Chapter 3: Processes

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- Process Concept
- Process Scheduling
- Operations on Processes
- Interprocess Communication
- Examples of IPC Systems
- Communication in Client-Server Systems

Objectives

- To introduce the notion of a process -- a program in execution, which forms the basis of all computation

- To describe the various features of processes, including scheduling, creation and termination, and communication

- To explore interprocess communication using shared memory and message passing

- To describe communication in client-server systems

Process Concept

- An operating system executes a variety of programs:
    - Batch system – **jobs**
    - Time-shared systems – **user programs** or **tasks**
- Textbook uses the terms *job* and *process* almost interchangeably
- **Process** – a program in execution; process execution must progress in sequential fashion
- Multiple parts
    - The program code, also called **text section**
    - Current activity including **program counter**, processor registers
    - **Stack** containing temporary data
        4 Function parameters, return addresses, local variables
    - **Data section** containing global variables
   - **Heap** containing memory dynamically allocated during run time

Process Concept (Cont.)

- Program is *passive* entity stored on disk (**executable file**), process is *active*

    - Program becomes process when executable file loaded into memory

- Execution of program started via GUI mouse clicks, command line entry of its name, etc

- One program can be several processes

    - Consider multiple users executing the same program
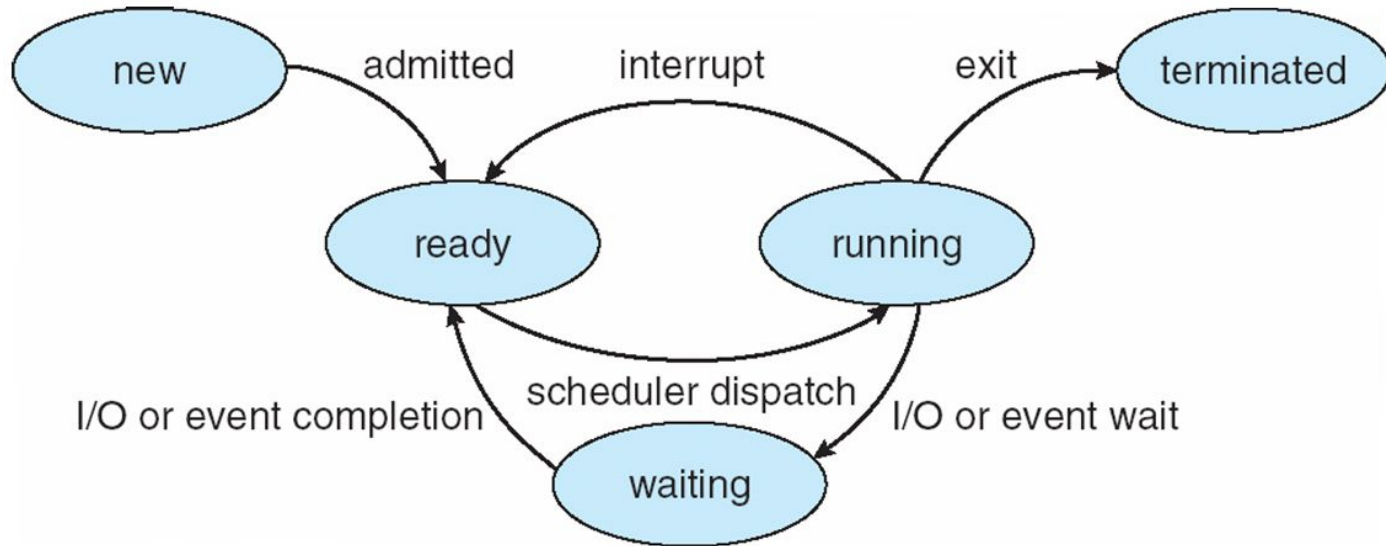
Process in Memory
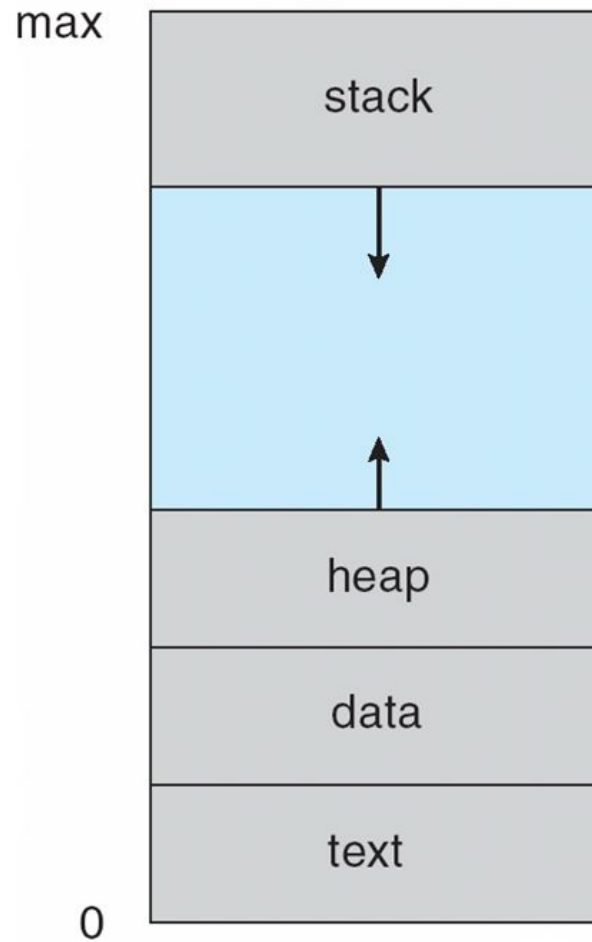
Process State

- As a process executes, it changes **state**
  - **new**: The process is being created
  - **running**: Instructions are being executed
  - **waiting**: The process is waiting for some event to occur
  - **ready**: The process is waiting to be assigned to a processor
 - **terminated**: The process has finished execution

Process Control Block (PCB)

Information associated with each process

(also called **task control block**)

- Process state – running, waiting, etc

- Program counter – location of instruction to next execute

- CPU registers – contents of all process-centric registers

- CPU scheduling information- priorities, scheduling queue pointers

- Memory-management information – memory allocated to the process

- Accounting information – CPU used, clock time elapsed since start, time limits

- I/O status information – I/O devices allocated to process, list of open files

| process state |
| process number |
| program counter |
| registers |
| memory limits |
| list of open files |
| • • • |
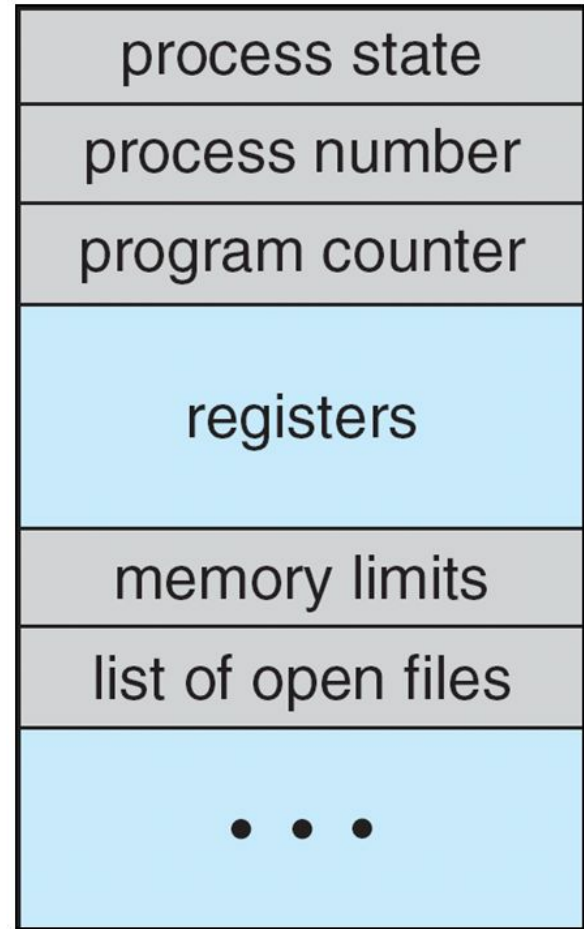
CPU Switch From Process to Process
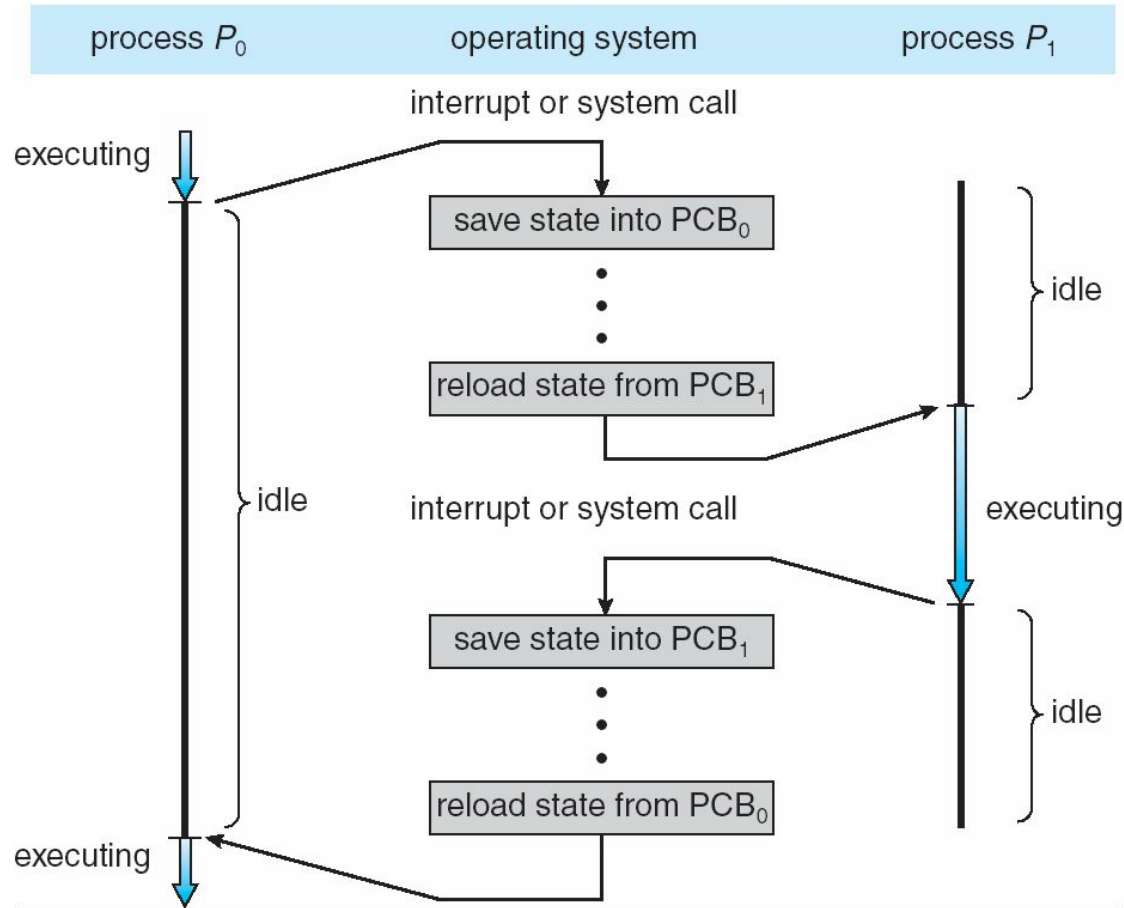
Threads

- So far, process has a single thread of execution

- Consider having multiple program counters per process

  - Multiple locations can execute at once

    4 Multiple threads of control -> **threads**

- Must then have storage for thread details, multiple program counters in PCB

- See next chapter

Process Representation in Linux

Represented by the C structure `task_struct`

```
pid t_pid; /* process identifier */
long state; /* state of the process */
unsigned int time_slice /* scheduling information */
struct task_struct *parent; /* this process's parent */
struct list_head children; /* this process's children */
struct files_struct *files; /* list of open files */
struct mm_struct *mm; /* address space of this process */
```
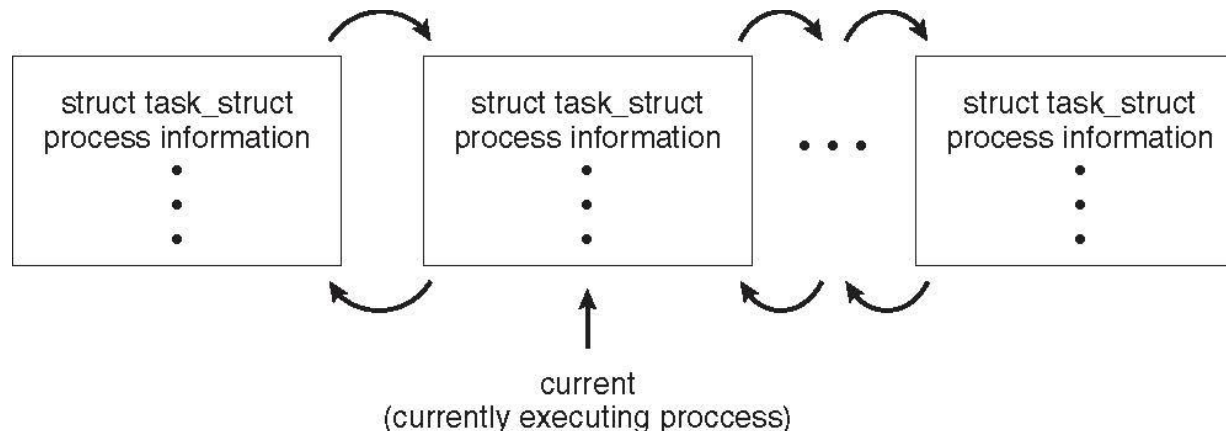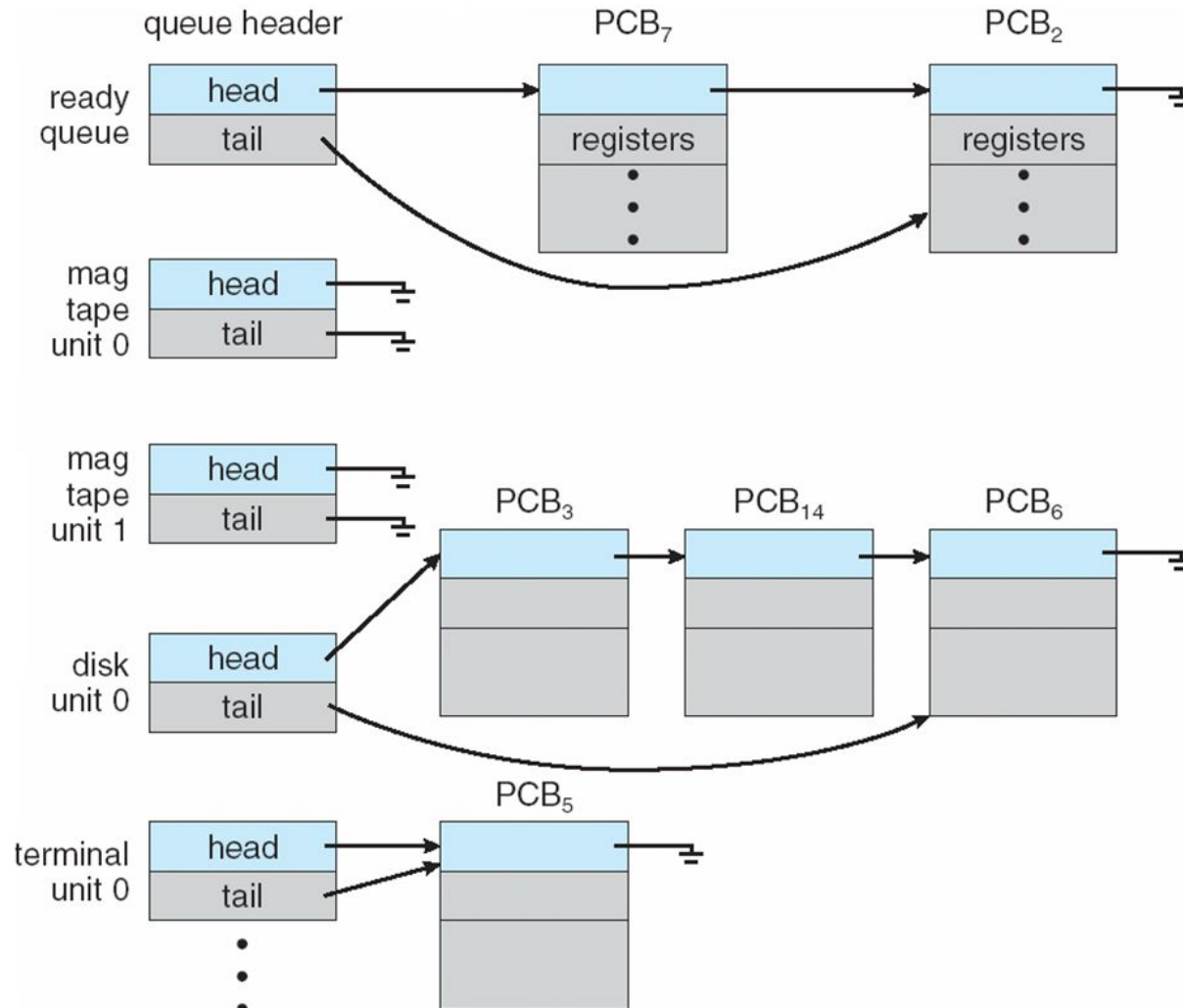
Process Scheduling

- Maximize CPU use, quickly switch processes onto CPU for time sharing

- **Process scheduler** selects among available processes for next execution on CPU

- Maintains **scheduling queues** of processes

  - **Job queue** – set of all processes in the system

  - **Ready queue** – set of all processes residing in main memory, ready and waiting to execute

  - **Device queues** – set of processes waiting for an I/O device

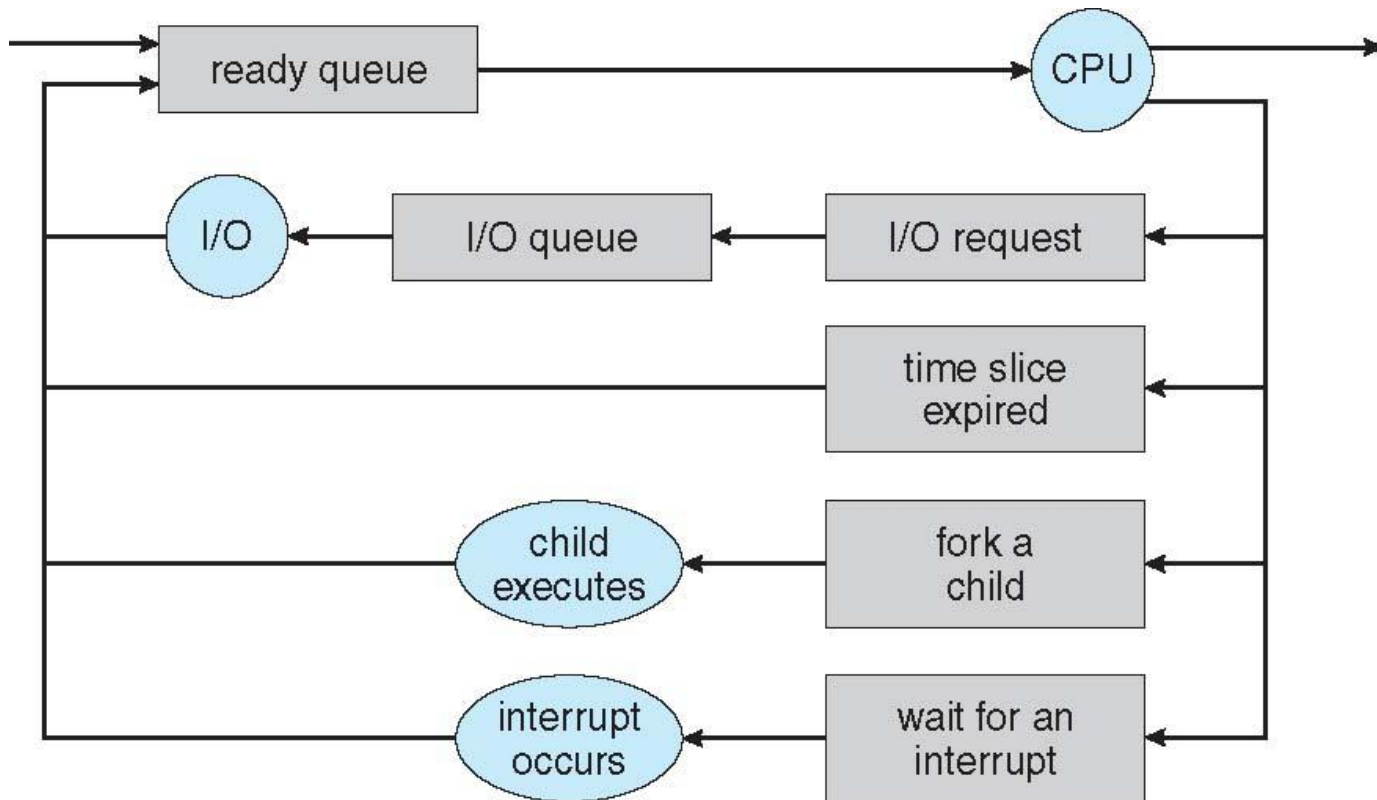 - Processes migrate among the various queues

Ready Queue And Various I/O Device Queues

Representation of Process Scheduling

- **Queueing diagram** represents queues, resources, flows

Schedulers

- **Short-term scheduler** (or **CPU scheduler**) – selects which process should be executed next and allocates CPU
  - Sometimes the only scheduler in a system
  - Short-term scheduler is invoked frequently (milliseconds) $\Rightarrow$ (must be fast)
- **Long-term scheduler** (or **job scheduler**) – selects which processes should be brought into the ready queue
  - Long-term scheduler is invoked infrequently (seconds, minutes) $\Rightarrow$ (may be slow)
  - The long-term scheduler controls the **degree of multiprogramming**
- Processes can be described as either:
  - **I/O-bound process** – spends more time doing I/O than computations, many short CPU bursts
  - **CPU-bound process** – spends more time doing computations; few very long CPU bursts
- Long-term scheduler strives for good *process mix*

Addition of Medium Term Scheduling

- **Medium-term scheduler** can be added if degree of multiple programming needs to decrease
  - Remove process from memory, store on disk, bring back in from disk to continue execution: **swapping**
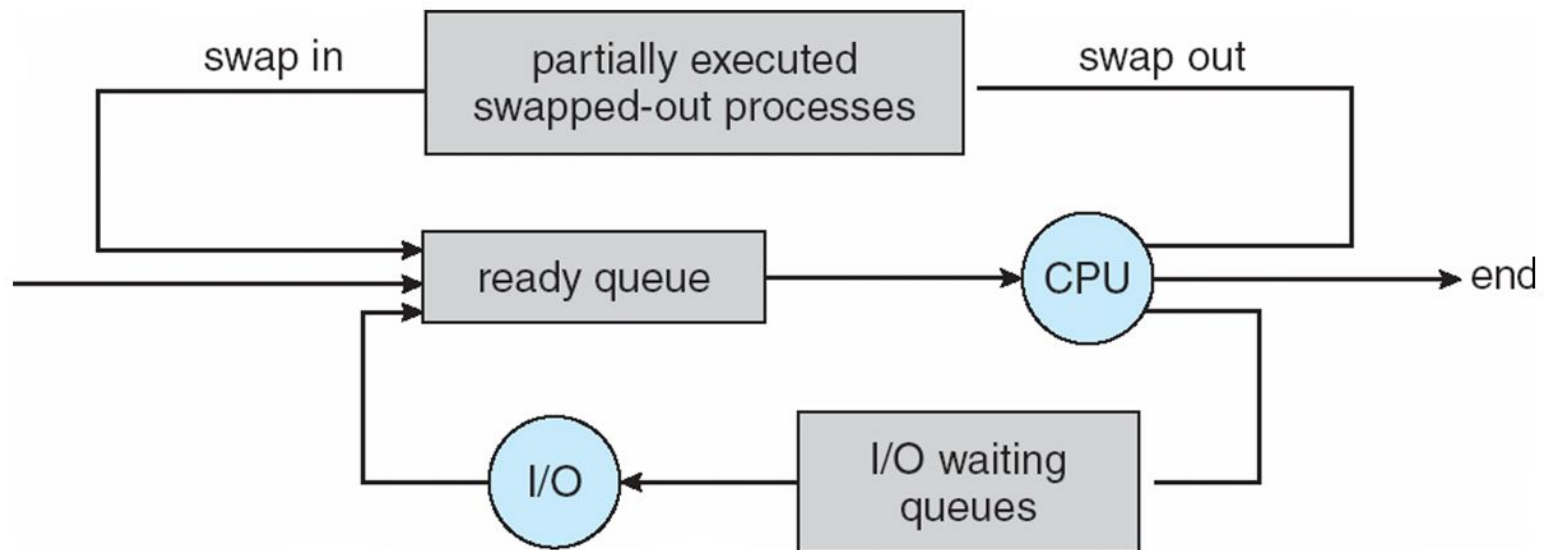
Multitasking in Mobile Systems

- Some mobile systems (e.g., early version of iOS) allow only one process to run, others suspended

- Due to screen real estate, user interface limits iOS provides for a

    - Single **foreground** process- controlled via user interface

    - Multiple **background** processes– in memory, running, but not on the display, and with limits

    - Limits include single, short task, receiving notification of events, specific long-running tasks like audio playback

- Android runs foreground and background, with fewer limits

    - Background process uses a **service** to perform tasks

    - Service can keep running even if background process is suspended

    - Service has no user interface, small memory use

Context Switch

- When CPU switches to another process, the system must **save the state** of the old process and load the **saved state** for the new process via a **context switch**

- **Context** of a process represented in the PCB

- Context-switch time is overhead; the system does no useful work while switching

  - The more complex the OS and the PCB ☐ the longer the context switch

- Time dependent on hardware support

  - Some hardware provides multiple sets of registers per CPU ☐ multiple contexts loaded at once

Operations on Processes

- System must provide mechanisms for:
    - process creation,
    - process termination,
   - and so on as detailed next

Process Creation

- **Parent** process create **children** processes, which, in turn create other processes, forming a **tree** of processes

- Generally, process identified and managed via a **process identifier** (**pid**)

- Resource sharing options
  - Parent and children share all resources
  - Children share subset of parent's resources
  - Parent and child share no resources

- Execution options
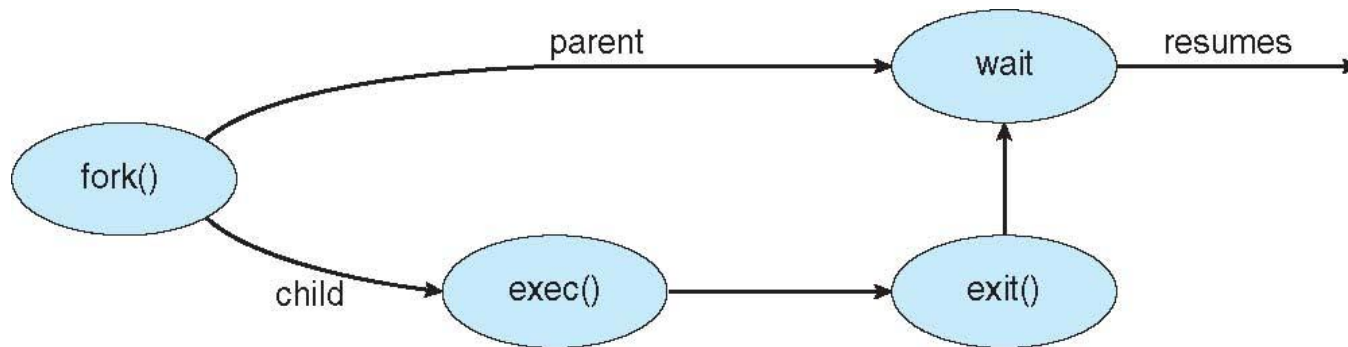 - Parent and children execute concurrently
  - Parent waits until children terminate
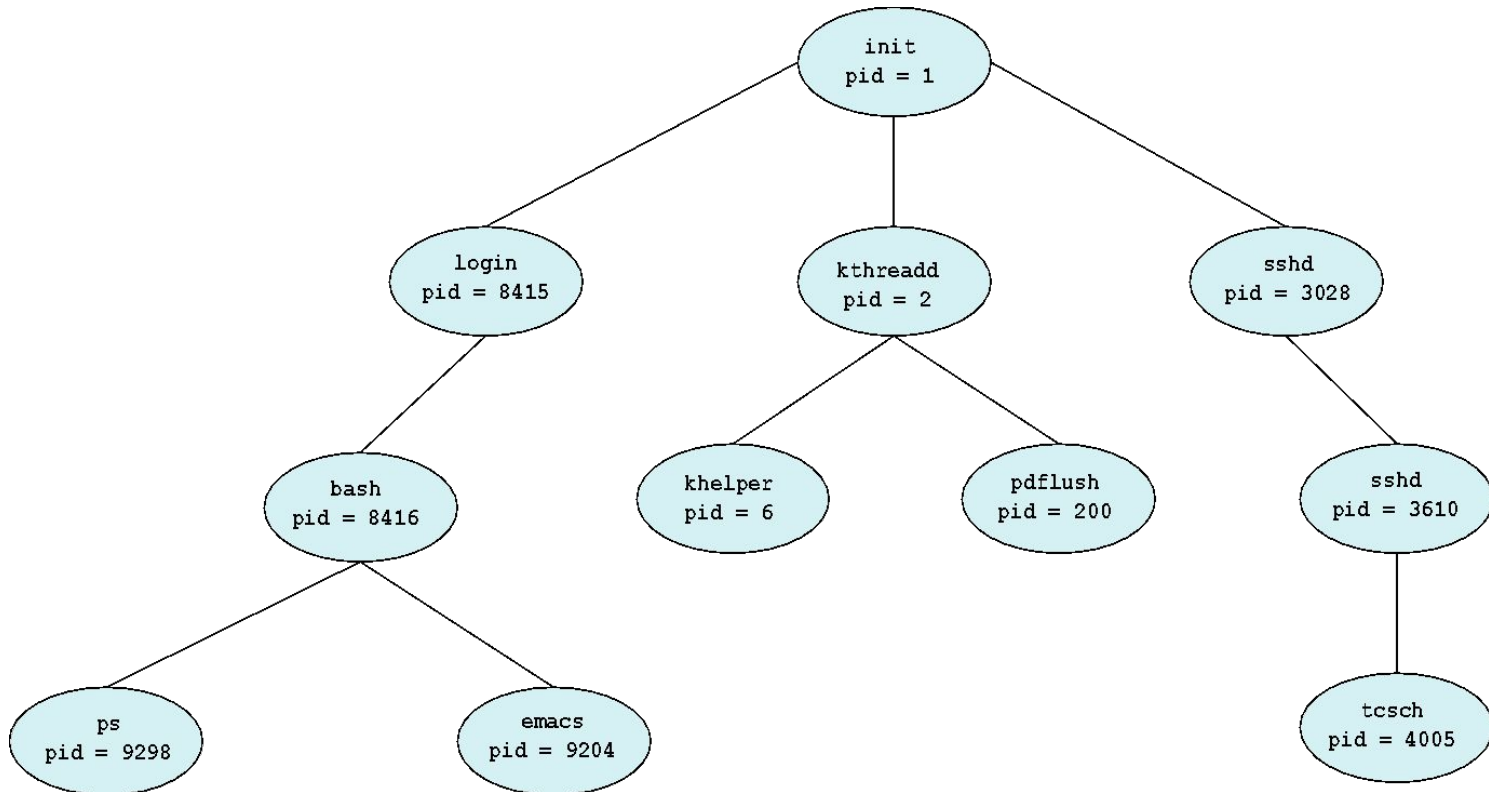
Process Creation (Cont.)

- Address space
  - Child duplicate of parent
  - Child has a program loaded into it
- UNIX examples
  - **fork()** system call creates new process
  - **exec()** system call used after a **fork()** to replace the process' memory space with a new program

```c
#include <sys/types.h>
#include <stdio.h>
#include <unistd.h>

int main()
{
pid_t pid;

    /* fork a child process */
    pid = fork();

    if (pid < 0) { /* error occurred */
        fprintf(stderr, "Fork Failed");
        return 1;
    }
    else if (pid == 0) { /* child process */
        execlp("/bin/ls","ls",NULL);
    }
    else { /* parent process */
        /* parent will wait for the child to complete */
        wait(NULL);
        printf("Child Complete");
    }

    return 0;
}
```

```c
#include <stdio.h>
#include <windows.h>

int main(VOID)
{
STARTUPINFO si;
PROCESS_INFORMATION pi;

    /* allocate memory */
    ZeroMemory(&si, sizeof(si));
    si.cb = sizeof(si);
    ZeroMemory(&pi, sizeof(pi));

    /* create child process */
    if (!CreateProcess(NULL, /* use command line */
      "C:\\WINDOWS\\system32\\mspaint.exe", /* command */
      NULL, /* don't inherit process handle */
      NULL, /* don't inherit thread handle */
      FALSE, /* disable handle inheritance */
      0, /* no creation flags */
      NULL, /* use parent's environment block */
      NULL, /* use parent's existing directory */
      &si,
      &pi))
    {
      fprintf(stderr, "Create Process Failed");
      return -1;
    }
    /* parent will wait for the child to complete */
    WaitForSingleObject(pi.hProcess, INFINITE);
    printf("Child Complete");

    /* close handles */
    CloseHandle(pi.hProcess);
    CloseHandle(pi.hThread);
}
```

Process Termination

- Process executes last statement and then asks the operating system to delete it using the **exit()** system call.
  - Returns status data from child to parent (via **wait()**)
  - Process' resources are deallocated by operating system
- Parent may terminate the execution of children processes using the **abort()** system call. Some reasons for doing so:
  - Child has exceeded allocated resources
  - Task assigned to child is no longer required
  - The parent is exiting and the operating systems does not allow a child to continue if its parent terminates

Process Termination

- Some operating systems do not allow child to exists if its parent has terminated. If a process terminates, then all its children must also be terminated.

    - **cascading termination.** All children, grandchildren, etc. are terminated.

    - The termination is initiated by the operating system.

- The parent process may wait for termination of a child process by using the **wait()** system call. The call returns status information and the pid of the terminated process

    ```
    pid = wait(&status);
    ```

- If no parent waiting (did not invoke **wait()**) process is a **zombie**

- If parent terminated without invoking **wait**, process is an **orphan**

Multiprocess Architecture – Chrome Browser

- Many web browsers ran as single process (some still do)
  - If one web site causes trouble, entire browser can hang or crash
- Google Chrome Browser is multiprocess with 3 different types of processes:
  - **Browser** process manages user interface, disk and network I/O
  - **Renderer** process renders web pages, deals with HTML, Javascript. A new renderer created for each website opened
    - 4 Runs in **sandbox** restricting disk and network I/O, minimizing effect of security exploits
  - **Plug-in** process for each type of plug-in

Each tab represents a separate process

Interprocess Communication

- Processes within a system may be *independent* or *cooperating*

- Cooperating process can affect or be affected by other processes, including sharing data

- Reasons for cooperating processes:
  - Information sharing
  - Computation speedup
  - Modularity
  - Convenience

- Cooperating processes need **interprocess communication** (**IPC**)

- Two models of IPC
  - **Shared memory**
  - **Message passing**
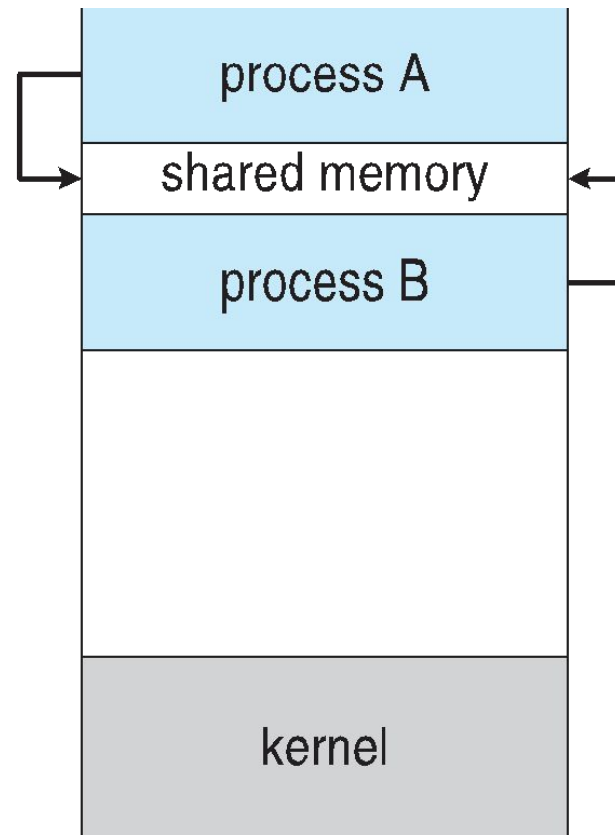
(a) Message passing.   (b) shared memory.
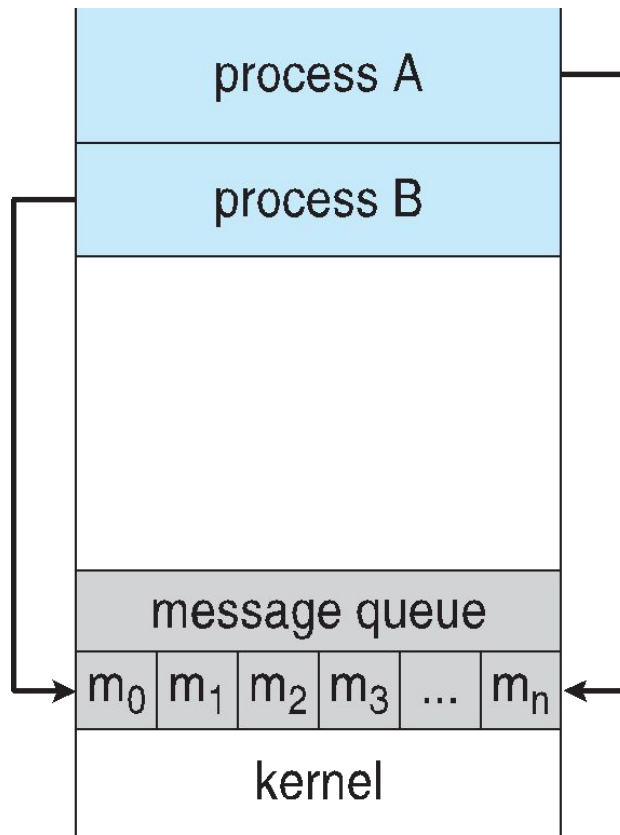
(a)

(b)

Cooperating Processes

- ***Independent*** process cannot affect or be affected by the execution of another process

- ***Cooperating*** process can affect or be affected by the execution of another process

- Advantages of process cooperation

  - Information sharing

  - Computation speed-up

  - Modularity

  - Convenience

Producer-Consumer Problem

- Paradigm for cooperating processes, *producer* process produces information that is consumed by a *consumer* process

    - **unbounded-buffer** places no practical limit on the size of the buffer

    - **bounded-buffer** assumes that there is a fixed buffer size

Bounded-Buffer – Shared-Memory Solution

- Shared data

```
#define BUFFER_SIZE 10
typedef struct {

. . .

} item;


item buffer[BUFFER_SIZE];
int in = 0;
int out = 0;
```

- Solution is correct, but can only use BUFFER_SIZE-1 elements

Bounded-Buffer – Producer

```
item next_produced;
while (true) {
        /* produce an item in next produced */
        while (((in + 1) % BUFFER_SIZE) == out)
                ; /* do nothing */
        buffer[in] = next_produced;
        in = (in + 1) % BUFFER_SIZE;
}
```

Bounded Buffer – Consumer

```
item next_consumed;
while (true) {
    while (in == out)
        ; /* do nothing */
    next_consumed = buffer[out];
    out = (out + 1) % BUFFER_SIZE;


    /* consume the item in next consumed */
}
```

Interprocess Communication – Shared Memory

- An area of memory shared among the processes that wish to communicate

- The communication is under the control of the users processes not the operating system.

- Major issues is to provide mechanism that will allow the user processes to synchronize their actions when they access shared memory.

- Synchronization is discussed in great details in Chapter 5.

Interprocess Communication – Message Passing

- Mechanism for processes to communicate and to synchronize their actions

- Message system – processes communicate with each other without resorting to shared variables

- IPC facility provides two operations:
  - **send**(*message*)
  - **receive**(*message*)

- The *message* size is either fixed or variable

Message Passing (Cont.)

- If processes *P* and *Q* wish to communicate, they need to:
  - Establish a ***communication link*** between them
  - Exchange messages via send/receive
- Implementation issues:
  - How are links established?
  - Can a link be associated with more than two processes?
  - How many links can there be between every pair of communicating processes?
  - What is the capacity of a link?
  - Is the size of a message that the link can accommodate fixed or variable?
  - Is a link unidirectional or bi-directional?

Message Passing (Cont.)

- Implementation of communication link
    - Physical:
        - Shared memory
        - Hardware bus
        - Network
    - Logical:
        - Direct or indirect
        - Synchronous or asynchronous
        - Automatic or explicit buffering

Direct Communication

- Processes must name each other explicitly:
    - **send** (*P, message*) – send a message to process P
    - **receive**(*Q, message*) – receive a message from process Q
- Properties of communication link
    - Links are established automatically
    - A link is associated with exactly one pair of communicating processes
    - Between each pair there exists exactly one link
    - The link may be unidirectional, but is usually bi-directional

Indirect Communication

- Messages are directed and received from mailboxes (also referred to as ports)
  - Each mailbox has a unique id
  - Processes can communicate only if they share a mailbox
- Properties of communication link
  - Link established only if processes share a common mailbox
  - A link may be associated with many processes
  - Each pair of processes may share several communication links
  - Link may be unidirectional or bi-directional

Indirect Communication

- Operations
    - create a new mailbox (port)
    - send and receive messages through mailbox
    - destroy a mailbox
- Primitives are defined as:

    send(*A, message*) – send a message to mailbox A

    receive(*A, message*) – receive a message from mailbox A

Indirect Communication

- Mailbox sharing
    - P_1, P_2, and P_3 share mailbox A
    - P_1, sends; P_2 and P_3 receive
    - Who gets the message?

- Solutions
    - Allow a link to be associated with at most two processes
    - Allow only one process at a time to execute a receive operation
    - Allow the system to select arbitrarily the receiver. Sender is notified who the receiver was.

Synchronization

- Message passing may be either blocking or non-blocking

- **Blocking** is considered **synchronous**

  - **Blocking send** -- the sender is blocked until the message is received

  - **Blocking receive** -- the receiver is blocked until a message is available

- **Non-blocking** is considered **asynchronous**

  - **Non-blocking send** -- the sender sends the message and continue

  - **Non-blocking receive** -- the receiver receives:

    - A valid message, or

    - Null message

- Different combinations possible

  - If both send and receive are blocking, we have a **rendezvous**

Synchronization (Cont.)

- Producer-consumer becomes trivial

```
message next_produced;
while (true) {
        /* produce an item in next produced */
send(next_produced);
}

message next_consumed;
while (true) {
   receive(next_consumed);

   /* consume the item in next consumed */
}
```

Buffering

- Queue of messages attached to the link.

- implemented in one of three ways

  1. Zero capacity – no messages are queued on a link.
      Sender must wait for receiver (rendezvous)

  2. Bounded capacity – finite length of n messages
      Sender must wait if link full

  3. Unbounded capacity – infinite length
      Sender never waits

Examples of IPC Systems - POSIX

- POSIX Shared Memory
    - Process first creates shared memory segment
      shm_fd = shm_open(name, O CREAT | O RDWR, 0666);
    - Also used to open an existing segment to share it
    - Set the size of the object

    ftruncate(shm fd, 4096);

    - Now the process could write to the shared memory

      sprintf(shared memory, "Writing to shared memory");

```c
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <fcntl.h>
#include <sys/shm.h>
#include <sys/stat.h>

int main()
{
/* the size (in bytes) of shared memory object */
const int SIZE = 4096;
/* name of the shared memory object */
const char *name = "OS";
/* strings written to shared memory */
const char *message_0 = "Hello";
const char *message_1 = "World!";

/* shared memory file descriptor */
int shm_fd;
/* pointer to shared memory obect */
void *ptr;

    /* create the shared memory object */
    shm_fd = shm_open(name, O_CREAT | O_RDWR, 0666);

    /* configure the size of the shared memory object */
    ftruncate(shm_fd, SIZE);

    /* memory map the shared memory object */
    ptr = mmap(0, SIZE, PROT_WRITE, MAP_SHARED, shm_fd, 0);

    /* write to the shared memory object */
    sprintf(ptr,"%s",message_0);
    ptr += strlen(message_0);
    sprintf(ptr,"%s",message_1);
    ptr += strlen(message_1);

    return 0;
}
```

```c
#include <stdio.h>
#include <stdlib.h>
#include <fcntl.h>
#include <sys/shm.h>
#include <sys/stat.h>

int main()
{
/* the size (in bytes) of shared memory object */
const int SIZE = 4096;
/* name of the shared memory object */
const char *name = "OS";
/* shared memory file descriptor */
int shm_fd;
/* pointer to shared memory obect */
void *ptr;

    /* open the shared memory object */
    shm_fd = shm_open(name, O_RDONLY, 0666);

    /* memory map the shared memory object */
    ptr = mmap(0, SIZE, PROT_READ, MAP_SHARED, shm_fd, 0);

    /* read from the shared memory object */
    printf("%s",(char *)ptr);

    /* remove the shared memory object */
    shm_unlink(name);

    return 0;
}
```

Examples of IPC Systems - Mach

- Mach communication is message based

    - Even system calls are messages

    - Each task gets two mailboxes at creation- Kernel and Notify

    - Only three system calls needed for message transfer

        msg_send(), msg_receive(), msg_rpc()

    - Mailboxes needed for commuication, created via

        port_allocate()

    - Send and receive are flexible, for example four options if mailbox full:

        4 Wait indefinitely

        4 Wait at most n milliseconds

   4 Return immediately

        4 Temporarily cache a message

Examples of IPC Systems – Windows

- Message-passing centric via **advanced local procedure call** (**LPC**) facility

  - Only works between processes on the same system

  - Uses ports (like mailboxes) to establish and maintain communication channels

  - Communication works as follows:

    4 The client opens a handle to the subsystem's **connection port** object.

    4 The client sends a connection request.

    4 The server creates two private **communication ports** and returns the handle to one of them to the client.

    4 The client and server use the corresponding port handle to send messages or callbacks and to listen for replies.
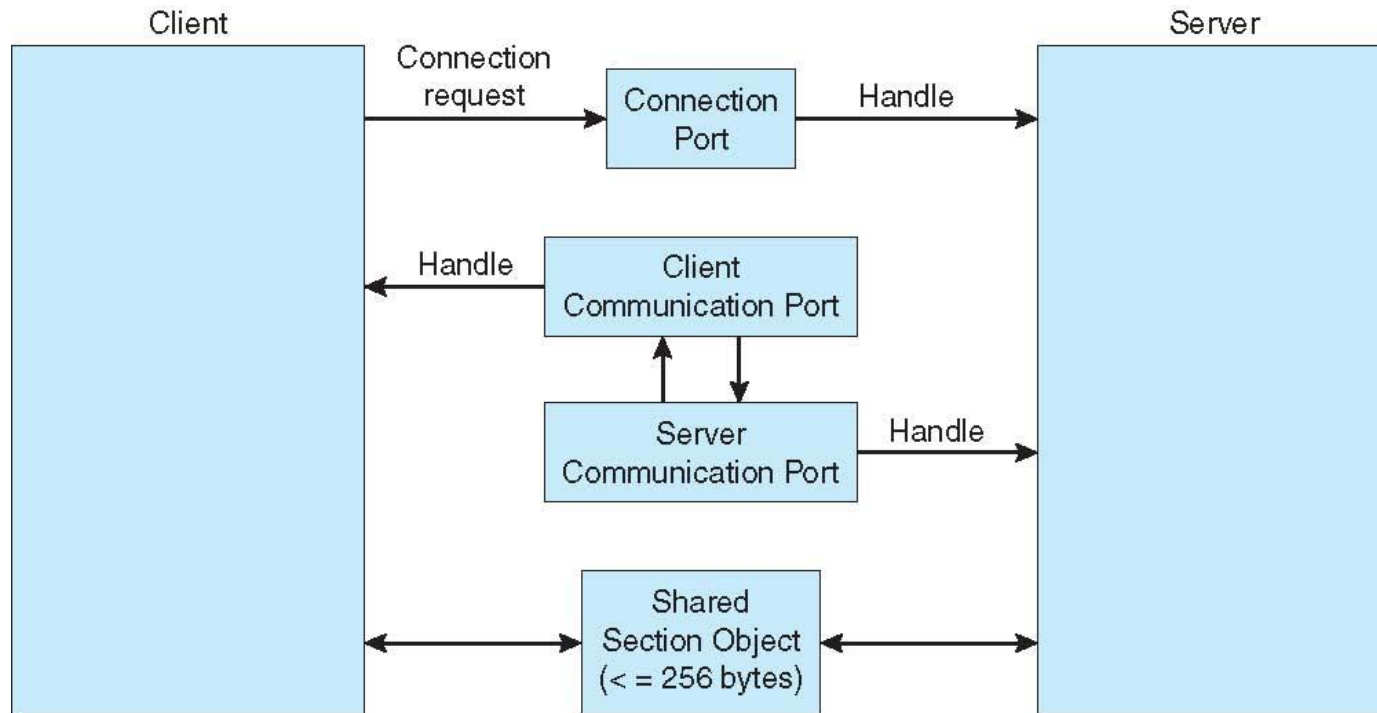
Communications in Client-Server Systems

- Sockets

- Remote Procedure Calls

- Pipes

- Remote Method Invocation (Java)

Sockets

- A **socket** is defined as an endpoint for communication

- Concatenation of IP address and **port** – a number included at start of message packet to differentiate network services on a host

- The socket **161.25.19.8:1625** refers to port **1625** on host **161.25.19.8**

- Communication consists between a pair of sockets

- All ports below 1024 are *well known*, used for standard services

- Special IP address 127.0.0.1 (**loopback**) to refer to system on which process is running

Socket Communication

host X
(146.86.5.20)

socket
(146.86.5.20:1625)

web server
(161.25.19.8)
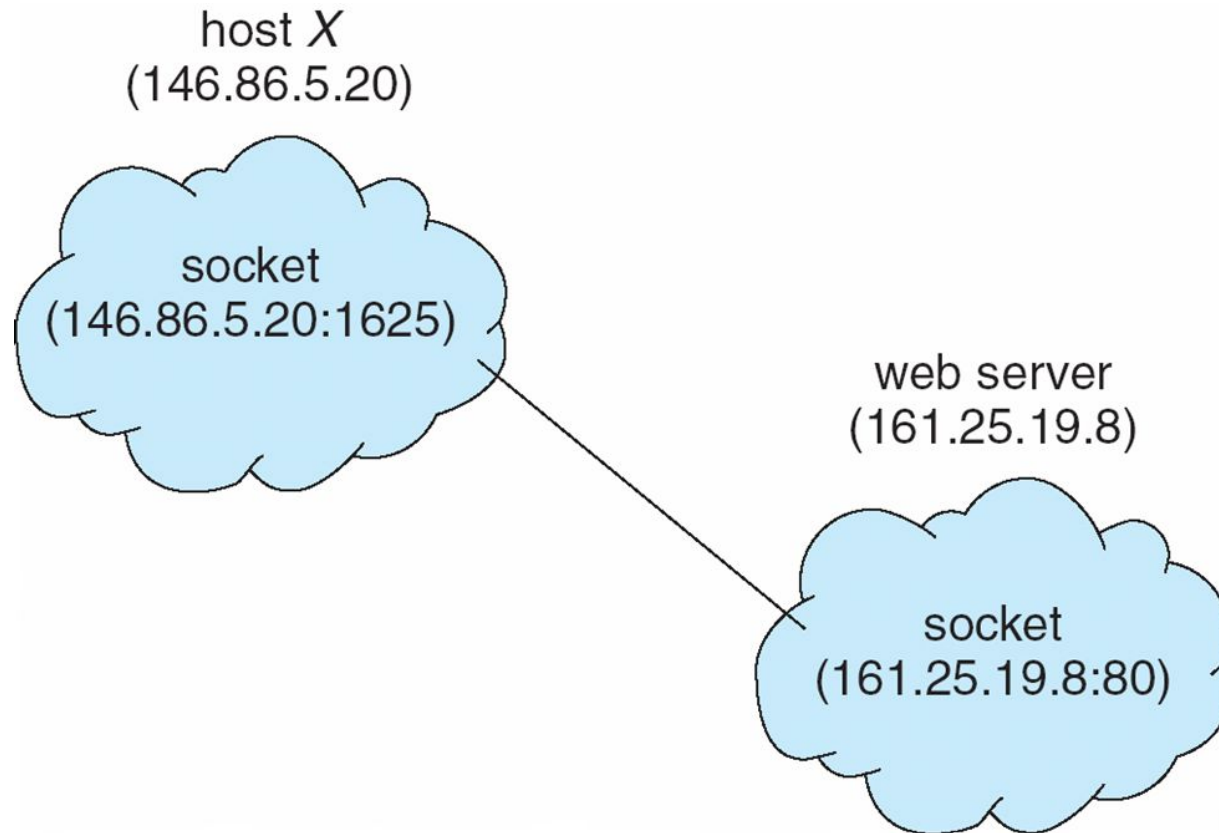
socket
(161.25.19.8:80)

Sockets in Java

- Three types of sockets

  - **Connection-oriented** (**TCP**)

  - **Connectionless** (**UDP**)

  - **MulticastSocket** class– data can be sent to multiple recipients

- Consider this "Date" server:

```java
import java.net.*;
import java.io.*;

public class DateServer
{
   public static void main(String[] args) {
      try {
         ServerSocket sock = new ServerSocket(6013);

         /* now listen for connections */
         while (true) {
            Socket client = sock.accept();

            PrintWriter pout = new
              PrintWriter(client.getOutputStream(), true);

            /* write the Date to the socket */
            pout.println(new java.util.Date().toString());

            /* close the socket and resume */
            /* listening for connections */
            client.close();
         }
      }
      catch (IOException ioe) {
         System.err.println(ioe);
      }
   }
}
```

Remote Procedure Calls

- Remote procedure call (RPC) abstracts procedure calls between processes on networked systems

  - Again uses ports for service differentiation

- **Stubs** – client-side proxy for the actual procedure on the server

- The client-side stub locates the server and **marshalls** the parameters

- The server-side stub receives this message, unpacks the marshalled parameters, and performs the procedure on the server

- On Windows, stub code compile from specification written in **Microsoft Interface Definition Language** (**MIDL**)

Remote Procedure Calls (Cont.)

- Data representation handled via **External Data Representation** (**XDL**) format to account for different architectures
    - **Big-endian** and **little-endian**
- Remote communication has more failure scenarios than local
    - Messages can be delivered *exactly once* rather than *at most once*
- OS typically provides a rendezvous (or **matchmaker**) service to connect client and server

Execution of RPC
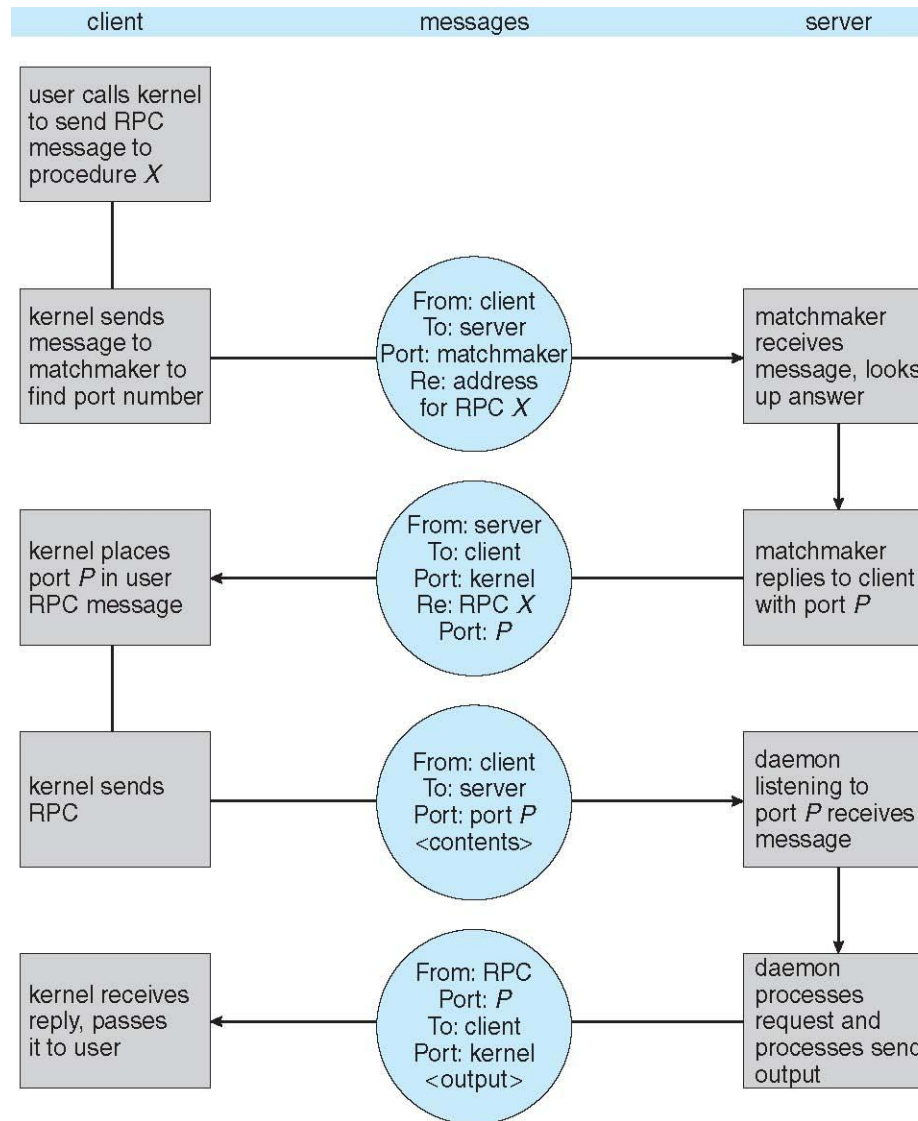
Pipes

- Acts as a conduit allowing two processes to communicate

- Issues:

  - Is communication unidirectional or bidirectional?

  - In the case of two-way communication, is it half or full-duplex?

  - Must there exist a relationship (i.e., *parent-child*) between the communicating processes?

  - Can the pipes be used over a network?

- Ordinary pipes – cannot be accessed from outside the process that created it. Typically, a parent process creates a pipe and uses it to communicate with a child process that it created.

- Named pipes – can be accessed without a parent-child relationship.

Ordinary Pipes

- Ordinary Pipes allow communication in standard producer-consumer style

- Producer writes to one end (the **write-end** of the pipe)

- Consumer reads from the other end (the **read-end** of the pipe)

- Ordinary pipes are therefore unidirectional

- Require parent-child relationship between communicating processes
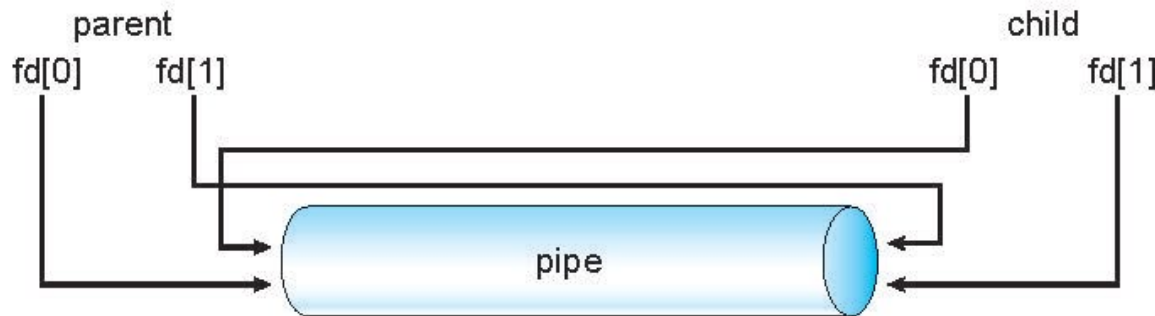
- Windows calls these **anonymous pipes**

- See Unix and Windows code samples in textbook

Named Pipes

- Named Pipes are more powerful than ordinary pipes

- Communication is bidirectional

- No parent-child relationship is necessary between the communicating processes

- Several processes can use the named pipe for communication

- Provided on both UNIX and Windows systems

End of Chapter 3