

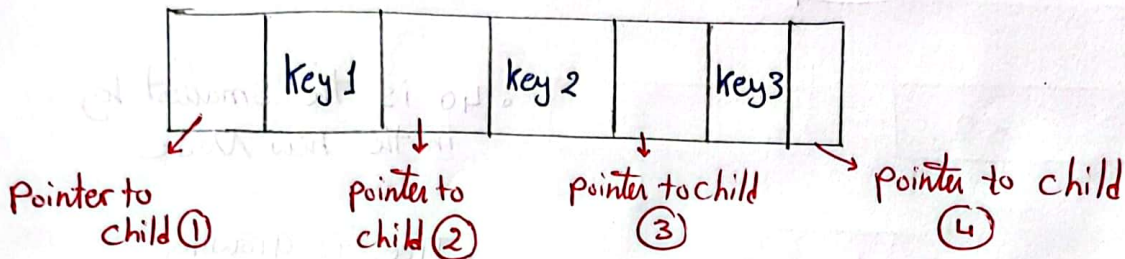
Insertion of B-tree Using time trade space offs

Structure of it

m-way \rightarrow m is the number of pointer to child

example 4-Way

Keys = $\boxed{m-1} \rightarrow 4-1 = \boxed{3} \rightarrow$ key in each node



Rules

- ① We have to at least fill the half in every node
ex: if degree is (m) we have to fill $\lceil \frac{m}{2} \rceil$
 - ② leaf nodes must be at same level
 - ③ Root can have minimum ② children
 - ④ We fill in Bottom up way
when we split we take the smallest key in the new node and then it becomes the parent
- It is similar to the search except it puts the new keys and splits into two nodes if needed

Insertion

Insert the following keys : 10, 20, 40, 50, 60, 70, 80, 30, 35,

child pointers = $\boxed{4} = \boxed{m}$ 5, 15

keys in node = $4-1 = \boxed{3}$

next page \rightarrow

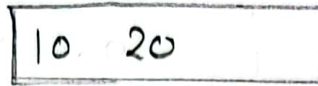
Nowran

①



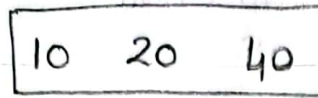
// ADD 10

②



// ADD 20

③



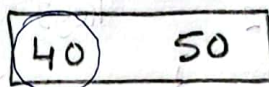
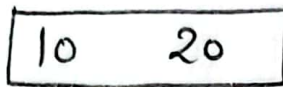
50

no space for 50 so we will split this node into ② → Create new one

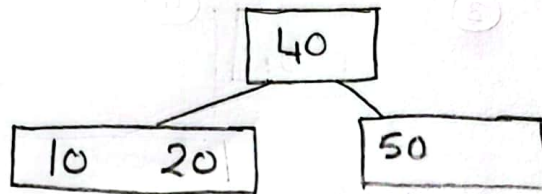
④



• we will now Balance



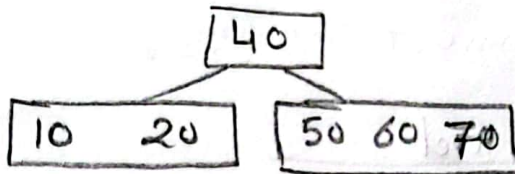
• 40 is the smallest key in the new Node



• Tree is growing upwards

• splitting while creating a B-tree

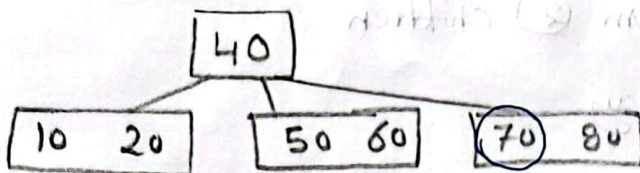
⑤



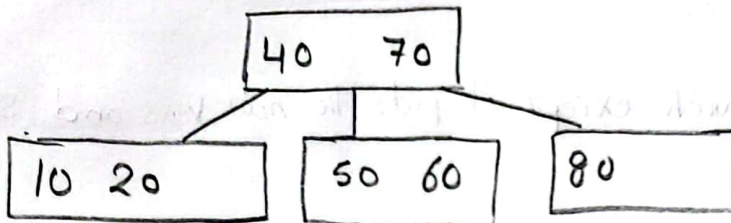
// ADD 60, 70

80

no space

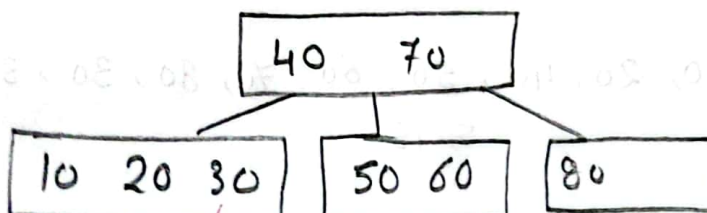


• 70 is the smallest key in the New Node



• Balancing

⑥



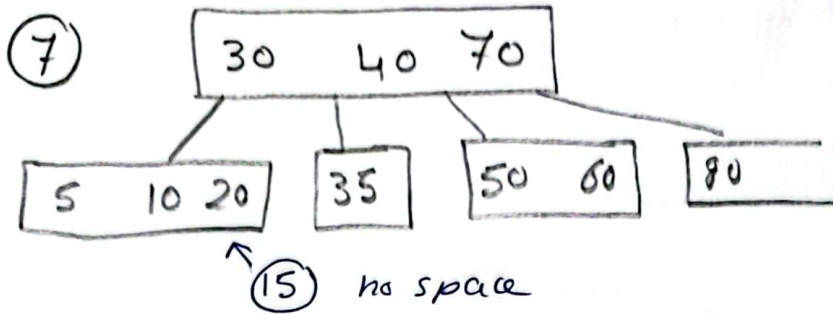
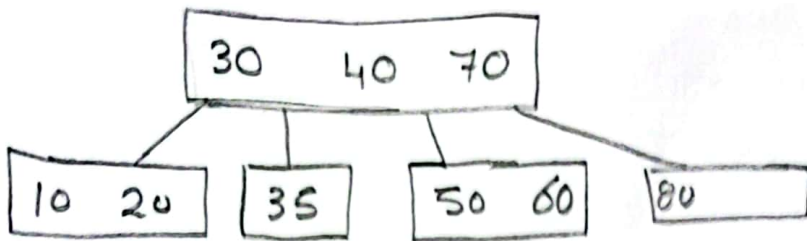
// ADD 30

will be smallest key in new node

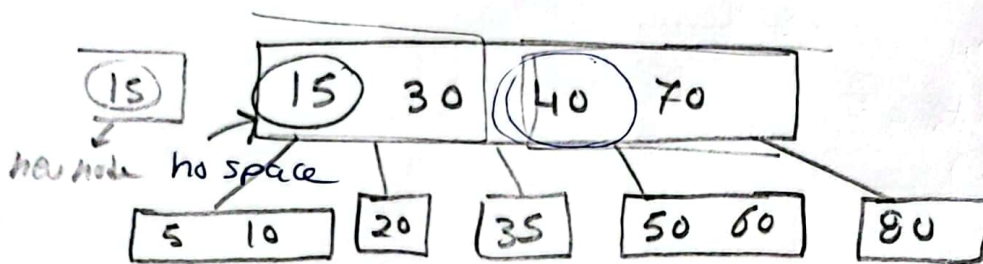
35

no space

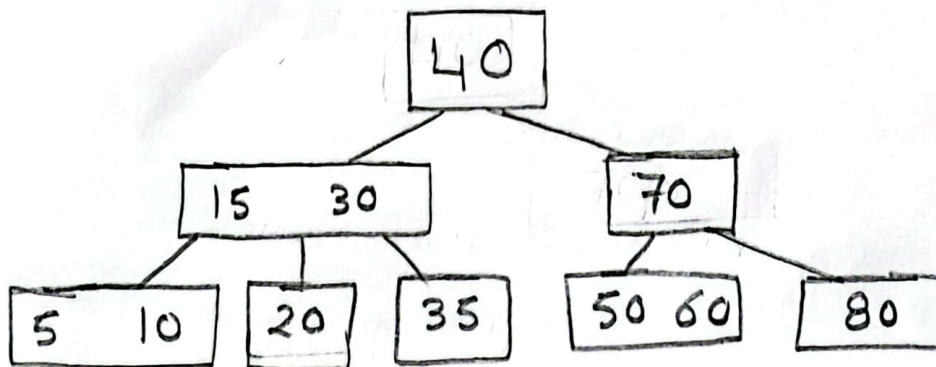
Nouran Hassan



• ADD (5)



Final Result :



• All leafs are on the same level