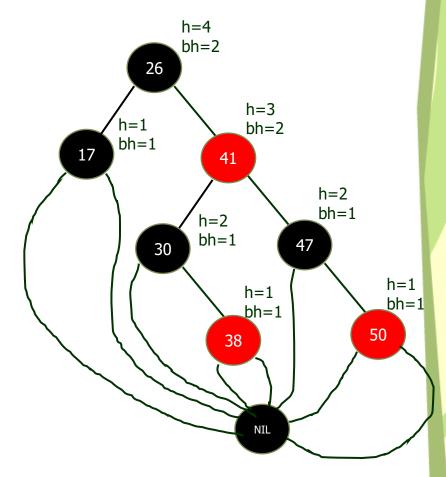
- A red-black tree is a binary search tree + 1 bit per node: an attribute color, which is either red or black.
- All leaves are empty (nil) and colored black.
 - A single sentinel, nil[T], is used for all the leaves of red-black tree T.
 - Color of [nil[T]](i.e. leaf node) is black.
 - The root's parent is also nil[T].
- All other attributes of binary search trees are inherited by red-black trees (i.e. key, left, right, and p as parent).
- We don't care about the key in nil[T].

• Properties:

- 1. Every node is either red or black.
- 2. The root is always black.
- 3. Every leaf (nil[T]) is black.
- 4. If a node is red, then both its children are black. (Hence no two reds in a row on a simple path from the root to a leaf is allowed.)
- 5. For each node, all paths from the node to descendant leaves contain the same number of black nodes.

Height of a red-black tree

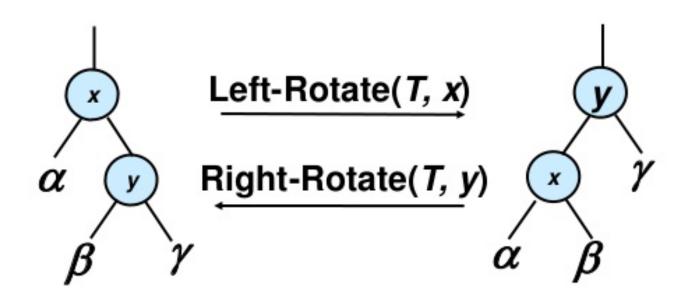
- Height of a node is the number of edges in a longest path to a leaf.
- Black-height of a node x: bh(x) is the number of black nodes (including nil[T]) on the path from x to leaf, not counting x. By property 5, black-height is well defined.
- So, a red-black tree of height h has black height >= h/2. Height of a red-black tree with n nodes is $h \le 2 \log_2(n + 1)$



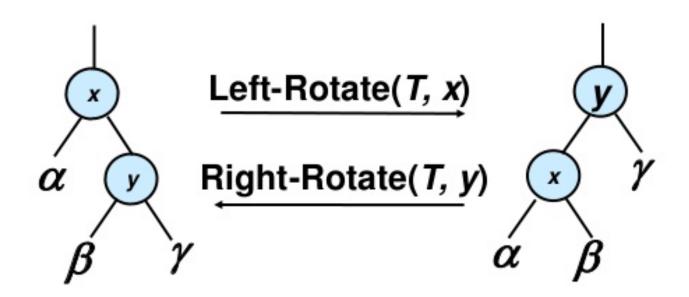
Rotations

- The basic tree-restructuring operation.
- Needed to maintain red-black trees as balanced binary search trees.
- Changes the local pointer structure. (Only pointers are changed.)
- Won't upset the binary-search-tree property.
- Have both left rotation and right rotation. They are inverses of each other.
- A rotation takes a red-black-tree and a node within the tree.

Rotations

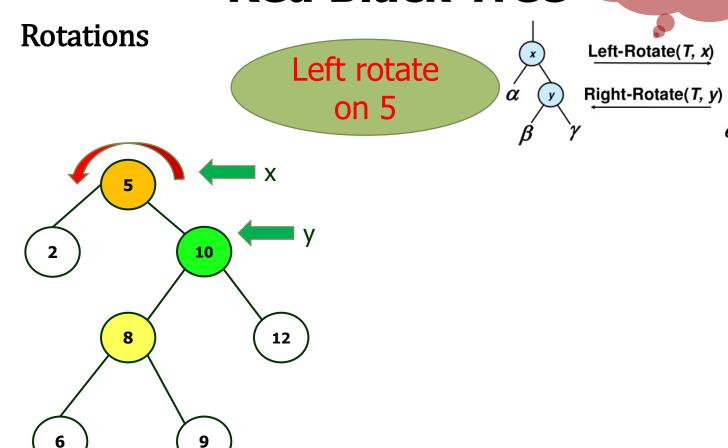


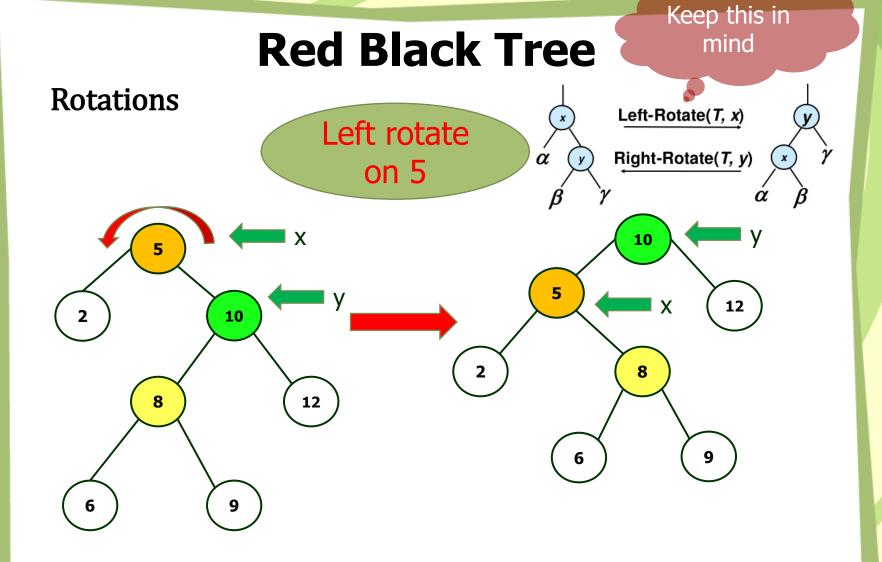
Rotations



Let's learn the game of changing pointer positions

Keep this in mind





Keep this in **Red Black Tree** mind **Rotations** Left-Rotate(T, x) Left rotate Right-Rotate(T, y) α on 5 10 У 12 2 10 8 12 8 6 9 Please note down the observations

Rotations

```
LEFT-ROTATE(T, x)
y \leftarrow right[x] //Set y.
right[x] \leftarrow left[y] //Turn y.s left subtree into x.s right subtree.
if left[y] \neq nil[T]
    then p[left[y]] \leftarrow x
p[y] \leftarrow p[x] //Link x.s parent
                                                             Left-Rotate(T, x)
if p[x] = nil[T]
   then root[T] \leftarrow y
                                                             Right-Rotate(T, y)
    else if x = left[p[x]]
            then left[p[x]] \leftarrow y
            else right[p[x]] \leftarrow y
left[y] \leftarrow x //Put x on y.s left.
p[x] \leftarrow y
```

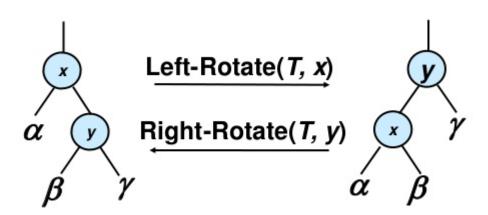
Rotations

The pseudocode for LEFT-ROTATE assumes that

right[x] = nil[T], and

root.s parent is nil[T].

Pseudocode for RIGHT-ROTATE is symmetric: exchange *left* and *right* everywhere.

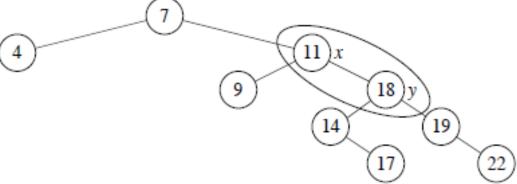


Rotations

```
RIGHT-ROTATE (T, x)
y \leftarrow left[x] //Set y.
left[x] \leftarrow right[y] //Turn y.s left subtree into x.s right subtree.
if right[y] \neq nil[T]
    then p[right[y]] \leftarrow x
p[y] \leftarrow p[x] //Link x.s parent
                                                             Left-Rotate(T, x)
if p[x] = nil[T]
   then root[T] \leftarrow y
                                                            Right-Rotate(T, y)
    else if x = right[p[x]]
           then right[p[x]] \leftarrow y
           else left[p[x]] \leftarrow y
right[y] \leftarrow x //Put x on y.s right.
p[x] \leftarrow y
```

Rotations (Example)

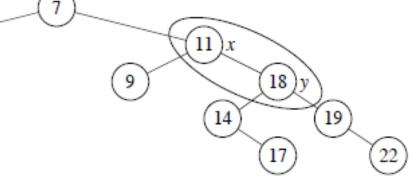
Demonstrate of left rotation that maintains in-order ordering of keys.



Rotations (Example)

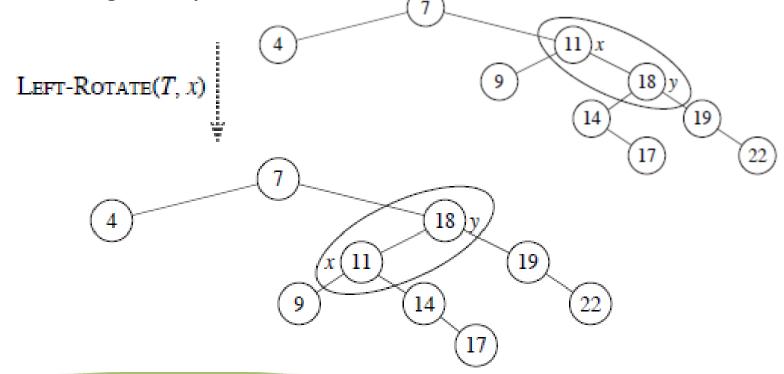
Demonstrate of left rotation that maintains in-order ordering of keys.

Left-Rotate(T, x)



Rotations (Example)

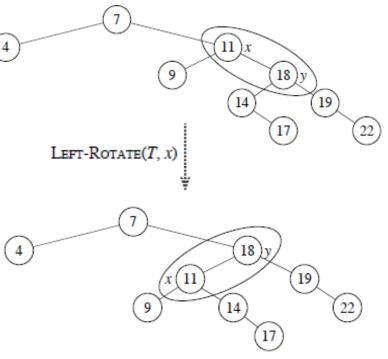
Demonstrate of left rotation that maintains in-order ordering of keys.



Rotations (Example)

Demonstrate of left rotation that maintains in-order ordering of keys.

- Before rotation: keys of x's left (4 subtree $\leq 11 \leq$ keys of y's left subtree $\leq 18 \leq$ keys of y's right subtree.
- Rotation makes y's left subtree into x's right subtree.
- After rotation: keys of x's left subtree $\leq 11 \leq$ keys of x's right subtree $\leq 18 \leq$ keys of y's right subtree.
- Time complexity: O(1) for both LEFT-ROTATE and RIGHT-ROTATE, since a constant number of pointers are modified.



Insertion:

Start by doing regular binary-search-tree insertion:

Insertion:

```
RB-INSERT(T, z)
      y \leftarrow nil[T]
      x \leftarrow root[T]
      while x \neq nil[T]
         do y \leftarrow x
           if key[z] < key[x]
            then x \leftarrow left[x]
            else x \leftarrow right[x]
      p[z] \leftarrow y
      if y = nil[T]
         then root[T] \leftarrow z
         else if key[z] < key[y]
                  then left[y] \leftarrow z
                  else right[y] \leftarrow z
      left[z] \leftarrow nil[T]
      right[z] \leftarrow nil[T]
      color[z] \leftarrow \text{RED}
       RB-INSERT-FIXUP(T, z)
```

Insertion:

```
RB-INSERT(T, z)
      y \leftarrow nil[T]
      x \leftarrow root[T]
       while x \neq nil[T]
         do y \leftarrow x
           if key[z] < key[x]
            then x \leftarrow left[x]
            else x \leftarrow right[x]
      p[z] \leftarrow y
      if y = nil[T]
         then root[T] \leftarrow z
         else if key[z] < key[y]
                  then left[y] \leftarrow z
                  else right[y] \leftarrow z
      left[z] \leftarrow nil[T]
       right[z] \leftarrow nil[T]
      color[z] \leftarrow \text{RED}
       RB-Insert-Fixup(T, z)
```

Simply Binary Search Tree Implementation

Insertion:

- RB-INSERT ends by coloring the new node z red.
- Then it calls RB-INSERT-FIXUP to maintain the properties of a red-black Tree.

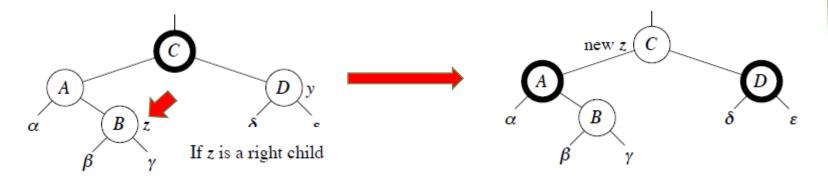
Which property might be violated?

1. Every node is either red or black.	OK
2. The root is always black.	If z is the root, then there's a violation. Otherwise, OK.
3. Every leaf (nil[T]) is black.	OK
4. If a node is red, then both its children are black.	If p[z] is red, there's a violation: both z and p[z] are red.
5. For each node, all paths from the node to descendant leaves contain the same number of black nodes.	OK

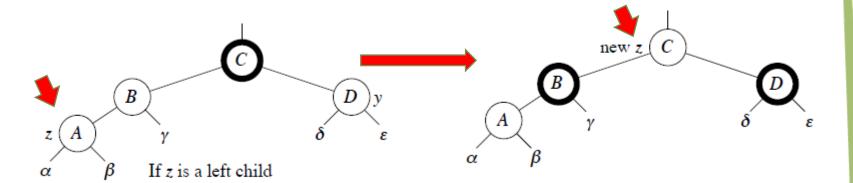
Insertion (RB-INSERT-FIXUP CASE 1) (y is red and z is a right child)

if
$$color[y] = \text{RED}$$

then $color[p[z]] \leftarrow \text{BLACK}$ //Case 1
 $color[y] \leftarrow \text{BLACK}$ //Case 1
 $color[p[p[z]]] \leftarrow \text{RED}$ //Case 1
 $z \leftarrow p[p[z]]$ //Case 1



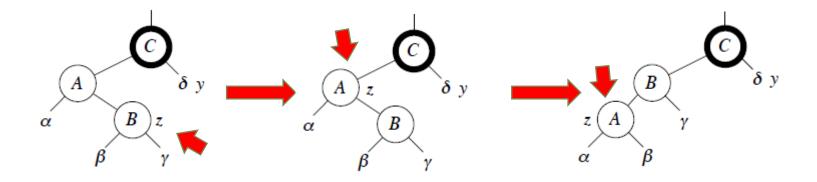
Insertion (RB-INSERT-FIXUP CASE 1) (y is red and z is a left child)



Insertion (RB-INSERT-FIXUP CASE 2 (y is black, z is a right child)

if
$$z = right[p[z]]$$

then $z \leftarrow p[z]$ //Case 2
LEFT-ROTATE(T, z) //Case 2

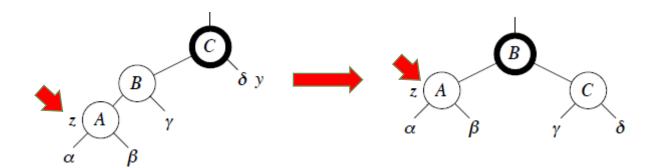


Insertion (RB-INSERT-FIXUP CASE 3 (y is black, z is a left child)

```
color[p[z]] \leftarrow BLACK // Case 3

color[p[p[z]]] \leftarrow RED // Case 3

RIGHT-ROTATE(T, p[p[z]]) // Case 3
```



Insertion:

```
RB-Insert-Fixup(T, z)
    while color[p[z]] = RED
       do if p[z] = left[p[p[z]]]
          then y \leftarrow right[p[p[z]]]
               if color[y] = RED
                 then
                    Apply Case 1
                 else if z = right[p[z]]
                        then
                               Apply Case 2
                        Apply Case 3
           else (same as then clause with .right. and .left. exchanged)
    color[root[T]] \leftarrow BLACK
```

Insertion:

```
RB-Insert-Fixup(T, z)
                                                       Once we apply
    while color[p[z]] = RED
                                                            Case 2
      do if p[z] = left[p[p[z]]]
                                                      immediately, we
         then y \leftarrow right[p[p[z]]]
                                                        apply Case 3
              if color[y] = RED
                then
                   Apply Case 1
                else if z = right[p[z]]
                      then
                             Apply Case 2
                      Apply Case 3
          else (same as then clause with .right. and .left. exchanged)
    color[root[T]] \leftarrow BLACK
```

```
Insertion:
                                                      Once we apply
                                                            Case 2
RB-Insert-Fixup(T, z)
                                                     immediately, we
     while color[p[z]] = RED
                                                       apply Case 3
       do if p[z] = left[p[p[z]]]
           then y \leftarrow right[p[p[z]]]
                if color[y] = RED
                  then color[p[z]] \leftarrow BLACK
                                                                          //Case 1
                        color[y] \leftarrow BLACK
                                                                          //Case 1
                        color[p[p[z]]] \leftarrow \text{RED}
                                                                         //Case 1
                                                                          //Case 1
                        z \leftarrow p[p[z]]
                  else if z = right[p[z]]
                                                                          //Case 2
                         then z \leftarrow p[z]
                                                                         //Case 2
                               LEFT-ROTATE(T, z)
                                                                          // Case 3
                        color[p[z]] \leftarrow BLACK
                                                                          // Case 3
                        color[p[p[z]]] \leftarrow \text{RED}
                        RIGHT-ROTATE(T, p[p[z]])
                                                                          // Case 3
           else (same as then clause with .right. and .left. exchanged)
     color[root[T]] \leftarrow BLACK
```

Insertion (Example):

Insert the following elements into an empty RB-Tree.

[11, 2, 14, 1, 7, 15, 5, 8, 4]



Insertion (Example 1):

Insert the following elements into an empty RB-Tree.

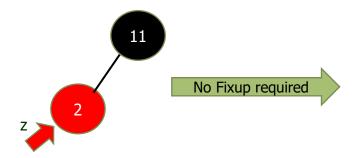
[11, 2, 14, 1, 7, 15, 5, 8, 4]



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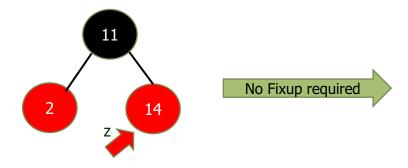
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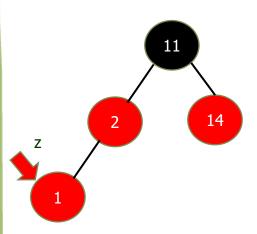
[11, 2, 14, 1, 7, 15, 5, 8, 4]



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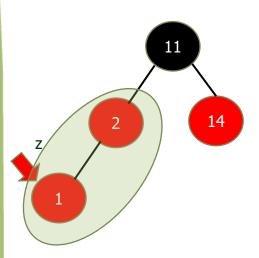
[11, 2, 14, 1, 7, 15, 5, 8, 4]



Insertion (Example 1):

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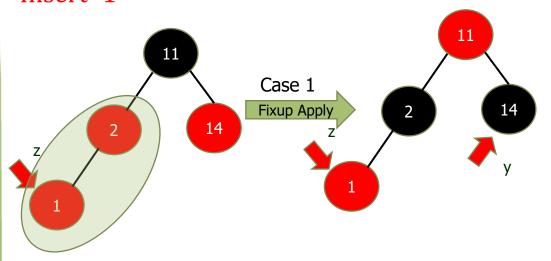
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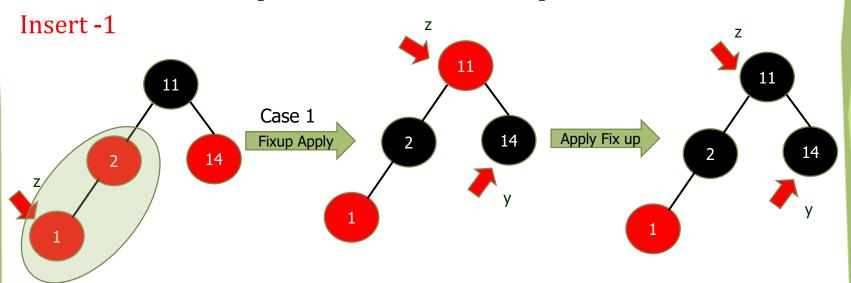
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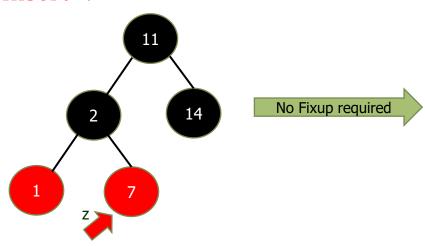
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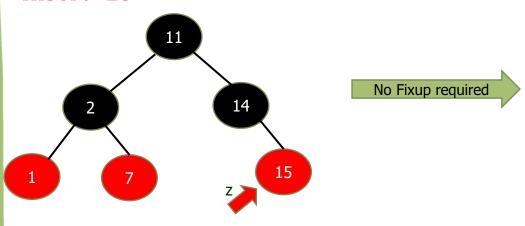
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Insertion (Example 1):

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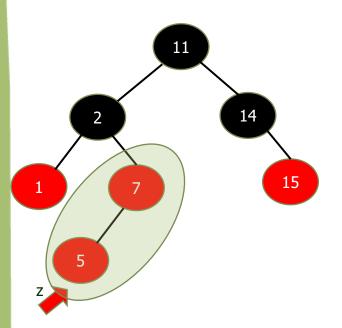
[11, 2, 14, 1, 7, 15, 5, 8, 4]



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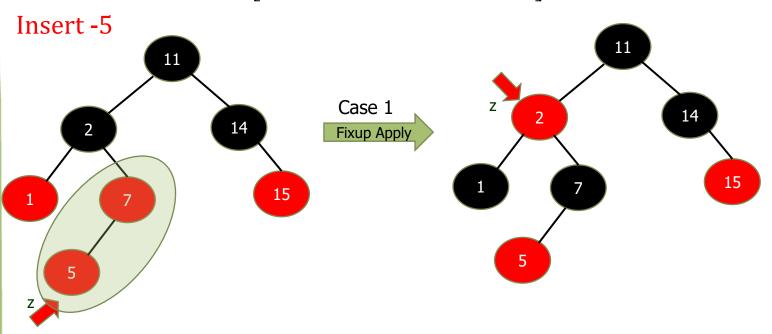
[11, 2, 14, 1, 7, 15, 5, 8, 4]



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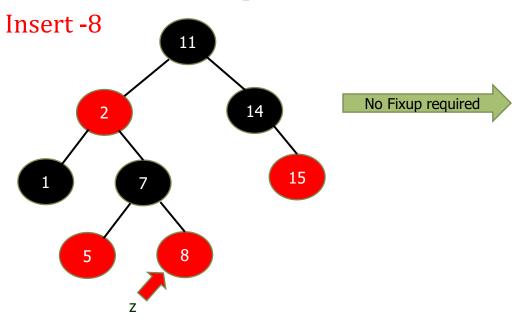
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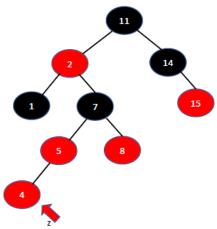
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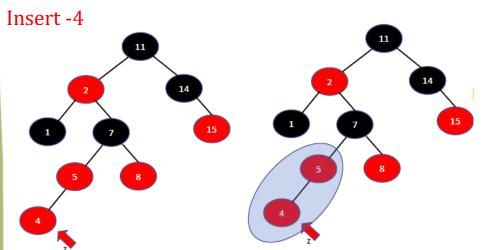
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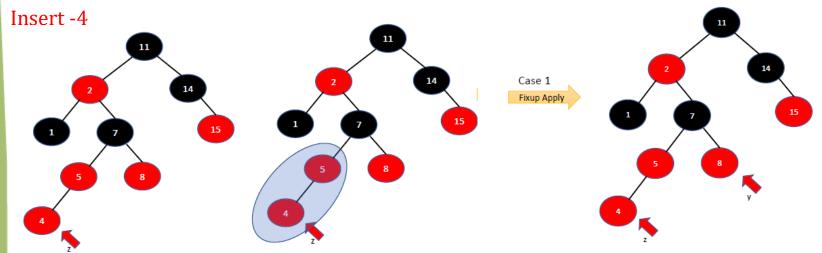
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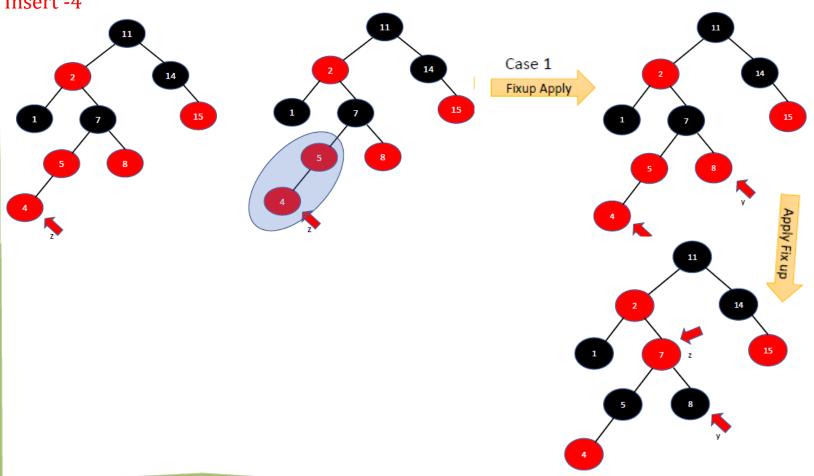
[11, 2, 14, 1, 7, 15, 5, 8, 4]



Red Black Tree Insertion (Example 1):

Insert the following elements into an empty RB-Tree.

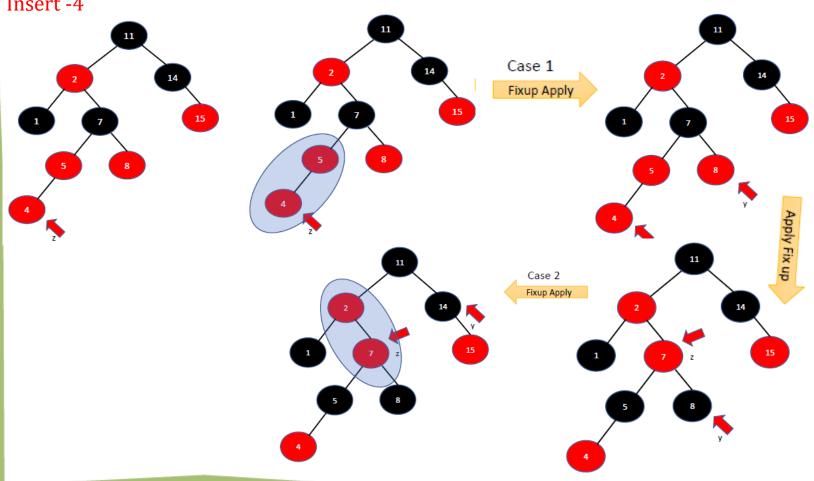
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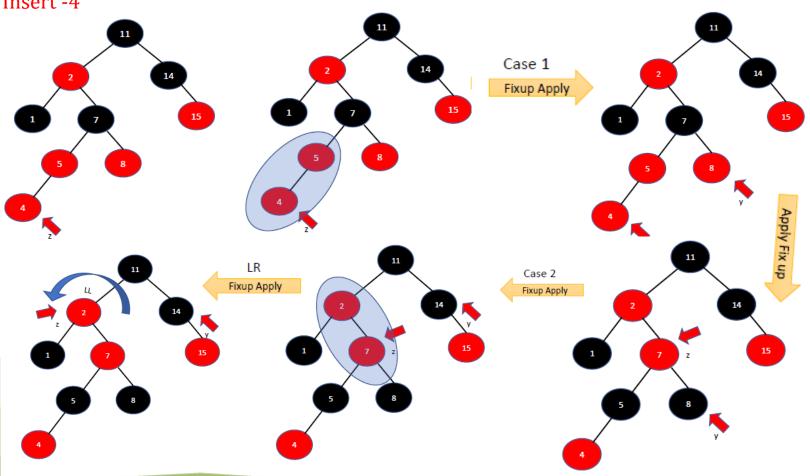


Red Black Tree Insertion (Example 1):

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[11, 2, 14, 1, 7, 15, 5, 8, 4]

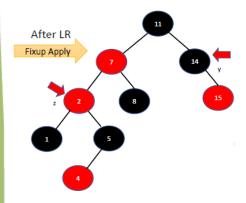




Insertion (Example 1): Red Black Tree

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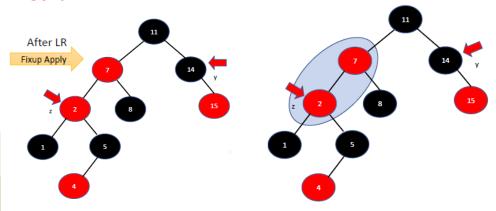
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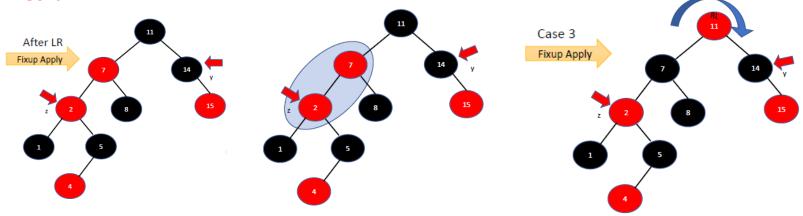
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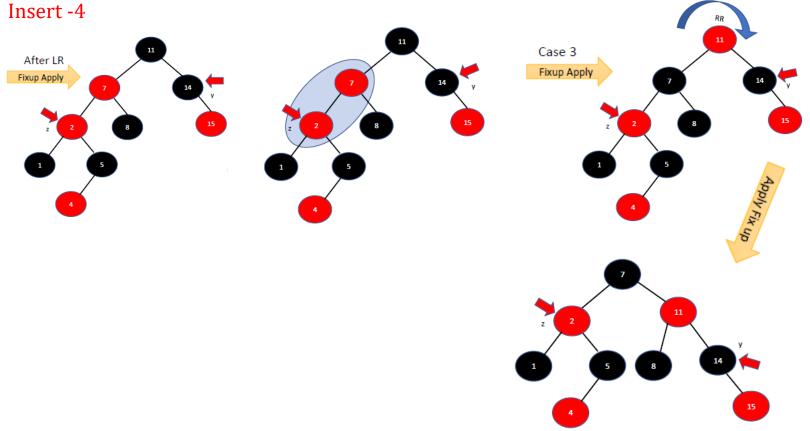


Red Black Tree Insertion (Example 1):

Insert the following elements into an empty RB-Tree.

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Insertion (Example 2):

Insert the following elements into an empty RB-Tree.

[50, 40, 30, 45, 20, 5]



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Insert the following elements into an empty RB-Tree.

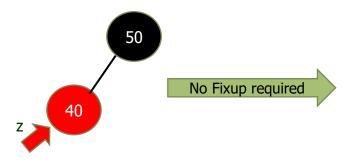
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Insertion (Example 2):

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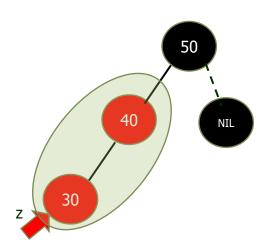
[50, 40, 30, 45, 20, 5]



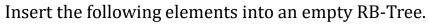
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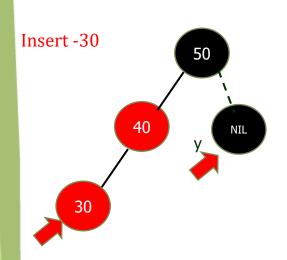
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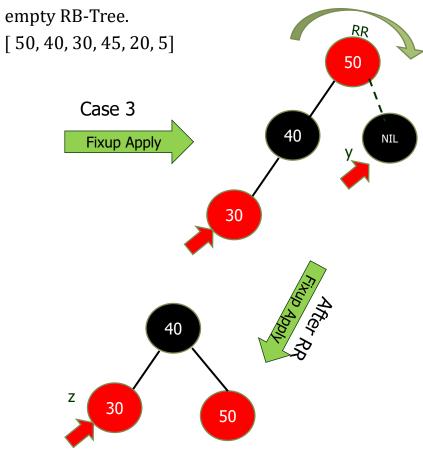
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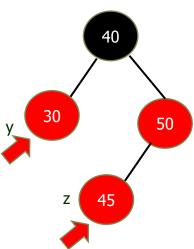




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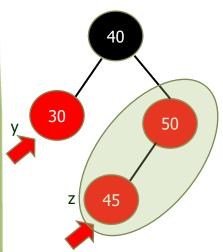
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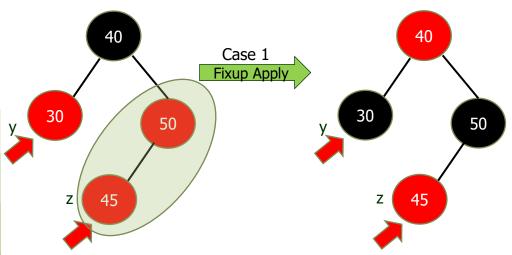
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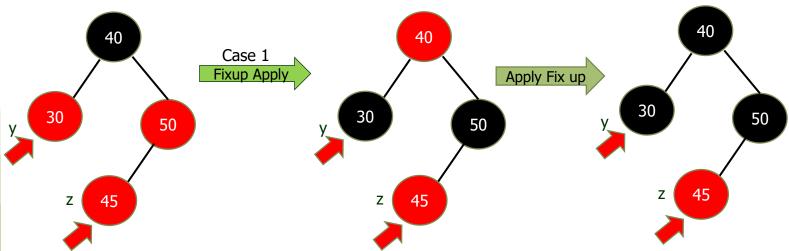
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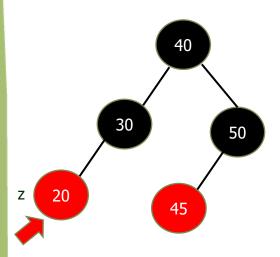
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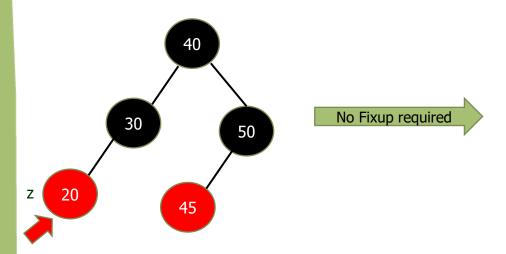
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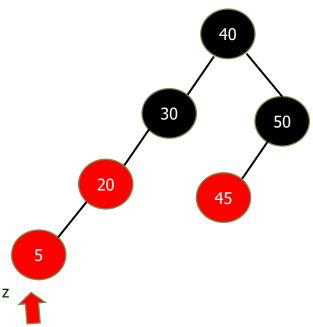
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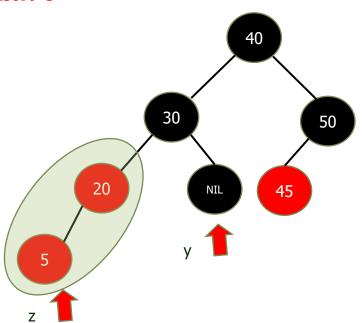
[50, 40, 30, 45, 20, 5]



Insertion (Example 2):

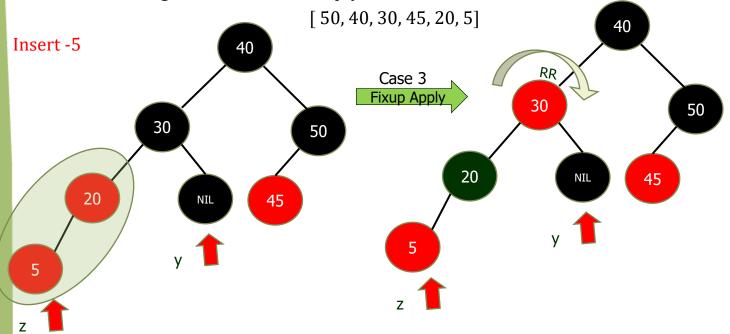
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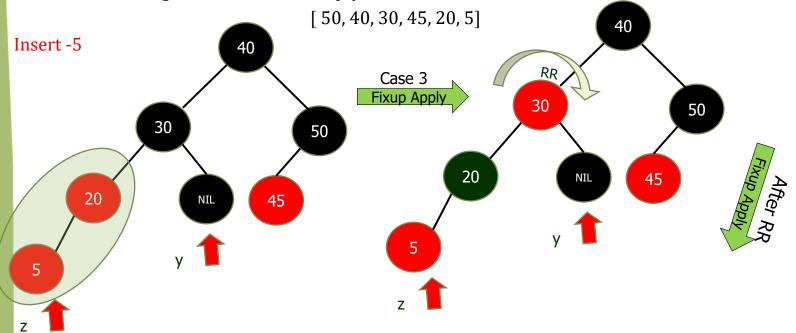
Insertion (Example 2):

Insert the following elements into an empty RB-Tree.



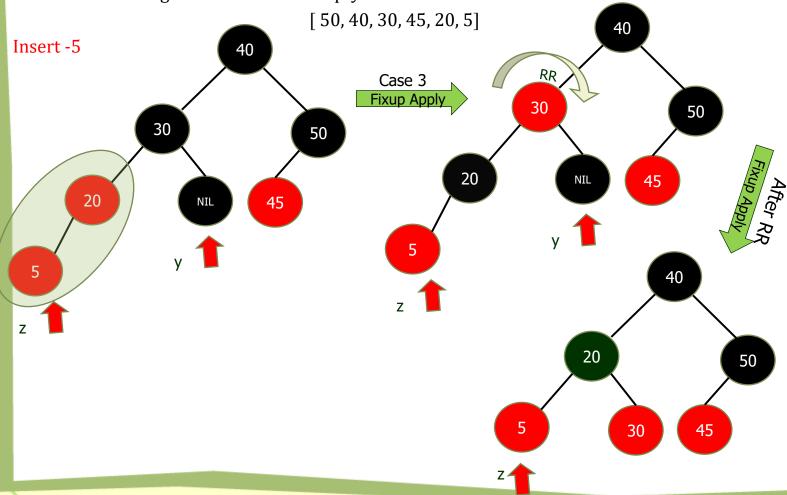
Insertion (Example 2):

Insert the following elements into an empty RB-Tree.



Insertion (Example 2):

Insert the following elements into an empty RB-Tree.



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

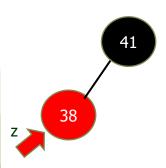
[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

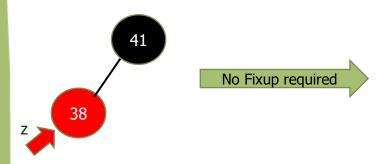
[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

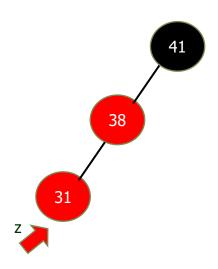
[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

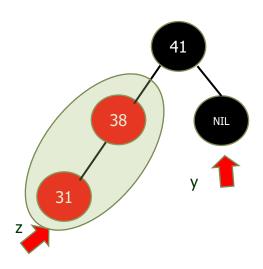
[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

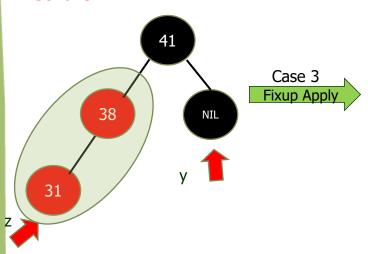
[41, 38, 31, 12, 19, 8]



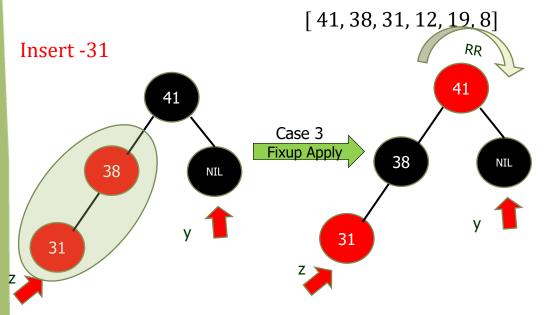
Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

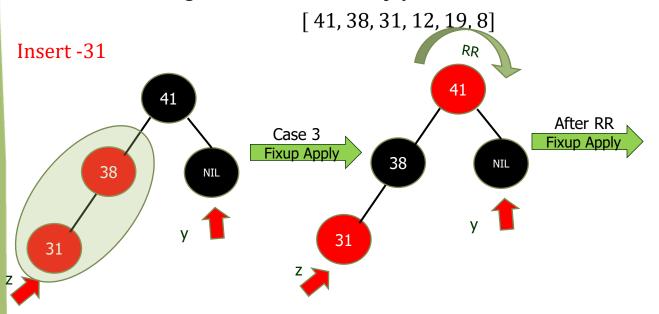
[41, 38, 31, 12, 19, 8]



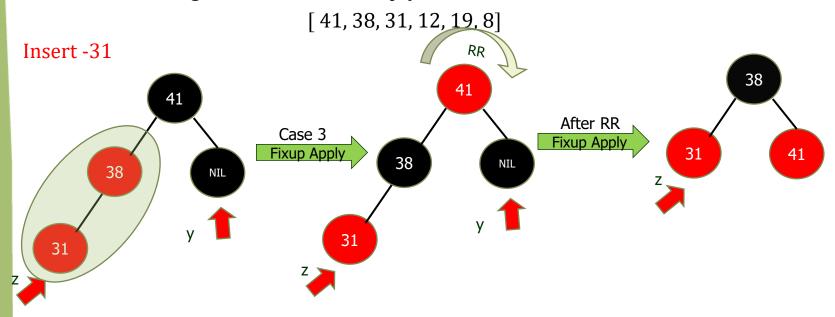
Insertion (Example 3):



Insertion (Example 3):



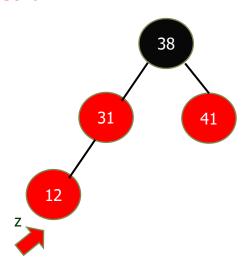
Insertion (Example 3):



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

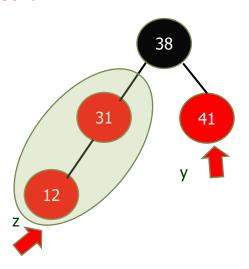
[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

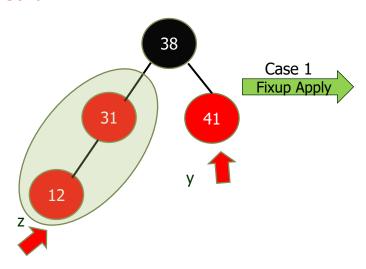
[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

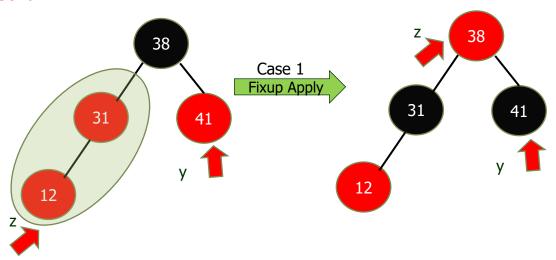
[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

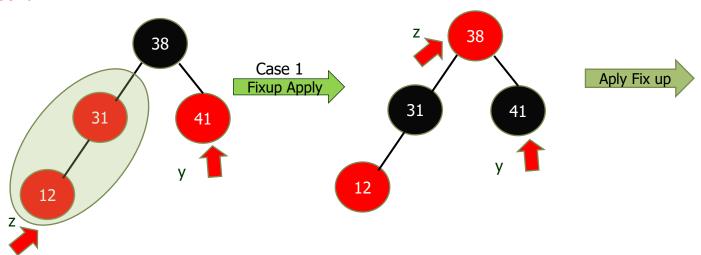
[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

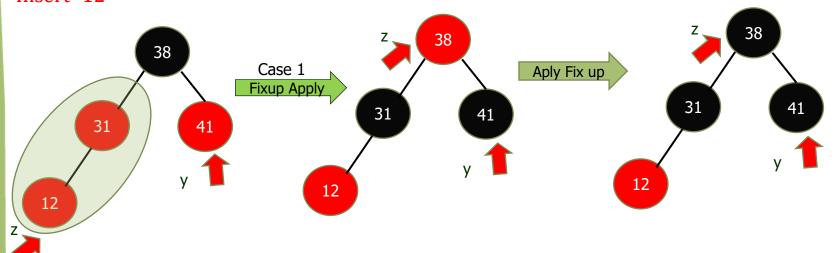
[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

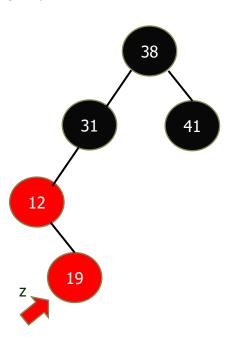
[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

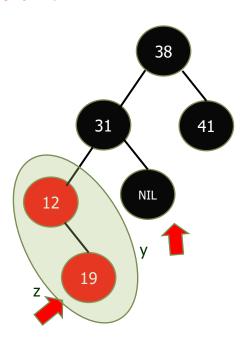
[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

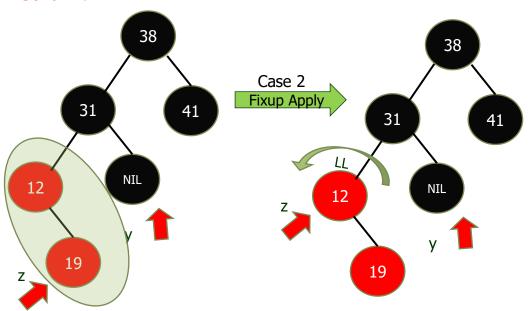
[41, 38, 31, 12, 19, 8]



Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

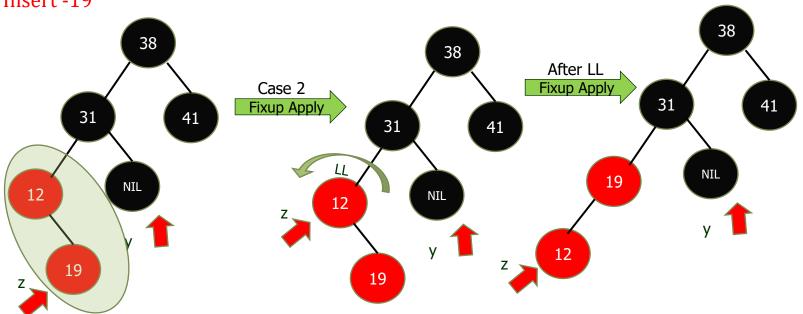
[41, 38, 31, 12, 19, 8]



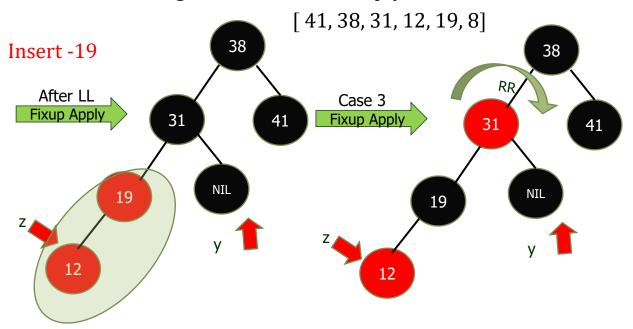
Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

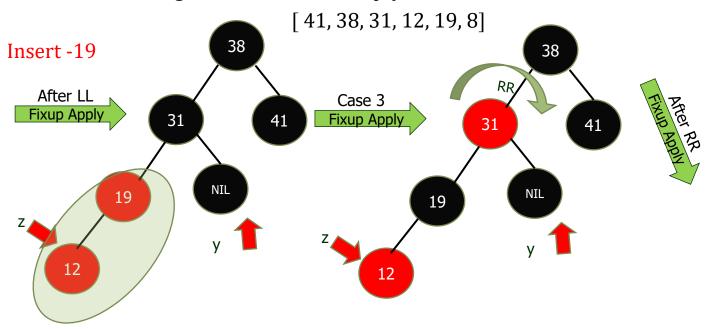
[41, 38, 31, 12, 19, 8]



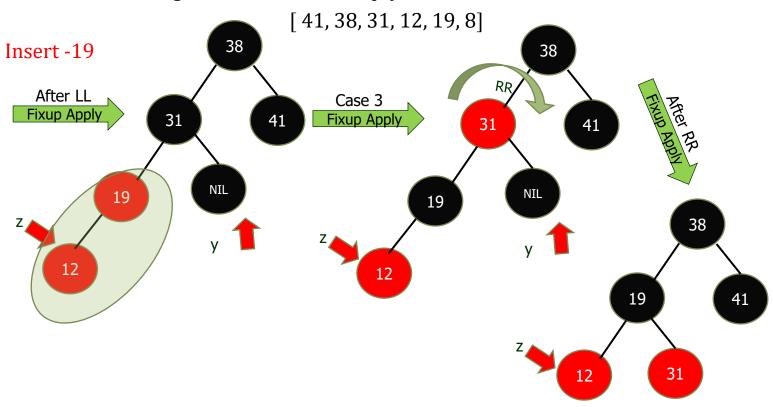
Insertion (Example 3):



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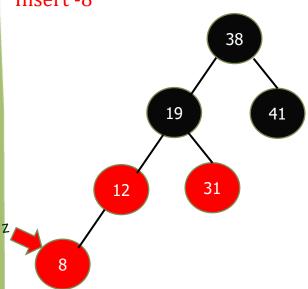


Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

[41, 38, 31, 12, 19, 8]

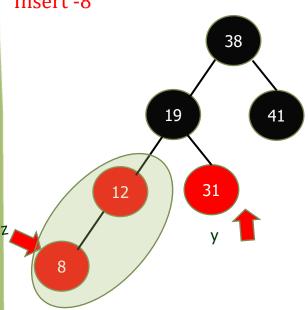




Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

[41, 38, 31, 12, 19, 8]

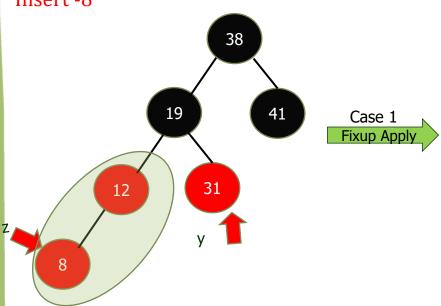


Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

[41, 38, 31, 12, 19, 8]

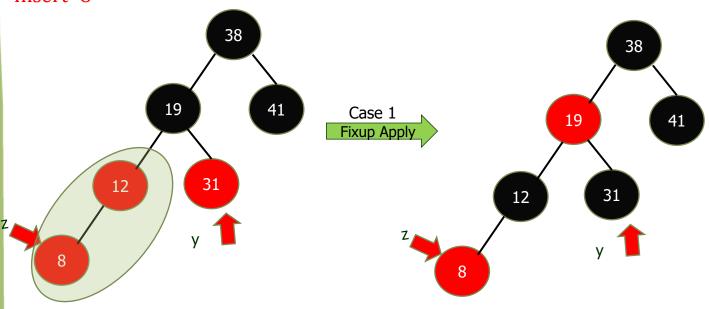




Insertion (Example 3):

Insert the following elements into an empty RB-Tree.

[41, 38, 31, 12, 19, 8]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

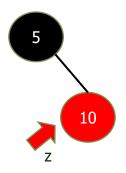
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

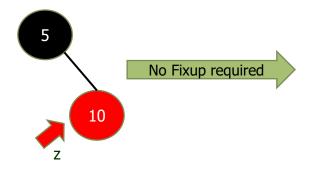
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

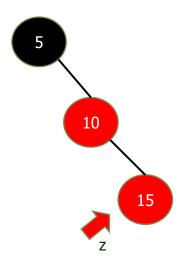
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

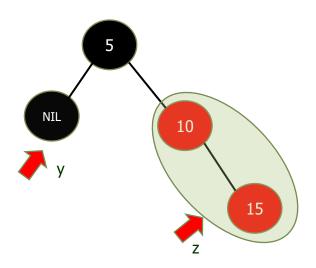
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

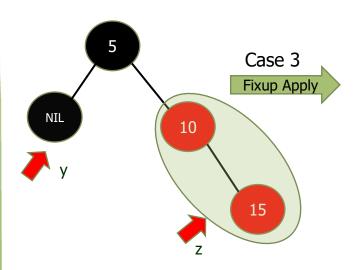
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

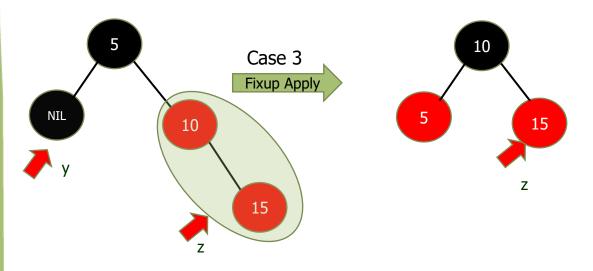
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

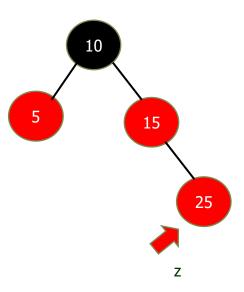
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

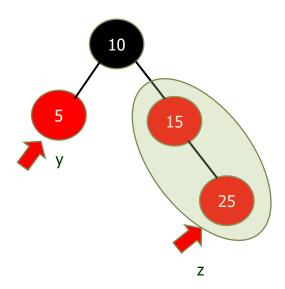
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

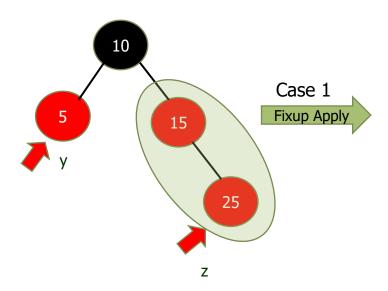
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

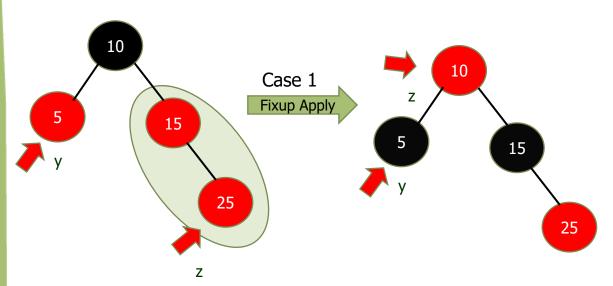
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

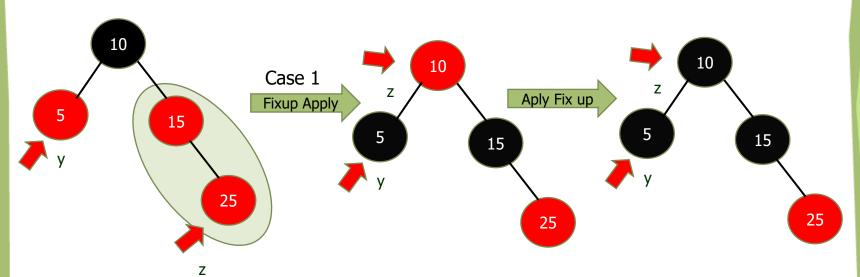
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

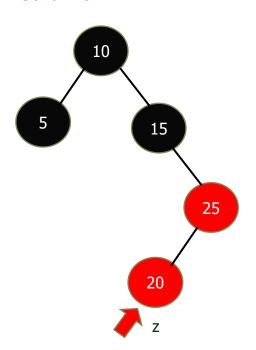
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

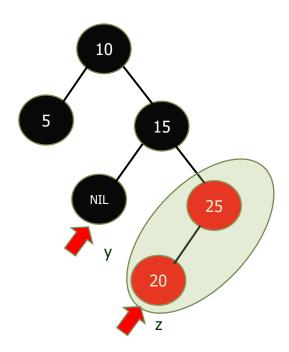
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

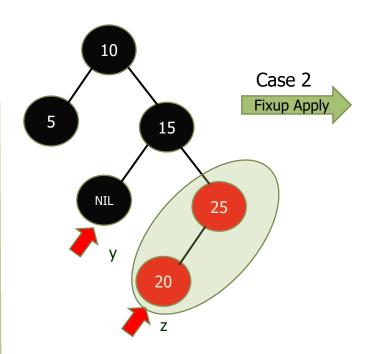
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

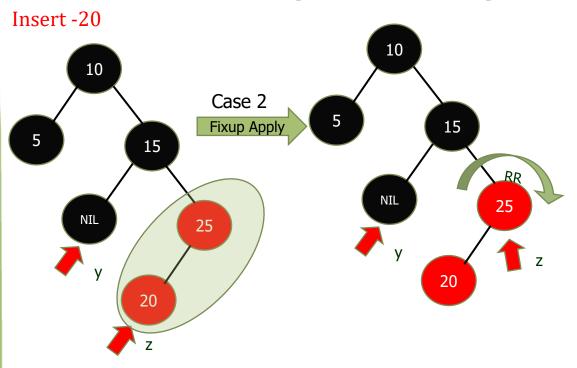
Insert the following elements into an empty RB-Tree.

[5, 10, 15, 25, 20, 30]



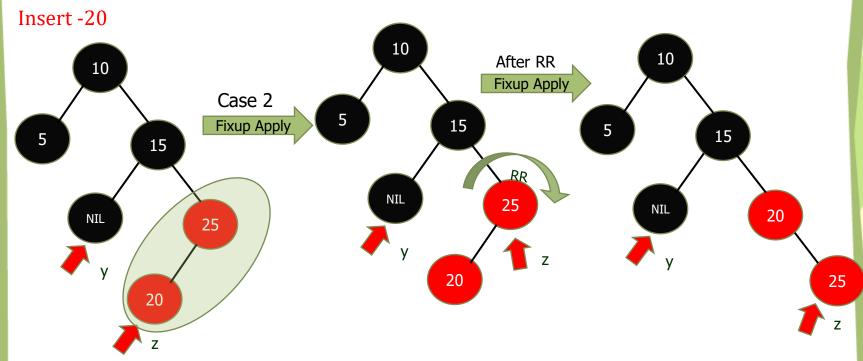
Insertion (Example 4):

Insert the following elements into an empty RB-Tree.



Insertion (Example 4):

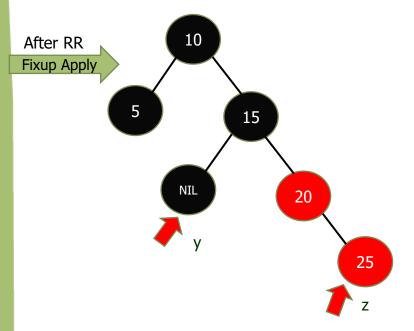
Insert the following elements into an empty RB-Tree.



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

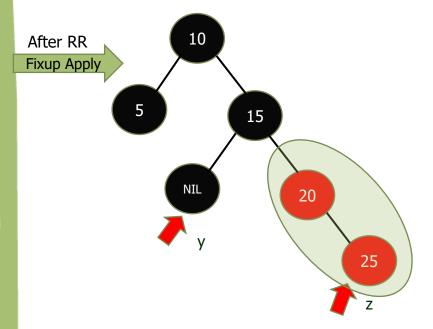
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

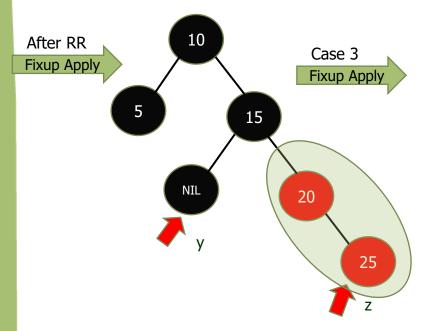
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

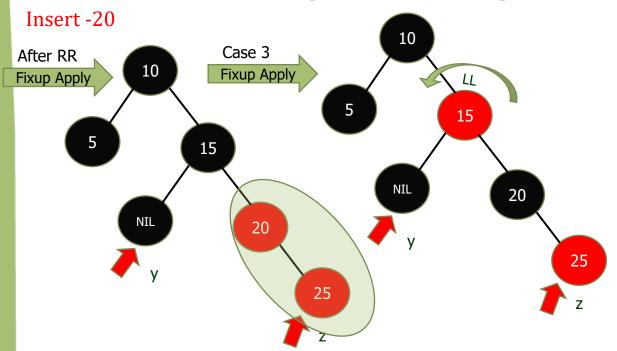
Insert the following elements into an empty RB-Tree.

[5, 10, 15, 25, 20, 30]



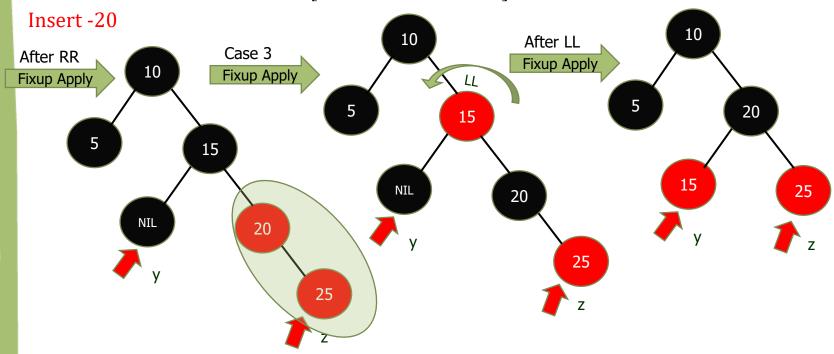
Insertion (Example 4):

Insert the following elements into an empty RB-Tree.



Insertion (Example 4):

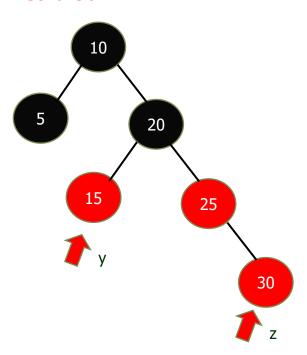
Insert the following elements into an empty RB-Tree.



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

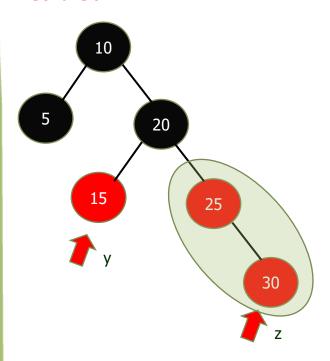
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

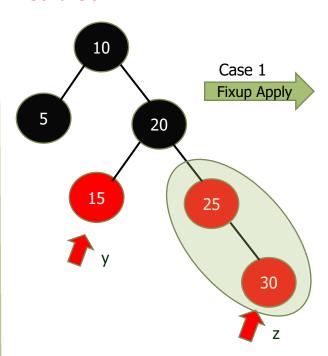
[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

[5, 10, 15, 25, 20, 30]



Insertion (Example 4):

Insert the following elements into an empty RB-Tree.

[5, 10, 15, 25, 20, 30]

