

Merge Sort Implementation - Explanation

Overview

This is an **efficient recursive merge sort** algorithm that sorts an array in $O(n \log n)$ time.

Function Breakdown

1. `merge(array, start, middle, end, temp)`

Merges two sorted halves of an array into one sorted section.

```
def merge(array, start, middle, end, temp):  
    left, right, pos = start, middle + 1, start
```

- `left`: pointer for the left half (starts at `start`, goes to `middle`)
- `right`: pointer for the right half (starts at `middle + 1`, goes to `end`)
- `pos`: position in temp array where we'll place the next element

```
while left <= middle and right <= end:  
    if array[left] <= array[right]:  
        temp[pos] = array[left]  
        left += 1  
    else:  
        temp[pos] = array[right]  
        right += 1  
    pos += 1
```

- Compare elements from both halves
- Put the smaller element into `temp`
- Move the appropriate pointer forward
- This continues until one half is exhausted

```
while left <= middle:  
    temp[pos] = array[left]  
    left, pos = left + 1, pos + 1
```

- Copy any remaining elements from the left half

```
while right <= end:  
    temp[pos] = array[right]
```

```
right, pos = right + 1, pos + 1
```

- Copy any remaining elements from the right half

```
array[start:end + 1] = temp[start:end + 1]
```

- Copy the sorted elements from `temp` back to the original `array`
-

2. `merge_sort_recursive(array, start, end, temp)`

The recursive function that divides the array and calls merge.

```
def merge_sort_recursive(array, start, end, temp):  
    if start < end:
```

- **Base case:** if `start >= end`, there's only 1 element (or none), so it's already sorted

```
        middle = (start + end) // 2
```

- Find the middle point to divide the array in half

```
        merge_sort_recursive(array, start, middle, temp)  
        merge_sort_recursive(array, middle + 1, end, temp)
```

- **Recursively sort** the left half (from `start` to `middle`)
- **Recursively sort** the right half (from `middle + 1` to `end`)

```
        merge(array, start, middle, end, temp)
```

- **Merge** the two sorted halves together
-

3. `merge_sort(array)`

Main entry point that sets up the sort.

```
def merge_sort(array):  
    if len(array) <= 1:  
        return array
```

- If array has 0 or 1 element, it's already sorted

```
temp = [0] * len(array)
```

- Create a temporary array of the same size (allocated once for efficiency)

```
merge_sort_recursive(array, 0, len(array) - 1, temp)  
return array
```

- Call the recursive function starting with the full array (indices 0 to n-1)
- Return the now-sorted array

4. `main()`

Handles user input and output.

```
def main():  
    nums = list(map(int, input("Enter Input : ").split()))
```

- Reads a line of space-separated numbers
- Converts them to integers
- Stores in a list

```
print(merge_sort(nums))
```

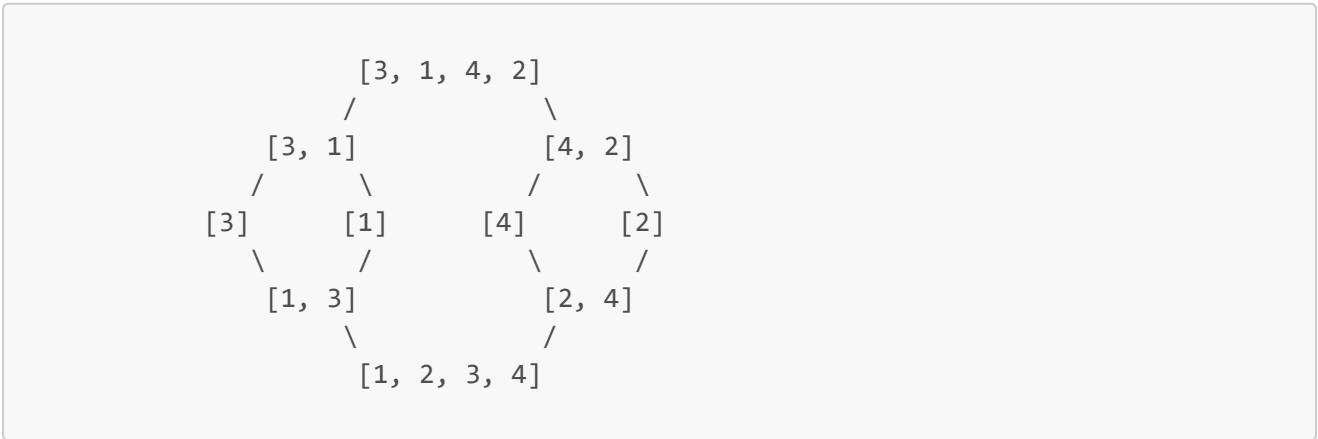
- Sorts the list and prints the result

How It Works (Example)

Input: `[3, 1, 4, 2]`

```
Step 1: Split into [3, 1] and [4, 2]  
Step 2: Split [3, 1] into [3] and [1]  
Step 3: Merge [3] and [1] → [1, 3]  
Step 4: Split [4, 2] into [4] and [2]  
Step 5: Merge [4] and [2] → [2, 4]  
Step 6: Merge [1, 3] and [2, 4] → [1, 2, 3, 4]
```

Visual Representation



Complexity Analysis

Aspect	Complexity	Explanation
Time (Best)	$O(n \log n)$	Always divides array in half
Time (Average)	$O(n \log n)$	Consistent performance
Time (Worst)	$O(n \log n)$	Even on reversed arrays
Space	$O(n)$	One temporary array
Stable	Yes	Equal elements maintain order
In-place	No	Requires extra space

Key Optimizations in This Implementation

- Single temp array:** Allocated once instead of creating many arrays during recursion
- Index-based operations:** Uses indices instead of array slicing to avoid copying overhead
- In-place updates:** Sorts the array in place after using temp for merging
- Efficient memory:** $O(n)$ space instead of $O(n \log n)$

When to Use Merge Sort

☑ Good for:

- Large datasets
- When you need stable sorting
- Linked lists (works great with $O(1)$ space)
- When consistent $O(n \log n)$ performance is required

✗ Not ideal for:

- Small arrays (insertion sort is faster)
 - Memory-constrained systems (requires $O(n)$ extra space)
 - When in-place sorting is required (use quicksort or heapsort)
-

Code

```
# ex01.py for Q3(Sort,Search)

def merge(array, start, middle, end, temp):
    left, right, pos = start, middle + 1, start
    while left <= middle and right <= end:
        if array[left] <= array[right]:
            temp[pos] = array[left]
            left += 1
        else:
            temp[pos] = array[right]
            right += 1
        pos += 1

    while left <= middle:
        temp[pos] = array[left]
        left, pos = left + 1, pos + 1

    while right <= end:
        temp[pos] = array[right]
        right, pos = right + 1, pos + 1

    array[start:end + 1] = temp[start:end + 1]

def merge_sort_recursive(array, start, end, temp):
    if start < end:
        middle = (start + end) // 2
        merge_sort_recursive(array, start, middle, temp)
        merge_sort_recursive(array, middle + 1, end, temp)
        merge(array, start, middle, end, temp)

def merge_sort(array):
    if len(array) <= 1:
        return array
    temp = [0] * len(array)
    merge_sort_recursive(array, 0, len(array) - 1, temp)
    return array

def main():
    nums = list(map(int, input("Enter Input : ").split()))
    print(merge_sort(nums))

if __name__ == "__main__":
    main()
```

