## Data Structures and Algorithms – MidTerm Exam

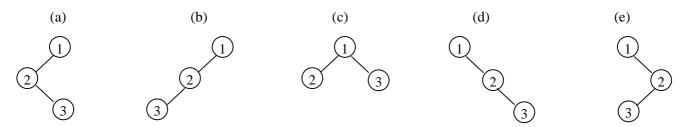
- 1. Let 'I' be "push", 'O' be "pop" and 'P' be "print" directly. To convert the infix expression a+(b\*c-d)/e to its postfix through a stack, the sequence of operations are \_\_\_\_\_PIIPIPOOIPOO \_\_\_\_\_ (For example: (a+b) is converted to ab+ by IPIPOO.) (5 points)
- 2. In a binary tree of N nodes, there are N+1 NULL pointers representing children. (2 points)
- 3. A sorting algorithm is *stable* if elements with equal keys are left in the same order as they occur in the input. Which of the following algorithms is/are stable? Answer: \_\_\_\_\_\_\_\_(a) (c) \_\_\_\_\_\_\_ (8 points) (a) insertion sort; (b) quick sort; (c) merge sort; (d) heap sort
- 4. The following routine removes duplicates from an array-based list A[0] ... A[N-1]. LastPosition is initially N-1.

```
for ( i = 0; i < LastPosition; i ++ ) {
    j = i + 1;
    while ( j < LastPosition )
        if ( A[i] == A[j] )        Delete(j);
        else    j ++;
}</pre>
```

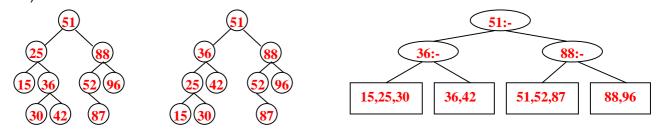
(a) What is the function of **Delete**? (3 points)

Delete A[ j ] by shifting A[ j+1 ] ... A[ LastPosition-1 ] to the left. LastPosition -- .

- (b)  $T_{worst}(N) = O(N^2)$ . (2 points)
- (c) Using linked list implementation,  $T_{worst}(N) = O(N^2)$ . (2 points)
- 5. Among the given trees, <u>d</u> has the same inorder and preorder traversal results, and <u>b</u> has the same postorder and inorder traversal results. (4 points)



6. Show the result of inserting { 51, 25, 36, 88, 42, 52, 15, 96, 87, 30 } into (a) an initially empty binary search tree; (b) an initially empty AVL tree; (c) an initially empty 2-3 tree. (30 points)



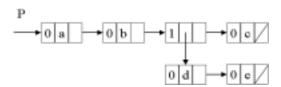
7. Please fill in the blanks in the programs. (12 points)

```
(a) Insertion for separate chaining hash table:
                                                          (b) Percolate down a max heap
void Insert( ElementType Key, HashTable H )
                                                          void PercolateDown(int p, PriorityQueue H)
{ Position Pos, NewCell;
                                                          { int child;
                                                            ElementType Tmp = H->Elements[ p ];
  List L;
  Pos = Find( Key, H );
                                                            for (; p * 2 <= H->Size; p = child) {
  if ( Pos == NULL ) {
                                                              child = p * 2;
    NewCell = malloc( sizeof( struct ListNode ) );
                                                              if ( child!=N-1&&H->Elements[child+1]>H->Elements [child] )
    L = H->TheLists[ Hash( Key, H->TableSize ) ];
                                                                child++;
    NewCell->Element = Key;
                                                              if ( H->Elements[ child ] > Tmp )
    NewCell->Next = L->Next
                                                                 H->Elements[p] = H->Elements[child];
     -->Next = NewCell
                                                               else
                                                                      break;
  }
}
                                                            H->Elements[p] = Tmp;
                                                         }
```

8. Assume that we represent trees using the list representation and that we define the node structure as:

TAG DATA LINK

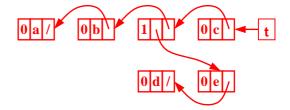
where LINK is a pointer pointing to the next element in the list; TAG is a field that holds the value of TRUE if the node is a link node in which DATA is a pointer pointing to the sublist, and a value of FALSE if the node is an atom node in which DATA is the data field. A sample tree is shown by the figure:



Please describe the function of the following program (7 points) and draw the resulting tree for the above example (5 points).

```
void r(GLNode_ptr p, GLNode_ptr *t )
{ GLNode_ptr temp1, temp2;
    if (! p) { *t=NULL; return; }
    if (! p->tag)
    { temp1 = p; p = p->link; temp1->link = NULL; }
    else
    { r ( p->data.sublist, t );
        temp1 = p; p = p->link; temp1->link = NULL; temp1->data.sublist = *t; }
    if ( p )
    { r ( p, t );
        temp2 = *t;
        while( temp2->link )        temp2 = temp2->link;
        temp2->link = temp1;
    }
    else    *t = temp1;
}
```

Reverse p, and t is the new head pointer.



9. Please write a C program to obtain the kth largest integer without destroying the original integer list. Your algorithm must have an average run time no worse than O(N log N). (20 points) int Find\_kth ( int A[ ], int N, int k ) /\* A[] stores the integer list; N is the size of the list; /\* and you are supposed to return the kth largest integer. \*/ Algorithm: Define a table[] and make table sort; /\* (+ quicksort or mergesort or heapsort) \*/ /\* Note: the list must be sorted in decreasing order \*/ Return A[ table[ k-1 ] ]. /\* Or if the list is sorted in increasing order \*/ Return A[ table[ N-k ] ]. A sample program – quicksort + table: Assume that Swap, Cutoff, and MAX\_SIZE are pre-defined. int Median3\_with\_table(int A[], int table[], int Left, int Right) { int Center = (Left + Right) / 2; if ( A[ table[ Left ] ] > A[ table[ Center ] ] ) Swap( &table[ Left ], & table[ Center ] ); if ( A[ table[ Left ] ] > A[ table[ Right ] ] ) Swap( & table[ Left ], & table[ Right ] ); if (A[table[Center]] > A[table[Right]]) Swap( & table[ Center ], & table[ Right ] ); /\* Invariant: A[ table[ Left ] ] <= A[ table[ Center ] ] <= A[ table[ Right ] ] \*/ Swap( & table[ Center ], & table[ Right - 1 ] ); /\* Hide pivot \*/ return A[ table[ Right - 1 ] ]; /\* Return pivot \*/ } void Qsort\_with\_table( int list[ ], int table[ ], int Left, int Right) { **int** i, j; int Pivot; if ( Left + Cutoff <= Right ) { Pivot = Median3\_with\_table( A, table, Left, Right ); i = Left; j = Right - 1;for(;;){ while( A[ table[ ++i ] ] < Pivot ) { } while( A[ table[ - -j ] ] > Pivot ) { } if(i < j)Swap( &table[ i ], &table[ j ] ); else break: Swap( &table[ i ], &table[ Right - 1 ] ); /\* Restore pivot \*/ Qsort( A, table, Left, i - 1 ); Qsort( A, table, i + 1, Right ); else /\* Do an insertion sort on the subarray \*/ InsertionSort( A + Left, table + Left, Right - Left + 1 );

}

```
void InsertionSort( int A[ ], int table[ ], int N )
{
    int j, P;
int Tmp;
    for (P = 1; P < N; P++) {
         Tmp = A[ table[ P ] ];
         for (j = P; j > 0 && A[table[j-1]] > Tmp; j--)
              A[ table[ j ] ] = A[ table[ j - 1 ] ];
         A[ table[ j ] ] = Tmp;
    }
}
int Find_kth ( int A[ ], int N, int k )
{ int i, table[ MAX_SIZE ];
   for (i = 0; i < N; i++)
     table[ i ] = i; /* initialize table */
   Qsort_with_table( A, table, 0, N-1);
   return A[table[N-k]];
}
```