

CS 291  
Homework 7

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Jingbo Wang  
jw6347

**Section 11.1, Exercise 2.b** Find a regular expression to describe each of the following languages.  
 $\{aa, ab, ac\}$

**Answer:**

$$\begin{aligned}\{aa, ab, ac\} &= \{a\}\{a, b, c\} \\ &= a(a + b + c)\end{aligned}$$

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**Section 11.1, Exercise 2.d** Find a regular expression to describe each of the following languages.  
 $\{a, aaa, aaaaa, \dots, a^{2n+1}, \dots\}$

**Answer:**

$$\begin{aligned}\{a, aaa, aaaaa, \dots, a^{2n+1}, \dots\} &= \{a\}\{\Lambda, aa, aaaa, \dots, a^{2n}, \dots\} \\ &= \{a\}\{\Lambda, aa, (aa)^2, \dots, (aa)^n, \dots\} \\ &= \{a\}\{aa\}^* \\ &= a(aa)^*\end{aligned}$$

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**Section 11.1, Exercise 2.e** Find a regular expression to describe each of the following languages.  
 $\{\Lambda, a, abb, abbbb, \dots, ab^{2n}, \dots\}$

**Answer:**

$$\begin{aligned}\{\Lambda, a, abb, abbbb, \dots, ab^{2n}, \dots\} &= \{\Lambda\} + \{a\}\{\Lambda, bb, \dots, b^{2n}, \dots\} \\ &= \{\Lambda\} + \{a\}\{\Lambda, bb, \dots, (bb)^n, \dots\} \\ &= \{\Lambda\} + \{a\}\{bb\}^* \\ &= \Lambda + a(bb)^*\end{aligned}$$

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**Section 11.1, Exercise 4.b** Find a regular expression for each of the following languages over the alphabet  $\{a, b\}$ .

Strings whose length is a multiple of 3.

**Answer:**

Strings whose length is a multiple of 1 over  $\{a, b\}$  is  $(a + b)$ .

Strings whose length is a multiple of 3 over  $\{a, b\}$  is  $(a + b)(a + b)(a + b) = (a + b)^3$ .

Therefore, the regular expression is  $((a+b)^3)^*$ .

*technically not part of regular expression*

**Section 11.1, Exercise 4.c** Find a regular expression for each of the following languages over the alphabet  $\{a, b\}$ .

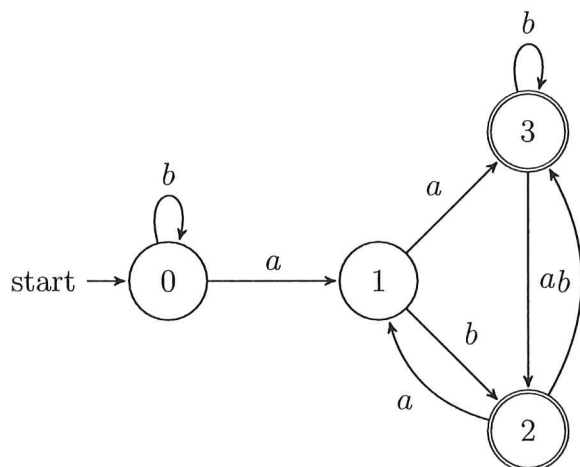
Strings containing the substring  $aba$ .

**Answer:**

$(a + b)^*$  describe all the string of one or more a's and one or more b's, and each string must have substring  $aba$ .

Therefore, the regular expression is  $(a + b)^*aba(a + b)^*$ .

**Section 11.2, Exercise 1** Write down the transition function for the following DFA.



**Answer:**

$$T(0, b) = 0,$$

$$T(0, a) = T(2, a) = 1,$$

$$T(1, a) = T(3, a) = 2,$$

$$T(1, b) = T(3, b) = T(2, b) = 3,$$

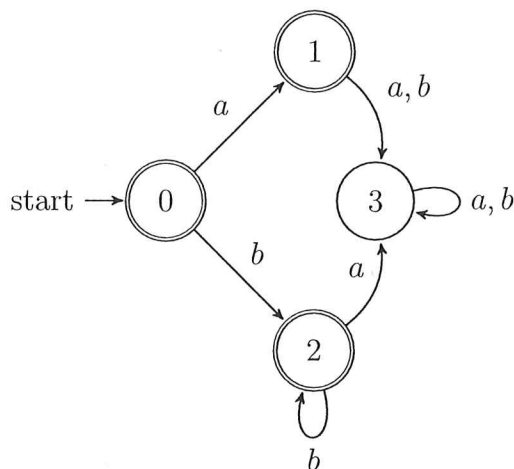
where 0 is the start and both 2 and 3 are final states.

	$T$	$a$	$b$
Start	0	1	0
	1	3	2
	2	1	3
Final	3	2	3

Section 11.2, Exercise 2.c Use your wits to construct a DFA for each of the following regular expressions.

$a + b^*$

Answer:

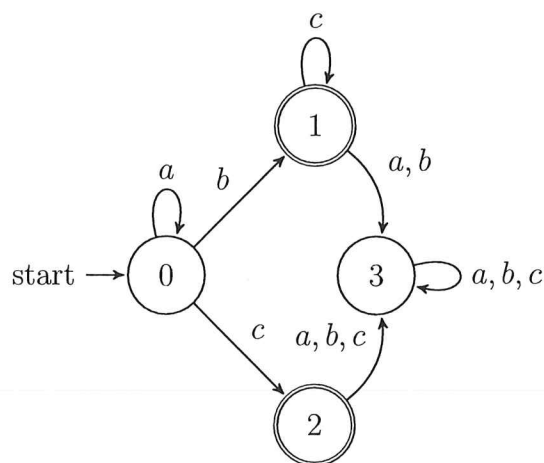


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Section 11.2, Exercise 2.f Use your wits to construct a DFA for each of the following regular expressions.

$a^*bc^* + ac$

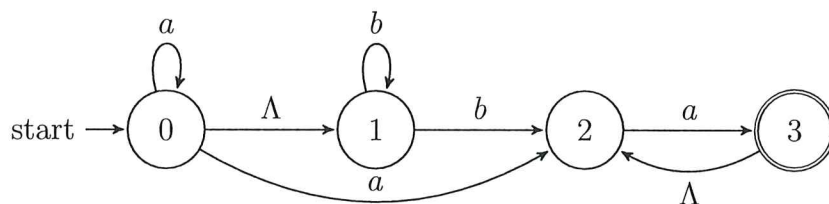
Answer:



close  
accepts aac  
which is not  
in the language

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Section 11.2, Exercise 4 Write down the transition function for the following NFA:



Answer:

$T(0, a) = 0,$   
 $T(0, \Lambda) = T(1, b) = 1,$   
 $T(1, b) = T(0, a) = T(3, \Lambda) = 2,$   
 $T(2, a) = 3,$   
 where 0 is the start and 3 is final state.

	$T$	$a$	$b$	$\Lambda$
Start	0	0	$\emptyset$	1
	1	$\emptyset$	2	$\emptyset$
	2	3	$\emptyset$	$\emptyset$
Final	3	$\emptyset$	$\emptyset$	2

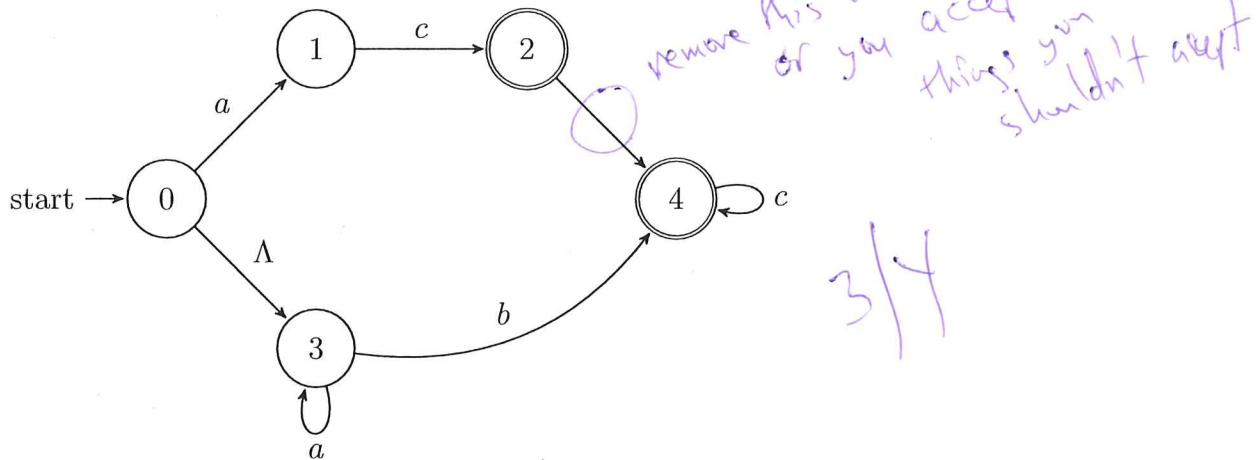
$\{0, 2\}$   
 $\{1, 2\}$   
 $\{2, 3\}$

Entries are given  
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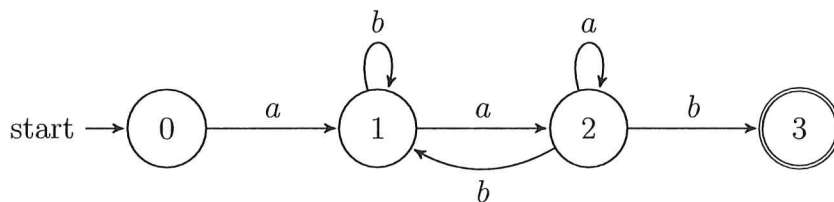
**Section 11.2, Exercise 5,a** Use your wits to construct an NFA for each of the following regular expressions.

$a^*bc^* + ac$

**Answer:**



**Section 11.2, Exercise 8,a** Given the following NFA:



Use algorithm (11.2.4) to find two regular expressions for the language accepted by the NFA as follows.

Delete state 1 before deleting state 2.

**Answer:**

$$\begin{aligned}
new(0, 2) &= old(0, 2) + old(0, 1)old(1, 1)^*old(1, 2) \\
&= \emptyset + ab^*a \\
&= ab^*a, \\
new(2, 2) &= old(2, 2) + old(2, 1)old(1, 1)^*old(1, 2) \\
&= a + bb^*a, \\
new(0, 3) &= old(0, 3) + old(0, 2)old(2, 2)^*old(2, 3) \\
&= \emptyset + ab^*a(a + bb^*a)^*b \\
&= ab^*a(a + bb^*a)^*b.
\end{aligned}$$

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The final DFA is  $ab^*a(a+bb^*a)^*b$

**Section 11.2, Exercise 8.b** Given the following NFA:

Delete state 2 before deleting state 1.

**Answer:**

$$\begin{aligned}
new(1, 3) &= old(1, 3) + old(1, 2)old(2, 2)^*old(2, 3) \\
&= \emptyset + aa^*b \\
&= aa^*b, \\
new(1, 1) &= old(1, 1) + old(1, 2)old(2, 2)^*old(2, 1) \\
&= b + aa^*b, \\
new(0, 3) &= old(0, 3) + old(0, 1)old(1, 1)^*old(1, 3) \\
&= \emptyset + a(b + aa^*b)^*aa^*b \\
&= a(b + aa^*b)^*aa^*b.
\end{aligned}$$

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The final DFA is  $a(b+aa^*b)aa^*b$ .

