### **Design and Analysis of Algorithm**

Greedy Methods
(Single Source shortest path,
Knapsack problem)

**Lecture - 62-64** 



#### **Overview**

- A greedy algorithm always makes the choice that looks best at the moment. (i.e. it makes a locally optimal choice in the hope that this choice will lead to a globally optimal solution).
- The objective of this section is to explores optimization problems that are solvable by greedy algorithms.

- In mathematics, computer science and economics, an optimization problem is the problem of finding the best solution from all feasible solutions.
- Algorithms for optimization problems typically go through a sequence of steps, with a set of choices at each step.
- Many optimization problems can be solved using a greedy approach.
- Greedy algorithms are simple and straightforward.

- A greedy algorithm always makes the choice that looks best at the moment.
- That is, it makes a locally optimal choice in the hope that this choice will lead to a globally optimal solution.
- Greedy algorithms do not always yield optimal solutions, but for many problems they do.
- This algorithms are easy to invent, easy to implement and most of the time provides best and optimized solution.

- Application of Greedy Algorithm:
  - A simple but nontrivial problem, the activityselection problem, for which a greedy algorithm efficiently computes a solution.
  - In combinatorics, (a branch of mathematics), a 'matroid' is a structure that abstracts and generalizes the notion of linear independence in vector spaces. Greedy algorithm always produces an optimal solution for such problems. Scheduling unit-time tasks with deadlines and penalties is an example of such problem.

- Application of Greedy Algorithm:
  - An important application of greedy techniques is the design of data-compression codes (i.e. Huffman code).
  - The greedy method is quite powerful and works well for a wide range of problems. They are:
    - Minimum-spanning-tree algorithms
      - (Example: Prims and Kruskal algorithm)
    - Single Source Shortest Path.

(Example: Dijkstra's and Bellman ford algorithm)

- Application of Greedy Algorithm:
  - A problem exhibits optimal substructure if an optimal solution to the problem contains within it optimal solutions to subproblems.
  - This property is a key ingredient of assessing the applicability of dynamic programming as well as greedy algorithms.
  - The subtleties between the above two techniques are illustrated with the help of two variants of a classical optimization problem known as knapsack problem. These variants are:
    - 0-1 knapsack problem (Dynamic Programming)
    - Fractional knapsack problem (Greedy Algorithm)

- Problem 5: Single source shortest path
  - It is a shortest path problem where the shortest path from a given source vertex to all other remaining vertices is computed.
  - Dijkstra's Algorithm and Bellman Ford Algorithm are the famous algorithms used for solving single-source shortest path problem.

- Problem 5: Single source shortest path
  - Dijkstra's Algorithm
    - Dijkstra Algorithm is a very famous greedy algorithm.
    - It is used for solving the single source shortest path problem.
    - It computes the shortest path from one particular source node to all other remaining nodes of the graph.

- Problem 5: Single source shortest path
  - Dijkstra's Algorithm (Feasible Condition)
    - Dijkstra algorithm works
      - for connected graphs.
      - for those graphs that do not contain any negative weight edge.
      - To provides the value or cost of the shortest paths.
      - for directed as well as undirected graphs.

#### Problem 5: Single source shortest path

Dijkstra's Algorithm (Implementation)

The implementation of Dijkstra Algorithm is executed in the following steps-

#### • Step-01:

- In the first step. two sets are defined-
- One set contains all those vertices which have been included in the shortest path tree.
- In the beginning, this set is empty.
- Other set contains all those vertices which are still left to be included in the shortest path tree.
- In the beginning, this set contains all the vertices of the given graph.

- Problem 5: Single source shortest path
  - Dijkstra's Algorithm (Implementation)

The implementation of Dijkstra Algorithm is executed in the following steps-

#### • Step-02:

For each vertex of the given graph, two variables are defined as-

- Π[v] which denotes the predecessor of vertex `v'
- d[v] which denotes the shortest path estimate of vertex 'v' from the source vertex.

- Problem 5: Single source shortest path
  - Dijkstra's Algorithm (Implementation)

The implementation of Dijkstra Algorithm is executed in the following steps-

#### • Step-02:

Initially, the value of these variables is set as-

- The value of variable ' $\Pi$ ' for each vertex is set to NIL i.e.  $\Pi[v] = NIL$
- The value of variable 'd' for source vertex is set to 0 i.e. d[S]
   = 0
- The value of variable 'd' for remaining vertices is set to  $\infty$  i.e.  $d[v] = \infty$

- Problem 5: Single source shortest path
  - Dijkstra's Algorithm (Implementation)

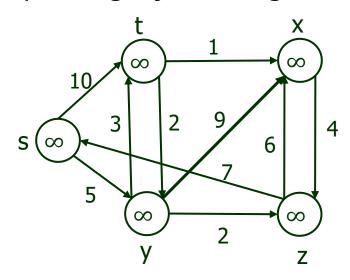
The implementation of Dijkstra Algorithm is executed in the following steps-

#### • Step-03:

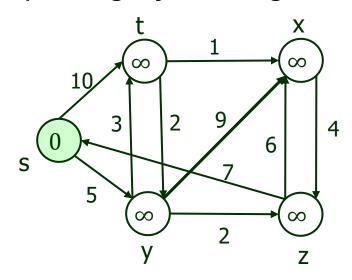
The following procedure is repeated until all the vertices of the graph are processed-

- Among unprocessed vertices, a vertex with minimum value of variable 'd' is chosen.
- Its outgoing edges are relaxed.
- After relaxing the edges for that vertex, the sets created in step-01 are updated.

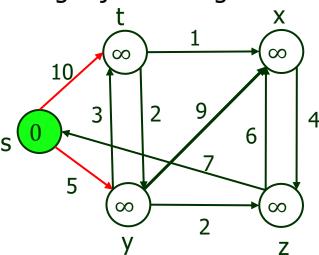
- Problem 5: Single source shortest path
  - Dijkstra's Algorithm



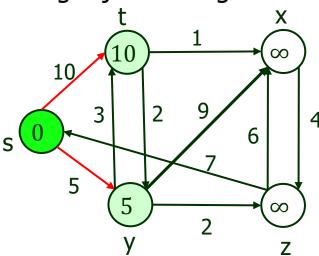
- Problem 5: Single source shortest path
  - Dijkstra's Algorithm



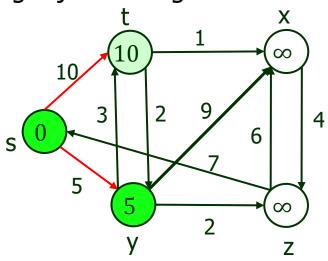
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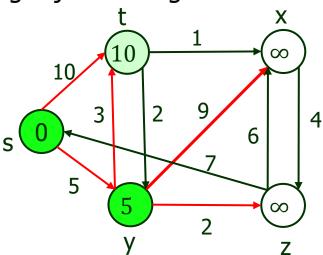
- Problem 5: Single source shortest path
  - Dijkstra's Algorithm



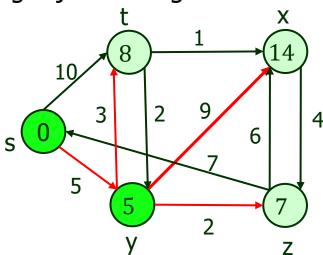
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  - Dijkstra's Algorithm



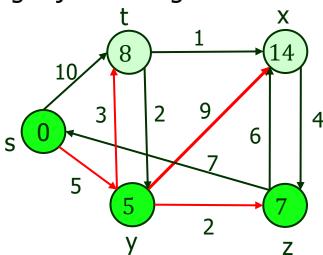
- Problem 5: Single source shortest path
  - Dijkstra's Algorithm



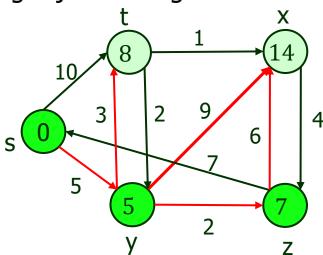
- Problem 5: Single source shortest path
  - Dijkstra's Algorithm



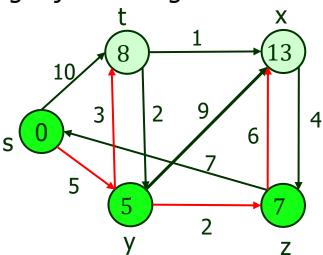
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  - Dijkstra's Algorithm



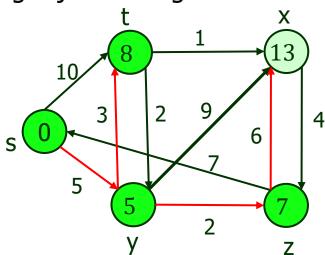
- Problem 5: Single source shortest path
  - Dijkstra's Algorithm



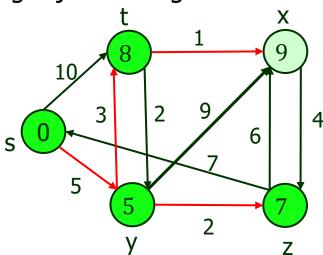
- Problem 5: Single source shortest path
  - Dijkstra's Algorithm



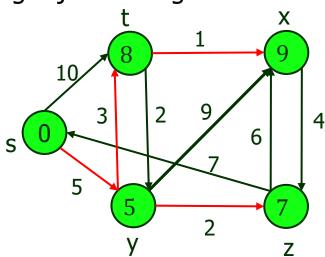
- Problem 5: Single source shortest path
  - Dijkstra's Algorithm



- Problem 5: Single source shortest path
  - Dijkstra's Algorithm

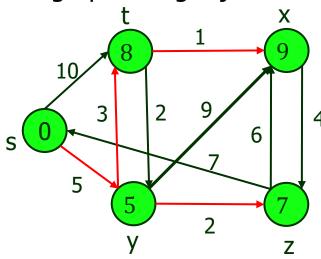


- Problem 5: Single source shortest path
  - Dijkstra's Algorithm



- Problem 5: Single source shortest path
  - Dijkstra's Algorithm

Example 1: Construct the Single source shortest path for the given graph using Dijkstra's Algorithm-



Hence the shortest path to all the vertex from s are:

$$s \rightarrow t = 8$$

$$s \rightarrow x = 9$$

$$s \rightarrow y = 5$$

$$s \rightarrow z = 7$$

- Problem 5: Single source shortest path
  - Dijkstra's Algorithm (Algorithm)

```
DIJKSTRA(G, w, s)

1 INITIALIZE-SINGLE-SOURCE(G, s)

2 S \leftarrow \emptyset

3 Q \leftarrow V[G]

4 while Q \neq \emptyset

5 do u \leftarrow EXTRACT-MIN(Q)

6 S \leftarrow S \in \{u\}

7 for each vertex v \in Adj[u]

8 do RELAX(u, v, w)
```

- Problem 5: Single source shortest path
  - Dijkstra's Algorithm (Algorithm)

```
INITIALIZE-SINGLE-SOURCE(G, s)

1 for each vertex v V[G]

2 do d[v] \leftarrow \infty

3 \Pi[v] \leftarrow NIL

4 d[s] \leftarrow 0

RELAX(u, v, w)

1 if d[v] > d[u] + w(u, v)

2 then d[v] \leftarrow d[u] + w(u, v)

3 \Pi[v] \leftarrow u
```

- Problem 5: Single source shortest path
  - Dijkstra's Algorithm (Complexity)

#### CASE-01: (IN CASE OF COMPLETE GRAPH)

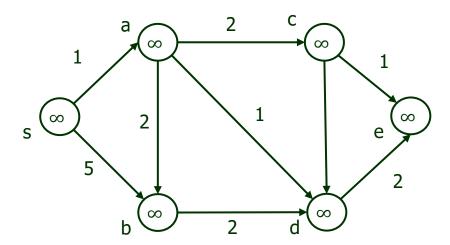
- A[i,j] stores the information about edge (i,j).
- Time taken for selecting i with the smallest dist is O(V).
- For each neighbor of i, time taken for updating dist[j] is O(1) and there will be maximum V neighbors.
- Time taken for each iteration of the loop is O(V) and one vertex is deleted from Q.
- Thus, total time complexity becomes  $O(V^2)$ .

- Problem 5: Single source shortest path
  - Dijkstra's Algorithm (Complexity)

#### **CASE-02:**

- With adjacency list representation, all vertices of the graph can be traversed using BFS in O(V + E) time.
- In min heap, operations like extract-min and decrease-key value takes  $O(\log V)$  time.
- So, overall time complexity becomes  $O(E + V) \times O(\log V)$  which is  $O((E + V) \times \log V) = O(E \log V)$
- This time complexity is reduced to  $O(E + V \log V)$  using Fibonacci heap.

- Problem 5: Single source shortest path
  - Dijkstra's Algorithm (Self Practice)



- Problem 5: Single source shortest path
- Bellman Ford Algorithm
  - Bellman-Ford algorithm solves the single-source shortest-path problem in the general case in which edges of a given digraph can have negative weight as long as G contains no negative cycles.
  - Like Dijkstra's algorithm, this algorithm, uses the notion of edge relaxation without using greedy method. Again, it uses d[u] as an upper bound on the distance d[u, v] from u to v.

- Problem 5: Single source shortest path
- Bellman Ford Algorithm
  - The algorithm progressively decreases an estimate d[v] on the weight of the shortest path from the source vertex s to each vertex  $v \in V$  until it achieve the actual shortest-path.
  - The algorithm returns Boolean TRUE if the given digraph contains no negative cycles that are reachable from source vertex s otherwise it returns Boolean FALSE.

Problem 5: Single source shortest path
 Bellman Ford Algorithm (Negative Cycle Detection)

-10

Assume:

$$d[u] \le d[x] + 4$$

$$d[v] \le d[u] + 5$$

$$d[x] \le d[v] - 10$$

Adding:

$$d[u] + d[v] + d[x] \le d[x] + d[u] + d[v] - 1$$

Because it's a cycle, vertices on left are same as those on right. Thus we get  $0 \le -1$ ; a contradiction.

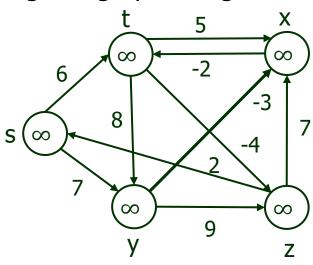
So for at least one edge (u, v), we have to check

$$d[v] > d[u] + w(u,v)$$

This is exactly what Bellman-Ford checks for.

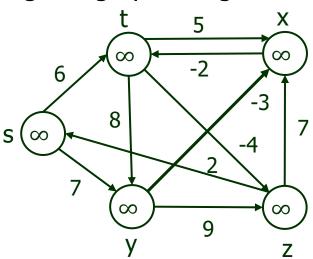
- Problem 5: Single source shortest path
- Bellman Ford Algorithm (Implementation)
  - Step 1: Start with the weighted graph.
  - Step 2: Choose the starting vertex by making the path value zero and assign infinity path values to all other vertices.
  - Step 3: Visit each edge and relax the path distances if they are inaccurate.
  - Step 4: Do step 3 'V' times because in the worst case a vertex's path length might need to be readjusted V times.
  - Step 5: After all vertices have their path lengths, check if a negative cycle is present or not.

- Problem 5: Single source shortest path
  - Bellman Ford Algorithm



- Problem 5: Single source shortest path
  - Bellman Ford Algorithm

Example 1: Construct the Single source shortest path for the given graph using Bellman ford Algorithm-



#### The edges are:

The cages arei	
1. (s,t)	6. (t,x)
2. (s,y)	7. (t,z)
3. (y,z)	8. (t,y)
4. (z,x)	9. (y,x)
5. (x.t)	10.(7.5)

- Problem 5: Single source shortest path
  - Bellman Ford Algorithm

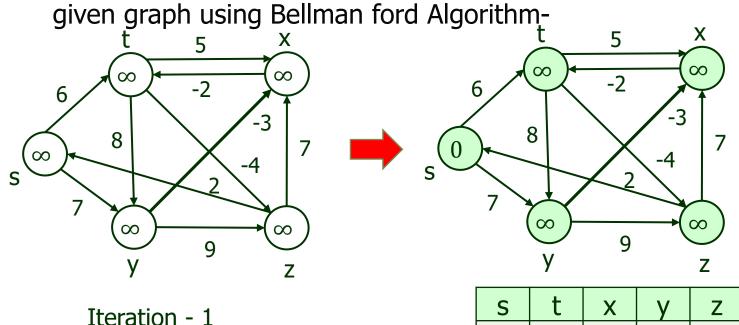
Example 1: Construct the Single source shortest path for the

 $\infty$ 

 $\infty$ 

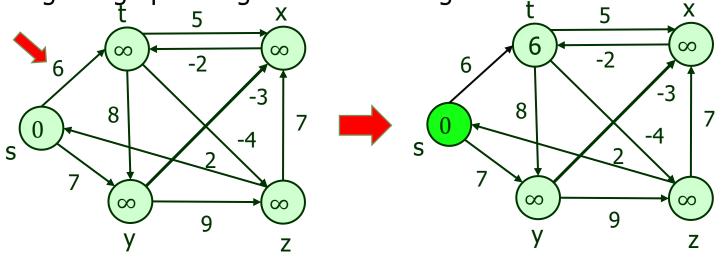
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- Problem 5: Single source shortest path
  - Bellman Ford Algorithm

Example 1: Construct the Single source shortest path for the given graph using Bellman ford Algorithm-

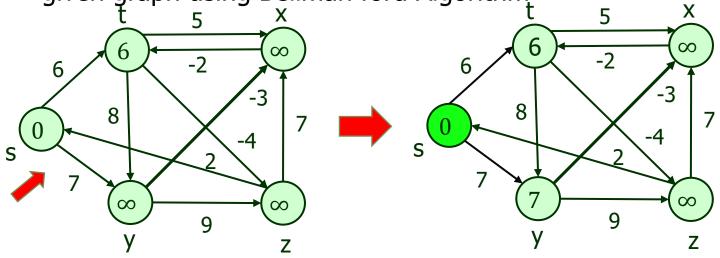


Iteration - 1

S	t	Х	У	Z
0	6	$\infty$	8	8

- Problem 5: Single source shortest path
  - Bellman Ford Algorithm

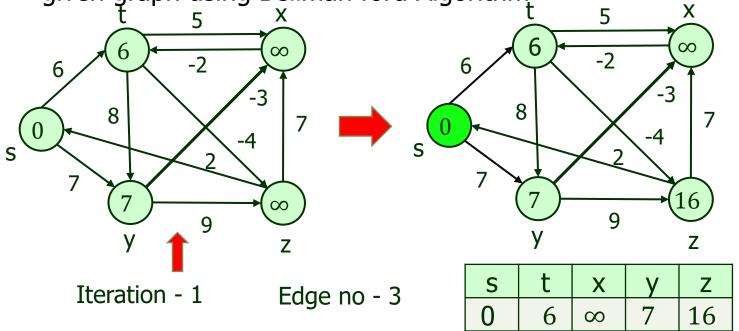
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Iteration - 1

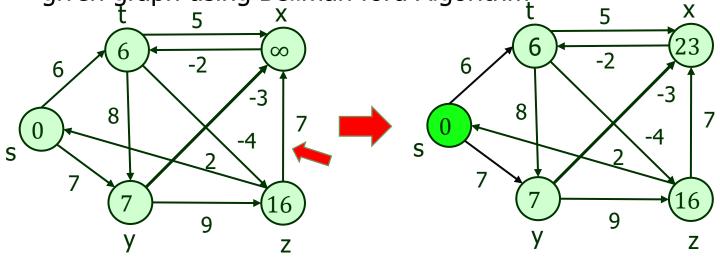
S	t	Х	У	Z
0	6	8	7	$\infty$

- Problem 5: Single source shortest path
  - Bellman Ford Algorithm



- Problem 5: Single source shortest path
  - Bellman Ford Algorithm

Example 1: Construct the Single source shortest path for the given graph using Bellman ford Algorithm-

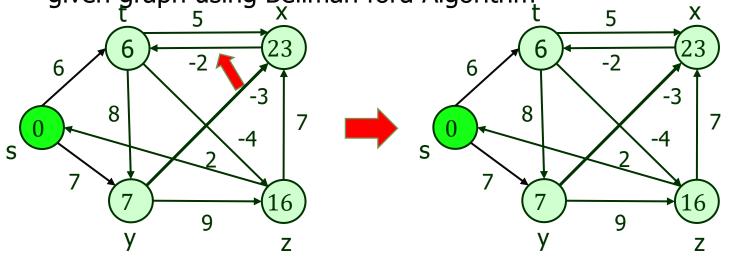


Iteration - 1

S	t	X	У	Z
0	6	23	7	16

- Problem 5: Single source shortest path
  - Bellman Ford Algorithm

Example 1: Construct the Single source shortest path for the given graph using Bellman ford Algorithm-

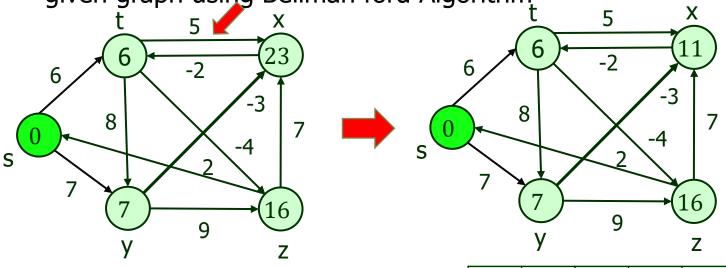


Iteration - 1

S	t	X	У	Z
0	6	23	7	16

- Problem 5: Single source shortest path
  - Bellman Ford Algorithm

Example 1: Construct the Single source shortest path for the given graph using Bellman ford Algorithm-

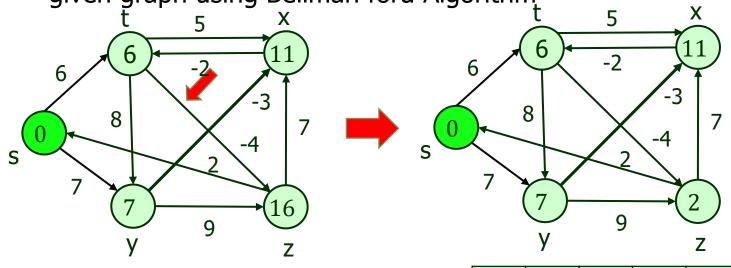


Iteration - 1

S	t	X	У	Z
0	6	11	7	16

- Problem 5: Single source shortest path
  - Bellman Ford Algorithm

Example 1: Construct the Single source shortest path for the given graph using Bellman ford Algorithm-

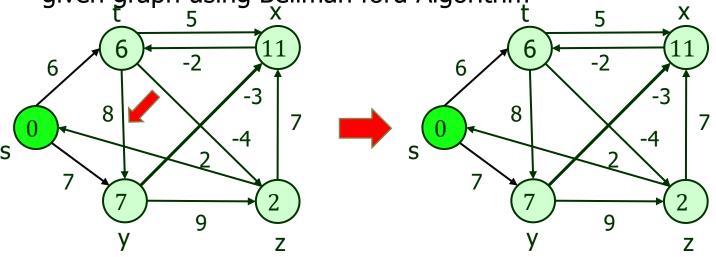


Iteration - 1

S	t	X	У	Z
0	6	11	7	2

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  - Bellman Ford Algorithm

Example 1: Construct the Single source shortest path for the given graph using Bellman ford Algorithm-

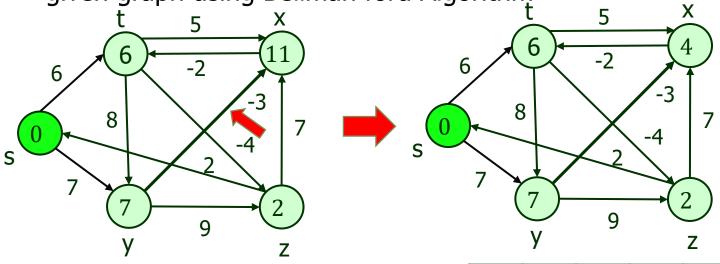


Iteration - 1

S	t	Х	У	Z
0	6	11	7	2

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  - Bellman Ford Algorithm

Example 1: Construct the Single source shortest path for the given graph using Bellman ford Algorithm-

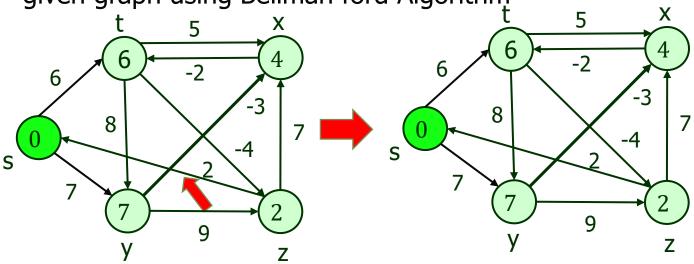


Iteration - 1

S	t	X	У	Z
0	6	4	7	2

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  - Bellman Ford Algorithm

Example 1: Construct the Single source shortest path for the given graph using Bellman ford Algorithm-

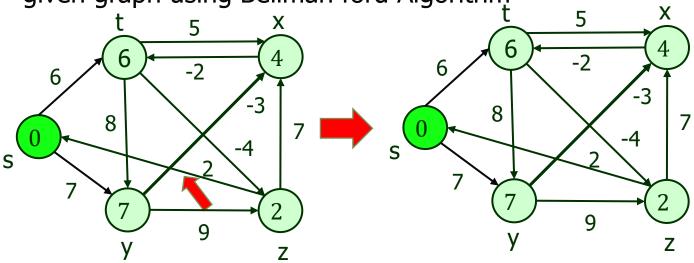


Iteration - 1

S	t	Х	У	Z
0	6	4	7	2

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  - Bellman Ford Algorithm

Example 1: Construct the Single source shortest path for the given graph using Bellman ford Algorithm-

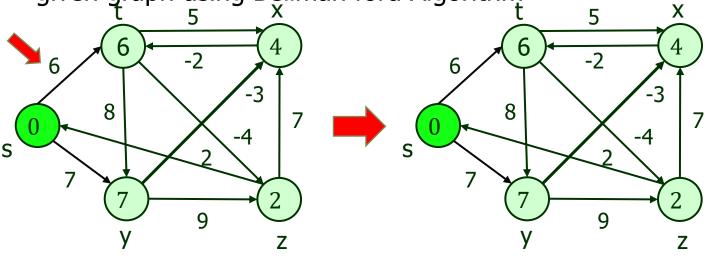


Iteration - 1

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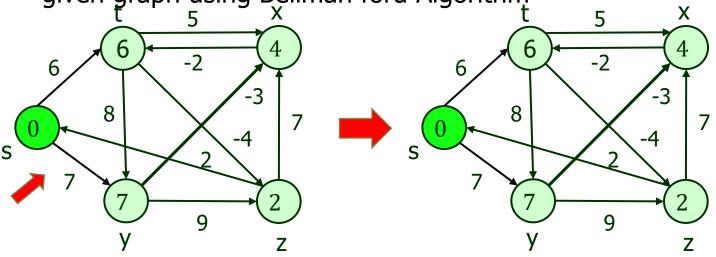


Iteration - 2

S	t	Х	У	Z
0	6	4	7	2

- Problem 5: Single source shortest path
  - Bellman Ford Algorithm

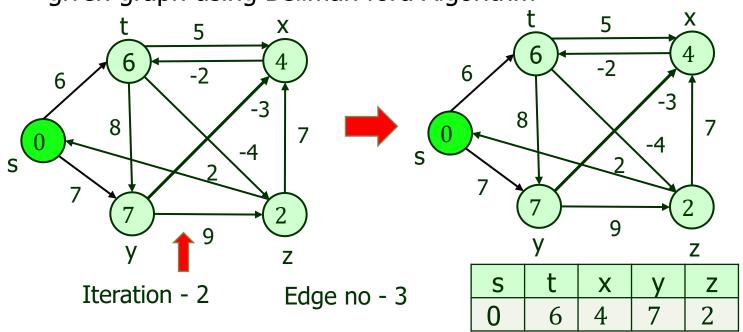
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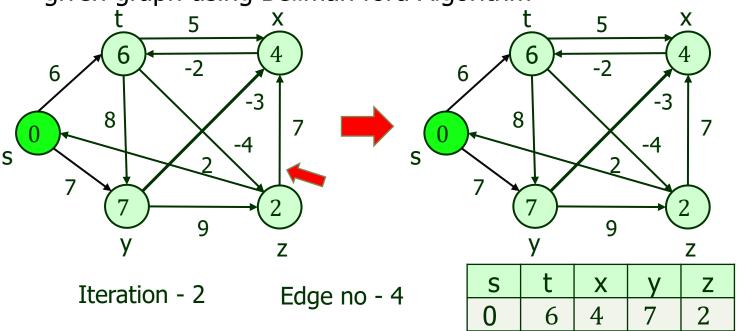
Iteration - 2

S	t	Х	У	Z
0	6	4	7	2

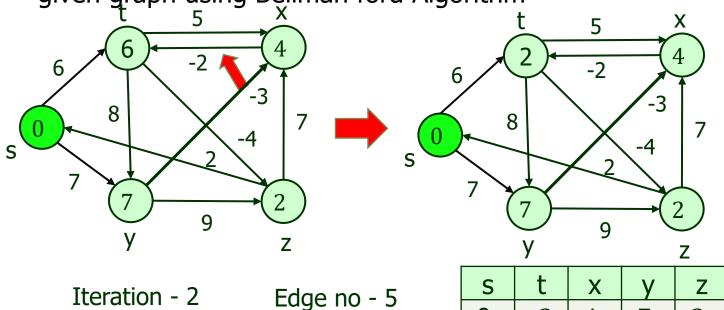
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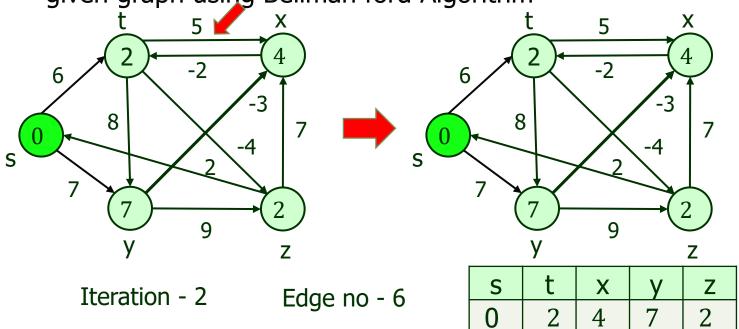
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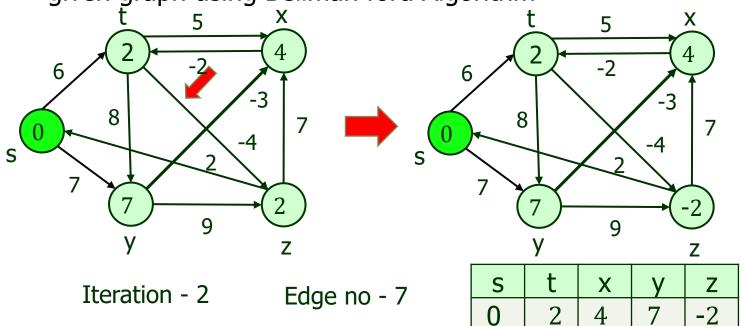
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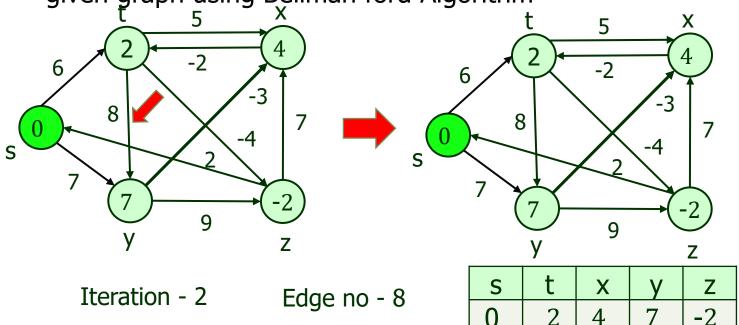
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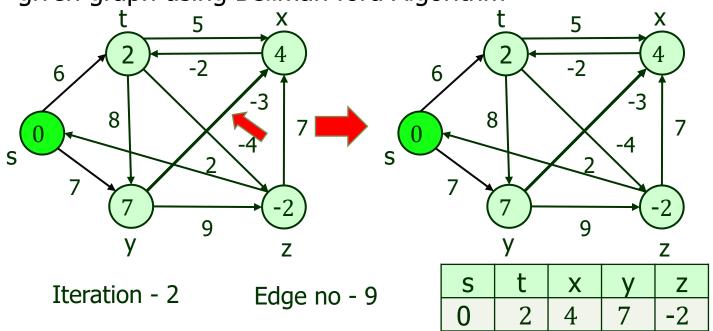
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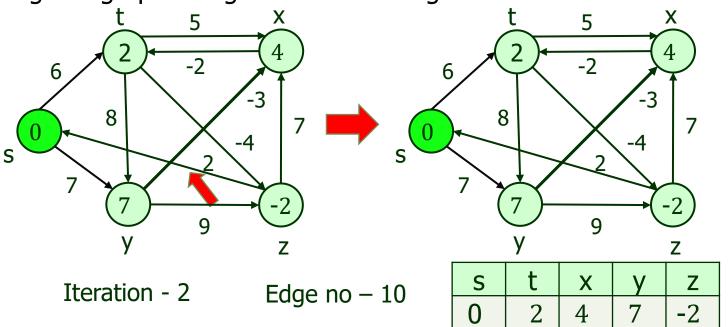
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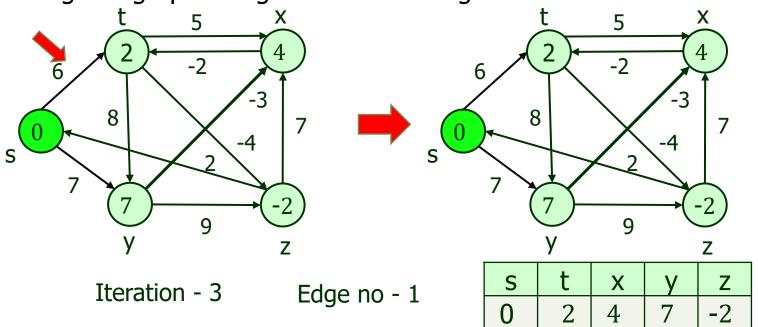
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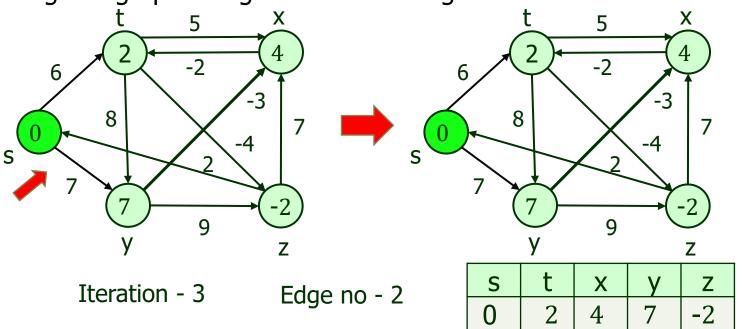
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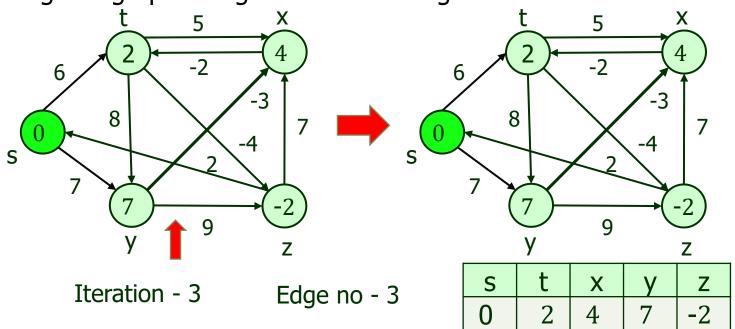
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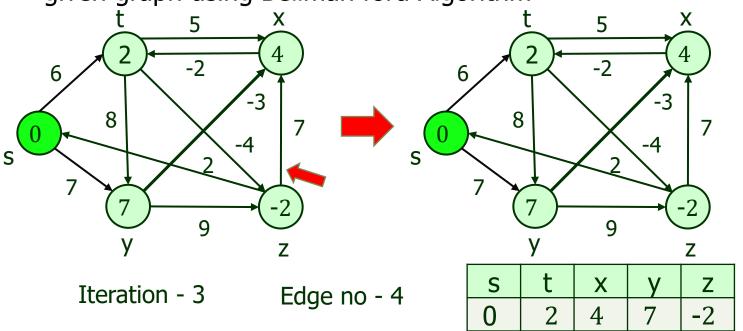
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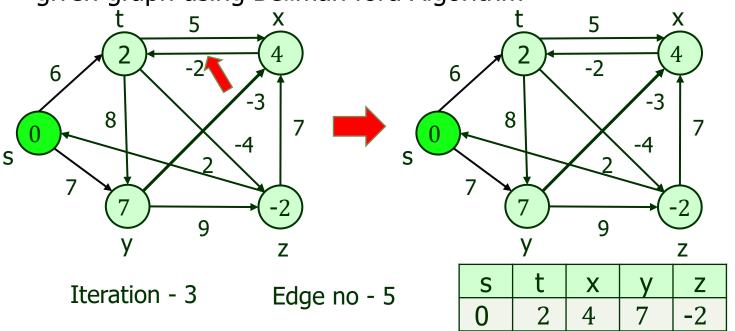
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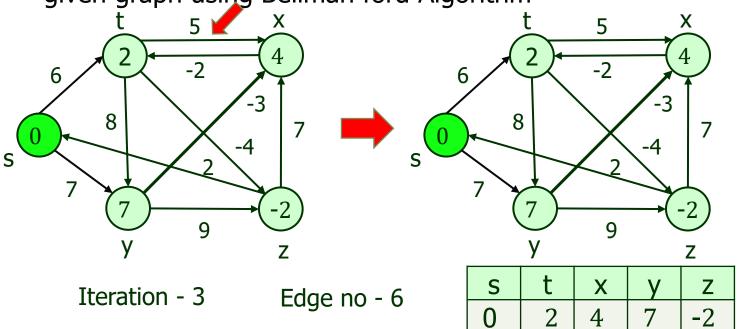
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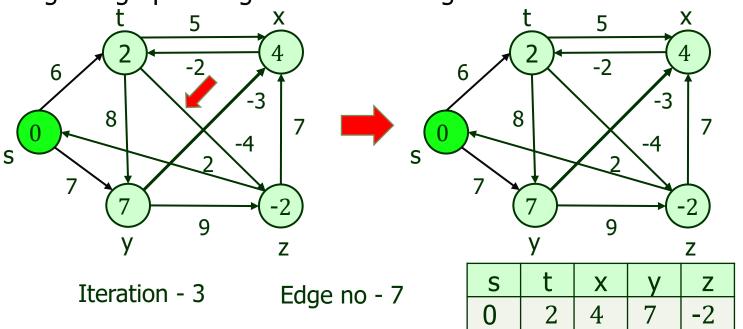
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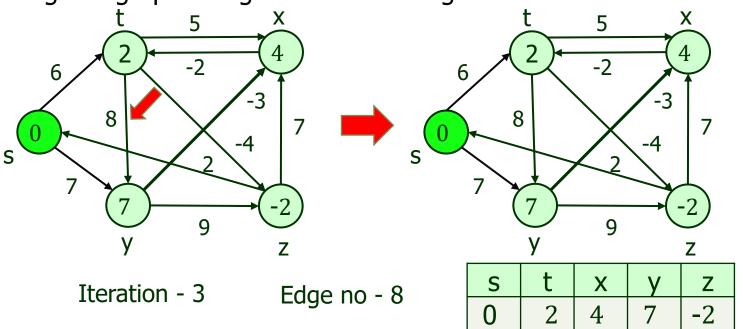
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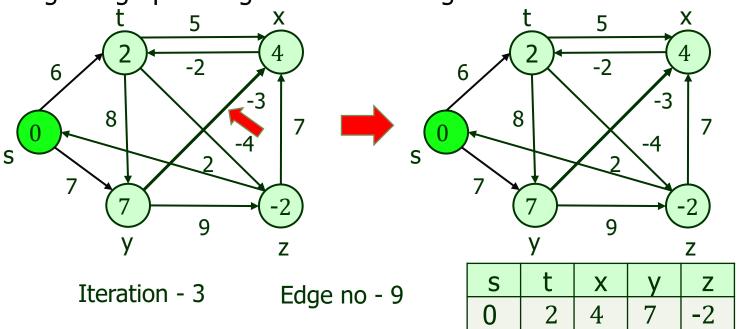
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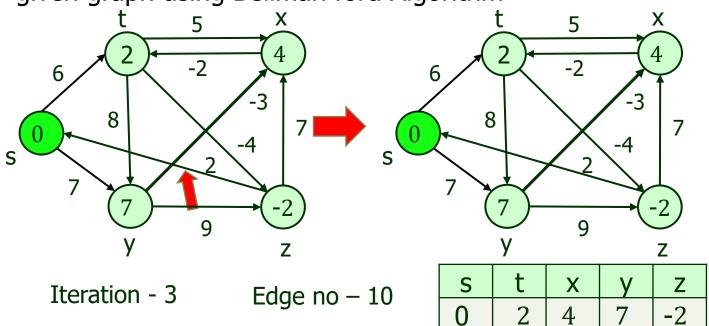
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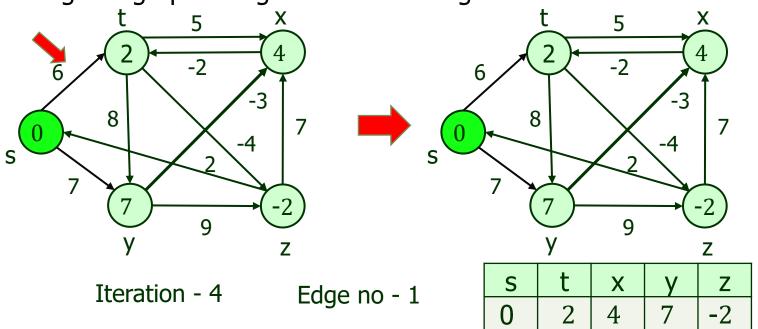
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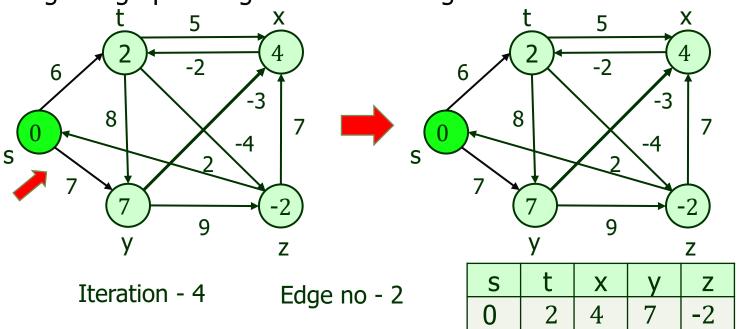
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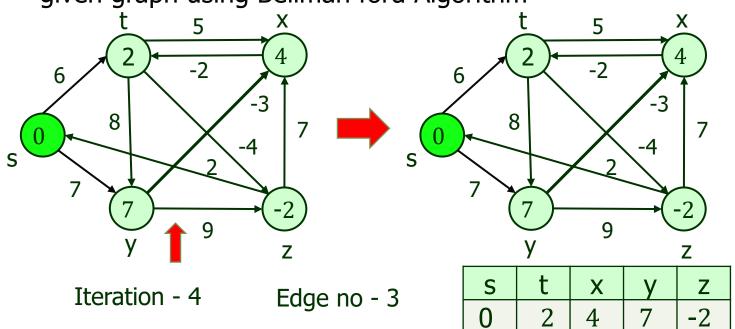
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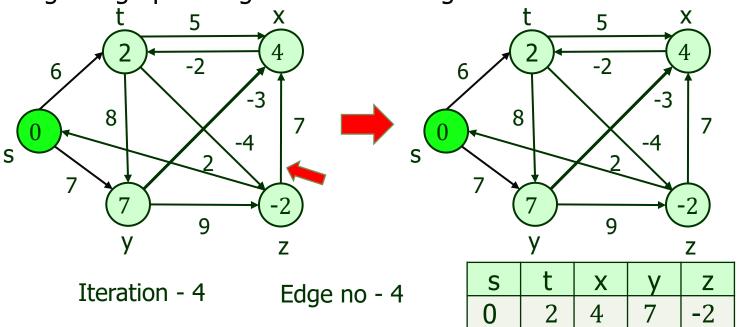
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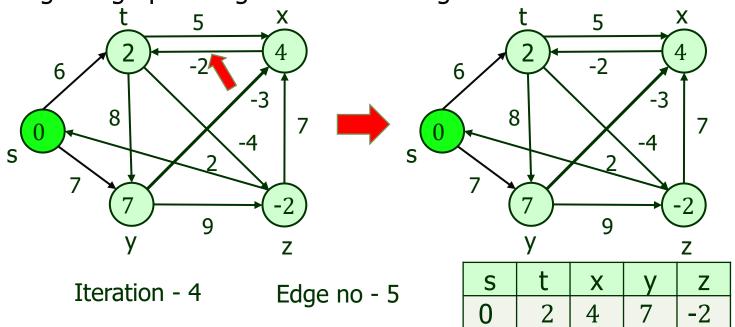
- Problem 5: Single source shortest path
  - Bellman Ford Algorithm



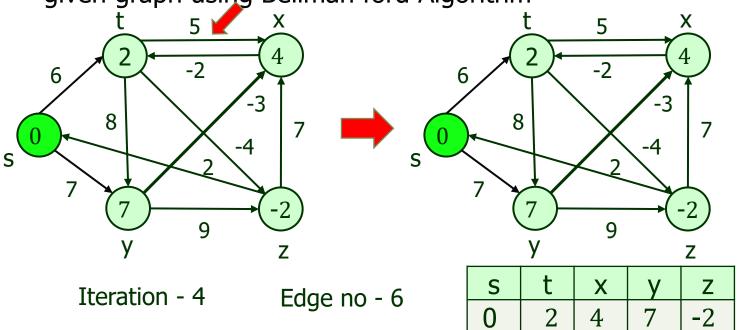
- Problem 5: Single source shortest path
  - Bellman Ford Algorithm



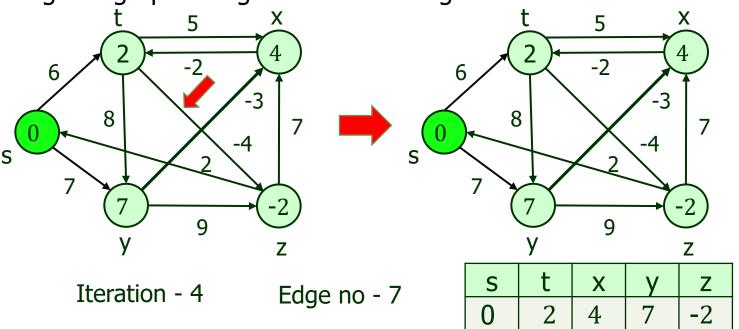
- Problem 5: Single source shortest path
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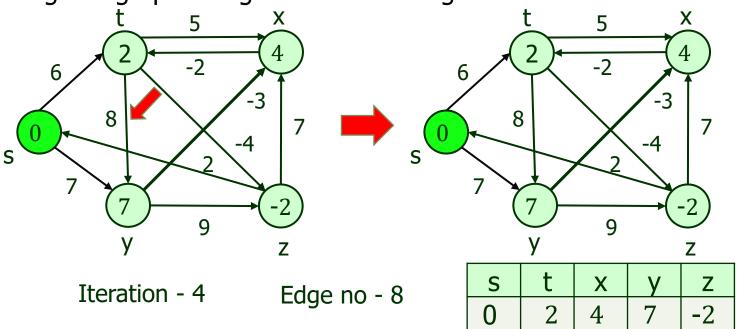
- Problem 5: Single source shortest path
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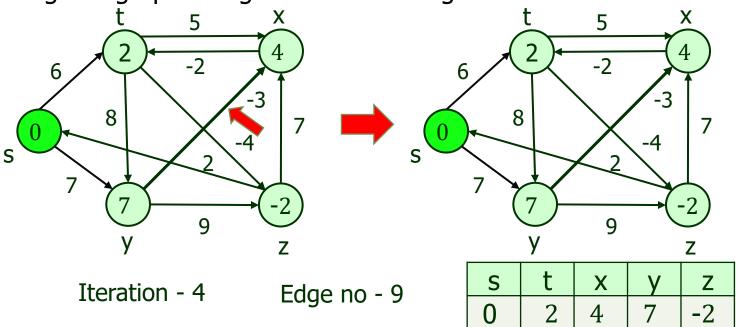
- Problem 5: Single source shortest path
  - Bellman Ford Algorithm



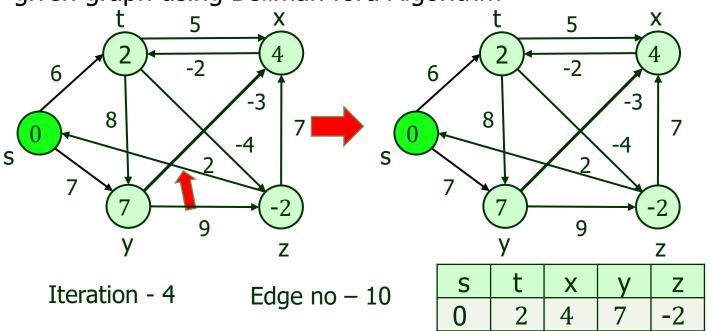
- Problem 5: Single source shortest path
  - Bellman Ford Algorithm



- Problem 5: Single source shortest path
  - Bellman Ford Algorithm



- Problem 5: Single source shortest path
  - Bellman Ford Algorithm



- Problem 5: Single source shortest path
  - Bellman Ford Algorithm

Example 1: Construct the Single source shortest path for the given graph using Bellman ford Algorithm-

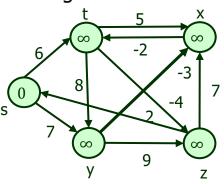
It was observed that after 3<sup>rd</sup> iteration there was no relaxation. So it was wise to stop after 3<sup>rd</sup> iteration and the final answer is :

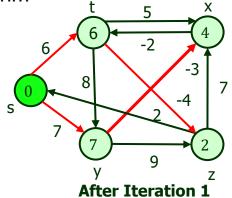
S	t	X	У	Z
0	2	4	7	-2

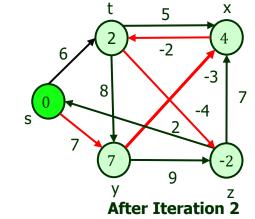
- Problem 5: Single source shortest path
  - Bellman Ford Algorithm

Example 1: Construct the Single source shortest path for the given graph

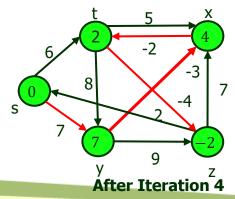
using Bellman ford Algorithm-







5 X 2 -2 -4 7 7 7 9 -2 7 After Iteration 3



- Problem 5: Single source shortest path
- Bellman Ford Algorithm (Algorithm)

```
INITIALIZE-SINGLE-SOURCE(G, s)
1 for each vertex v V[G]
         do d[v] \leftarrow \infty
         \Pi[V] \leftarrow NIL
4 d[s] \leftarrow 0
RELAX(u, v, w)
1 if d[v] > d[u] + w(u, v)
2 then d[v] \leftarrow d[u] + w(u, v)
3
         \Pi[V] \leftarrow U
```

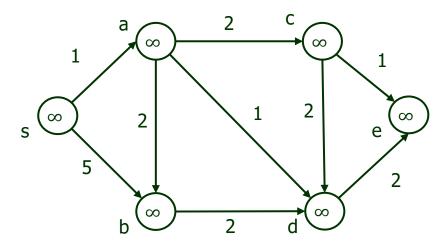
- Problem 5: Single source shortest path
- Bellman Ford Algorithm (Analysis)

```
BELLMAN-FORD(G, w, s)
   INITIALIZE-SINGLE-SOURCE(G, s) \rightarrow \Theta(V)
2 for i \leftarrow 1 to |V[G]| - 1
        do for each edge (u, v) \in E[G]
                                            O(VE)
             do RELAX(u, v, w)
  for each edge (u, v) \in E[G]
        do if d[v] > d[u] + w(u, v)
                                             O(E)
             then return FALSE
8 return TRUE
```

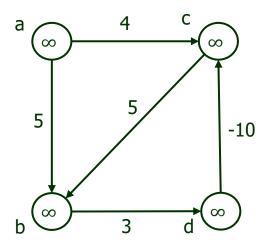
- Problem 5: Single source shortest path
- Bellman Ford Algorithm (Algorithm)

```
BELLMAN-FORD(G, w, s)
   INITIALIZE-SINGLE-SOURCE(G, s)
2 for i \leftarrow 1 to |V[G]| - 1
        do for each edge (u, v) \in E[G]
3
            do RELAX(u, v, w)
5 for each edge (u, v) \in E[G]
        do if d[v] > d[u] + w(u, v)
6
            then return FALSE
8 return TRUE
```

- Problem 5: Single source shortest path
  - Bellman Ford Algorithm(Self Practice)



- Problem 5: Single source shortest path
  - Bellman Ford Algorithm(Self Practice)



- Problem 5: Knapsack Problem
   Problem definition:
  - The Knapsack problem is an "combinatorial optimization problem", a topic in mathematics and computer science about finding the optimal object among a set of object.
  - Given a set of items, each with a weight and a profit, determine the number of each item to include in a collection so that the total weight is less than or equal to a given limit and the total profit is as large as possible.

- Problem 5: Knapsack Problem
   Versions of Knapsack:
  - Fractional Knapsack Problem:

Items are divisible; you can take any fraction of an item and it is solved by using **Greedy Algorithm.** 

0/1 Knapsack Problem:

Items are indivisible; you either take them or not and it is solved by using **Dynamic Programming(DP).** 

- Problem 5: Knapsack Problem
  - Fractional Knapsack Problem:

Given n objects and a knapsack with a capacity "M" (weight)

- Each object i has weight  $w_i$  and profit  $p_i$ .
- For each object i, suppose a fraction  $x_i$ ,  $0 \le x_i \le 1$ , can be placed in the knapsack, then the profit earned is  $p_i \cdot x_i$ .

### Problem 5: Knapsack Problem

#### **Fractional Knapsack Problem:**

The objective is to maximize profit subject to capacity constraints.

i.e.maximize

$$\sum_{i=1}^{n} p_i x_i$$
 -----(1)

Subject to

$$\sum_{i=1}^{n} w_i x_i \le M$$
 -----(2)

Where, 
$$0 \le x_i \le 1$$
,

$$p_i > 0$$
,

$$w_i > 0$$
.

A feasible solution is any subset  $\{x_1, x_2, x_3, \dots, x_n\}$  satisfying (1) & (2). An optimal solution is a feasible solution that maximize  $\sum_{i=1}^{n} p_i x_i$ .

#### Problem 5: Knapsack Problem

#### Fractional Knapsack Problem(Implementation)

Fractional knapsack problem is solved using greedy method in the following steps-

#### Step-01:

• For each item, compute its (profit / weight) ratio.(i.e  $p_i/x_i$ )

#### Step-02:

 Arrange all the items in decreasing order of their (profit / weight) ratio.

#### Step-03:

- Start putting the items into the knapsack beginning from the item with the highest ratio.
- Put as many items as you can into the knapsack.

#### Problem 5: Knapsack Problem

#### Fractional Knapsack Problem(Implementation)

Example 1: For the given set of items and knapsack capacity = 60 kg, find the optimal solution for the fractional knapsack problem making use of greedy approach.

Item	Weight	Profit
1	5	30
2	10	40
3	15	45
4	22	77
5	25	90

• Problem 5: Knapsack Problem

#### Fractional Knapsack Problem(Implementation)

Example 1: For the given set of items and knapsack capacity = 60 kg, find the optimal solution for the fractional knapsack problem making use of greedy approach.

Step-01: Compute the (profit / weight) ratio for each item-

#### Problem 5: Knapsack Problem

#### Fractional Knapsack Problem(Implementation)

Example 1: For the given set of items and knapsack capacity = 60 kg, find the optimal solution for the fractional knapsack problem making use of greedy approach.

Step-01: Compute the (profit / weight) ratio for each item(i.e. one unit

cost)-

Item	Weight	Profit	Ratio
(i)	$(w_i)$	$(p_i)$	$\binom{p_i}{w_i}$
1	5	30	6
2	10	40	4
3	15	45	3
4	22	77	3.5
5	25	90	3.6

#### Problem 5: Knapsack Problem

#### Fractional Knapsack Problem(Implementation)

Example 1: For the given set of items and knapsack capacity = 60 kg, find the optimal solution for the fractional knapsack problem making use of greedy approach.

Step-02: Sort all the items in decreasing order of their value / weight

ratio-

Item (i)	Weight $(w_i)$	Profit $(p_i)$	Ratio $\binom{p_i}{w_i}$
1	5	30	6
2	10	40	4
5	25	90	3.6
4	22	77	3.5
3	15	45	3

#### Problem 5: Knapsack Problem

#### Fractional Knapsack Problem(Implementation)

Example 1: For the given set of items and knapsack capacity = 60 kg, find the optimal solution for the fractional knapsack problem making use of greedy approach.

Item	Weight	Profit	Ratio
(i)	$(w_i)$	$(p_i)$	$(p_i/w_i)$
1	5	30	6
2	10	40	4
5	25	90	3.6
4	22	77	3.5
3	15	45	3

Knapsack weight	Items in Knapsack	Cost
60	Ø	0

#### Problem 5: Knapsack Problem

#### Fractional Knapsack Problem(Implementation)

Example 1: For the given set of items and knapsack capacity = 60 kg, find the optimal solution for the fractional knapsack problem making use of greedy approach.

Item	Weight	Profit	Ratio
(i)	$(w_i)$	$(p_i)$	$\left \binom{p_i}{w_i}\right $
1	5	30	6
2	10	40	4
5	25	90	3.6
4	22	77	3.5
3	15	45	3

Knapsack weight	Items in Knapsack	Cost
60	Ø	0

#### Problem 5: Knapsack Problem

#### Fractional Knapsack Problem(Implementation)

Example 1: For the given set of items and knapsack capacity = 60 kg, find the optimal solution for the fractional knapsack problem making use of greedy approach.

Item	Weight	Profit	Ratio
1	5	30	6
2	10	40	4
5	25	90	3.6
4	22	77	3.5
3	15	45	3

Knapsack weight	Items in Knapsack	Cost
60	Ø	0
55	1	30

#### Problem 5: Knapsack Problem

#### Fractional Knapsack Problem(Implementation)

Example 1: For the given set of items and knapsack capacity = 60 kg, find the optimal solution for the fractional knapsack problem making use of greedy approach.

Item	Weight	Profit	Ratio
(i)	$(w_i)$	$(p_i)$	$\left \binom{p_i}{w_i}\right $
1	5	30	6
2	10	40	4
5	25	90	3.6
4	22	77	3.5
3	15	45	3

Knapsack		
weight	Knapsack	Cost
60	Ø	0
55	1	30

#### Problem 5: Knapsack Problem

#### Fractional Knapsack Problem(Implementation)

Example 1: For the given set of items and knapsack capacity = 60 kg, find the optimal solution for the fractional knapsack problem making use of greedy approach.

Item	Weight	Profit	Ratio
(i)	$(w_i)$	$(p_i)$	$\left \binom{p_i}{w_i}\right $
1	5	30	6
2	10	40	4
5	25	90	3.6
4	22	77	3.5
3	15	45	3

Knapsack	Items in	
weight	Knapsack	Cost
60	Ø	0
55	1	30
45	1, 2	70

#### Problem 5: Knapsack Problem

#### Fractional Knapsack Problem(Implementation)

Example 1: For the given set of items and knapsack capacity = 60 kg, find the optimal solution for the fractional knapsack problem making use of greedy approach.

Item	Weight	Profit	Ratio
(i)	$(w_i)$	$(p_i)$	$\binom{p_i}{w_i}$
1	5	30	6
2	10	40	4
5	25	90	3.6
4	22	77	3.5
3	15	45	3

Knapsack		
weight	Knapsack	Cost
60	Ø	0
55	1	30
45	1, 2	70

#### Problem 5: Knapsack Problem

#### Fractional Knapsack Problem(Implementation)

Example 1: For the given set of items and knapsack capacity = 60 kg, find the optimal solution for the fractional knapsack problem making use of greedy approach.

Item	Weight	Profit	Ratio
(i)	$(w_i)$	$(p_i)$	$\binom{p_i}{w_i}$
1	5	30	6
2	10	40	4
5	25	90	3.6
4	22	77	3.5
3	15	45	3

Knapsack	Items in	
weight	Knapsack	Cost
60	Ø	0
55	1	30
45	1,2	70
20	1,2,5	160

#### Problem 5: Knapsack Problem

#### Fractional Knapsack Problem(Implementation)

Example 1: For the given set of items and knapsack capacity = 60 kg, find the optimal solution for the fractional knapsack problem making use of greedy approach.

Item	Weight	Profit	Ratio
(i)	$(w_i)$	$(p_i)$	$\left \binom{p_i}{w_i}\right $
1	5	30	6
2	10	40	4
5	25	90	3.6
4	22	77	3.5
3	15	45	3

Knapsack	Items in	
weight	Knapsack	Cost
60	Ø	0
55	1	30
45	1,2	70
20	1,2,5	160

# Problem 5: Knapsack Problem Fractional Knapsack Problem(Implementation)

#### Example 1:

Item	Weight	Profit	Ratio
(i)	$(w_i)$	$(p_i)$	$\left \binom{p_i}{w_i}\right $
1	5	30	6
2	10	40	4
5	25	90	3.6
4	22	77	3.5
3	15	45	3

Knapsack weight	Items in Knapsack	Cost
60	Ø	0
55	1	30
45	1,2	70
20	1,2,5	160
0	1,2,5,	220
<u> </u>	frac(4)	230

Now, Knapsack weight left to be filled is 20 kg but item-4 has a weight of 22 kg. Since in fractional knapsack problem, even the fraction of any item can be taken. So, knapsack will contain the fractional part of item 4.(20 out of 22)

Total cost of the knapsack =  $160 + (20/22) \times 77 = 160 + 70 = 230$  units

Problem 5: Knapsack Problem
 Fractional Knapsack Problem(Algorithm)

```
FRACTIONAL_KNAPSACK (v, w, M)

1. load = 0

2. i = 1

3. While (load < M) and i \le n

4. if w_i \le M - load

5. take \ all \ of \ item \ i

6. else \ take \ (M - load) / w_i \ of \ item \ i

7. add \ what \ was \ taken \ to \ load \ (load = load + w_i)

8. i=i+1
```

(Note: Assume that the items are already sorted on the basis of ratio)

- Problem 5: Knapsack Problem
   Fractional Knapsack Problem(Algorithm)
  - The main time taking step is the sorting of all items in decreasing order of their (profit / weight) ratio.
  - If the items are already arranged in the required order, then while loop takes O(n) time.
  - The average time complexity of Quick Sort is  $O(n \log n)$ .
  - Therefore, total time taken including the sort is  $O(n \log n)$ .

### Problem 5: Knapsack Problem

### Fractional Knapsack Problem(self practice)

Example 2: For the given set of items and knapsack capacity = 60 kg, find the optimal solution for the fractional knapsack problem making use of greedy approach.

Item	Weight	Profit
Α	1	5
В	3	9
С	2	4
D	2	8

#### Problem 5: Knapsack Problem

### Fractional Knapsack Problem(self practice)

Example 3: Find the optimal solution for the fractional knapsack problem making use of greedy approach. Consider-

```
n = 3

M = 20 kg

(w1, w2, w3) = (18, 15, 10)

(p1, p2, p3) = (25, 24, 15)
```

