

# Mini-Project: Sorting Algorithms

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January 13, 2025

## 1 Introduction

Sorting is one of the most classically studied families of algorithms, because they are among the modules essential for the good running of more advanced algorithms. The general principle of a sorting algorithm is to order (in ascending order for example) the objects of a collection of data (values), according to a comparison criterion (key – for us, values and keys are here confused: these are the elements of an array of integers). We generally carry out sorting by using an "in-place" approach: the sorted values are stored in the same array as the initial values (which therefore becomes an input-output parameter).

## 2 Sorting Algorithms

### 2.1 Bubble Sort

#### 2.1.1 Algorithm

The **Bubble Sort** function in C sorts an array of integers in ascending order using the bubble sort algorithm. It iterates through the array multiple times, comparing and swapping adjacent elements if they are in the wrong order. The process is repeated until the array is sorted. An optimization is included with a **sorted flag** that allows the function to exit early if no swaps are made during a pass, indicating that the array is already sorted. This reduces unnecessary iterations and improves efficiency.

```

void bubble_sort(int arr[], int n) {
    for (int i = 0; i < n; i++) {
        int sorted = 1;
        for (int j = 0; j < n - i - 1; j++) {
            if (arr[j] > arr[j + 1]) {
                swap(&arr[j], &arr[j + 1]);
                sorted = 0;
            }
        }
        if (sorted) {
            return;
        }
    }
}

```

Listing 1: Bubble Sort implementation

### 2.1.2 Complexity

Bubble Sort has a worst-case and average-case time complexity of  $O(n^2)$ , where  $n$  is the number of elements to be sorted. The best-case time complexity is  $O(n)$  when the array is already sorted. The space complexity is  $O(1)$  as it is an in-place sorting algorithm.

### 2.1.3 Experimental Study

To study the real cost of the algorithms, we tested them on arrays of integers of increasing size  $n$  filled randomly. Time measurement was done several times (5 times) for a given array size. The values of  $n$  chosen were 100, 1000, 10000, and 100000.

### 2.1.4 Results

## 2.2 Gnome Sort

### 2.2.1 Algorithm

The **Gnome Sort** function in C sorts an array of integers in ascending order by iteratively comparing each element with the next one and swapping them if they are out of order, moving the index backward after a swap to recheck previous elements, and moving it forward otherwise, continuing this process until the entire array is sorted, ensuring that each element is in its correct position by the end of the algorithm.

```

void gnome_sort(int arr[], int n) {
    int i = 0;

    while (i < n - 1) {
        if (arr[i] <= arr[i + 1]) {
            i++;
        } else {
            swap(&arr[i], &arr[i + 1]);
            if (i > 0) {
                i--;
            } else {
                i++;
            }
        }
    }
}

```

Listing 2: Gnome Sort implementation

### 2.2.2 Complexity

Gnome Sort has a worst-case and average-case time complexity of  $O(n^2)$ , where  $n$  is the number of elements to be sorted. The best-case time complexity is  $O(n)$  when the array is already sorted. The space complexity is  $O(1)$  as it is an in-place sorting algorithm.

### 2.2.3 Experimental Study

To study the real cost of the algorithms, we tested them on arrays of integers of increasing size  $n$  filled randomly. Time measurement was done several times (5 times) for a given array size. The values of  $n$  chosen were 100, 1000, 10000, and 100000.

### 2.2.4 Results

## 2.3 Radix Sort

### 2.3.1 Algorithm

The **Sort Aux** function in C sorts an array of integers based on a specific digit (determined by the parameter ‘i’) using a counting sort approach, where it first counts the occurrences of each digit, then calculates the cumulative count, places the elements into an output array in sorted order according to the current digit, and finally copies the sorted elements back into the original array, ensuring that the array is partially sorted by the specified digit.

```

void sort_aux(int arr[], int n, int i) {
    int *output = (int *)malloc(n * sizeof(int));
    int count[10] = {0};

    for (int j = 0; j < n; j++) {
        count[key(arr[j], i)]++;
    }
    for (int j = 1; j < 10; j++) {
        count[j] += count[j - 1];
    }
    for (int j = n - 1; j >= 0; j--) {
        output[count[key(arr[j], i)] - 1] = arr[j];
        count[key(arr[j], i)]--;
    }
    for (int j = 0; j < n; j++) {
        arr[j] = output[j];
    }

    free(output);
}

```

Listing 3: Sort Aux implementation

The **Radix Sort Helper** function in C sorts an array of integers by iterating through each digit from the least significant to the most significant, calling the **Sort Aux** function to sort the array based on the current digit, while the **Radix Sort** function determines the maximum number of digits in the largest number of the array and then calls **Radix Sort Helper** to perform the sorting process, ensuring the entire array is sorted by all digits.

```

void radix_sort_helper(int arr[], int n, int k) {
    // k represents the maximum number of digits
    for (int i = 0; i < k; i++) {
        sort_aux(arr, n, i);
    }
}

void radix_sort(int arr[], int n) {
    // maximum number of digits is the number of digits of the
    // largest number
    int imax = 0;
    for (int i = 1; i < n; i++) {
        if (arr[i] > arr[imax]) {
            imax = i;
        }
    }

    int k = (int)log10(arr[imax]) + 1;
    radix_sort_helper(arr, n, k);
}

```

Listing 4: Radix Sort implementation

### 2.3.2 Complexity

Radix Sort has a time complexity of  $O(d \cdot (n + k))$ , where  $d$  is the number of digits in the largest number,  $n$  is the number of elements, and  $k$  is the range of the digit values. The space complexity is  $O(n + k)$ .

### 2.3.3 Experimental Study

To study the real cost of the algorithms, we tested them on arrays of integers of increasing size  $n$  filled randomly. Time measurement was done several times (5 times) for a given array size. The values of  $n$  chosen were 100, 1000, 10000, and 100000.

### 2.3.4 Results

## 2.4 Quick Sort

### 2.4.1 Algorithm

The **Partition** function in C rearranges the elements of an array around a pivot element (chosen as the first element of the array segment) such that all elements less than or equal to the pivot are on its left and all elements greater than the pivot are on its right, by iteratively moving indices from both ends towards the center, swapping elements as needed, and finally placing the pivot in its correct position, returning the index of the pivot.

```
int partition(int arr[], int low, int high) {
    int p = arr[low];
    int i = low;
    int j = high;

    while (i < j) {
        while (arr[i] <= p && i <= high - 1) {
            i++;
        }
        while (arr[j] > p && j >= low + 1) {
            j--;
        }
        if (i < j) {
            swap(&arr[i], &arr[j]);
        }
    }
    swap(&arr[low], &arr[j]);
    return j;
}
```

Listing 5: Partition implementation

The **Quick Sort Helper** function in C recursively sorts an array of integers by partitioning the array around a pivot element and then sorting the subarrays on either side of the pivot, while the **Quick Sort** function serves as a wrapper that initializes the recursive process by calling **Quick Sort Helper** with the full range of the array.

```
void quick_sort_helper(int arr[], int low, int high) {
    if (low < high) {
        int pi = partition(arr, low, high);

        quick_sort_helper(arr, low, pi - 1);
        quick_sort_helper(arr, pi + 1, high);
    }
}

void quick_sort(int arr[], int n) { quick_sort_helper(arr, 0, n - 1); }
```

Listing 6: Quick Sort implementation

### 2.4.2 Complexity

Quick Sort has a worst-case time complexity of  $O(n^2)$ , which occurs when the pivot selection is poor. However, the average-case and best-case time complexity is  $O(n \log n)$ . The space complexity is  $O(\log n)$  due to the recursive stack space.

### 2.4.3 Experimental Study

To study the real cost of the algorithms, we tested them on arrays of integers of increasing size  $n$  filled randomly. Time measurement was done several times (5 times) for a given array size. The values of  $n$  chosen were 100, 1000, 10000, and 100000.

### 2.4.4 Results

## 2.5 Heap Sort

### 2.5.1 Algorithm

The **Heapify** function in C maintains the heap property of a binary heap by comparing a node with its children and swapping it with the largest child if necessary, then recursively applying the same process to the affected subtree, ensuring that the subtree rooted at the given node becomes a valid max-heap.

```

void heapify(int arr[], int n, int i) {
    int largest = i;
    int left = 2 * i + 1; // Left child
    int right = 2 * i + 2; // Right child

    if (left < n && arr[left] > arr[largest])
        largest = left;

    if (right < n && arr[right] > arr[largest])
        largest = right;

    if (largest != i) {
        swap(&arr[i], &arr[largest]);
        heapify(arr, n, largest);
    }
}

```

Listing 7: Heapify implementation

The **Heap Sort** function in C sorts an array of integers in ascending order by first building a max heap from the array, then repeatedly extracting the maximum element from the heap and placing it at the end of the array, and finally re-heapifying the remaining elements until the entire array is sorted.

```

void heap_sort(int arr[], int n) {
    // Build max heap
    for (int i = n / 2 - 1; i >= 0; i--)
        heapify(arr, n, i);

    // Extract elements from heap one by one
    for (int i = n - 1; i > 0; i--) {
        swap(&arr[0], &arr[i]);
        heapify(arr, i, 0);
    }
}

```

Listing 8: Heap Sort implementation

### 2.5.2 Complexity

Heap Sort has a time complexity of  $O(n \log n)$  for all cases (worst, average, and best). The space complexity is  $O(1)$  as it is an in-place sorting algorithm.

### 2.5.3 Experimental Study

To study the real cost of the algorithms, we tested them on arrays of integers of increasing size  $n$  filled randomly. Time measurement was done several times (5 times) for a given array size. The values of  $n$  chosen were 100, 1000, 10000, and 100000.

#### **2.5.4 Results**

### **3 Conclusion**

In this project, we put into practice and tested some sorting algorithms, studied their complexity, and compared theoretical complexity with the evaluation of running cost. The experimental results matched the theoretical expectations, confirming the efficiency of the studied sorting algorithms.