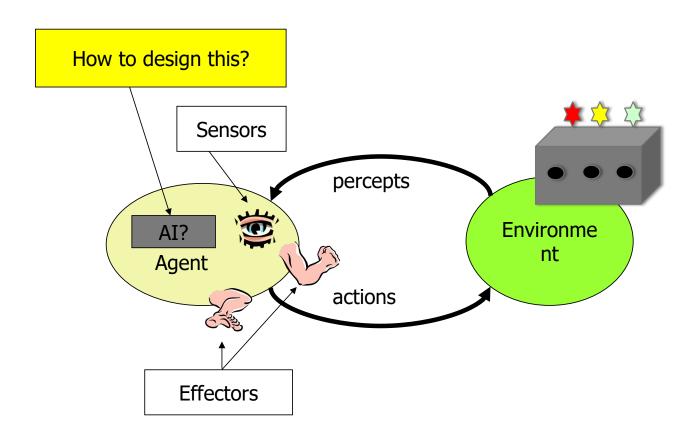
CSCI 561 - Foundation for Artificial Intelligence

DISCUSSION SECTION (WEEK 2)

PROF WEI-MIN SHEN SHEN@ISI.EDU

WHAT IS "PROBLEM SOLVING"? WHAT IS "SEARCH"?

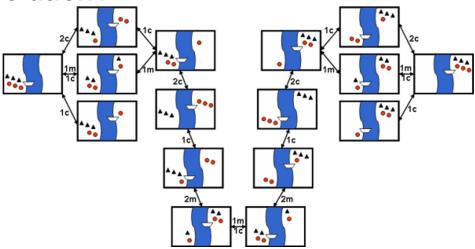


ESSENTIALS OF SEARCH

- How to represent a "problem"?
 - How to construct a Search Tree/Graph?
 - Nodes, Goals, Initials, Links
- How to find a solution "systematically" or "optimally" in your representation?
 - Use the <u>uninformed</u> algorithms you learned
 - Use the <u>informed</u> algorithms you learned

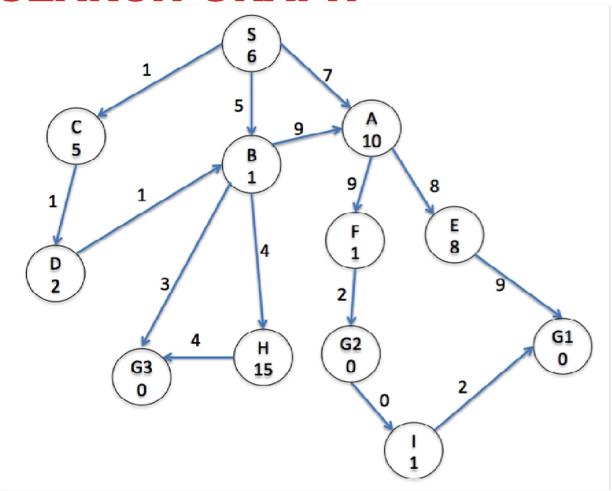
MISSIONARIES AND CANNIBALS

Did you find that there was much search involved in finding a solution?



Why do people have a hard time solving this problem?

SEARCH GRAPH



GRAPH SEARCH

```
function GRAPH-SEARCH(problem) return a solution or failure
  frontier ← MAKE-QUEUE(MAKE-NODE(problem.INITIAL-STATE))
 explored set ← empty
  loop do
       if EMPTY?(frontier) then return failure
       node ← REMOVE-FIRST(frontier)
       if problem.GOAL-TEST applied to node.STATE succeeds
              then return SOLUTION(node)
       explored set \leftarrow INSERT(node, explored set)
       for each new_node in EXPAND(node, problem) do
               if NOT(MEMBER?(new_node, frontier)) and
                NOT(MEMBER?(new node, explored set))
                then frontier ← INSERT(new node, frontier)
```

GRAPH SEARCH

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function GRAPH-SEARCH(problem) return a solution or failure
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if problem.GOAL-TEST applied to node.STATE succeeds
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for each new_node in EXPAND(node, problem) do

if NOT(MEMBER?(new_node, frontier)) and
NOT(MEMBER?(new_node, explored_set))
then frontier ← INSERT(new_node, frontier)
```

```
How to modify this algorithm to become the following algorithms? (important!)

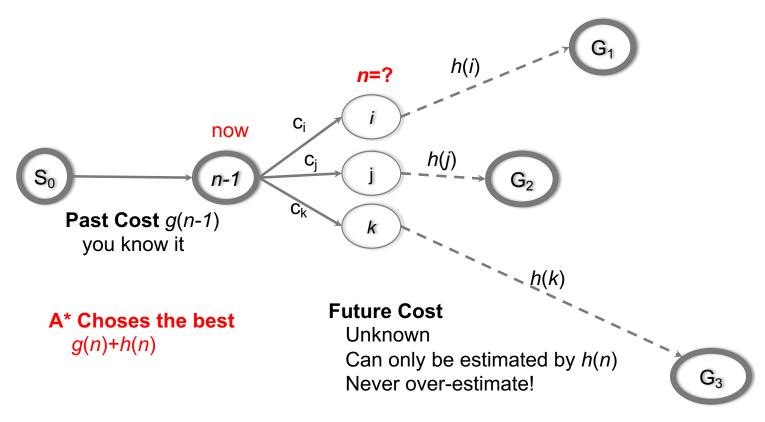
BFS

DFS

UCS

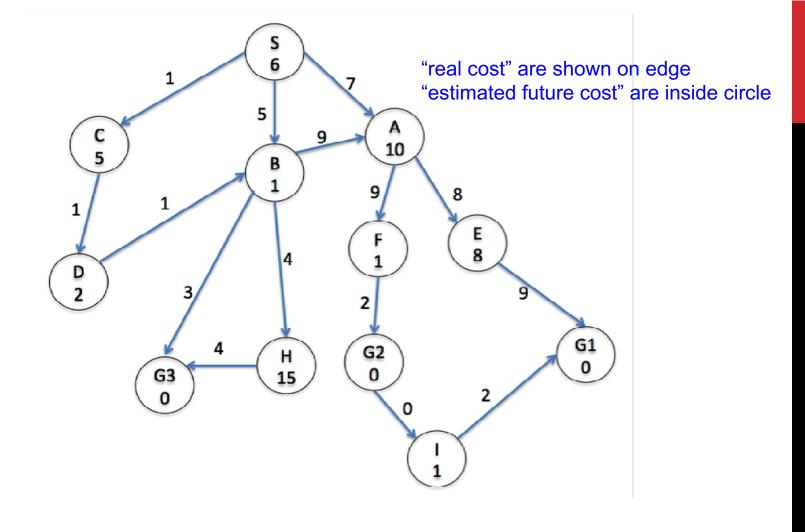
A*
```

A* = BEST-FIRST (PAST + ESTIMATED FUTURE)



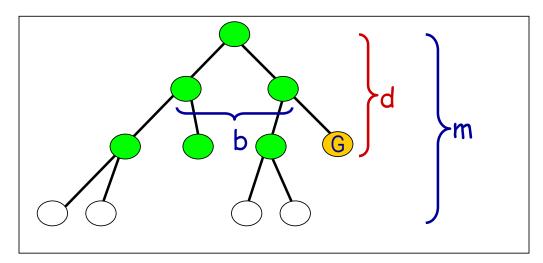
Note: Uniform-cost search uses only g(n), no h(n).

Is it good to have h(x)=0 for all x?



TIME COMPLEXITY OF BREADTH-FIRST SEARCH

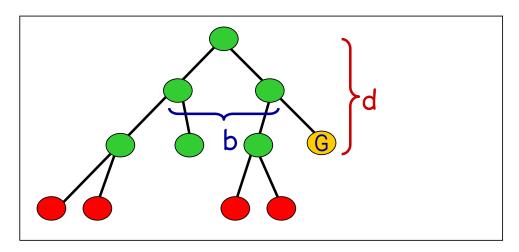
- Illustrates when goal check is done when node is selected for expansion
- If a goal node is found on depth **d** of the tree, all nodes up till that depth are created and examined (note: and the children of nodes at depth d are created and queued, but not yet examined).



Thus: O(bd+1)

SPACE COMPLEXITY OF BREADTH-FIRST SEARCH

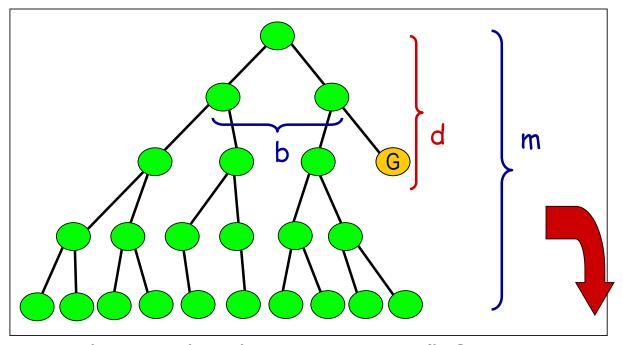
- Illustrates when goal check is done when node is selected for expansion
- Largest number of nodes in FRONTIER is reached on the level d+1 just beyond the goal node.



- QUEUE contains all nodes. (Thus: 4).
- In General: b^{d+1} − b ~ b^{d+1}

TIME COMPLEXITY OF DEPTH-FIRST SEARCH

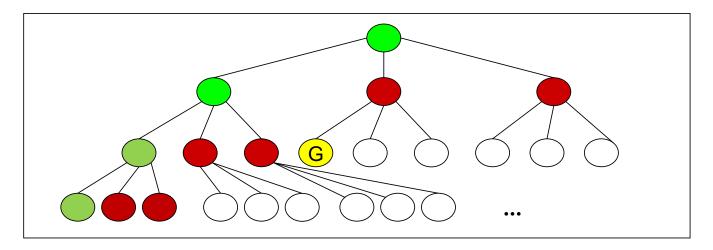
- In the worst case:
 - the (only) goal node may be on the right-most branch,



• Time complexity = $b^m + b^{m-1} + ... + 1 = O(b^m)$

SPACE COMPLEXITY OF DEPTH-FIRST

- Largest number of nodes in FRONTIER is reached in bottom leftmost node.
- Example: m = 3, b = 3:



- FRONTIER contains all nodes. Thus: 6.
- In General FRONTIER contains ((b-1) * m)
- Order: O(m*b)

SEARCH IN AI APPLICATIONS

Is search involved in these Al applications? If so, in what part or parts of the application?

 Building a driverless car that will drive down a roadway. (Leave aside the search involved in route planning.)

- Building a system like Siri.
- Text-to-speech synthesis

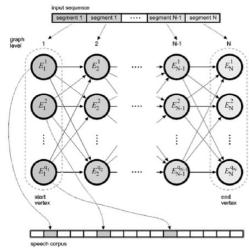
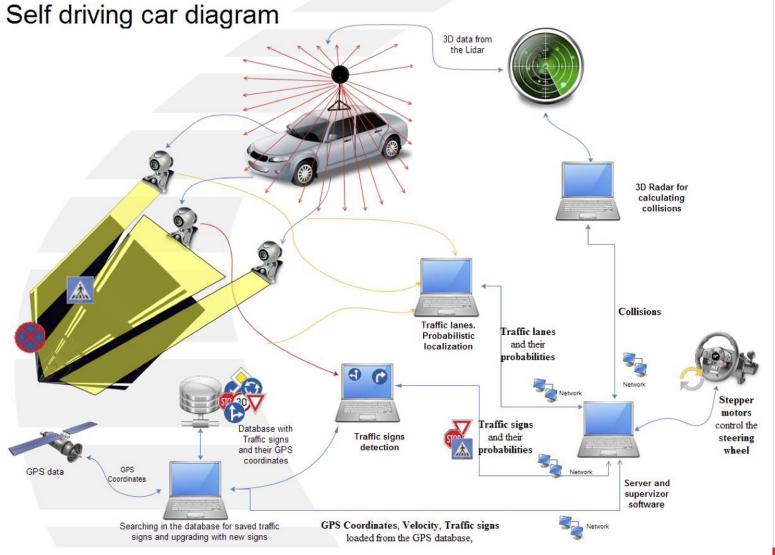


Figure 1. Structure of the graph for finding the optimal speech unit sequence; E_1^i are the graph initial-level vertices, E_N^i are the graph final-level vertices.



WHAT YOU SHOULD KNOW

- What is the difference between uninformed and informed search? Which ones are optimal?
- What are the advantages and disadvantages of depth-first search?
- Be familiar with the differences between search strategies shown in Figure 3.21

Criterion	Breadth- First	Uniform- Cost	Depth- First	Depth- Limited	Iterative Deepening	Bidirectional (if applicable)
Complete?	Yes ^a	$\operatorname*{Yes}^{a,b}O(b^{1+\lfloor C^*/\epsilon floor})$	No (1m)	No	Yes ^a	$\operatorname{Yes}^{a,d}$
Time			$O(b^m)$	$O(b^\ell)$	$O(b^d)$	$O(b^{d/2})$
Space	$O(b^d)$	$O(b^{1+\lfloor C^*/\epsilon floor})$	O(bm)	$O(b\ell)$	O(bd)	$O(b^{d/2})$
Optimal?	Yes^c	Yes	No	No	Yes ^c	$\mathrm{Yes}^{c,d}$

WANT MORE?

BigO and complexity:

https://apelbaum.wordpress.com/2011/05/05/big-o/