

HW1

Sicheng Wang

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1 PROBLEM 3

1.1 DHS chapter8, Pb1.

Suppose a decision tree has a path with repeated splits s_1 and s_2 , where $s_1 = s_2 = (f, t)$.

Say that s_1 is prior to s_2 (s_1 is closer to the root).

Suppose at split s_1 , this path goes to the left branch, which means all the data subject to this path have feature f less than threshold t .

Then, at split s_2 , these data will again all fall into the left path. So s_2 has an empty right branch. Therefore, removing s_2 from this path will have no impact on the decision result.

Similarly, if this path goes to the right branch at s_1 , removing s_2 from this path will have no impact on the decision result.

Consequently, an equivalent decision tree only with distinct splits on each path can be constructed by iteratively removing all the repeated splits which is closer to the leaves.

1.2 DHS chapter8, 2.(a)

For a B -way split, if this feature is numeric, then this split can be treated as cutting the feature value range into B sections. In this case, it can be represented with $B - 1$ binary splits, from low to high. If the feature is categorical, this split can be represented by $B - 1$ 'equals' split. Therefore a B -way split, the same function can be achieved by replacing it with $B - 1$ binary splits, which will also produce B possible leaves, each of which corresponding to one of the child nodes of the B -way split.

That is to say, given an arbitrary decision tree, all of its binary non-binary split nodes can be replaced with binary splits. By doing so, this new decision tree is a binary tree because it contains only binary splits.

1.3 DHS chapter8, 2.(b)

For a B -way split, there are B possible outcomes, since one binary split will only introduce one more outcome, there needs to be at least $B - 1$ nodes. which will take at least $\log_2(B - 1)$ levels (by splitting a numerical feature space in a binary search fashion). As for the upper bound, there can be as many as $B - 1$ splits in a row, in a root-leaf path (no repeating split), which can take as much as $B - 1$ levels.

1.4 DHS chapter8, 2.(c)

If the split is limited to a single feature: As demonstrated in b), the lower bound for number of nodes is $B - 1$. The upper bound case is $B - 1$, when all thresholds are encountered from low to high in a path (all the other branches are leaf nodes). If the split is not limited to a single feature: the lower bound is $B - 1$. And the upper bound is $2^B - 1$, in which case the decision tree is a complete tree.

2 PROBLEM 4

2.1 (a)

$$\begin{aligned}\Delta i(N) &= i(N) - P_L i(N_L) - (1 - P_L) i(N_R) \\ &= -\sum_j \left[\frac{N_j}{N} \log_2 \left(\frac{N_j}{N} \right) \right] + \frac{N_L}{N} \sum_j \left[\frac{N_{jL}}{N_L} \log_2 \left(\frac{N_{jL}}{N_L} \right) \right] + \frac{N - N_L}{N} \sum_j \left[\frac{N_j - N_{jL}}{N - N_L} \log_2 \left(\frac{N_j - N_{jL}}{N - N_L} \right) \right] \\ &= -\frac{1}{N} \sum_j \left([N_j (\log_2 N_j - \log_2 N) - N_{jL} (\log_2 N_{jL} - \log_2 N_L) - (N_j - N_{jL}) (\log_2 (N_j - N_{jL}) - \log_2 (N - N_L))] \right) \\ &= -\frac{1}{N} \sum_j \left([N_j \log_2 N_j - N_{jL} \log_2 N_{jL} - (N_j - N_{jL}) \log_2 (N_j - N_{jL})] \right) \\ &\quad + \frac{1}{N} \sum_j \left([N_j \log_2 N - N_{jL} \log_2 N_L - (N_j - N_{jL}) \log_2 (N - N_L)] \right)\end{aligned}$$

The first part is a convex function, which is maximized to 0 when $N_{jL} = 0$ or $N_{jL} = N_j$, depending on N_L because $\sum_i N_{jL} = N_L$.

The second part is a convex function, which is minimized to 1 when $N_L = \frac{N}{2}$.

Finally, this equation is maximized to 1 when the given data is separated to two subsets with equal size, and there are no labels appear in both subsets.

In another way, "yes/no" query means a binary split, the data points decoded by this point either go to the left branch or the right branch. This information takes up to 1 bit to encode (maximum of 1 if they are split equally and one label never goes to two branches, otherwise it's less informative).

2.2 (b)

The decrease in entropy impurity provided by a single B -way split can never be greater than $\log_2 B$ bit.

3 PROBLEM 5

To find the normal equations solution, minimize

$$e(a, b) = \sum_i [(y_i - ax_i - b)^2]$$

Take the partial derivative of this equation:

$$\frac{\partial e(a, b)}{\partial a} = \sum_i [(y_i - ax_i - b)x_i]$$

$$\frac{\partial e(a, b)}{\partial b} = \sum_i [(y_i - ax_i - b)]$$

To find the minimal, let $\frac{\partial e(a, b)}{\partial a} = \frac{\partial e(a, b)}{\partial b} = 0$ Then,

$$\begin{cases} \sum_i (x_i^2)a + \sum_i x_i b - \sum_i x_i y_i = 0 \\ \sum_i x_i a + nb - \sum_i y_i = 0 \end{cases}$$

Solve this equation,

$$\begin{cases} a = \frac{n \sum_i (x_i y_i) - (\sum_i x_i)(\sum_i y_i)}{n \sum_i (x_i^2) - (\sum_i x_i)^2} \\ b = \frac{\sum_i (y_i)(\sum_i x_i)^2 - \sum_i (x_i y_i) \sum_i (x_i)}{n \sum_i (x_i^2) - (\sum_i x_i)^2} \end{cases}$$

4 PROBLEM 7

Suppose two convex hulls are both linearly separable and intersected.

Because they intersect, there must exist a point that is inside both convex hulls.

Let $x = [x_1 \ x_2 \ \dots \ x_n]$ be a point in the intersecting area, $\alpha^{(1)} = [\alpha_1 \ \alpha_2 \ \dots \ \alpha_n]$ for this point represented with the first convex hull's vectors, and $\alpha^{(2)} = [\alpha_1 \ \alpha_2 \ \dots \ \alpha_n]$ for this point represented with the second convex hull's vectors.

$$\alpha^{(1)}x = \alpha^{(2)}x$$

Because these two convex hulls are linearly separable, then we should have:

$\alpha^{(1)}x > 0$ if x is within the first convex hull

and

$\alpha^{(2)}x < 0$ if x is within the second convex hull

This conflicts with the first equation. So two convex hulls cannot be both linearly separable and intersected.

5 PROBLEM 8

$$\begin{aligned} & \nabla_A \text{tr}(ABA^T C) \\ &= \nabla_A \text{tr}(ABA^T C)(A^T \text{ treated as constant}) + (\nabla_{A^T} \text{tr}(A^T CAB))^T (A \text{ treated as constant}) \\ &= (BA^T C)^T + ((CAB)^T)^T \\ &= C^T AB^T + CAB \end{aligned}$$