

# UNIT-7

## Number Theoretic Algorithms:

Small easier unit so read everything instead of imp only ✓

### Number Theoretic Notations:

i) Divisibility and Divisors: The notation  $d|a$  (read "d divides a") means that  $a=kd$  for some integer  $k$ . If  $d|a$ , then we say that 'a' is a multiple of  $d$ . Every integer divides 0. If  $a>0$  and  $d|a$ , then  $|d| \leq |a|$ . If  $d|a$  and  $d \geq 0$ , we say that  $d$  is a divisor of  $a$ . A divisor of an integer 'a' is at least 1 but not greater than  $|a|$ .

For example: The divisors of 24 are 1, 2, 3, 4, 6, 8, 12 and 24.

trivial divisors of 24: 1 and 24

non-trivial divisors of 24: 2, 3, 4, 6, 8 and 12

→ 1 and the no. itself

→ other remaining ones than that of trivial

### Prime and Composite Numbers:

An integer  $a>1$  whose only divisors are the trivial divisors (i.e. 1 and 'a' itself only) is said to be a prime number or simply prime.

An integer  $a>1$  that is not prime is said to be a composite number. The integer 1 is said to be a unit and is neither prime nor composite. Similarly, the integer 0 and all negative integers are neither prime nor composite.

### Common Divisors and Greatest Common Divisors:

If 'd' is a divisor of 'a' and also divisor of 'b', then 'd' is a common divisor of 'a' and 'b'. Note that 1 is a common divisor of any two integers. An important property of common divisors is that:

→ if  $d|a$  and  $d|b$  then,  $d|(a+b)$  and  $d|(a-b)$ .

→ if  $d|a$  and  $d|b$  then,  $d|(ax+by)$

The greatest common divisor  $\gcd(a, b)$  is the largest of the common divisors of  $a$  and  $b$ . We define  $\gcd(0, 0)$  to be 0.

→ For any integers  $a$  and  $b$  if  $d|a$  and  $d|b$  then,  $d|\gcd(a, b)$ .

→ For all integers  $a$  and  $b$  and any non-negative integer  $n$ ,

$$\gcd(an, bn) = n \gcd(a, b)$$

→ For all positive integers  $n, a$ , and  $b$ , if  $n|ab$  and  $\gcd(a, n)=1$ , then  $n|b$ .



### Relatively Prime Integers:

Two integers  $a, b$  are said to be relatively prime if their only common divisor is 1, that is if  $\gcd(a, b) = 1$ .

→ For any integers  $a, b$ , and  $p$ , if both  $\gcd(a, p) = 1$  and  $\gcd(b, p) = 1$  then,  $\gcd(ab, p) = 1$ .

→ For all primes  $p$  and all integers  $a, b$  if  $p \nmid ab$ , then  $p \nmid a$  or  $p \nmid b$  or both.

### Euclid's Algorithm for solving Modular Linear Equations: [Imp]

Euclidean algorithm is an efficient method for computing the greatest common divisor of two numbers.

#### Algorithm:

EUCLID( $a, b$ )

{ if  $b = 0$   
  return  $a$ ;

else

  return EUCLID( $b, a \bmod b$ )

}

Analysis: Since it is recursive algorithm so we need to find their recurrence relation. Since every time the problem is divided into two parts one is  $b$  and another is  $a \bmod b$ .

Thus the size of sub-problem =  $\frac{n}{2}$

Dividing and merging time = constant =  $O(1)$

Thus, recurrence relation is,  $T(n) = T(n/2) + O(1)$

By, solving this we get,  $T(n) = O(\log n)$ .

Example: Find  $\gcd(30, 21)$  by Euclid's Algorithm.

#### Solution

EUCLID(30, 21)

= EUCLID(21,  $30 \bmod 21$ )

= EUCLID(21, 9)

= EUCLID(9,  $21 \bmod 9$ )

= EUCLID(9, 3)

= EUCLID(3,  $9 \bmod 3$ )

= EUCLID(3, 0)

Since  $b = 0$ , so return  $a$

= 3

∴  $\gcd(30, 21) = 3$ .



## \* Extended Euclid's Algorithm for solving Modular Linear Equation: [Imp]

Extended Euclid's Algorithm is an extension to Euclid's algorithm which computes the coefficients of Bezout's identity, (which are integers  $x$  and  $y$ ) in addition to the greatest common divisors of integers  $a$  and  $b$  such that:

$$d = \gcd(a, b) = ax + by$$

where,  $x$  and  $y$  may be zero or negative.

### Algorithm:

EXTENDED-EUCLID( $a, b$ )

{ if  $b = 0$

return ( $a, 1, 0$ );

( $d', x', y'$ )  $\leftarrow$  EXTENDED-EUCLID( $b, a \bmod b$ )

( $d, x, y$ )  $\leftarrow$  ( $d', y', x' - \text{floor}(a/b)y'$ )

return ( $d, x, y$ );

}

Analysis: Same as Euclid's algorithm i.e.,  $T(n) = O(\log n)$ .

Example: Find GCD(161, 28) and value of  $x$  and  $y$  by using extended Euclidean's algorithm.

Solution:

We have,  $a = 161$ ,  $b = 28$  and  $\gcd(a, b) = ax + by$ .

Let's define following three equations;

$$r = r_1 - q * r_2$$

$$x = x_1 - q * x_2$$

$$y = y_1 - q * y_2$$

Consider  $a = r_1$  and  $b = r_2$ .

formulas that will be used to calculate  $r, x, y$  in below table

always initialize  
 $x_1 = 1, x_2 = 0$   
&  $y_1 = 0, y_2 = 1$

remainder

$q$	$r_1$	$r_2$	$r$	$x_1$	$x_2$	$x$	$y_1$	$y_2$	$y$
5	161	28	21	1	0	1	0	1	-5
1	28	21	7	0	1	-1	1	-5	6
3	21	7	0	1	-1	4	-5	6	-23
	7	0		-1	4		6	-23	

$$x = -1, y = 6$$

Since:  $ax + by = \gcd(a, b)$

$$\text{or, } \{161 * (-1)\} + (28 * 6) = 7$$

$$\text{or, } 7 = 7$$

마지막으로  
step #4  
 $\rightarrow r = 0$   
so  
last  
step  
shift only



## ⊗ Miller-Rabin Randomized Primality Test:

This test algorithm test determines whether a given number is prime or not.

### Algorithm:

/\* It returns false if  $n$  is composite and returns true if  $n$  is probably prime.  $k$  is an input parameter that determines accuracy level. Higher value of  $k$  indicates more accuracy. \*/

bool IsPrime(int  $n$ , int  $k$ )

1) Handle base cases for  $n < 3$

2) If  $n$  is even, return false.

3) Find an odd number  $d$  such that  $n-1$  can be written as  $d \cdot 2^r$ .  
Note that since  $n$  is odd,  $(n-1)$  must be even and  $r$  must be greater than 0.

4) Do following  $k$  times

if (millerTest( $n, d$ ) == false)  
return false.

5) Return true.

bool millerTest(int  $n$ , int  $d$ )

1) Pick a random number  $a$  in range  $[2, n-2]$

2) Compute:  $x = \text{pow}(a, d) \% n$

3) If  $x == 1$  or  $x == n-1$ , return true.

// Below loop mainly runs ' $r-1$ ' times.

4) Do following while  $d$  doesn't become  $n-1$ .

a)  $x = (x * x) \% n$ .

b) If  $(x == 1)$  return false.

c) If  $(x == n-1)$  return true.

Example: Input:  $n=13, k=2$ .

1). Compute  $d$  and  $r$  such that  $d \cdot 2^r = n-1$ ,  
 $d=3, r=2$ .

2) Call millerTest  $k$  times.

1st Iteration:

1) Pick a random number ' $a$ ' in range  $[2, n-2]$

Suppose  $a=4$

2). Compute:  $x = \text{pow}(a, d) \% n$

$$x = 4^3 \% 13 = 12$$

3). Since  $x = (n-1)$ , return true.

2nd Iteration:

1) Pick a random number ' $a$ ' in range  $[2, n-2]$

Suppose  $a=5$

2) Compute:  $x = \text{pow}(a, d) \% n$

$$x = 5^3 \% 13 = 8$$

3)  $x$  neither 1 nor 12.

4) Do following  $(r-1)=1$  times

$$a) x = (x * x) \% 13 = (8 * 8) \% 13 = 12$$

b) Since  $x = (n-1)$ , return true.

Since both iterations return true, we return true.



## ⊗ Chinese Remainder Theorem:

The Chinese remainder theorem states that if one knows the remainders of the Euclidean division of an integer  $n$  by several integers, then one can determine uniquely the remainder of the division of  $n$  by the product of these integers, under the condition that the divisors are pairwise co-prime.

Statement: If  $m_1, m_2, \dots, m_k$  are pairwise relatively prime positive integers, and if  $a_1, a_2, \dots, a_k$  are any integers, then the simultaneous congruences

$$x \equiv a_1 \pmod{m_1},$$

$$x \equiv a_2 \pmod{m_2}, \dots$$

$$x \equiv a_k \pmod{m_k} \text{ have a solution, and the solution is unique modulo } m.$$

Here we need to calculate  $x$  with the help of following formulas:

$$x = (M_1 x_1 a_1 + M_2 x_2 a_2 + \dots + M_k x_k a_k) \pmod{M}$$

$$M = m_1 \cdot m_2 \cdot \dots \cdot m_k$$

$$M_j = \frac{M}{m_j}$$

$$\& M_j x_j = 1 \pmod{m_j}$$

Example: Solve following congruences by using Chinese remainder theorem.

$$x \equiv 1 \pmod{5}$$

$$x \equiv 1 \pmod{7}$$

$$x \equiv 3 \pmod{11}$$

Solution:

we have,  $a_1 = 1, a_2 = 1, a_3 = 3$

$$m_1 = 5, m_2 = 7, m_3 = 11$$

$$x = (M_1 x_1 a_1 + M_2 x_2 a_2 + M_3 x_3 a_3) \pmod{M}$$

$$M = m_1 \cdot m_2 \cdot m_3 = 5 \times 7 \times 11 = 385$$

$$M_1 = \frac{M}{m_1} = \frac{385}{5} = 77$$

$$M_2 = \frac{M}{m_2} = \frac{385}{7} = 55$$



$$M_3 = \frac{M}{m_3} = \frac{385}{11} = 35$$

Now,  $M_1 x_1 = 1 \pmod{5}$

or,  $77x_1 = 1 \pmod{5}$

or,  $2x_1 = 1 \pmod{5}$

or,  $x_1 = 3$

Since we have to put value of  $x_1$  in such a way that (product of 2 and  $x_1$ )  $\pmod{5}$  becomes equal to 1.  
for eg. let we put 1 then,  $(2 \times 1) \pmod{5} = 2$   
Now we put 3 then,  $2 \times 3 = 6 \pmod{5} = 1$   
So,  $x_1 = 3$

$M_2 x_2 = 1 \pmod{7}$

or,  $55x_2 = 1 \pmod{7}$

or,  $6x_2 = 1 \pmod{7}$

or,  $x_2 = 6$

$M_3 x_3 = 1 \pmod{11}$

or,  $35x_3 = 1 \pmod{11}$

or,  $2x_3 = 1 \pmod{11}$

or,  $x_3 = 6$

Random multiply	mod 7 value
$6 \times 1 = 6$	6 ( $\neq 1$ )
$6 \times 2 = 12$	5 ( $\neq 1$ )
$6 \times 3 = 18$	4 ( $\neq 1$ )
$6 \times 4 = 24$	3 ( $\neq 1$ )
$6 \times 5 = 30$	2 ( $\neq 1$ )
$6 \times 6 = 36$	1 ( $= 1$ )

Hence,  $x_2 = 6$ . on putting 6 and doing mod we got 1

$$\begin{aligned} \therefore x &= (M_1 x_1 a_1 + M_2 x_2 a_2 + M_3 x_3 a_3) \pmod{M} \\ &= (77 \times 3 \times 1 + 55 \times 6 \times 1 + 35 \times 6 \times 3) \pmod{385} \\ &= (231 + 330 + 630) \pmod{385} \\ &= 1191 \pmod{385} \\ &= 36 \text{ Ans.} \end{aligned}$$

Also, we can test the solution as;

$36 \pmod{5} = 1$

$36 \pmod{7} = 1$

$36 \pmod{11} = 3$