a) What are the properties of Entity called as: - Attributes CH-2:- DB DESIGN AND ER MODEL: Data Modeling: - A model is a representation of reality that retains only carefully objected essential D details.

Logical org. of data for optimum Pylo. extraction & data Manipulation. Why Model ? · To understand & identify exertal data elements.

To produce a representation that can be transformed into a scheme Data Hodeling is to be able to use DB to store data, it should be designed in an efficient manner. The 1st step in designing a DB is data modeling. It enables a DB designer to execute a model that represents the way in which infortikely to be augarized in the DB. There are 2 major categories methodologies used to create a data Model: · ER approach & · Object model > lun pui Jocus is the design of the DB → BB Design: · Foundation of a successful info. System should promote:--data integrity - Prevent data redundancies Lanomalies · Must yield a DB that is effecient in its puovision of data access. It > DB duign Strategius: · Top-down design (e.g:-ER modeling)
Identify entitles Idata onto Define attributes / data elements for each entity. · Dor Bottom-up design (e.g. - normatigation) Identify attributes
Group them together to define entitles

· Centralized design:

-Bruell no. of Jobj. & procedures

-Brugle design process 3 Supports 3 · Decentralized design: 2 Bottom-3 -large no. of entities with complex relations & oft. up design -multiple parallel design of subsystems & aggregation 3 3 > Steps in disign Process-1) Requirements analysis (where the org. is presently working & where they want to work ii) Conceptual DB design ] - format of data to othere AT) Logical DB design ] in future M) Schema refinement (otr defined) v) Physical DB design (now actually data stored in physical servers) vi) Appla & Security design > Feasibility Analysis: (Is the design actually possible in real life?) · Technological Jessibility: - What How, I sow & additional would be needed!
- What is available in-house? what has to be punchased? - How will the new system be integrated? · Openational leasibility: (day to day activities) - who will design the dystem? -,, ,, do training on help-dush supposet? - Can the available penson provide the time · Lepected Doost of the overall project? - 3/w, H/w, appla development, Staff-Have - Hidden cost (unforcescen) Benefits
How doon expected? Steps to Design CR Diagrams: Step1: - Identify the obtoing & weak dull ty bets Step 2: - Identify the Relevant Attributes The next oftep is to get all the attributes that are next applicable for each entity out. Do this work by considering each entity out in mind & also the type of attributes.

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3

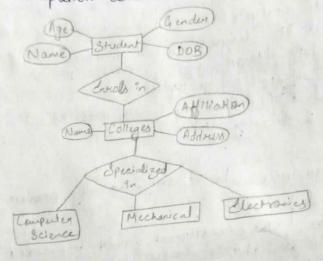
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Step 33- Identify the Relationship Sets In this oslep we need to find all the meaningful relations tup dets among possible entity osets. This oslep is very toicky, as redundant relationships may lead to complicated design & in twen a bad implementation.

Step 4 :- Indentify the Cardinality (no. of tuples) Ratio & Participation Constraints.

Sample:



# Design Issues:

1. Use of entity dets Vs. Attributes:

- · Choice mainly depends on the sto. of the enterprise being modeled,
- · What is an dutity & Affribute?
- · When to represent a value as an attribute?
- · 11 11 11 11 11 a departe intity det?
- · Representing as a operande entity out allows details to be added later.
- 2. It is difficult to examine of an obj. can be best expressed by an entity bet on sulationship det.
  - To understand & determine the eight use, the user need to destionate a relationship set for describing an action that occurs in b/w the
- Generally, the relationships described in the DB are binary relationships described in the DB are binary relationships can be represented by a everal binary relationships.

Frendeg: we can create & represent a ternary relationship

"parent" that may relate to a child, his father, as well as his mother. The condinality nation can become an affective measure in the placement of the relationship attributes. So, it is better to associate the attributes of one-to-one on one-to-many relationship sets with #Weak dutity sets: . It is one that can only exist when owned by another one. Example: - a ROOM can only exist in a BUILDING. On the other hand, a TIRE might be considered as a strong entity beog it also exist without being attached to a CAR. · Unlike strong entity, a weak entity does not have any pt. key. It instead has a partial discumenator key.

· A weak entity is represented by a double rectangle.

The combination of discriminator of prickey of the strong entity of the makes it possible to uniquely identify all entitles of the weak entity oset. the weak entity set. Dri. Key of Weak Entity = lits own discurrentor + Pri. Key of Strong Entity Set Example: - In a University, a course is a strong entity, & a course offering can be modeled as a weak entity. Would be demester & section no. (4) # extentded Extentives: As the complexity of data is increasing, it became more I more difficult to use the traditional ER model for DB modeling. Hence Dome improvements on enhancements were made to the existing IR Hodel to make it able to handle the complex applu better. · Generalization · Specialization · Hagregation These is used for data abstraction in which abstraction uchanism is used to wide details of a set of obj.

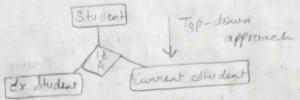
Generalization:

It is a bottom-up approach in which I lower level entities combine to form a higher level entity. In this, the higher-level entity can also combine with other lower level entity to make further higher level entity.

Saving A Convert

Decialization: opp. to generalization.

It is a top-down approach in which one higher level entity can be broken down into 2 lower-level entity. In spect this, a higher level entity may not have any lower-level entity sets, it's possible.



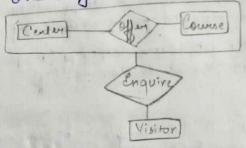
A process when relation blu 2 entitles is treated as a dingle entity

It is a process in which a dingle entity alone is not able to

make sense in a relationship so the relationship of 2 entitles

acts as one entity.

-



# Relational DB:

RDBMS is used to manage Relational DB. Relational DB is a collection of organized oset of tables related to each other, I from which data can be accessed easily. There DB is the most bused DB Tuple: row/record

Attribute: column

All ribute Domain: When an attribute is defined in a welation (table), it is defined to hold only a certain type of values, which is

-> Relation Schema: It describes the other. of the relation, with the name of the relation (manue of table), its attributes & their names > Relational Integrity Constraints: every relation has done conditions that must hold for it to be a valid relation. These conditions are called Relational Integrity Constraints. There are 3 main integrity constraints-· Key constraints. · Domain constraints · Referential integrity constraints. · Key Constraints: In a relation with a key attribute, no Etuples can have identical values for key attributes. A key attribute ear not have NULL values. - Pri- key ; Unique key · Domain Constraints & Attributes have specific values in real-world scenario. eg: - age can only be tre int. The Every attribute is bound to have a specific range of values. - Check constraint. · Referential Integrity Constraints: Works on the concept of Foreign keys. If a relation refers to a key attribute of a diff. or same relation, then that key element much exist. # Relational Algebra Operators: - Wit theory of with & Base of SQL) · SELECT Operations works on tuples Notation: - op(r) p is called the odection predicate dxample: - OCITY= "AMBALA" (STUDENT) o = Select operator Unarry T= relation op evalory · PROJECT Operation: works on attributes Notation: - TTAI, A2, ..., AK (r) where A1, A2 are attribute names & ris the relation name. TT = Project operator Example: TTCITY, CLASS (STUDENT)

Duplicates are nemoved

· UNION Operation: Binary operator Notation: rus

Fay rus to be valid.

1. I do nust have the same only ( dance no. of attributes)

2. The attribute domains must be compatible

red are 2 relations.

example: - To find all the customers with either an account or aloan. 11 customer-name (depositor) UTTenstomer-name (borrower)

Duplicates not allowed.

· SET DIFFERENCE Operation:

Notation: r-s

3. This opd, must be taken blw compatible relations. 8 & & must have the same ority.

attribute domains of r 2 d must be compatible.

> All the values of r which are not present in os.

· C'ARTESIAN PRODUCT:

Notation: TXS

Applying this on 2 relations that is one two dets of tuples, it will take every tuple one by one from the left det (relation) & will pair it up with all the tupled in the bright det (relation)

· KENAME Operator :-

Allows us to name, & 00, to englar to, the results of Helational - algebra expressions:

Allows us to sujer to a relation by more than one name \* xample :- PX(E)

secturns the expression E under the name X.