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Problem 03: Discuss time complexity of insertion, deletion, and search in Balanced binary search tree (BST). add simulation steps for every operation.

Solution:

Time Complexity between BST insertion, deletion and search is given below:

Operation	Best case	Average case	Worst case
Insertion	$O(1)$	$O(\log_2 n)$	$O(\log_2 n)$
Deletion	$O(1)$	$O(\log_2 n)$	$O(\log_2 n)$
Search	$O(1)$	$O(\log_2 n)$	$O(\log_2 n)$

Simulation steps for every operations

Insertion:

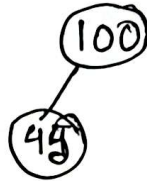
Step 1: After insert root "100" the tree will be:



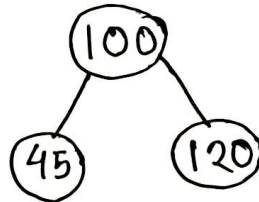
Step 2: It will have two child element, Right and Left. ~~The~~ if the element is greater than the Root then the element will be insert in right; Otherwise the element will be insert in left.

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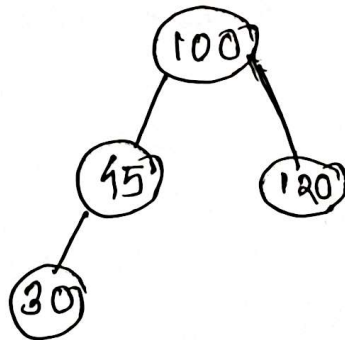
Now, Insert "45", it is smaller than "100", so, It will be inserted in left side of root node. ( $100 > 45$ ).



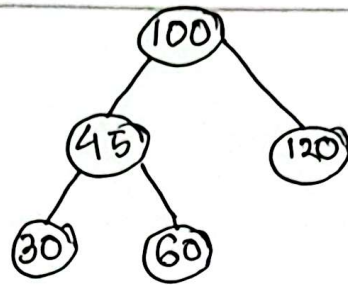
Step 3 Insert "120", Since "120" is greater than "100". So, It will be inserted in right node of root.



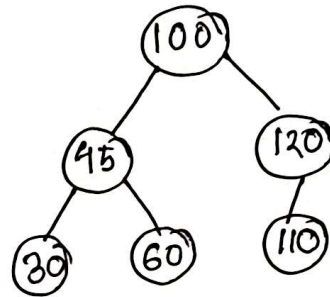
Step 4 Insert "30". "30" will be inserted in the left sub tree, as "100" has already a left node. So, "30" will compare with "45", as "45" is greater than "30" [ $45 > 30$ ], "30" will be "45"'s left node.



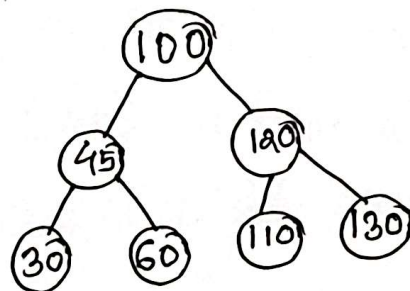
Step 5 Insert "60". "60" is smaller than "100". ( $100 > 60$ ). It will be inserted in "100" left sub-tree. as "45" is smaller than "60". So, "60" will be inserted in the right of "45".



Step 6: Insert "110"; "110" is greater than "100" ( $110 > 100$ ). So, "110" will be insert in the right sub tree. and "110" also ~~greater~~ smaller than "120". ( $120 < 110$ ). So, "110" will be "120"'s left node.



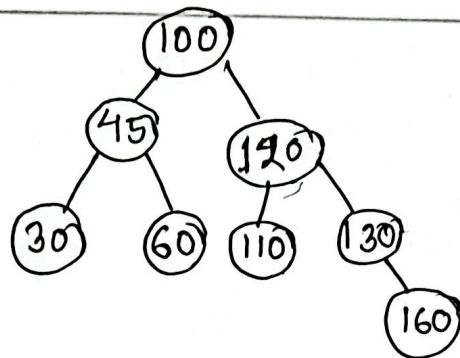
Step 7: Insert "130". "130" is greater than "100" ( $100 < 130$ ), and "130" is also greater than "120" ( $130 > 120$ ). So, "130" will insert in right sub-tree of "100" and "120".



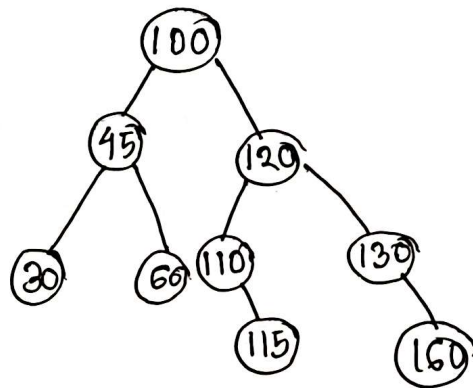
Step 8: Insert "160". "160" is greater than "100", It is also greater than "120" and also greater than "130", [ $100 < 120 < 130 < 160$ ]; So, Here "160" will be "100" and "120"'s right sub-tree and "130"'s right child.



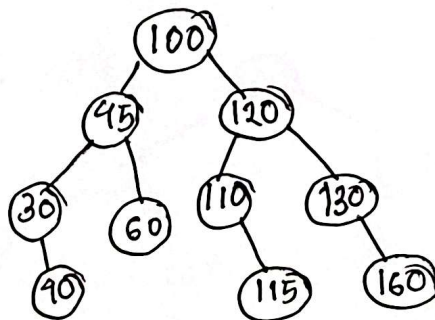
(9)



Step 9: 09% Insert '125', '125' will be insert is the same way. it will be insert is '110's right child.  $[100 < 120 > 110 < 130]$

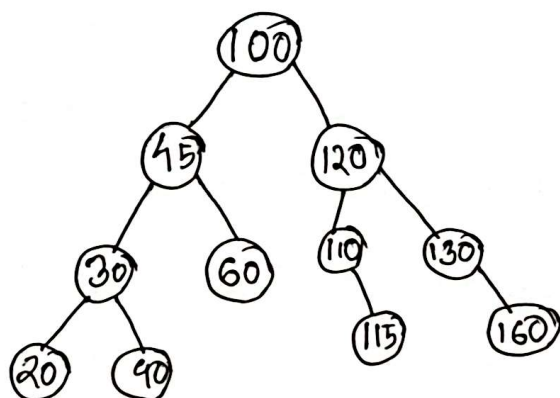


Step 10: 10% Insert "40"; "40" is smaller than "100" ( $100 > 40$ ) also smaller than "45" and "30". so, it will be insert in left subtree of "100", "45" and and greater than "30" ( $30 < 40$ ). so, "40" will be "100", "45" left subtree but "30"'s right node.  $[100 > 45 > 30 < 40]$

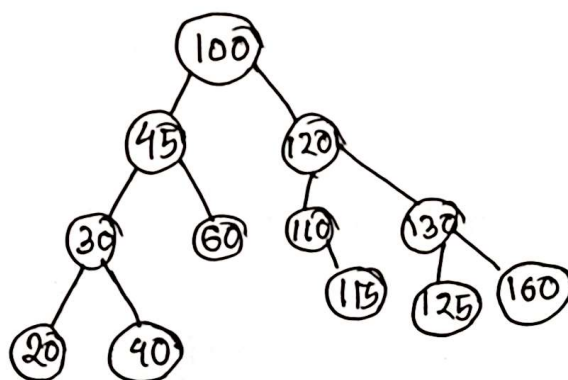


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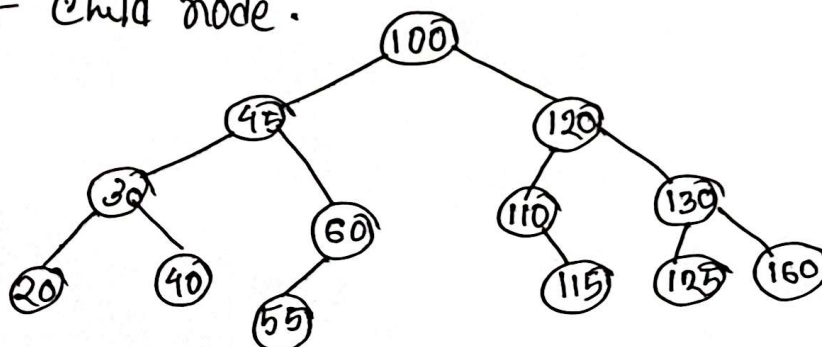
Step 10: Insert "20"; "20" is smaller than "100", "45", and "30". It will be inserted in "100"'s left subtree, "45"'s left subtree, and "30"'s left child. [ $100 > 45 > 30 > 20$ ]



Step 11: Insert "125"; ( $100 < 120 < 130 > 125$ ), so, "125" will be "100"'s right subtree and "130"'s left child.



Step 12: Insert "55", ( $100 > 45 < 60 > 55$ ). So, "55" will be "100"'s left subtree, "45"'s right subtree, and "60" left child node.



Deletion:

Step: 01: ~~Delete~~ Delete "115"; "115" is compare with the root "100". Though  $115 > 100$ , so, "115" go to "100"'s Right sub-tree.

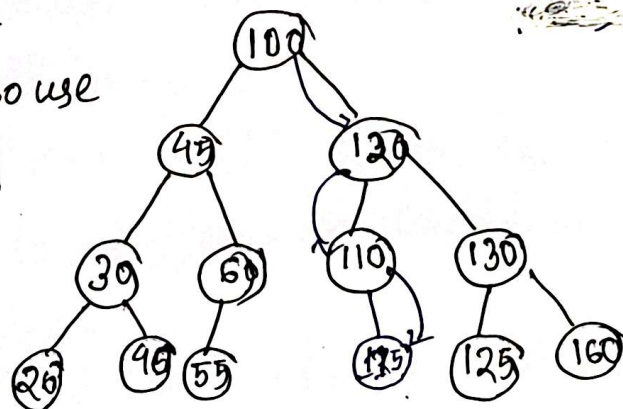
Step: 02: "100"'s Right child node is "120". Here "115" is compare with "120", Though  $115 < 120$ , so, It will go to "120" left child node to compare.

Step: 03: "120"'s left child node is "110". Here "115" is compare with "110", Though  $110 < 115$ , so, It will go to "110"'s ~~left child node~~ right node to compare.

Step: 04: "110"'s right node is "115", it will be compared with the "115", ( $115 = 115$ ). So, 115 will be delete, the "110"'s right node's pointer will be "NULL".

The tree

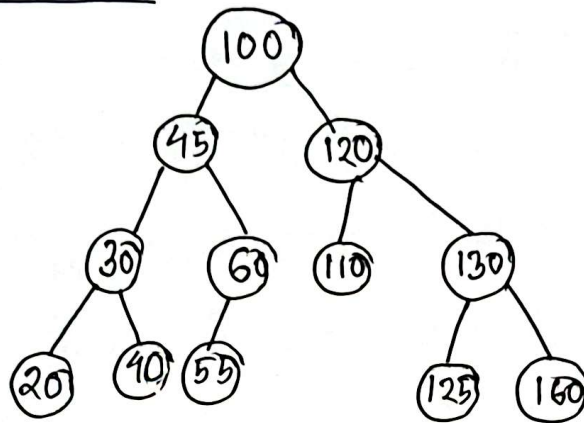
[Here we also use Searching]



[100 < 115]  
[120 > 115]  
[115 > 100]



### The new tree



Before Deletion, the tree in Pre-Order traversal:  
100, 45, 30, 20, 40, 60, 55, 120, 110, 130, 125, 160

After Deletion, the tree is Pre-Order traversal:  
100, 45, 30, 20, 40, 60, 55, 120, 110, 130, 125, 160

### Search:

Process: If we want to search any element,

Step 1: the search element will compare with the root. if search element and ~~key~~ target matched then we will return root.

Step 2: Otherwise, we will search the element in left or right based on two conditions.

The conditions:

(i) if  $(\text{root} \rightarrow \text{target} < \text{key})$ , then search in left;

(ii) if  $(\text{root} \rightarrow \text{target} > \text{key})$ , then search in right;

### Simulation :

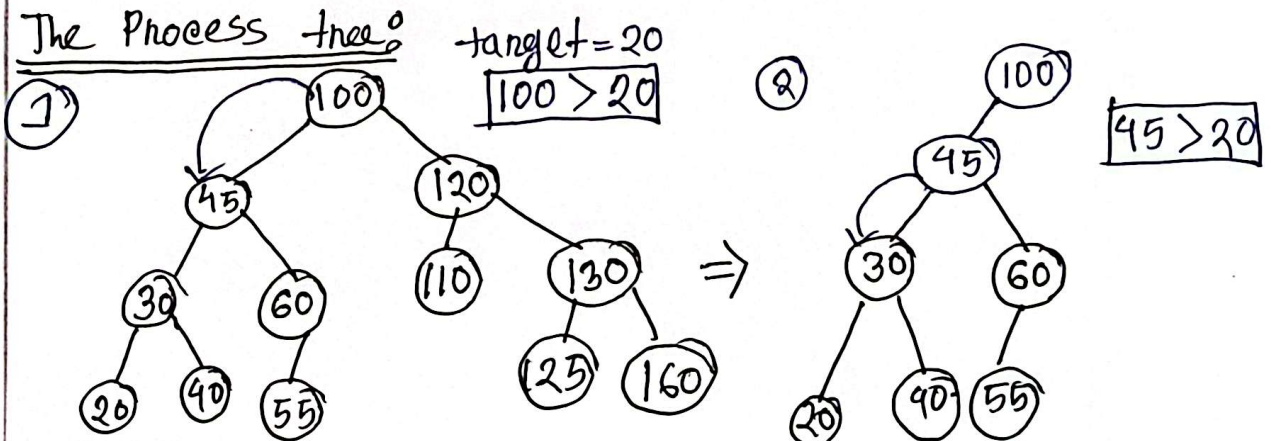
Step 01 : target = "20"; Here, 20 is compare with 100, Though  $(100 > 20)$ , So, the search will start in 100's left sub-tree. and go to "100"'s left node to compare.

Step 02 : "100" left child node is "45". 45 will be compare with "20". Though  $(20 < 45)$ . So, the search will start in 45's left sub-tree. And go to "45"'s left node to compare.

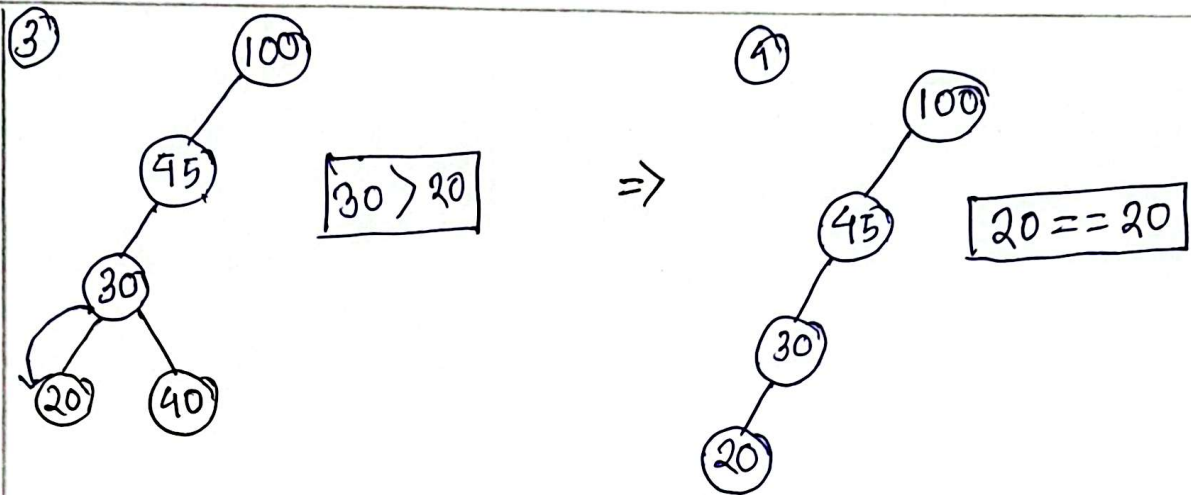
Step 03 : "45" left child node is "30", so "30" will be compare with "20". Though  $(20 < 30)$ ; the search will start in left sub-tree. And go to "30"'s left child node to compare.

Step 04 : "30"'s left child is "20", so, "20" will compare with target = "20", though  $(20 = 20)$ . So, then it will return the index of node "20".

### The Process tree :







Here in this, Searching function we only search of only half of the sub-tree, so, other side we don't check or search.

This is how the BST (Binary Search tree)'s insertion, deletion and search worked.