

Chapter 4

Relational Algebra



fit@hcmus

KHOA CÔNG NGHỆ THÔNG TIN
TRƯỜNG ĐẠI HỌC KHOA HỌC TỰ NHIÊN

Content

1. Introduction
2. Relational Algebra Operations From Set Theory
3. Select Operation
4. Projection Operation
5. CARTESIAN Operations
6. Join Operation
7. Division Operation
8. Additional Relational Operations
9. Update Operations

1. Introduction

□ Update on the relation “KHOA”

- ❑ Add a new Khoa ‘Hóa học’ into the relation
- ❑ Move “CNTT” to the room “B12”

MÃKHOA	TÊNKHOA	NĂMTL	PHÒNG	ĐIỆNTHOAI	TRƯỞNGKHOA	NGÀYNHẬNCHỨC
CNTT	Công nghệ thông tin	1995	B12	0838123456	002	20/02/2005
VL	Vật lý	1976	B21	0838223223	005	18/09/2003
SH	Sinh học	1980	B31	0838454545	004	11/10/2000
HH	Hóa học	1980	B41	NULL	007	15/10/2001

- ❑ Query: name of the faculty found from the year 1980

TÊNKHOA
Sinh học
Hóa học

1. Introduction (cont.)



Two types of operations:

- ❑ Make change the data (update): add new, delete and modify
- ❑ Don't make change the data (extract data): query



Query Language – QL:

- ❑ Allow to extract or update data stored in a database schema



Types of Relational Query:

- ❑ Relational Algebra

❑ A sequence of relational algebra operations forms a **relational algebra expression**, whose result will also be a relation that represents the result of a database query (or retrieval request).

- ❑ Relational Calculus

❑ Considered to be a nonprocedural language. This differs from relational algebra, where we must write a sequence of operations to specify a retrieval request; hence relational algebra can be considered as a procedural way of stating a query



SQL (Structured Query Language)

1. Introduction (cont.)

□ Algebra

- ❑ Operator (toán tử)
- ❑ Operand (toán hạng)

□ Arithmetic (số học)

- ❑ Operators: +, -, *, /
- ❑ Operands or variables: x, y, z
- ❑ Constant
- ❑ Expression
 - ❑ $(x+7) / (y-3)$
 - ❑ $(x+y)*z$ and/or $(x+7) / (y-3)$

1. Introduction (cont.)

Relational Algebra

- Operands – variables are relations
 - ? Or set (tập hợp)
- Operators are relational operations
 - ? Relational Algebra Operations From Set Theory
 - Union \cup (hội)
 - Intersect \cap (giao)
 - Difference set/ Minus – (trừ)
 - ? Unary Relational Operations
 - Selection σ (chọn)
 - Projection π (chiếu)
 - ? Additional relational operations
 - Cartesian product \times (Tích Cartesian)
 - Join \bowtie (kết)

1. Introduction (cont.)

In Algebra relation, a query is expressed through a **relational algebra expression**:

- ❑ The result of a retrieval is a new relation, which may have been formed from one or more relations.
- ❑ The algebra operations thus produce new relations, which can be further manipulated using operations of the same algebra
- ❑ It is formed by a sequence of relational algebra operations
- ❑ the result of the previous operation is input to the next operation.

1. Introduction (cont.)

	Algebra	Relational algebra
Operand	<ul style="list-style-type: none">- Variable : x, y, z, \dots- Constant : $150, \dots$	<ul style="list-style-type: none">- Schema : $NhanVien, \dots$- Relation/Instances/ tuples : t, v, \dots
Operator	<ul style="list-style-type: none">- Operations between operands to form a new value: $+, -, *, /, \dots$	<ul style="list-style-type: none">- Operations between relations to form a new relation: $\text{selection } \sigma, \text{ union } \cup, \dots$
Expression	<ul style="list-style-type: none">- Sequence of algebra operators- Return a new value $(x+7) / (y-3)$ $(x+y)*z \text{ and/or } (x+7) / (y-3)$	<ul style="list-style-type: none">- Sequence of relational algebra (query)- Return a new relation $\pi_{MANV} (NHANVIEN)$

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2. Relational Algebra Operations From Set Theory

Operations

- ❑ Union $r \cup s$
- ❑ Intersect $r \cap s$
- ❑ Difference set/ Minus $r - s$

Union Compatibility (khả hợp)

- ❑ The operand relations $R(A_1, A_2, \dots, A_n)$ and $S(B_1, B_2, \dots, B_n)$:
- ❑ Two are compatible if
 - ❑ R and S have the same number of attributes (e.g. n)
 - ❑ and $\text{DOM}(A_i) = \text{DOM}(B_i)$, $1 \leq i \leq n$

The resulting relation for $R \cup S$, $R \cap S$, or $R - S$ has the same attribute names as the *first* operand relation R (by convention).

2. Union

Example

SINHVIEN	TENSV	NGSINH	PHAI
	Tung	12/08/1955	Nam
	Hang	07/19/1968	Nu
	Nhu	06/20/1951	Nu
	Hung	09/15/1962	Nam

GIAOVIEN	TENGV	NG_SINH	GIOITINH
	Trinh	04/05/1986	Nu
	Khang	10/25/1983	Nam
	Phuong	05/03/1958	Nu
	Minh	02/28/1942	Nam
	Chau	12/30/1988	Nu

Degree n=3

$\text{DOM}(\text{TENSV}) = \text{DOM}(\text{TENGV})$

$\text{DOM}(\text{NGSINH}) = \text{DOM}(\text{NG_SINH})$

$\text{DOM}(\text{PHAI}) = \text{DOM}(\text{GIOITINH})$

2. Union

- Given two compatible relations **r** and **s**
- Union of **r** and **s**

- Notation: $r \cup s$
- The result of the union operation includes tuples in **r** or in **s**, or in both relations (duplicated tuples are eliminated)

$$r \cup s = \{ t / t \in r \vee t \in s \}$$

- Example

r	A	B
	α	1
	α	2
	β	1

s	A	B
	α	2
	β	3

$r \cup s$	A	B
α	1	
α	2	
β	1	
α	2	
β	3	

2. Union - example

SinhVien	
HOTEN	DIACHI
Đinh Bá Tiên	119 Công Quỳnh, Tp HCM
Nguyễn Thanh Tùng	222 Nguyễn Văn Cừ, Tp HCM
Lê Quỳnh Như	291 Hồ Văn Huê, Tp HCM

GiaoVien	
HOTEN	DIACHI
Đinh Bá Tiên	119 Công Quỳnh, Tp HCM
Trần Thanh Tâm	553 Mai Thị Lựu, Tp HCM

SinhVien \cup GiaoVien	
HOTEN	DIACHI
Đinh Bá Tiên	119 Công Quỳnh, Tp HCM
Nguyễn Thanh Tùng	222 Nguyễn Văn Cừ, Tp HCM
Trần Thanh Tâm	553 Mai Thị Lựu, Tp HCM

2. Intersect



Given two compatible relations r and s



Intersect of r and s

- ❑ Notation: $r \cap s$
- ❑ The result relation of the intersect operation includes tuples that are both in r and s

$$r \cap s = \{ t / t \in r \wedge t \in s \}$$

r	A	B
α	1	
α	2	
β	1	

s	A	B
α	2	
β	3	

$r \cap s$	A	B
α	2	

2. Intersect (cont.) - Example

SinhVien	
HOTEN	DIACHI
Đinh Bá Tiên	119 Công Quỳnh, Tp HCM
Nguyễn Thanh Tùng	222 Nguyễn Văn Cừ, Tp HCM
Lê Quỳnh Như	291 Hồ Văn Huê, Tp HCM

GiaoVien	
HOTEN	DIACHI
Đinh Bá Tiên	119 Công Quỳnh, Tp HCM
Trần Thanh Tâm	553 Mai Thị Lựu, Tp HCM

SinhVien \cap GiaoVien	
HOTEN	DIACHI
Đinh Bá Tiên	119 Công Quỳnh, Tp HCM

2. Minus

- Given two compatible relations **r** and **s**
- Minus of **r** and **s**

- Notation: $r - s$
- The result of the minus operation includes tuples **in r** and **not in s**

$$r - s = \{ t / t \in r \wedge t \notin s \}$$

r	A	B
α	1	
α	2	
β	1	

s	A	B
α	2	
β	3	

r - s	A	B
α	1	
β	1	

2. Minus - example

SinhVien

HOTEN	DIACHI
Đinh Bá Tiên	119 Công Quỳnh, Tp HCM
Nguyễn Thanh Tùng	222 Nguyễn Văn Cừ, Tp HCM
Lê Quỳnh Như	291 Hồ Văn Huê, Tp HCM

GiaoVien

HOTEN	DIACHI
Đinh Bá Tiên	119 Công Quỳnh, Tp HCM
Trần Thanh Tâm	553 Mai Thị Lựu, Tp HCM

SinhVien – GiaoVien

HOTEN	DIACHI
Nguyễn Thanh Tùng	222 Nguyễn Văn Cừ, Tp HCM
Lê Quỳnh Như	291 Hồ Văn Huê, Tp HCM

2. Characteristics

- Notice that both union and intersection are *commutative operations* (*tính giao hoán*); that is

$$r \cup s = s \cup r$$

$$r \cap s = s \cap r$$

- Both union and intersection can be treated as n-ary operations applicable to any number of relations as both are *associative operations* (*tính kết hợp*); that is

$$r \cup (s \cup t) = (r \cup s) \cup t$$

$$r \cap (s \cap t) = (r \cap s) \cap t$$

- The minus operation is *not commutative*; that is, in general

$$R - S \neq S - R$$

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3. Select Operation

- ❑ Is used to select a *subset* of the tuples from a relation that satisfy a selection condition P. It is a filter that keeps only those tuples that satisfy a qualifying condition – those satisfying the condition are selected while others are discarded
- ❑ Denoted by
$$\sigma_P(r)$$

σ: sigma
- ❑ P is an expression consisting of clauses as:
 - ❑ <attribute> <comparison> <constant>
 - ❑ <attribute> <comparison> <comparision>
 - ❑ <comparison>: < , > , ≤ , ≥ , ≠ , =
 - ❑ Clauses are connected by logical connections \wedge , \vee , \neg

3. Select Operation (cont.)

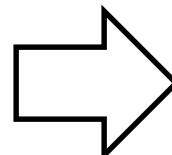
□ The result relation:

- ❑ Attributes: the same attributes of the relation r
- ❑ Tuples: less or equal of r

$$\sigma_{(A=B) \wedge (D>5)}(r)$$

□ Example

r	A	B	C	D
	α	α	1	7
α	β	5	7	
β	β	12	3	
β	β	23	10	



A	B	C	D
α	α	1	7
β	β	23	10

3. Select Operation (cont.)

- The select is *commutative operations*; that is

$$\sigma_{p_1}(\sigma_{p_2}(r)) = \sigma_{p_2}(\sigma_{p_1}(r)) = \sigma_{p_1 \wedge p_2}(r)$$

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4. Projection Operation

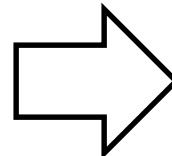
- This operation selects certain columns from a relation r and discards the other columns
- Notation

$$\pi_{A_1, A_2, \dots, A_k}(r)$$

- The result is a relation
 - ? having k attributes
 - ? Number of tuples are less than r
- Example

r	A	B	C
	α	10	1
	α	20	1
	β	30	1
	β	40	2

$$\pi_{A,C}(r)$$



$\pi_{A,C}(r)$	A	C
	α	1
	β	1
	β	2

4. Projection Operation (cont.)

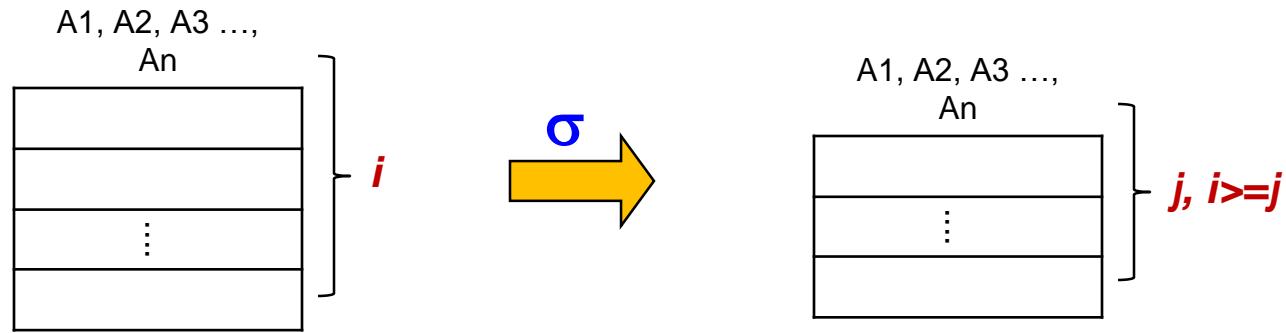
- The projection is *not commutative operations*

$$\pi_{X,Y}(r) = \pi_X(\cancel{\pi_Y(r)})$$

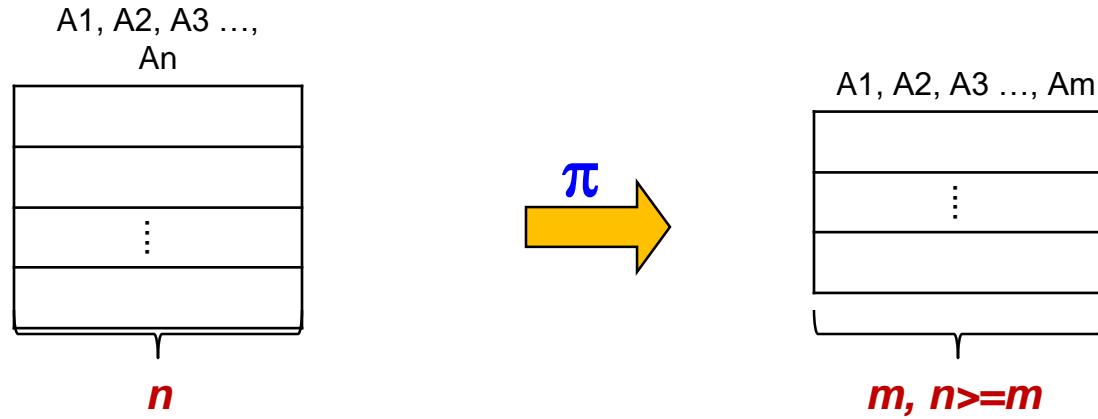
$$\pi_{A_1, A_2, \dots, A_n}(\pi_{A_1, A_2, \dots, A_m}(r)) = \pi_{A_1, A_2, \dots, A_n}(r), \text{ với } n \leq m$$

4. Select vs Projection

Select



Projection



4. Projection – example 1

- Retrieve the name and salary of female teachers

$$\pi_{\text{HOTEN, LUONG}} (\sigma_{\text{PHAI}='Nữ'} (\text{GIAOVIEN}))$$

4. Project Operation – example 2

- List out the codes/IDs of teachers who belong to the dept. “HTTT” or participate in the topic code “001”

$$\pi_{\text{MAGV}} (\sigma_{\text{MABM}='HTTT'} (\text{GIAOVIEN})) \cup \pi_{\text{MAGV}} (\sigma_{\text{MAĐT}='001'} (\text{TG_ĐETAI}))$$

4. Project Operation – example 3

- Retrieve the code/ID of the deans who are also the project managers

$$\pi_{\text{TRUONGKHOA}}(\text{KHOA}) \cap \pi_{\text{GVCNĐT}}(\text{ĐETAI})$$

4. Project Operation – example 3

- Retrieve the names of the jobs that started between 01/01/2007 and 01/08/2007

$$\sigma_{(NGAYBD \geq '1/1/2007' \wedge NGAYBD \leq '1/8/2007')} (CONGVIEC)$$

Quiz #1

Given the schema **PHANCONG(MÃNV, MÃĐA, THOIGIAN)** and the query “*Cho biết danh sách mã nhân viên vừa có tham gia đề án số 1 vừa có tham gia đề án số 2*”. Choose the correct answer:

- A. $\Pi_{MÃNV}(\sigma_{MÃĐA = 1 \text{ AND } MÃĐA = 2} (PHÂNCÔNG))$
- B. $\Pi_{MÃNV} ((\sigma_{MÃĐA = 1} (PHÂNCÔNG) \cap (\sigma_{MÃĐA = 2} (PHÂNCÔNG)))$
- C. $\Pi_{MÃNV} (\sigma_{MÃĐA = 1} (PHÂNCÔNG)) \cap \Pi_{MÃNV} (\sigma_{MÃĐA = 2} (PHÂNCÔNG))$
- D. B, C

4. Projection Operation - General

- ❑ Extend the project operation by adding arithmetical operations in attributes
- ❑ Notation $\pi_{F_1, F_2, \dots, F_n}(E)$, where:
 - ❑ E relation
 - ❑ F_1, F_2, \dots, F_n arithmetical expression containing:
 - ❑ Constants
 - ❑ Attributes of E

4. Projection Operation - General

Example

- Retrieve the name and salary of teachers after increasing 10%

$$\pi_{\text{HOTEN, LUONG}*1.1}(\text{GIAOVIEN})$$

4. Sequence of Operations

□ Combining algebra operations

? Combining expressions

$$\pi_{A_1, A_2, \dots, A_k}(\sigma_P(r))$$
$$\sigma_P(\pi_{A_1, A_2, \dots, A_k}(r))$$

? Execute step by step

■ B1

$$\sigma_P(r)$$

■ B2

$$\pi_{A_1, A_2, \dots, A_k}(\text{result of B1})$$

Rename the resulting relation

4. Assignment

❑ Is used to retrieve a result relation from an algebra operation

❑ Notation: \leftarrow

❑ Example:

❑ B1

$$s \leftarrow \sigma_p(r)$$

❑ B2

$$KQ \leftarrow \pi_{A1, A2, \dots, Ak}(s)$$

4. Rename Operation



We may want to apply several relational algebra operations one after the other. Either we can write the operations as a single **relational algebra expression** by nesting the operations, or we can apply one operation at a time and create **intermediate result relations**. In the latter case, we must give names to the relations that hold the intermediate results.

- Given relation $r(B, C, D)$

$\rho_s(r)$: rename the relation r to s

- Rename attribute

$\rho_{x, c, d}(r)$: rename the attribute B to X

- Rename relation and attribute:

$\rho_{s(x,c,d)}(r)$: rename the relation r to s , and the attribute B to X

4. Example 5

- Retrieve the code/ID and name of teachers who work in the dept. “HTTT”.

$\pi_{\text{MAGV, HOTEN}} (\sigma_{\text{MABM}=\text{'HTTT'}} (\text{GIAOVIEN}))$

□ C1: $\text{GV_HTTT} \leftarrow \sigma_{\text{MABM}=\text{'HTTT'}} (\text{GIAOVIEN})$

$\text{KQ} \leftarrow \pi_{\text{MAGV, HOTEN}} (\text{GV_HTTT})$

□ C2: $\text{KQ(MA, TEN)} \leftarrow \pi_{\text{MAGV, HOTEN}} (\text{GV_HTTT})$

$\rho_{\text{KQ(MA, TEN)}} (\pi_{\text{MAGV, HOTEN}} (\text{GV_HTTT}))$

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5. CARTESIAN Operations

- This operation is used to combine tuples from two relations in a combinatorial fashion.
- Notation

$$r \times s$$

- The result relation is a relation q with:
 - ? Each tuple of q is a combination between 1 tuple in r and 1 tuple in s
 - ? If the relation r has u tuples and s has v tuples, then q has $u \times v$ tuples
 - ? If r has n attributes and s has m attributes then q has $n + m$ attributes ($R^+ \cap S^+ = \emptyset$)

5. CARTESIAN Operations - Example

r	A	B
α	1	
β	2	

s	X	C	D
α	10	+	
β	10	+	
β	20	-	
γ	10	-	

unambiguous

$r \times s$	A	R.B	X	C	D
α	1		α	10	+
α	1		β	10	+
α	1		β	20	-
α	1		γ	10	-
β	2		α	10	+
β	2		β	10	+
β	2		β	20	-
β	2		γ	10	-

$$\rho_{(X,C,D)}(s)$$

5. CARTESIAN Operations

- Usually, the Cartesian operation is followed by the select operation

$r \times s$

A	R.B	S.B	C	D
α	1	α	10	+
α	1	β	10	+
α	1	β	20	-
α	1	γ	10	-
β	2	α	10	+
β	2	β	10	+
β	2	β	20	-
β	2	γ	10	-

$\sigma_{A=S.B} (r \times s)$

A	R.B	S.B	C	D
α	1	α	10	+
β	2	β	10	+
β	2	β	20	-

5. CARTESIAN Operations – Example 6



Retrieve the information of the dept. and the information of the lecturer who is the head of that department

TENBM	MABM	TRUONGBM	NGAYNHANCHUC	...
Hệ thống thông tin	HTTT	002	20/09/2004	
Công nghệ tri thức	CNTT			
Mạng máy tính	MMT	001	15/05/2005	

MAGV	HOTEN	NGSINH	MABM	PHAI	LUONG
001	Nguyễn Hoài An	15/02/1973	MMT	Nam	2000	
002	Trần Trà Dương	20/06/1960	HTTT	Nu	2500	
003	Nguyễn Ngọc Anh	11/05/1975	HTTT	Nu	2200	
004	Trương Nam Sơn	20/06/1959	VS	Nam	2300	

5. CARTESIAN Operations – Example 6

TENBM	MABM	TRUONGBM	NGAYNHANCHUC	GV	HOTEN	...
Hệ thống thông tin	HTTT	002	20/09/2004	002	Trần Trà Dương	
Mạng máy tính	MMT	001	15/05/2005	001	Trương Nam Sơn	

5. CARTESIAN Operations – Example 6

B1: Cartesian operation BOMON and GIAOVIEN

$$\text{BM_GV} \leftarrow (\text{GIÁOVIÊN} \times \text{BỘMÔN})$$

B2: select tuples that satisfy TRUONGBM = MAGV

$$\text{KQ} \leftarrow \sigma_{\text{TRUONGBM} = \text{MAGV}}(\text{BM_GV})$$

GIÁO VIÊN	MÃ GV	HỌ TÊN	NGÀY SINH	SỐ NHÀ
	001	Nguyễn Hoài An	15/02/1973	25/3
	002	Trần Trà Hương	20/06/1960	125
	003	Nguyễn Ngọc Ánh	11/05/1975	12/21

BỘ MÔN	MÃ BM	TÊN BM	PHÒNG	TRƯỞNG BM
	HTTT	Hệ thống thông tin	B13	002
	CNTT	Công nghệ tri thức	B15
	MMT	Mạng máy tính	B16	001

$\sigma_{\text{TRUONGBM}=MAGV}(\text{BM}_\text{GV})$									
MÃ GV	HỌ TÊN	NGÀY SINH	MÃ BM	TÊN BM	PHÒNG	TRƯỞNG BM
001	Nguyễn Hoài An	15/02/1973	HTTT	Hệ thống thông tin	B13	002
001	Nguyễn Hoài An	15/02/1973	CNTT	Công nghệ tri thức	B15	002
001	Nguyễn Hoài An	15/02/1973	MMT	Mạng máy tính	B16	001
....

5. CARTESIAN Operations – Example 7

☐ Retrieve the highest salary of the lecturers?

?] GIANGVIEN(MSGV, HOTEN, LUONG, ...)

HOTEN	...	LUONG	LUONG	...
Nguyễn Hoài An	...	2000	2000	...
Trần Trà Hương	...	2500	2500	...
Nguyễn Ngọc Anh	...	2200	2200	...

5. CARTESIAN Operations – Example 7

B1: Retrieve the salaries that are **not highest**

$$r1 \leftarrow (\pi_{LUONG} (GIAOVIEN))$$
$$r2 \leftarrow \sigma_{GIAOVIEN.LUONG < R1.LUONG} (GIAOVIEN \times r1)$$
$$r3 \leftarrow \pi_{R2.LUONG} (r2)$$

B2: Full salary set minuses the salary set in R3

$$KQ \leftarrow \pi_{LUONG} (GIAOVIEN) - r3$$

5. CARTESIAN Operations – Example 8

Retrieve the name of lecturers working the same department with ‘Trần Trà Hương’

Relation: GIAOVIEN

Attributes: HOTEN, MABM

Condition: HOTEN = ‘Trần Trà Hương’

What dept. “Trần Trà Hương” belongs to?

MABM	HOTEN
MMT	Nguyễn Hoài An
HTTT	Trần Trà Hương
HTTT	Nguyễn Ngọc Anh
VS	Trương Nam Sơn
...	

What lecturers belong to that dept.?

MABM	HOTEN
MMT	Nguyễn Hoài An
HTTT	Trần Trà Hương
HTTT	Nguyễn Ngọc Anh
VS	Trương Nam Sơn
...	

5. CARTESIAN Operations – Example 8

□ B1: Find the dept. of ‘Trần Trà Hương’

$$r1 \leftarrow \pi_{MABM, MAGV} (\sigma_{HOTEN='Trần Trà Hương'} (GIAOVIEN))$$

□ B2: Retrieve the name of lecturers of that dept.

$$r2 \leftarrow \sigma_{HOTEN <> 'Trần Trà Hương'} (GIAOVIEN)$$
$$r3 \leftarrow \sigma_{R1.MABM=R2.MABM} (r1 \times r2)$$
$$KQ \leftarrow \pi_{HOTEN} (r3)$$

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6. **Join Operation**
 1. Natural join (kết tự nhiên)
 2. Theta join (kết có điều kiện tổng quát)
 3. Equi join (kết bằng)
7. Division Operation
8. Additional Relational Operations
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6. Join Operation

- ❑ To identify and select related tuples from two relations
- ❑ Denoted by: $r \bowtie s$
 - ❑ $R(A_1, A_2, \dots, A_n)$ and $S(B_1, B_2, \dots, B_m)$
- ❑ The result relation is a relation q
 - ❑ Having $n + m$ attributes: $Q(A_1, A_2, \dots, A_n, B_1, B_2, \dots, B_m)$
 - ❑ Each tuple of q is a combination of 1 in r and 1 in s that satisfy join condition:
 - ❑ Join condition: $A_i \theta B_j$
 - ❑ A_i attribute of R, B_j attribute of S
 - ❑ A_i and B_j have the same value domain
 - ❑ θ is a comparison: $\neq, =, <, >, \leq, \geq$

6. Join Operation

□ Theta join

- ❑ Ký hiệu $r \bowtie_C s$
- ❑ C join condition

□ Equi join: when the condition C is a equal join

□ Natural join: when two attribute of the condition C have the same name:

- ❑ Denoted by: $r \bowtie s$ or $r * s$
- ❑ $R^+ \cap S^+ \neq \emptyset$
- ❑ The attributes of the result relation are in r or s , minus the attributes in the join condition

6. Join Operation

□ Theta join - Example

r	A	B	C
1	2	3	
4	5	6	
7	8	9	

s	D	E
3		1
6		2

$r \bowtie_{B < D} s$

A	B	C	D	E
1	2	3	3	1
1	2	3	6	2
4	5	6	6	2

$$r \bowtie_C s = \sigma_C(r \times s)$$

6. Join Operation

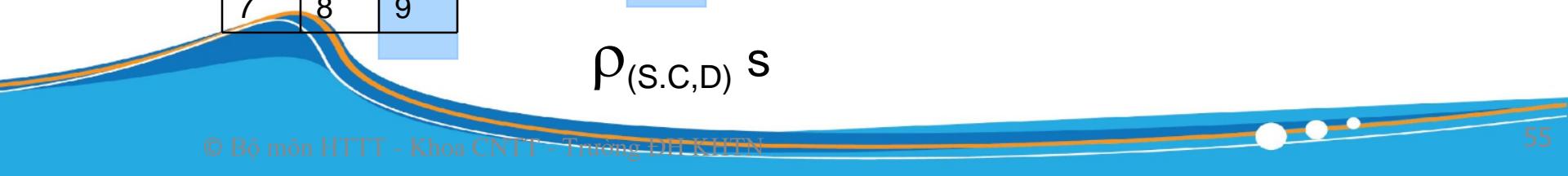
□ Equi join - Example

r	A	B	C
1	2	3	
4	5	6	
7	8	9	

s	D	E
3	1	
6	2	

 $r \bowtie_{C=D} s$

A	B	C	D	E
1	2	3	3	1
4	5	6	6	2



r	A	B	C
1	2	3	
4	5	6	
7	8	9	

s	s.C	D
3	1	
6	2	

 $r \bowtie_{C=s.C} s$

A	B	C	s.C	D
1	2	3	3	1
4	5	6	6	2

 $\rho_{(S.C,D)} s$

6. Join Operation

□ Natural join - Example

r	A	B	C
1	2	3	
4	5	6	
7	8	9	

s	C	D
3		1
6		2

$r \bowtie s$

A	B	C	D	
1	2	3	1	
4	5	6	2	

6. Join Operation – Example 10

□ Retrieve the lecturers who have the salary greater than the lecturer ‘Nguyễn Hoài An’

❑ Relation: **GIAOVIEN**

❑ Attribute: **LUONG**
GIAOVIEN(MAGV, HOTEN, LUONG, PHAI,
NGAYSINH,...)

$R1(LG) \leftarrow \pi_{LUONG} (\sigma_{HOTEN='Nguyễn Hoài An'} (GIAOVIEN))$

$KQ \leftarrow GIAOVIEN \bowtie_{LUONG > LG} R1$

$KQ(MAGV, HOTEN, LUONG, PHAI, NGAYSINH, \dots, LG)$

6. Join Operation – Example 11

For each lecturer, list out the information of the department that they are working for

Relation: **GIAOVIEN, BOMON**

GIAOVIEN(MAGV, HOTEN, LUONG, PHAI, ..., **MABM**, ...)

BOMON(**MABM**, TENBM, PHONG, DIENTHOAI, ...)

KQ \leftarrow GIAOVIEN \bowtie BOMON

KQ(MAGV, HOTEN, ..., **MABM**, TENBM, PHONG, ...))

6. Join Operation – Example 12

For each project, list out the information of the lecturers who manage that project

Relation: **ĐETAI, GIAOVIEN**

ĐETAI(MAĐT, TENĐT, KINHPHI, ..., **GVCNĐT**)

GIAOVIEN(**MAGV**, HOTEN, LUONG, PHAI, ...)

KQ \leftarrow ĐETAI $\bowtie_{GVCNĐT = MAGV}$ GIAOVIEN

KQ(MAĐT, TENĐT, KINHPHI, ..., **GVCNĐT, MAGV, HOTEN, ...**)

6. Join Operation – Example 13

- Retrieve the highest salary of the department 'HTTT'?

Quiz

Given the DB schema below:

- PHÒNG(Mã Phòng, Diện_tích, Giá_tiền, Loại)
- THUÊ(Mã Phòng, Mã KH, Ngày_nhân, Số_ngày)
- KHÁCH_HÀNG(Mã KH, Họ_tên, Điện_thoại, Giới_tính)

Find out the name of customers who rent the room ‘R101’:

- A) $\pi_{Họ_Tên}(\sigma_{Mã_Phòng='R101'}(KHÁCH_HÀNG * THUÊ))$
- B) $\pi_{Họ_Tên}(\sigma_{Mã_Phòng='R101'}(KHÁCH_HÀNG \bowtie KHÁCH_HÀNG.Mã_KH \leftrightarrow THUÊ.Mã_KH THUÊ))$
- C) $\pi_{Họ_Tên}(KHÁCH_HÀNG \bowtie KHÁCH_HÀNG.Mã_KH = THUÊ.Mã_KH \text{ AND } Mã_Phòng <> 'R101' THUÊ))$
- D) Chỉ có A và B đúng.
- E) Tất cả A, B và C đúng.

Complete Set of Relational Operations

□ The set of operations including **select** σ , **project** π , **union** \cup , **set difference** $-$, and **cartesian product** \times is called a complete set because any other relational algebra expression can be expressed by a combination of these five operations.

Example

- $r \cap s = r \cup s - ((r-s) \cup (s-r))$
- $r \bowtie_c s = \sigma_c(r \times s)$

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6. Join Operation
7. **Division Operation**
8. Additional Relational Operations
9. Update Operations

7. Division Operation

- To extract tuples in the relation r that satisfy all tuples in the relation s
- Denoted by $r \div s$
 - $r(Z) \text{ và } x(X)$
 - Z are attributes of r , X are attributes of s
 - $X \subseteq Z$
- The result relation of the division $t(Y)$
 - With $Y = Z - X$
 - t_0 is a tuple of t if all tuples $t_s \in s$, there exists a tuple $t_r \in r$ that satisfies 2 two below conditions:
 - $t_r(Y) = t_0$
 - $t_r(X) = t_s(X)$

R(Z)
X Y

S(X)

T(Y)

7. Division Operation - Example

 $r \div s$

r	A	B	C	D	E
•	α	a	α	a	1
•	α	a	γ	a	1
•	α	a	γ	b	1
•	β	a	γ	a	1
•	β	a	γ	b	3
•	γ	a	γ	a	1
•	γ	a	γ	b	1
•	γ	a	β	b	1

s	D	E
a	1	
b	1	

A	B	C
α	a	γ
γ	a	γ

7. Division Operation - Example

DANGKY	MASV	MANH
	SV01	MH01
	SV02	MH01
	SV03	MH01
	SV01	MH02
	SV02	MH02
	SV03	MH03
	SV01	MH03
	SV04	MH01

MONHOC	MANH
	MH01
	MH02
	MH03

DANGKY ÷ MONHOC	MASV
	SV01

7. Division Operation – Example 17

- Retrieve the codes/IDs of lecturers who participate all works of the project “001”

7. Division Operation – Example 18

- Retrieve the name of the projects that all lecturers of the dept. “Hệ thống thông tin” are working on

7. Division Operation

□ **Division operation** can be expressed through a complete set of relational algebra operations

$$Q1 \leftarrow \pi_Y(r)$$

$$Q2 \leftarrow Q1 \times s$$

$$Q3 \leftarrow \pi_Y(Q2 - r)$$

$$KQ \leftarrow Q1 - Q3$$

Recap of Relational Algebra Operations

TABLE 6.1 OPERATIONS OF RELATIONAL ALGEBRA

Operation	Purpose	Notation
SELECT	Selects all tuples that satisfy the selection condition from a relation R .	$\sigma_{\text{<SELECTION CONDITION>}}(R)$
PROJECT	Produces a new relation with only some of the attributes of R , and removes duplicate tuples.	$\pi_{\text{<ATTRIBUTE LIST>}}(R)$
THETA JOIN	Produces all combinations of tuples from R_1 and R_2 that satisfy the join condition.	$R_1 \bowtie_{\text{<JOIN CONDITION>}} R_2$
EQUIJOIN	Produces all the combinations of tuples from R_1 and R_2 that satisfy a join condition with only equality comparisons.	$R_1 \bowtie_{\text{<JOIN CONDITION>}} R_2$, OR $R_1 \bowtie_{(\text{<JOIN ATTRIBUTES 1>}), (\text{<JOIN ATTRIBUTES 2>})} R_2$
NATURAL JOIN	Same as EQUIJOIN except that the join attributes of R_2 are not included in the resulting relation; if the join attributes have the same names, they do not have to be specified at all.	$R_1 \bowtie_{\text{<JOIN CONDITION>}} R_2$, OR $R_1 \bowtie_{(\text{<JOIN ATTRIBUTES 1>}), (\text{<JOIN ATTRIBUTES 2>})} R_2$ OR $R_1 * R_2$
UNION	Produces a relation that includes all the tuples in R_1 or R_2 or both R_1 and R_2 ; R_1 and R_2 must be union compatible.	$R_1 \cup R_2$
INTERSECTION	Produces a relation that includes all the tuples in both R_1 and R_2 ; R_1 and R_2 must be union compatible.	$R_1 \cap R_2$
DIFFERENCE	Produces a relation that includes all the tuples in R_1 that are not in R_2 ; R_1 and R_2 must be union compatible.	$R_1 - R_2$
CARTESIAN PRODUCT	Produces a relation that has the attributes of R_1 and R_2 and includes as tuples all possible combinations of tuples from R_1 and R_2 .	$R_1 \times R_2$
DIVISION	Produces a relation $R(X)$ that includes all tuples $t[X]$ in $R_1(Z)$ that appear in R_1 in combination with every tuple from $R_2(Y)$, where $Z = X \cup Y$.	$R_1(Z) \div R_2(Y)$

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8. Additional Relational Operations
 - Aggregate function (Hàm kết hợp)
 - Grouping (gom nhóm)
 - Outer join (Kết ngoài)
9. Update Operations

8. Aggregate functions

- A type of request that cannot be expressed in the basic relational algebra is to specify mathematical **aggregate functions** on collections of values from the database.
- Examples of such functions include retrieving the average or total salary of all employees or the total number of employee tuples. These functions are used in simple statistical queries that summarize information from the database tuples.
- Common functions applied to collections of numeric values include **SUM**, **AVERAGE**, **MAXIMUM**, and **MINIMUM**. The **COUNT** function is used for counting tuples or values.

8. Aggregate functions

- Input: set of values or relation
- Output: single value

- ❑ AVG
- ❑ MIN
- ❑ MAX
- ❑ SUM
- ❑ COUNT

r	A	B
1	2	
3		4
1	2	
1	2	

$$\begin{aligned} \text{SUM}(B) &= 10 \\ \text{AVG}(A) &= 1.5 \\ \text{MIN}(A) &= 1 \\ \text{MAX}(B) &= 4 \\ \text{COUNT}(A) &= 4 \end{aligned}$$

8. Grouping

- Grouping the relation into groups based on a grouping condition
- Notation

$G_1, G_2, \dots, G_n \overline{\sum} F_1(A_1), F_2(A_2), \dots, F_n(A_n)(E)$

- [?] E relation/relational algebra expression
- [?] G₁, G₂, ..., G_n attributes defining groups/ grouping condition
- [?] F₁, F₂, ..., F_n aggregate functions
- [?] A₁, A₂, ..., A_n attributes

8. Grouping - Example

r	A	B	C
α	2	7	
α	4	7	
β	2	3	
γ	2	10	

 $\mathfrak{I}_{\text{SUM}(C)}(r)$

SUM_C
27

 $\textcolor{red}{A}\mathfrak{I}_{\text{SUM}(C)}(r)$

A	SUM_C
α	14
β	3
γ	10

8. Grouping – Example 19

- Retrieve the number of lecturers and their summary of salary

8. Grouping – Example 19

- Retrieve the number of lecturers and average of salary of each department

8. Grouping – Example 20

□ What is/are the name(s) of the faculty with the most teachers?

- ?] Calculate the number of lecturers according to faculties

$$r1 \leftarrow \pi_{MaKhoa, TenKhoa} \Sigma_{COUNT(HoTen)} (GIAOVIEN * BOMON * KHOA)$$

- ?] Get the highest number of lecturers

$$r2(MAX_COUNT) \leftarrow \Sigma_{MAX(COUNT_HoTen)} (r1)$$

- ?] Faculty with the most lecturers

$$KQ \leftarrow \pi_{MaKhoa, TenKhoa} (r1 \bowtie_{COUNT_HoTen = MAX_COUNT} r2)$$

8. Grouping – Example 21

- What are the subject and the number of projects of that subject?

Quiz

DB schema:

- PHÒNG(Mã Phòng, Diện_tích, Giá_tiền, Loại)
- THUÊ(Mã Phòng, Mã_KH, Ngày_nhận, Số_ngày)
- KHÁCH_HÀNG(Mã KH, Họ_tên, Điện_thoại, Giới_tính)

Retrieve the types of rooms that have never been rented :

- A. $\pi_{\text{Loại}} (\sigma_{\text{COUNT}(\text{Mã Phòng}) = 0} (\text{Loại} \setminus \text{COUNT}(\text{Mã Phòng})) \text{ (THUÊ * PHÒNG) })$
- B. $\pi_{\text{Mã Phòng}, \text{Loại}} (\text{KHÁCH_HÀNG} * \text{THUÊ} * \text{PHÒNG}) \div \pi_{\text{Mã Phòng}} (\text{PHÒNG})$
- C. $\pi_{\text{Mã Phòng}, \text{Loại}} (\text{PHÒNG}) \div \pi_{\text{Mã Phòng}} (\text{PHÒNG})$
- D. $\pi_{\text{Loại}} (\text{KHÁCH_HÀNG} * \text{THUÊ} * \text{PHÒNG}) - \pi_{\text{Loại}} (\text{PHÒNG})$
- E. All wrong

8. Outer Join

- Extend the JOIN operation to avoid information loss
 - ❑ Perform the Join operation
 - ❑ Get more tuples that do not satisfy the join condition
- 3 types:
 - ❑ Left outer join 
 - ❑ Right outer join 
 - ❑ Full outer join 

8. Outer Join – Example 22

□ Retrieve the name of the lecturer and the name of the department he/she is heading, if any

$R1 \leftarrow GIAOVIEN \bowtie_{MAGV=TRUONGBM} BOMON$

$KQ \leftarrow \pi_{HOTEN, TENBM}(R1)$

HOTEN	TENBM
Nguyễn Hoài An	Mạng máy tính
Trần Trà Hương	Hệ thống thông tin
Nguyễn Ngọc Ánh	null
...	...

8. Outer Join – Example 23

- Retrieve the dept. name and dept. head's name if any

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9. Update Operations

□ Date can be updated by:

- ▢ Insertion
- ▢ Deletion
- ▢ Update/Modify

□ Update Operations is expressed through the assignment:

$$r_{\text{new}} \leftarrow \text{relational operations } r_{\text{old}}$$

9. Insertion

□ Denoted by

$$r_{\text{new}} \leftarrow r_{\text{old}} \cup E$$

- ❑ r relation
- ❑ E relational algebra expression

□ Example

- ❑ Assign lecturer code “001” to participate in work 4 of the project “001” with an allowance level 2.

THAMGIAĐT \leftarrow THAMGIAĐT \cup ('001', '001', 4, 2)

9. Deletion

□ Denoted by

$$r_{\text{new}} \leftarrow r_{\text{old}} - E$$

- ❑ r relation
- ❑ E relational algebra expression

□ Example

- ❑ Remove works of the lecturer “001”

$$\text{THAMGIAĐT} \leftarrow \text{THAMGIAĐT} - \sigma_{\text{MAGV}='001'}(\text{THAMGIAĐT})$$

9. Update/ Modify

□ Denoted by

$$r_{\text{new}} \leftarrow \pi_{F_1, F_2, \dots, F_n} (r_{\text{old}})$$

- ❑ r relation
- ❑ F_i expression that calculates the new value for the updating attribute

□ Example

- ❑ Increase the allowance for all works up to 1.5 times

$$\text{THAMGIAĐT} \leftarrow \pi_{\text{MAGV}, \text{MAĐT}, \text{STT}, \text{PHUCAP}^*1.5} (\text{THAMGIAĐT})$$