

# AVL Tree

Class 22

# BST Implementation

```
class Tree_Node
{
    Object data;
    Tree_node* left_child;
    Tree_node* right_child;
};
```

- the data has a key of a comparable type (i.e.,  $<$  is defined on data.key)
- the data in a node is greater than any value in its left subtree
- the data in a node is less than any value in its right subtree
- simplifying assumption: there are no duplicate values in the tree

## Find

```
1  bool find(const Comparable& key, Tree_Node* root) const
2  {
3      if (root == nullptr)
4      {
5          return false;
6      }
7      if (key < root->data.key)
8      {
9          return find(key, root->left);
10     }
11     if (key > root->data.key)
12     {
13         return find(key, root->right);
14     }
15     return true;
16 }
```

# AVL Tree

- in 1962 Adelson-Velskii and Landis proposed a height-balanced tree
- use a BST with all its requirements
- additionally, insist that the height of any node's left and right subtrees can differ by at most 1
- more formally
  1. an empty BST is height-balanced
  2. a non-empty BST with root  $T$  and children  $T_L$  and  $T_R$  is height-balanced iff
    - 2.1  $|\text{height}(T_L) - \text{height}(T_R)| < 2$ , and
    - 2.2  $T_L$  and  $T_R$  are height-balanced recursively

# AVL Tree

- height-balancing guarantees that for *find*,  $T(n) \in O(\lg n)$
- but how to accomplish height-balancing?
- insert and delete must be modified to maintain the balance
- each tree node must have a field for balance factor

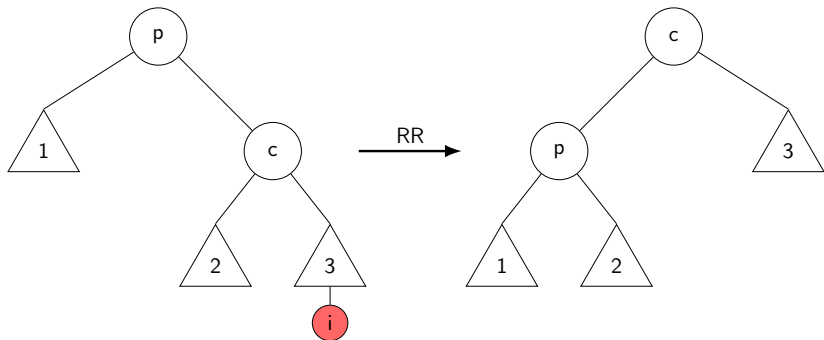
```
class AVL_node
{
    Object data;
    int balance;
    AVL_node* left_child;
    AVL_node* right_child;
};
```

# AVL Implementation

- maintain **height-balanced** BST
- every insert or remove can **potentially** unbalance the tree
- an unbalanced tree requires a **rotation** to rebalance
- there are 4 types of rotations
  1. RR
  2. LL
  3. LR
  4. RL

example

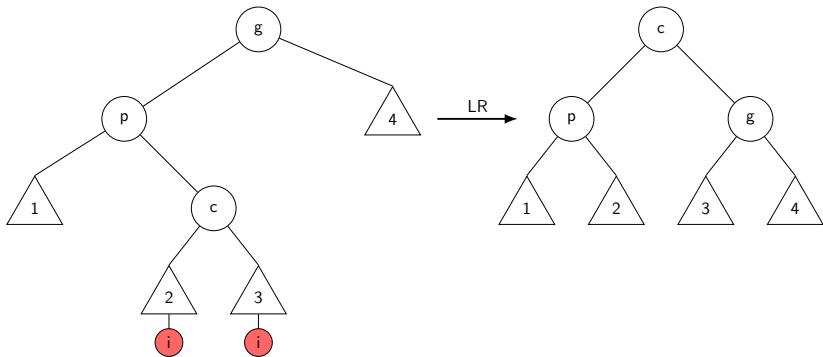
RR



- *p* is the lowest unbalanced node
- *c* is *p*'s right child
- tree 3 contains the insert that caused the tree to become unbalanced, depicted by *i*
- LL is the mirror image of this



## LR



- g is the lowest unbalanced node
- p is g's left child
- c is p's right child
- the insert that caused the tree to become unbalanced is either in 2 or 3, or is c itself