

# Reviewing Induction

## Principle of Mathematical Induction

Suppose  $P(k)$  is a property about integers. If:

- 1  $P(1)$  is true, and
- 2 The statement, “For all integers  $k \geq 1$ , if  $P(k)$  is true, then  $P(k + 1)$  is true,” is true;

then  $P(k)$  is true for all  $k \geq 1$ .

- Two parts to a proof by induction: The *base case* and the *inductive step*.
- PMI also works if we start at  $k = 0$ , or  $k = 5$ , etc.
- Notice: Proving the inductive step involves proving a universal conditional statement! Choose an arbitrary element of the domain satisfying the hypothesis. (Choose arbitrary  $k \geq 1$  s.t.  $P(k)$  is true, show  $P(k+1)$  must also be true.)

Rewriting the statement so it is a statement about integers:

**Theorem** (Already proved by contradiction)

$P(k)$

For all  $k \geq 1$ , if a graph  $G$  contains a walk of length  $k$  connecting  $u$  and  $v$  for any distinct  $u, v \in V(G)$ , then  $G$  also contains a path connecting  $u$  and  $v$ .

*base* { Proof (by induction): For the base case, consider  $k=1$ . Let  $G$  be an arbitrary walk of length 1 connecting  $u$  and  $v$  for some distinct  $u, v \in V(G)$ . Since a walk of length 1 cannot repeat vertices, this walk is a path. Thus  $G$  contains a  $u-v$  path.

For the inductive step, let  $k \geq 1$  be an integer, and assume that if  $G$  contains a walk of length  $k$  connecting  $u$  and  $v$  for any distinct  $u$  and  $v$  in  $V(G)$ , then  $G$  also contains a  $u-v$  path. (Inductive hypothesis)

For arbitrary vertices  $x, y \in V(G)$ , suppose  $W$  is an  $x$ - $y$  walk of length  $k+1$  in  $G$ . We must show that  $G$  contains an  $x$ - $y$  path.

Let  $xa$  be the first edge of  $W$ . If  $a=y$ , then  $xa$  is an  $x$ - $y$  path. Otherwise, removing  $xa$  from  $W$  leaves an  $a$ - $y$  walk of length  $k$ .

By the inductive hypothesis, there is an  $a$ - $y$  path in  $G$ ,  $P$ . We consider two cases:

Case 1:  $P$  does not contain  $x$ . Then adding  $xa$  to  $P$  results in an  $x$ - $y$  path in  $G$ .

Case 2:  $P$  does contain  $x$ . Then the subpath of  $P$  connecting  $x$  and  $y$  is an  $x$ - $y$  path in  $G$ .

Therefore  $G$  contains an  $x$ - $y$  path, and the inductive step is complete.

(Therefore by the Principle of Mathematical Induction, the theorem is true.)

# Important: Avoid the induction trap

- In the inductive step, we started with a walk of length  $k + 1$ , and we went and found a walk of length  $k$  inside it.
- You CANNOT start with the walk of length  $k$ , and build up the walk of length  $k + 1$  from there!
- Why not? Because we have to verify the result for an *arbitrary* walk of length  $k + 1$ , not a walk of length  $k + 1$  that was built by adding an edge to a walk of length  $k$ .
- In this case, it is true that all walks of length  $k + 1$  can be built by adding an edge to a walk of length  $k$ , but that won't always be true!

**Always START from the larger object, and FIND the smaller object inside it.**