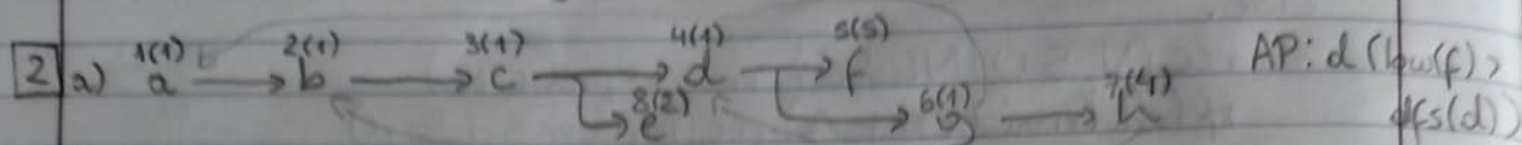
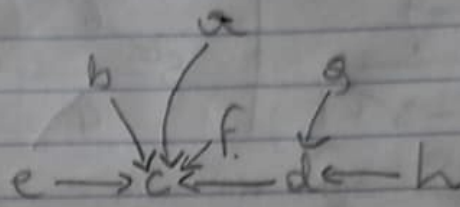
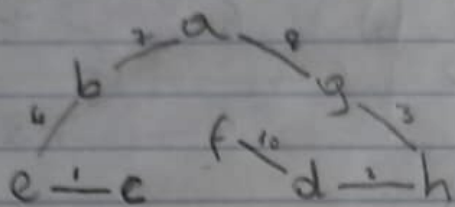


1) $a=1, b=2, c=0 \quad \log_2 1 = 0 \rightarrow O(\log n)$



b)



3) a) Func $\text{path}(G, \text{vis}, \text{node}, \text{target}, \text{len})$:
 if $\text{node} = \text{target}$: $\text{max_len} = \max(\text{max_len}, \text{len})$
 else:
 for node, v in G :
 $\text{vis}[v] = \text{True}$
 $\text{path}(G, \text{vis}, v, \text{target}, \text{len} + 1)$
 $\text{vis}[v] = \text{False}$

b) É mostrar que, p.e., o caminho mais longo de s a d é s, a, b, c, t, d , que não está incluído em S .

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4) a) Func $\text{TD}(p, i, j)$:
 if $\text{dp}[i, j]$ has a value: return $\text{dp}[i, j]$
 if $i = j$: $\text{dp}[i, j] = 0$
 else:
 $\text{dp}[i, j] = +\infty$
 for $k = i \rightarrow j - 1$:
 $\text{dp}[i, j] = \min(\text{dp}[i, j], \text{TD}(p, i, k) + \text{TD}(p, k + 1, j) + p[i, k, j])$
 return $\text{dp}[i, j]$

b) Func $\text{BU}(p, a, b)$:
 for $k = 1 \rightarrow n$: $\text{dp}[k, k] = 0$
 for $j = 2 \rightarrow n$: for $i = j - 1 \rightarrow 1$:
 $\text{dp}[i, j] = +\infty$
 for $k = i \rightarrow j - 1$:
 $\text{dp}[i, j] = \dots$
 return $\text{dp}[a, b]$