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Tutorial 2: Introduction to Finite Field F_2

0. Let us take a number $x = 100 + (\text{ID} \bmod 100)$ and $y = 200 + (\text{MyKAD} \bmod 100)$

An addition is an exclusive-or operation

 $x = 187$ and $y = 239$.

$$x = 187_{10} = \textcolor{blue}{1011} \ 1011_2 = x^7 + x^5 + x^4 + x^3 + x + 1 \quad \text{and}$$

$$y = 239_{10} = \textcolor{blue}{1110} \, 1111_2 = x^7 + x^6 + x^5 + x^3 + x^2 + x + 1,$$

$$\begin{array}{r} x + y = \textcolor{blue}{1011} \ 1011_2 + \textcolor{blue}{1110} \ 1111_2 = \textcolor{blue}{1110} \ 1111 \\ \textcolor{blue}{1011} \ 1011 \\ \hline \textcolor{blue}{0101} \ 0100 \end{array}$$

$$\begin{aligned} \text{In mathematical notation } x + y &= (x^7 + x^5 + x^4 + x^3 + x + 1) + (x^7 + x^6 + x^5 + x^3 + x^2 + x + 1) \\ &= 2x^7 + x^6 + 2x^5 + x^4 + 2x^3 + x^2 + x + 2 \\ &= 0x^7 + x^6 + 0x^5 + x^4 + 0x^3 + x^2 + x + 0 \pmod{2} \\ &= x^6 + x^4 + x^2 + x \end{aligned}$$

1. Multiplication is a convolution modulo 2.

$$\begin{aligned} x \cdot y &= 11101111_2 \cdot 10111011_2 \\ &= 11101111_2 \\ &\quad 11101111_2 \\ &\quad\quad 11101111_2 \\ &\quad\quad\quad 11101111_2 \\ &\quad\quad\quad\quad 11101111_2 \\ &\quad\quad\quad\quad\quad 11101111_2 \\ &= 110001011011001_2 \end{aligned}$$

2. Division gives quotient and remainder.

Let us divide by $283_{10} = 100011011_2 = x^8 + x^4 + x^3 + x + 1$

Quotient	$110001011011001 = 25,305_{10}$
1	$\begin{array}{r} 100011011 \\ \hline 10010000011001 \end{array}$
11	$\begin{array}{r} 100011011 \\ \hline 0011101111001 \end{array}$
11001	$\begin{array}{r} 100011011 \\ \hline 1100010101 \end{array}$
110011	$\begin{array}{r} 100011011 \\ \hline 100100011 \end{array}$
1100111	$\begin{array}{r} 100011011 \\ \hline 00111000 = 56_{10} \text{ Reminder} \end{array}$

3. An inverse is a good starting point to be a cryptographer.

Lets compute an inverse $x^5+x^4+x^3$ from previous number modulo $x^8+x^4+x^3+x+1$

Let us invoke an Extended Euclidean Algorithm

Extended Euclidean Algorithm								
i	$b =$	$a *$	q	+	r	u	v	$w=u-vq$
0	100011011	111000	1101		11	0	1	1101
1	111000	11	10111		1	1	1101	11110011

Line 0: Quotient 100011011 $w=u-vq=0-1\cdot1101=1101$
 1000 111000
 11011011
 1100 111000
 0111011
 1101 111000
 000011 Remainder

Line 1: Quotient 111000 $w=u-vq=1-1101\cdot10111$
 10000 11 =10111
 01000 10111
 10100 11 10111
 100 11110011
 10110 11
 10
 10111 11
 1 Remainder

Line 0: $v \cdot q = 1 \cdot 1101 = 1 \cdot (x^3 + x^2 + 1) = 1101$

- Always remember to check your answer whether $a \cdot a^{-1} \equiv 1 \pmod{p}$
4. Check your division: Another say another A.

Let us check that

$a^{-1} = 11110011$ is indeed an inverse of $a = 111000$ modulo 100011011 .

$$a \cdot a^{-1}$$

$$\begin{array}{r}
 = 111000 * 11110011 \\
 \begin{array}{r}
 11110011 \\
 11110011 \\
 11110011 \\
 11110011 \\
 \hline
 1011011001 \\
 100011011 \\
 \hline
 011101111
 \end{array}
 \end{array}$$