

**AIM :**

Implement DFS algorithm use on undirected graph and develop a recursive algorithm for searching all the vertices of graph OR tree data structure

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**Introduction****Depth First Search**

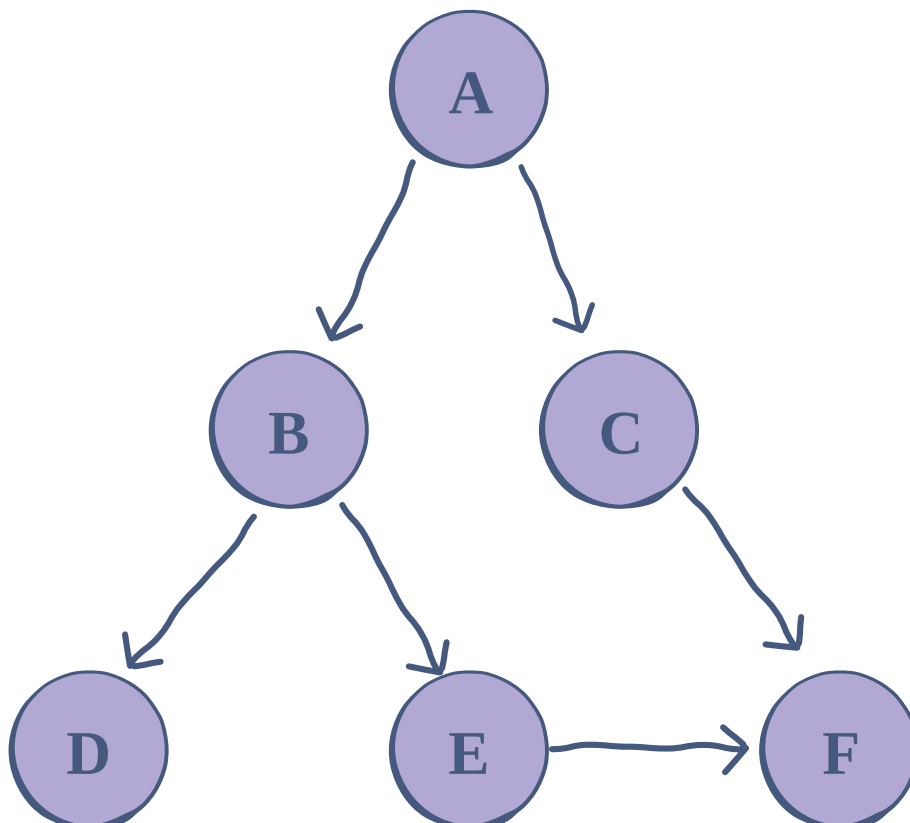
**Depth-First-Search(DFS)** is an algorithms for tree traversal on graph or tree data structure it can be implemented easily recursion and data structures like dictionaries and sets

**The Algorithm**

1. pick any node . if it is unvisited , mark it as visited and recur on all its adjacent nodes
2. Repeat until all the nodes are visited , or the node to be searched is found

**Implementation**

Consider this graph implemented in the code below:



In [3]:

```
#using a python dictionary to act as an adjacency list
graph = {
    'A': ['B', 'C'],
    'B': ['D', 'E'],
    'C': ['F'],
    'D': [],
    'E': ['F'],
    'F': [],
}

visited = set() #Set to keep track of visited nodes

def dfs(visited, graph, node):
    if node not in visited:
        print (node)
        visited.add(node)
        for neighbour in graph[node]:
            dfs(visited, graph, neighbour)

#driver code
dfs(visited, graph, 'A')
```

A  
B  
D  
E  
F  
C

### Explanation

- **lines 2-9:** The illustrated graph is represented using an **adjacency list**- an easy way to do it in python is to use a dictionary data structure. Each vertex has a list of its adjacent nodes stored
- **line 11:** `visited` is a set that is used to keep track of visited nodes.
- **line 21:** The `dfs` function is called and is passed the `visited` set, the `graph` in the form of dictionary and `A` which is the starting node
- **Lines 13-18:** `dfs` follows the algorithm describe above
  - it first checks if the current node is unvisited - if yes, it is appended in the `visited` set
  - Then for each neighbor of the current node, the `dfs` function is invoked again
  - The base case is invoked when all the nodes are visited the function then returns

### Time Complexity

since all the nodes and vertices are visited the average time complexity for DFS on graph is  $O(V + E)$  where  $V$  is the number of vertices and  $E$  is the number of edges. In case of DFS on tree the time complexity is  $O(V)$ , where  $V$  is the number of nodes

average time complexity because a set's operation has an average time complexity of  $O(1)$ . If we used a list, the complexity would be higher.

In [ ]: