

Experiment #6

Interrupts in Personal Computers



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Abstract

This experiment aims at understanding and expanding the concept of interruption. Students will learn how to use and call Interrupt Service Routines (ISR) based on MS Windows XP OS. Besides, creating a new ISR using debug and TASM will be a main task of the experiment.

PART I Theoretical Introduction

The 8086 interrupts can be classified into three types. These are

1. Predefined interrupts
2. User-defined software interrupts
3. User-defined hardware interrupts

The interrupt vector address for all the 8086 interrupts are determined from a table stored in locations 00000H through 003FFH. The starting addresses for the service routines for the interrupts are obtained by the 8086 using this table. Four bytes of the table are assigned to each interrupt: two bytes for IP and two bytes for CS. The table may contain up to 256 8-bit vectors. If fewer than 256 interrupts are defined in the system, the user need only provide enough memory for the interrupt pointer table for obtaining the defined interrupts. The interrupt address vector (contents of IP and CS) for all the interrupts of the 8086 assigns every interrupt a type code for identifying the interrupt. There are 256 type codes associated with 256 table entries. Each entry consists of two addresses, one for storing the IP contents and the other for storing the CS contents. Each 8086-interrupt physical address vector is 20 bits wide and is computed from the 16-bit contents of IP and CS. For obtaining an interrupt address vector, the 8086 calculates two addresses in the pointer table where IP and CS are stored for a particular interrupt type. For example, for the interrupt type nn (instruction $INT\ nn$), the table address for $IP=4 \times nn$ and the table address for $CS=4 \times nn + 2$. For servicing the 8086's non-maskable interrupt (NMI pin), the 8086 assigns the type code 2 to this interrupt. The 8086 automatically executes the $INT2$ instruction internally to obtain the interrupt address vector as follows:

Address for IP = $4 \times 2 = 00008H$

Address for CS = $4 \times 2 + 2 = 0000AH$

The 8086 loads the values of IP and CS from the 20-bit physical address 00008H and 0000AH in the pointer table. The user must store the desired 16-bit values of IP and CS in these locations. Similarly, the IP and CS values for other interrupts are calculated. The 8086-interrupt pointer table layout is shown in Figure 1.

Interrupt type code		20-bit Memory Address
	↓	↓
0	IP	00000H
	CS	00002H
1	IP	00004H
	CS	00006H
2		00008H
		0000AH
.		.
.		.
.		.
.		.
.		.
255	IP	003FEH
	CS	00400H

Figure 1 Interrupt Vector Table

Interrupt Vector Table			
<div> <div>080H</div> <div>32-255 User defined</div> <div>14-31 Reserved</div> </div>			
040H	Coprocessor error	16	
03CH	Unassigned	15	
038H	Page fault	14	
034H	General protection	13	
030H	Stack seg overrun	12	
02CH	Segment not present	11	
028H	Invalid task state seg	10	
024H	Coproc seg overrun	9	
020H	Double fault	8	
01CH	Coprocessor not avail	7	
018H	Undefined Opcode	6	
014H	Bound	5	
010H	Overflow (INTO)	4	
00CH	1-byte breakpoint	3	
008H	NMI pin	2	
004H	Single-step	1	
000H	Divide error	0	
4 bytes			

The interrupt vector table is located in the first 1024 bytes of memory at addresses 000000H through 0003FFH.

There are 256 4-byte entries (segment and offset in real mode).

Seg high	Seg low	Offset high	Offset low
Byte 3	Byte 2	Byte 1	Byte 0

Figure 2 Interrupt Vector Table

In response to an interrupt, the 8086 pushes flags, CS, and IP onto the stack, clears TF and IF flags, and then loads IP and CS from the pointer table using the type code. Interrupt service routine should be terminated with the IRET (Interrupt Return) instruction which pops the top three stack words into IP, CS, and flags, thus returning to the right place in the main program. The 256 interrupt type codes are assigned as follows;

- Types 0 to 4 are for the predefined interrupts.
- Types 5 to 31 are reserved by Intel for future use.
- Types 32 to 255 are available for maskable interrupts.

Our focus in this experiment is on **software interrupts**.

PART II Pre-Lab

1. Briefly explain the following terms:
 - ISR [Interrupt Service Routine]
 - IVT [Interrupt Vector Table]
2. What's the address of CS & IP of the ISR of INT 62H?

PART III Practices

3.1 PRACTIC I: Activating the Interrupt Service Routine Manually

You have notices that Interrupt Vector table contains addresses for Interrupt Service Routines (ISPs). These Routines can be activated by executing their code. One way to do so is using “INT n¹” Assembly instruction. These instructions activate the ISP its address loaded in vector number “n”.

Step 1: Open Command Dos Terminal and execute “Debug”.

Step 2: Assemble the following code to offset 100.

```
-INT 0
```

Step 3: Run the program using G (Make sure the IP value is 100).

TASKS:

1. Explain what does “INT 0” do?

3.2 PRACTICE II: Activating the Interrupt Service Routine Automatically

Sometimes, ISP can be called automatically as a response of a specific situation or operation. For example, ISP 0 is activated whenever there is a division over zero.

Step 1: Run the Debug program.

Step2: Assemble the following program

```
MOV AX, 1234  
MOV CL,0  
DIV CL  
INT 3
```

Step 3: Run the program using G.

TASKS:

1. Explain the result (Notice Figure 3)?
2. INT 3 used to set break point and stops the executions. If this instruction was removed, will there be any changes in the result?

¹ n: Interrupt number

```

C:\ Command Prompt
Microsoft Windows XP [Version 5.1.2600]
(C) Copyright 1985-2001 Microsoft Corp.

C:\Documents and Settings\abedl_karim>DEBUG
-A
1362:0100 MOV AX,1234
1362:0103 MOV CL,0
1362:0105 DIV CL
1362:0107 INT 3
1362:0108
-G
Divide overflow
C:\DOCUME~1\ABEDL_~1>

```

Figure 3 Activating ISP Automatically

3.3 PRACTICE III: Activating Interrupt Service Routine Manually (Another Way)

What happened in the first practice was calling the ISP using an assembly instruction. This instruction actually refers to the Interrupt Vector table each time it is executed. It goes there and gets the address for the specified interrupt number, all what it needs is getting the Segment address and the offset.

What we are trying to do in this practice is changing the **CS** and **IP** to the address of the ISP number 0 and executing using G. This supposed to work as it works with “**INT 0**” instruction.

Step 1: Run the debug program.

Step 2: Run the R command and check the content of the registers.

Finding Out the Address for ISP # 0

Step 3: The vector table is located in the address **0:0** on the memory, Display the content of memory at this address (Hint: use D command).

Step 4: The address of interrupt service routine corresponding to INT 0 is located the range 0-3 as shown in table below

3	Segment (high)
2	Segment (low)
1	Offset (high)

0

Offset (low)

Find out the values for ISP # 0 address.

(P.S.: The values of segment and offset below are not necessary the same on your PC)

Step 5: We can run the interrupt service routine by making the CS register points to the code segment that the interrupt service routine located in it and making the IP register points to the start address of the interrupt service routine. (Hint: Use R command)

Step 6: Use G to execute the program start at address **CS:IP**.

TASKS:

1. You should notice that the results of this practice and Practice I are the Same.
2. What is the routine address of INT F stored in the Interrupt Vector Table?

```

C:\DOCUMENTS\ABEDL_~1>DEBUG
-D 0:0
0000:0000 68 10 A7 00 BB 01 70 00-16 00 8D 03 8B 01 70 00 h.....p.....p.
0000:0010 8B 01 70 00 B9 06 0C 02-40 07 0C 02 FF 03 0C 02 ..p.....0.....
0000:0020 46 07 0C 02 0A 04 0C 02-3A 00 8D 03 54 00 8D 03 F.....:..T...
0000:0030 6E 00 8D 03 88 00 8D 03-A2 00 8D 03 FF 03 0C 02 n.....l...
0000:0040 A9 08 0C 02 A4 09 0C 02-AA 09 0C 02 5D 04 0C 02 .....1...
0000:0050 B0 09 0C 02 0D 02 DB 02-C4 09 0C 02 8B 05 0C 02 .....
0000:0060 0E 0C 0C 02 14 0C 0C 02-1F 0C 0C 02 AD 06 0C 02 .....?.....
0000:0070 AD 06 0C 02 A4 F0 00 F0-37 05 0C 02 94 2C 00 C0 .....
-R CS
CS 1362
:00A7
-R IP
IP 0100
:1068
-g
Divide overflow
C:\DOCUMENTS\ABEDL_~1>_

```

Figure 4 Activating the ISP Manually (Another Way)

3.4 PRACTICE IV: Creating Your Own ISP

Now, to change the INT response. What about replacing the code of ISP 0? or maybe creating a new code and replace the vector values with its address? We can write our routine and make it run when dived by zero occur by putting the address (segment and offset) of our new routine in memory location 0-3.

Out task: is to rewrite the interrupt routine for INT 0 to display AAAA when a divide by zero occur.

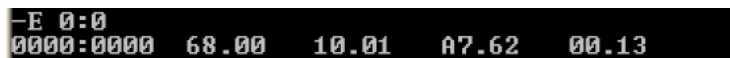
Step 1: Run the debug program.

Step 2: Assemble the following program to offset 100

```
MOV CX,4
L1: MOV DL,41
MOV AH,6
INT 21
LOOP L1
INT 3
```

Notice that the first instruction should be at address CS:0100.

Step 3: Change the address of interrupt service routine corresponding to INT 0 To be CS:0100 (Hint: Use E command) (Notice Figure 5).



```
-E 0:0
0000:0000  68.00  10.01  A7.62  00.13
```

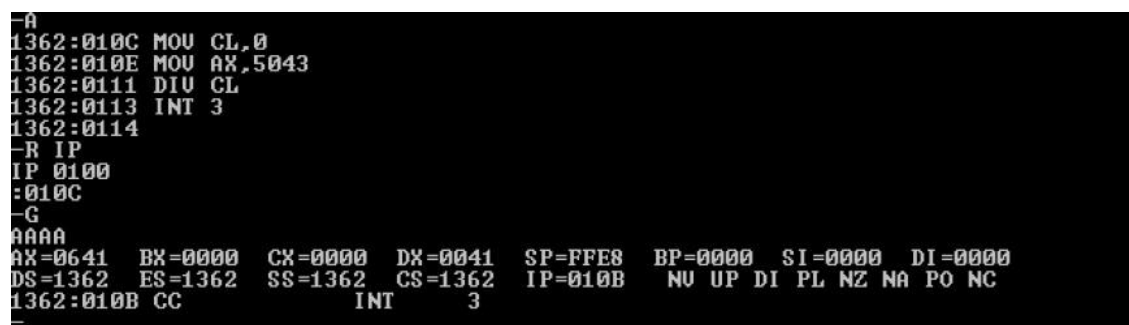
Figure 5 Changing the Values of INT 0 Vector Entry

Step 4: Now we can test the new interrupt routine by making a program that makes a division by zero (e.g. Practice II)

Code Example:

```
MOV CL,0
MOV AX,5043
DIV CL
INT 3
```

Notice that when we execute the program AAAA is displayed instead of divide overflow (Notice Figure 6).



```
-A
1362:010C MOV CL,0
1362:010E MOV AX,5043
1362:0111 DIV CL
1362:0113 INT 3
1362:0114
-R IP
IP 0100
:010C
-G
AAAA
AX=0641 BX=0000 CX=0000 DX=0041 SP=FFES BP=0000 SI=0000 DI=0000
DS=1362 ES=1362 SS=1362 CS=1362 IP=010B  NU UP DI PL NZ NA PO NC
1362:010B CC          INT      3
```

Figure 6 Test the new INT 0

3.4 PRACTICE V: Creating Your Own ISP Using TASM

In this practice we would work on installing a new interrupt service routine for interrupt number 62.

Setp1: Write down the following code and save it to an Assembly file.

```
.MODEL TINY
.CODE
.STACK 100H

JMP INSTALL
MESS DB "WELCOME TO COOLEST LAB ON EARTH ENCS411 :D $"

MYINT PROC

    MOV BX,OFFSET MESS
    MOV DL,44 ;column 20 (# of mess chars)

L1:
    MOV AL,CS:[BX]
    CMP al,'$'
    JZ L2

    PUSH BX

    MOV BH,0 ;page number 0
    MOV AH,2 ;set the curser position
    MOV DH,20 ;row 0
    INC DL
    INT 10H ;to access video memory

    MOV AH,9H ;Write character and attribute at cursor position

    MOV BH,0 ;page number 0
    MOV BL,84H ;color of the character
    MOV CX,1 ;number of times to write the character
    INT 10H ;video memory access

    POP BX
    INC BX
    JMP L1

L2:
    IRET
MYINT ENDP

INSTALL:
;install new interrupt vector
    MOV AH,25H ;set interrupt vector
    MOV AL,62H ;desired interrupt vector
    MOV DX,CS
    MOV DS,DX
    MOV DX,OFFSET MYINT
    INT 21H

; Terminate & stay resident

    MOV AX,3100H
    INT 21H

END
```

Step2: Compile and build this ASM file and execute it on MDA-8086 kit. (How? Review Exp#1 Intro. To MDA Kit)

Step 3: To test the code

- Open debug.
- Execute the instruction "INT 62". What is the result?

Another way to test the code:

- Write the following code to an assembly file

```
.MODEL SMALL
.STACK 100
.CODE
    INT 62H
    mov AH,4CH
    INT 21H
END
```

- Compile and run it. What is the result?

TASKS:

1. Explain what does this code do (Line by line)?
2. What does INT 21 do?