

1

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[]:

l1 to l3 cache

Binary code

fetch line by line

decode

execute

In detail by naveen notes, where?

Text

compiled code, loaded from memory

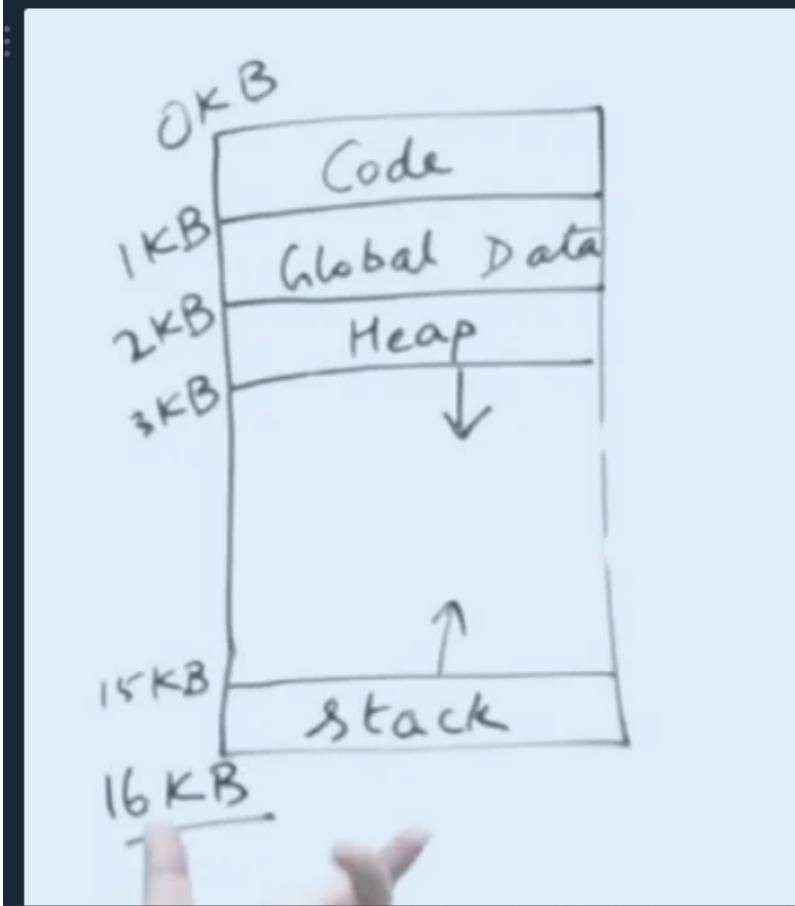
Data

Global and static data

Stack

local variables, function args and return values

This representation of this is called the address space of the process



stack pointer

current line that is being executed,

frame pointer

current function ends, next function start

#Components of OS 1. kernel

kernel

1. A kernel is that part of the operating system which interacts directly with the hardware
2. Heart of OS/Core component
3. Very first part of OS to load on start-up Apps interact with kernel : using system calls (process is moved to kernel stack from user stack)

0.0.1 Process management

1. Scheduling
2. Creating & deleting both user and system process
3. Suspending and resuming processes
4. process synchronization

5. process communication

0.0.2 Memory management

1. Allocating and deallocating memory space as per need
2. Keeping track of which part of memory are currently being used and by which process.

0.0.3 File management

1. Creating and deleting files
2. Creating and deleting directories to organize files
3. Mapping files into secondary storage

1 IO management

1. Buffering (data copy between two devices)
 1. Within one job.
 2. Youtube video buffering
2. caching
 1. Memory caching
 2. Web caching etc
3. spooling
 1. Within differing speed two jobs
 2. Print spooling and mail spooling

(printer store the data in some area and then print it slowly)

2 User space

Where application software runs, apps don't have privileged access to the underlying hardware. It interacts with kernel

Monolithic kernel

Linux, Unix, MS-DOS

All functions are in kernel itself

Bulky in size

Memory required to run is high

Less reliable, one module crashes -> whole kernel is down

High performance as communication is fast

Less user mode, kernel mode switching overheads

Micro kernel

Symbian OS

Only major functions are in kernel 1. Process mgmt 2. Memory mgmt

Rest in User space File mgmt. IO mgmt.

smaller in size

More Reliable and More stable

Performance is slow (Overhead switching b/w user mode and kernel mode)

Hybrid kernel

MacOS, Windows

same as above except

IO in kernel

system calls

calls to os to get resources

mmap for available memory

scanf is wrapper over read()

Transitions from US to KS done by

1. hardware interrupt

urgent interrupt or non urgent

keyboard and mouse, creates an interrupt

2. software interrupts

traps/exception/faults

special kind of software interrupt

division by 0

Illegal memory access

1. Degree of multi-programming: The number of processes in the memory
2. LTS controls degree of multi-programming