

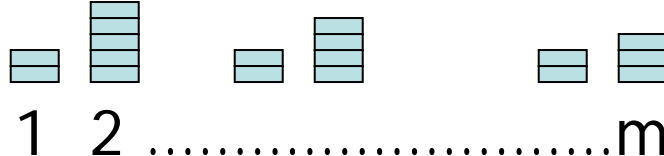
L2 Norm Estimation



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Lecture 2

L2 Norm Estimation

Vector x : 

- A stream is a sequence of updates $(i, 1)$
 $x_i = x_i + 1$
- Want to estimate $\|x\|_2$ up to $1 \pm \epsilon$
- Last week, we have seen how to do that for $\|x\|_0$:
 - Space: $(1/\epsilon + \log m)^{O(1)}$
 - Technique:
 - Linear sketches $\text{Sum}_S(x) = \sum_{i \in S} x_i$ for “random” sets S
 - (Somewhat messy) estimator involving median
- Today: three methods for estimating $\|x\|_2$
 - Alon-Matias-Szegedy
 - Johnson-Lindenstrauss
 - Median-based
 - Really cute and simple
 - Need in future lectures
 - Generalizes to L_p norm for $p < 2$

Why L₂ norm ?

- Database join (on A):
 - All triples (Rel1.A, Rel1.B, Rel2.B)
s.t. Rel1.A=Rel2.A
- Self-join: if Rel1=Rel2
- Size of self-join:

$$\sum_{\text{val of A}} \text{Rows}(\text{val})^2$$
- Updates to the relation
increment/decrement
Rows(val)

Rel1		Rel2	
A	B	A	B
Lec1	distinct	Lec1	distinct
Lec1	elements	Lec1	elements
Lec1	norm	Lec1	norm
Lec2	L2	Lec2	L2
Lec2	norm	Lec2	norm
....		



A	Rel1.B	Rel2.B
Lec1	distinct	distinct
Lec1	distinct	elements
Lec1	distinct	norm
Lec1	elements	distinct
Lec1	elements	elements
	

Algorithm I: AMS

Alon-Matias-Szegedy'96

- Choose $r_1 \dots r_m$ to be i.i.d. r.v., with
 $\Pr[r_i=1]=\Pr[r_i=-1]=1/2$

- Maintain

$$Z = \sum_i r_i x_i$$

under increments/decrements to x_i

- Algorithm A:

$$Y = Z^2$$

- “Claim”: Y “approximates” $\|x\|_2^2$ with “good” probability

Analysis

- The expectation of $Z^2 = (\sum_i r_i x_i)^2$ is equal to

$$E[Z^2] = E[\sum_{i,j} r_i x_i r_j x_j] = \sum_{i,j} x_i x_j E[r_i r_j]$$

- We have
 - For $i \neq j$, $E[r_i r_j] = E[r_i] E[r_j] = 0$ – term disappears
 - For $i = j$, $E[r_i r_j] = 1$

- Therefore

$$E[Z^2] = \sum_i x_i^2 = \|x\|_2^2$$

(unbiased estimator)

Analysis, ctd.

- The second moment of $Z^2 = (\sum_i r_i x_i)^2$ is equal to the expectation of $Z^4 = (\sum_i r_i x_i) (\sum_i r_i x_i) (\sum_i r_i x_i) (\sum_i r_i x_i)$
- This can be decomposed into a sum of
 - $\sum_i (r_i x_i)^4 \rightarrow \text{expectation} = \sum_i x_i^4$
 - $6 \sum_{i < j} (r_i r_j x_i x_j)^2 \rightarrow \text{expectation} = 6 \sum_{i < j} x_i^2 x_j^2$
 - Terms involving **single** multiplier $r_i x_i$ (e.g., $r_1 x_1 r_2 x_2 r_3 x_3 r_4 x_4$)
 $\rightarrow \text{expectation} = 0$

Total: $\sum_i x_i^4 + 6 \sum_{i < j} x_i^2 x_j^2$

- The variance of Z^2 is equal to

$$\begin{aligned}
 E[Z^4] - E^2[Z^2] &= \sum_i x_i^4 + 6 \sum_{i < j} x_i^2 x_j^2 - (\sum_i x_i^2)^2 \\
 &= \sum_i x_i^4 + 6 \sum_{i < j} x_i^2 x_j^2 - \sum_i x_i^4 - 2 \sum_{i < j} x_i^2 x_j^2 \\
 &= 4 \sum_{i < j} x_i^2 x_j^2 \\
 &\leq 2 (\sum_i x_i^2)^2
 \end{aligned}$$

Analysis, ctd.

- We have an estimator $Y=Z^2$
 - $E[Y] = \sum_i x_i^2$
 - $\sigma^2 = \text{Var}[Y] \leq 2 (\sum_i x_i^2)^2$
- Chebyshev inequality^{Wiki} :
$$\Pr[|E[Y]-Y| \geq c\sigma] \leq 1/c^2$$
- Algorithm B:
 - Maintain $Z_1 \dots Z_k$ (and thus $Y_1 \dots Y_k$), define $Y' = \sum_i Y_i / k$
 - $E[Y'] = k \sum_i x_i^2 / k = \sum_i x_i^2$
 - $\sigma'^2 = \text{Var}[Y'] \leq 2k(\sum_i x_i^2)^2 / k^2 = 2 (\sum_i x_i^2)^2 / k$
- Guarantee:
$$\Pr[|Y' - \sum_i x_i^2| \geq c (2/k)^{1/2} \sum_i x_i^2] \leq 1/c^2$$
- Setting c to a constant and $k=O(1/\epsilon^2)$ gives $(1 \pm \epsilon)$ -approximation with const. probability

Comments

- Only needed 4-wise independence of $r_1 \dots r_m$
 - Can generate such vars from $O(\log m)$ random bits
- What we did:
 - Maintain a “linear sketch” vector $\mathbf{Z}=[Z_1 \dots Z_k] = \mathbf{R} \mathbf{x}$
 - Estimator for $\|\mathbf{x}\|_2^2$: $(Z_1^2 + \dots + Z_k^2)/k = \|\mathbf{R}\mathbf{x}\|_2^2 / k$
 - “Dimensionality reduction”: $\mathbf{x} \rightarrow \mathbf{R}\mathbf{x}$
 - ... but the tail somewhat “heavy”
 - Reason: only used second moment of the estimator

Algorithm II: Dim. Reduction (JL)

Interlude: Normal Distribution

- Normal distribution $N(0,1)$:
 - Range: $(-\infty, \infty)$
 - Density: $f(x) = e^{-x^2/2} / (2\pi)^{1/2}$
 - Mean=0, Variance=1
- Basic facts:
 - If X and Y independent r.v. with normal distribution, then $X+Y$ has normal distribution^{WIKI}
 - $\text{Var}(cX) = c^2 \text{Var}(X)$
 - If X, Y independent, then $\text{Var}(X+Y) = \text{Var}(X) + \text{Var}(Y)$

A different linear sketch

- Instead of ± 1 , let r_i be i.i.d. random variables from $N(0,1)$
- Consider

$$Z = \sum_i r_i x_i$$

- We still have that $E[Z^2] = \sum_i x_i^2 = \|x\|_2^2$, since:

- $E[r_i] E[r_j] = 0$
- $E[r_i^2] = \text{variance of } r_i, \text{ i.e., } 1$

- As before we maintain $\mathbf{Z} = [Z_1 \dots Z_k]$ and define

$$Y = \|\mathbf{Z}\|_2^2 = \sum_j Z_j^2 \quad (\text{so that } E[Y] = k \|x\|_2^2)$$

- We show that there exists $C > 0$ s.t. for small enough $\epsilon > 0$

$$\Pr[| Y - k \|x\|_2^2 | > \epsilon k \|x\|_2^2] \leq \exp(-C \epsilon^2 k)$$

Proof

- See the attached notes,
by Ben Rossman and Michel Goemans