CKY Parser

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Parsing Algorithms

Top-down vs. bottom-up:

- Top-down: (goal-driven): from the start symbol down.
- Bottom-up: (data-driven): from the symbols up.

Naive vs. dynamic programming:

- Naive: enumerate everything.
- Backtracking: try something, discard partial solutions.
- Dynamic programming: save partial solutions in a table.

Examples:

- •CKY: bottom-up dynamic programming.
- Early parsing: top-down dynamic programming.

CKY (Cocke-Kasami-Younger)

- One of the earliest recognition and parsing algorithms
- ■The standard version of CKY can only recognize languages defined by context-free grammars in Chomsky Normal Form (CNF).
- It is also possible to extend the CKY algorithm to handle some grammars which are not in CNF
 - Harder to understand
- Based on a "dynamic programming" approach:
 - Build solutions compositionally from sub-solutions
- Uses the grammar directly.

Chomsky Normal Form (CNF)

A context-free grammar where the right side of each production rule is restricted to be either two non terminals or one terminal. Production can be one of following formats:

$$\circ A \rightarrow \alpha$$

$$\circ$$
 A \rightarrow BC

Any CFG can be converted to a weakly equivalent grammar in CNF

CNF Conversion

Three main conditions:

1) Hybrid rules:

2) Unit productions:

$$A \rightarrow B$$

3) Long productions:

$$A \rightarrow BCD$$

CNF Conversion

1) Hybrid rule conversion:

- Replace all terminals with dummy non-terminals
- E.g. INF-VP \rightarrow to VP •To \rightarrow to , INF-VP \rightarrow To VP

2) Unit productions:

- Rewrite RHS with RHS of all derivable non-unit productions
- If $A \Rightarrow B$ and $B \rightarrow w$, then add $A \rightarrow w$

3) Long productions:

- Introduce new non-terminals and spread over rules
- A → B C D
 - $A \rightarrow ED, E \rightarrow BC$

Consider the grammar G given by:

$$S \rightarrow AB$$

$$S \rightarrow XB$$

$$T \rightarrow AB$$

$$T \rightarrow XB$$

$$X \rightarrow AT$$

$$A \rightarrow a$$

$$B \rightarrow b$$

Is w = aaabbb in L(G)?

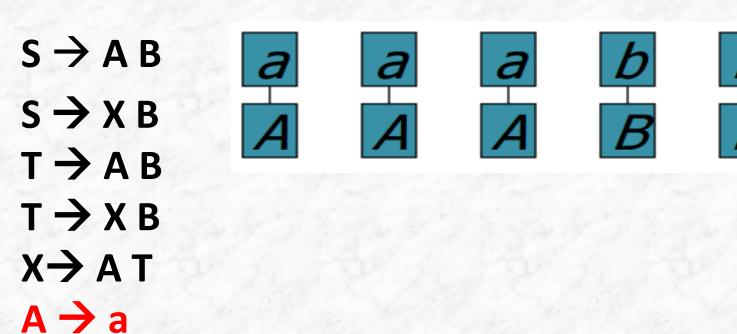
Now look at aaabbb:

$$S \rightarrow AB$$
 $S \rightarrow XB$
 $T \rightarrow AB$
 $T \rightarrow XB$
 $X \rightarrow AT$
 $A \rightarrow a$

 $B \rightarrow b$

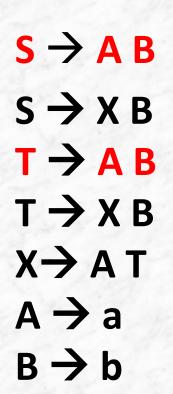


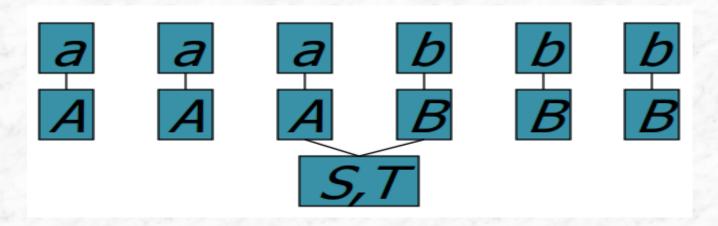
1) Write variables for all length 1 substrings.



 $B \rightarrow b$

2) Write variables for all length 2 substrings.

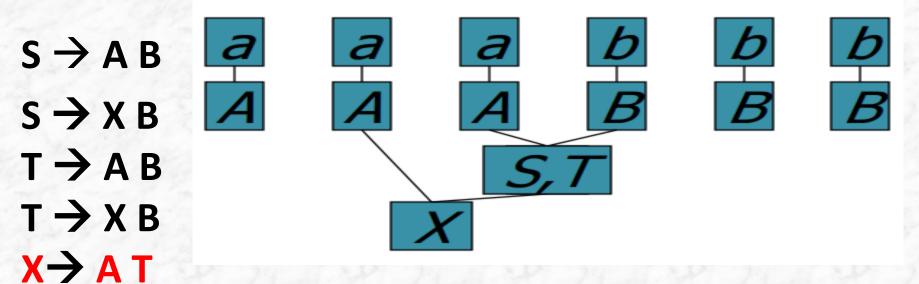




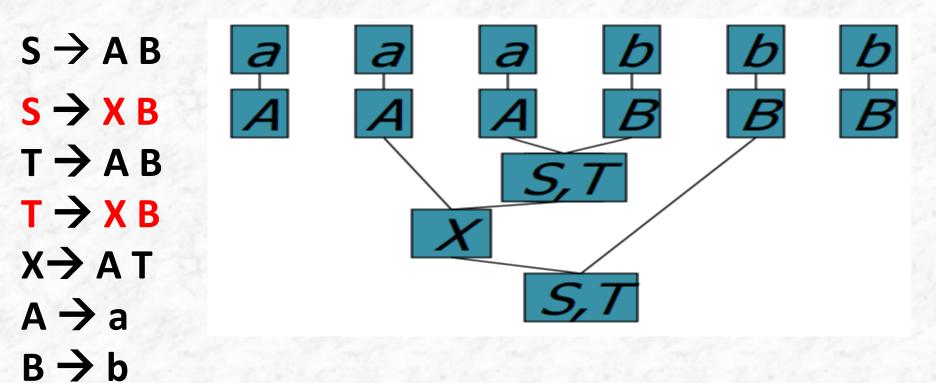
3) Write variables for all length 3 substrings.

 $A \rightarrow a$

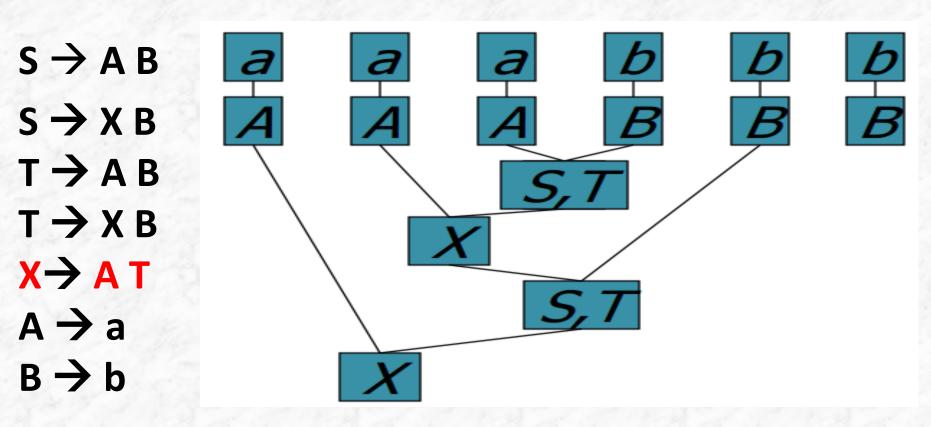
 $B \rightarrow b$



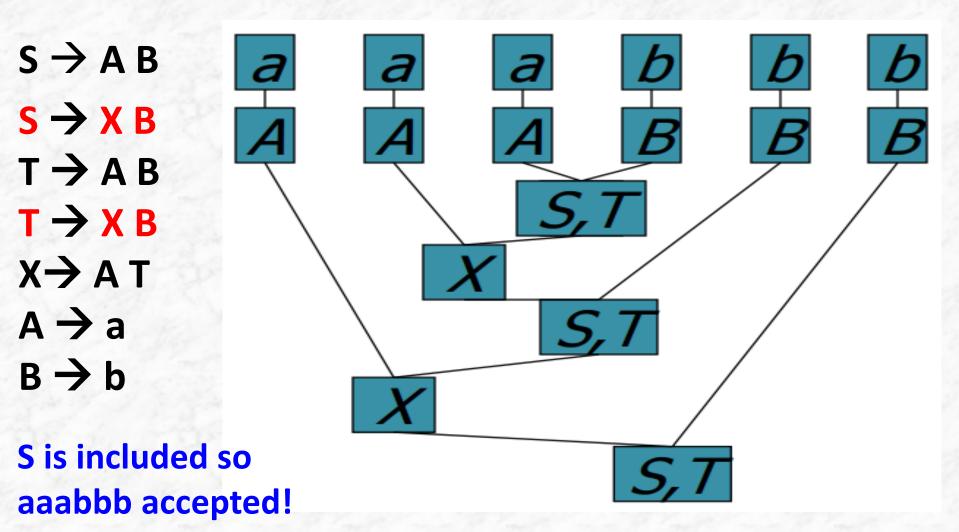
3) Write variables for all length 4 substrings.



3) Write variables for all length 5 substrings.



3) Write variables for all length 6 substrings.



The CKY Algorithm

function CKY (word w, grammar P) returns table

```
for i \leftarrow from 1 to LENGTH(w) do
  table[i-1, i] \leftarrow {A | A \rightarrow w<sub>i</sub> \in P }
for j \leftarrow from 2 to LENGTH(w) do
  for i \leftarrow from j-2 down to 0 do
        for k \leftarrow i + 1 to j - 1 do
           table[i,j] \leftarrow table[i,j] \cup {A | A \rightarrow BC \in P,
                   B \in table[i,k], C \in table[k,j]
```

If the start symbol $S \in \text{table}[0,n]$ then $w \in L(G)$

The table chart used by the algorithm:

/	1	2	3	4 b	5 b	6 6
0			1	24	624	36 24
1						
2						
3						
4						
5				240		34 24

Build an n+1 x n+1 matrix, where n = number of words in input

The table chart used by the algorithm:

	1	2	3	4	5	6
	a	a	a	b	b	b
0	[0,1]		4	24		98 AL
1	The same of the sa	[1,2]				
2			[2,3]			
3	The San			[3,4]	Verille.	
4				300	[4,5]	
5				240	24	[5,6]

Illustrate the numbering of cells: [i,j]'s represent spans

The table chart used by the algorithm:

-	1	2	3	4	5	6
	a	a	a	b	b	b
0	[0,1] A		150	24		26
1		[1,2] A				
2	$S \rightarrow AB$		[2,3]			
3	$S \rightarrow XB$ $T \rightarrow AB$			[3,4]		
4	$T \rightarrow X B$ $X \rightarrow A T$ $A \rightarrow a$				[4,5]	
5	$B \rightarrow b$					[5,6]

AA get nothing

The table chart used by the algorithm:

	1	2	3	4	5	6
	a	a	a	<u></u>	b	b
0	[0,1] A			24		
1		[1,2] A				
2	$S \rightarrow AB$		[2,3] A			
3	$S \rightarrow XB$ $T \rightarrow AB$ $T \rightarrow XB$			[3,4]	Series Contractions of the Contraction of the Contr	
4	$X \rightarrow A T$ $A \rightarrow a$				[4,5]	
5	$B \rightarrow b$					[5,6]

'a' is labeled A

The table chart used by the algorithm:

<u> </u>	1	2	3	4	5	6
	a	<u>a</u>	a	b	b	b
0	[0,1] A			24	6719	
1		[1,2] A				
2	$S \rightarrow AB$		[2,3] A			
3	$S \rightarrow X B$ $T \rightarrow A B$ $T \rightarrow X B$			[3,4] B		
4	Т → Х Б Х → А Т А → а				[4,5]	
5	$B \rightarrow b$					[5,6]

'b' is labeled B

The table chart used by the algorithm:

j	1	2	3	4	5	6
'	a	a	a	b	b	b
0	[0,1] A		4	24		
1		[1,2] A				
2	$S \rightarrow AB$		[2,3] A			
3	$S \rightarrow XB$ $T \rightarrow AB$			[3,4] B	No May	
4	T → X B X→ A T A → a				[4,5]	
5	$B \rightarrow b$					[5,6]

Found **S** and **T** [2,3],[3,4]

S->AB T->AB

The table chart used by the algorithm:

/	1	2	3	4 6	5 6	6
0	[0,1] A				8 / Ling	\$6 //L
1		[1,2] A				
2	$S \rightarrow AB$		[2,3] A	[2,4] S,T		
3	$S \rightarrow XB$ $T \rightarrow AB$			[3,4] B		
4	T → X B X→ A T A → a				[4,5]	
5	$B \rightarrow b$					[5,6]

Found **S** and **T** [2,3],[3,4]

S->AB T->AB

The table chart used by the algorithm:

/	1	2	3 =	4 6	5 6	6 <u>b</u>		
0	[0,1] A					900		
1		[1,2] A		[1,4] X				
2	$S \rightarrow AB$		[2,3] A	[2,4] S,T				
3	$S \rightarrow XB$ $T \rightarrow AB$ $T \rightarrow XB$			[3,4] B				
4	Т → Х Б Х → А Т А → а				[4,5]			
5	$B \rightarrow b$					[5,6]		

Test all filled cell with column = the row of current cell Failed (AS)

The table chart used by the algorithm:

/	1	2	3	4	5 6	6 <u> </u>
0	[0,1] A		26-36	249	626	3674
1		[1,2] A		[1,4] X		
2			[2,3] A	[2,4] S,T		
3				[3,4] B		
4					[4,5]	
5						[5,6]

Test all filled cell with column = the row of current cell Failed (AX)

The table chart used by the algorithm:

	1	2	3	4	5	6
	a	a	a	b	b	b
0	[0,1] A			25	126	
1		[1,2] A		[1,4] X		
2	$S \rightarrow AB$		[2,3] A	[2,4] S,T		
3	$S \rightarrow XB$ $T \rightarrow AB$ $T \rightarrow XB$			[3,4] B		
4	Т → Д Б Х → Д Т Д → а				[4,5] B	
5	$B \rightarrow b$					[5,6]

'b' is labeled B

The table chart used by the algorithm:

-	1	2	3	4	5	6
	a	a	a	b	b	b
0	[0,1] A			25	24	1900 July
1		[1,2] A		[1,4] X		
2	$S \rightarrow AB$		[2,3] A	[2,4] S,T		
3	$S \rightarrow X B$ $T \rightarrow A B$ $T \rightarrow X B$			[3,4] B		
4	Т → Х Б Х → А Т А → а				[4,5] B	
5	$B \rightarrow b$					[5,6]

'BB not found

The table chart used by the algorithm:

/	1	2	3	4 b	2 0	6 <mark>6</mark>
0	[0,1] A		100	24	6214	26
1		[1,2] A		[1,4] X		
2	$S \rightarrow AB$		[2,3] A	[2,4] S,T		
3	$S \rightarrow XB$ $T \rightarrow AB$			[3,4] B		
4	T → X B X→ A T A → a				[4,5] B	
5	$B \rightarrow b$					[5,6]

SB or TB not found

The table chart used by the algorithm:

j	1	2	3	4	5	6
	a	a	a	b	b	b
0	[0,1] A		14 30	25		900
1		[1,2] A		[1,4] X	[1,5] T,S	
2	$S \rightarrow AB$		[2,3] A	[2,4] S,T		
3	$S \rightarrow X B$ $T \rightarrow A B$ $T \rightarrow X B$			[3,4] B		
4	X → A T A → a				[4,5] B	
5	$B \rightarrow b$					[5,6]

'XB' produce T from [1,4],[4,5]

The table chart used by the algorithm:

-	1	2	3	4	5	6
	a	a	a	b	b	b
0	[0,1] A			26	[0,5] X	967/4
1		[1,2] A		[1,4] X	[1,5] T,S	
2	$S \rightarrow AB$		[2,3] A	[2,4] S,T		
3	$S \rightarrow XB$ $T \rightarrow AB$ $T \rightarrow XB$			[3,4] B		
4	X → A T A → a				[4,5] B	
5	$B \rightarrow b$					[5,6]

'AT' produce X from [0,1], [1,5]

The table chart used by the algorithm:

	and the second	for the state of the state of		The second second		the state of the s
j	1	2	3	4	5	6
	a	a	a	b	b	b
0	[0,1] A			24	[0,5] X	76/24
1		[1,2] A		[1,4] X	[1,5] T,S	
2	$S \rightarrow A B$ $S \rightarrow X B$		[2,3] A	[2,4] S,T		
3	$T \rightarrow AB$ $T \rightarrow XB$			[3,4] B		
4	$X \rightarrow A T$ $A \rightarrow a$				[4,5] B	
5	$B \rightarrow b$					[5,6] B

'b' is labeled B

The table chart used by the algorithm:

j	1	2	3	4	5	6
	a	6	a	b	b	b
0	[0,1] A			240	[0,5] X	76/2/2
1		[1,2] A		[1,4] X	[1,5] T,S	
2	$S \rightarrow A B$ $S \rightarrow X B$		[2,3] A	[2,4] S,T		
3	$T \rightarrow AB$ $T \rightarrow XB$			[3,4] B		
4	$X \rightarrow A T$ $A \rightarrow a$				[4,5] B	
5	$B \rightarrow b$					[5,6] B

BB not found

The table chart used by the algorithm:

		Contract of the second				
	1	2	3	4	5	6
	<u>a</u>	a	a	b	b	b
0	[0,1] A			24	[0,5] X	76/2/4
1		[1,2] A		[1,4] X	[1,5] T,S	
2	$S \rightarrow A B$ $S \rightarrow X B$		[2,3] A	[2,4] S,T		
3	$T \rightarrow AB$ $T \rightarrow XB$			[3,4] B		
4	$X \rightarrow AT$ $A \rightarrow a$				[4,5] B	
5	$B \rightarrow b$					[5,6] B

'TB' not found

The table chart used by the algorithm:

/	1 00	2	3	4 6	5 b	6 6
0	[0,1] A				[0,5] X	[0,6] S,T
1		[1,2] A		[1,4] X	[1,5] T,S	
2	$S \rightarrow AB$		[2,3] A	[2,4] S,T		
3	$S \rightarrow X B$ $T \rightarrow A B$ $T \rightarrow X B$			[3,4] B		
4	$X \rightarrow A T$ $A \rightarrow a$				[4,5] B	
5	$B \rightarrow b$					[5,6] B

XB produce S,T from [0,5],[5,6]

The table chart used by the algorithm:

/ /	1 a	2	3	4 b	5 b	6 b
0	[0,1] A			240	[0,5] X	[0,6] S,T
1		[1,2] A		[1,4] X	[1,5] T,S	
2			[2,3] A	[2,4] S,T	68 3	
3				[3,4] B		
4					[4,5] B	
5			V-314			[5,6] B

Found S node: [0,5] [5,6] Recognition algorithm returns True when a root node is found in [0,n]

Parsing results

- We keep the results for every wij in a table.
- Note that we only need to fill in entries up to the diagonal.
- Every entry in the table T[i,j] can contains up to
 r=|N| symbols (the size of the non-terminal set).
- •We then want to find T[0,n,S] = true.