



Today's agenda

↳ Reverse an array

↳ Reverse a given Part of array.

↳ Rotate array by K.

↳ greater than itself.

↳ Two Sum.



AlgoPrep



Q) Reverse array

In Given array of length n , Reverse the whole array.

ex: arr[5]: {^{0 1 2 3 4}
10 20 30 40 50}
 \downarrow
50 40 30 20 10

arr[8]: {^{0 1 2 3 4 5 6 7}
10 20 30 40 50 60 70 80}
 \downarrow
80 70 60 \downarrow 50 40 30 20 10

arr[8]: {^{0 1 2 3 4 5 6 7}
10 20 30 40 50 60 70 80}
 $\downarrow \uparrow \downarrow \uparrow \downarrow \uparrow \downarrow \uparrow$
80 70 20 50 40 20 20 10

Swap(0, 7)

Swap(1, 6)

Swap(2, 5)

Swap(3, 4)



arr[9]: {⁰ 10 ¹ 20 ² 30 ³ 40 ⁴ 50 ⁵ 60 ⁶ 70 ⁷ 80 ⁸ 90}
90 80 70 60 50 40 30 20 10



Swap(0, 8)
↓↓
↓↓

Swap(1, 7)

```
int main() {
    // Input - arr?
    reverse(arr);
    for (int i=0; i<arr.length; i++) {
        System.out.print(arr[i]);
    }
}
```

Public static void reverse (int [] arr){

arr.length

int sp = 0;

int ep = arr.length - 1;

T.C: O(n)

while (sp < ep) {

int temp = arr[sp];

arr[sp] = arr[ep];

arr[ep] = temp;

sp++;

ep--;



SP = 42
28 4

CP = 75
34 8

→ while () {

int temp = arr[SP];
arr[SP] = arr[CP];
arr[CP] = temp;
SP++;
CP--;

arr[8]: {
1 2 3 4 5 6 7
8 7 6 5 4 3 2 1}

Swap (4, 7)

Swap (5, 6)

Swap (2, 5)

Swap (3, 4)



AlgoPrep



Q) Reverse a Part of array

Given N array element and $[s, e]$, reverse the array from $[s, e]$ index.

$[3, 7]$

$\text{arr}[10] = \{ -3 \ 4 \ 2 \ 8 \ 3 \ 9 \ 6 \ 2 \ 8 \ 10 \}$

$\{ -3 \ 4 \ 2 \ [2 \ 6 \ 9 \ 3 \ 8] 8 \ 10 \}$

$\text{swap}(3, 7)$
 $\text{swap}(4, 6)$

Public static void reverse (int[] arr, int s, int e){

Big O notation
worst case scenario

int SP = s;
int EP = e;

T.C: $O(n)$

S.C: $O(1)$ $\rightarrow O(1)$
space complexity ignore input

```
while (SP < EP) {
    int temp = arr[SP];
    arr[SP] = arr[EP];
    arr[EP] = temp;
    SP++;
    EP--;
}
```

3



Q) Rotate the array \rightarrow {google, meta, amazon}

In Given N elements, Rotate array from last to first by K times.

$k=3$ $\text{arr}[7]: \{ \underline{\underline{3}} \ \underline{-2} \ \underline{1} \ \underline{4} \ \underline{6} \ \underline{9} \ \underline{8} \}$

\downarrow 1st rot.

{ $\underline{\underline{8}}$ } $3 \ -2 \ 1 \ 4 \ 6 \ 9 \ 3$

\downarrow 2nd rot.

{ $\underline{\underline{9}}$ } $8 \ 3 \ -2 \ 1 \ 4 \ 6 \ 3$

\downarrow 3rd rot.

{ $\underline{\underline{6}}$ } $9 \ 8 \ 3 \ -2 \ 1 \ 4 \ 3$

$k=3$ $\text{arr}[7]: \{ \underline{\underline{3}} \ \underline{-2} \ \underline{1} \ \underline{4} \ \underline{6} \ \underline{9} \ \underline{8} \}$



{ $\underline{\underline{6}}$ } $9 \ 8 \ 3 \ -2 \ 1 \ 4 \ 3$

$k=4$ $\text{arr}[7]: \{ \underline{\underline{3}} \ \underline{-2} \ \underline{1} \ \underline{4} \ \underline{6} \ \underline{9} \ \underline{8} \}$



$4 \ 6 \ 9 \ 8 \ 3 \ -2 \ 1$



K=3

arr[7]: { 3 1 -2 1 4 6 9 8 }



Reverse the array

{ 3 1 -2 1 4 6 9 8 }



Reverse the first K elements

{ 8 9 6 4 1 -2 3 }



Reverse the remaining elements

{ 6 9 8 3 -2 1 4 }





II Pseudo code

```
int main() {  
    input → arr[N], K  
    K = K * N
```

// Step 1: Reverse the whole array.

```
reverse (arr, 0, N-1);
```

// Step 2: Reverse the first K elements.

```
reverse (arr, 0, K-1);
```

// Step 3: Reverse the remaining elements

```
reverse (arr, K, N-1);
```

T.C: O(N)

S.C: O(1)

}

```
Public static void reverse (int [] arr, int s, int e){
```

```
    int SP = s;
```

```
    int EP = e;
```

```
    while (SP < EP) {  
        int temp = arr[SP];  
        arr[SP] = arr[EP];  
        arr[EP] = temp;  
        SP++;  
        EP--;
```

}



$K=1$

$\text{arr}[4]: \{ \underset{\downarrow \text{arr+1}}{4} \underset{0}{3} \underset{1}{7} \underset{2}{10} \underset{3}{3} \}$

10 4 3 7

$\downarrow \text{arr+2}$

7 10 4 3

$\downarrow \text{arr+3}$

3 7 10 4

$\downarrow \text{arr+4}$

4 3 7 10

$\downarrow \text{arr+5}$

10 4 3 7

$\downarrow \text{arr+6}$

7 10 4 3

$\downarrow \text{arr+7}$

3 7 10 4

$\rightarrow K > N \rightarrow K = K - N ??$

$K=10$

$\text{arr}[4]: \{ \underset{\downarrow \text{arr+1}}{4} \underset{0}{3} \underset{1}{7} \underset{2}{10} \underset{3}{3} \}$

$\downarrow \text{arr+1}$

$\downarrow \text{arr+2}$

$\downarrow \text{arr+3}$

$\downarrow \text{arr+4}$

$\downarrow \text{arr+5}$

$\downarrow \text{arr+6}$

$\downarrow \text{arr+7}$

$\downarrow \text{arr+8}$

$\downarrow \text{arr+9}$

$\downarrow \text{arr+10}$



effective no. of rotation: $K \% N$

N

7

K

24

effective rotation

$$24 - 7 = 17 - 7 = 10 \cdot 7 : 3$$

$\{ \dots \}$

$$\hookrightarrow 24 \% 7 = 3$$

7

31

$$31 \% 7 = 3$$

7

3

$$3 \% 7 = 3$$



AlgoPrep

Break till 10:40 Pm



Q) Given N array elements, count total no. of elements having atleast 1 element greater than itself.

ex: $\text{arr}[7]: \{ -4^0, -3^1, 1^2, 9^3, 3^4, 9^5, 4^6 \}$
↳ ans = 5

$\text{arr}[8]: \{ 3^0, 4^1, 11^2, 8^3, 2^4, 10^5, 9^6, 11^7 \}$
↳ ans = 6

$\text{arr}[5]: \{ 7^0, 7^1, 7^2, 7^3, 7^4 \}$
↳ ans = 0

$\text{arr}[6]: \{ 1^0, 2^1, 2^2, 3^3, 3^4, 4^5 \}$
↳ ans = 5

Obs1: man elements of the array are not valid.

Obs2: non-man elements are always valid.

→ find occurrence of man element. → mancount

↳ ans = no. of Elements - mancount.



arr[8]: { 3 4 2 8 11 10 9 11 }
0 1 2 2 4 5 6 7

11P Suedo Code

man: 3 4 8 11

manCount = 0 2

int Countgreater (int arr[N]) {

int man = arr[0];

for (int i=1; i < N; i++) {
if (arr[i] > man) { man = arr[i]; }
}

T.C: O(2N) ≈ O(N)

S.C: O(1)

int manCount = 0;
for (int i=0; i < N; i++) {
if (arr[i] == man) { manCount++; }
}

return N - manCount;

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Q) Two sum

Given N array elements, check if there exists a pair (i, j) such that $\text{arr}[i] + \text{arr}[j] = k$ and $i \neq j$

Note: i and j are index values, k is given sum.

ex: $\text{arr}[7]: \{ 2 -1 0 3 2 5 7 \}$
 $K=8$ ↳ True

$\text{arr}[4]: \{ 1 3 -2 6 \}$
 $K=5$ ↳ False

$\text{arr}[5]: \{ 2 4 -3 7 10 \}$
 $K=8$ ↳ $\text{arr}[1] + \text{arr}[2] = 8 \rightarrow \text{True}$

$\text{arr}[6]: \{ 3 5 2 8 3 7 \}$
 $K=6$ ↳ True



$k=12$

$\text{arr}[5]: \{3 \ 5 \ 2 \ 7 \ 5\}$

$i \ j$

(0, 0)	(1, 0)	(2, 0)	(3, 0)	(4, 0)
(0, 1)	(1, 1)	(2, 1)	(3, 1)	(4, 1)
(0, 2)	(1, 2)	(2, 2)	(3, 2)	(4, 2)
(0, 3)	(1, 3)	(2, 3)	(3, 3)	(4, 3)
(0, 4)	(1, 4)	(2, 4)	(3, 4)	(4, 4)

$i \rightarrow 0 \downarrow 1 \downarrow 2 \downarrow 3 \downarrow 4$
 $j \rightarrow (1, 2, 3, 4) \quad (2, 3, 4) \quad (3, 4) \quad (4)$

Public static boolean TwoSum (int arr[], int k){

for (int i = 0; $i < \text{arr}^{\text{size}}$; i++) {

 for (j = i + 1; $j < N - 1$; j++) {

 if ($\text{arr}[i] + \text{arr}[j] == k$) { return true; }

 return false;

Time Complexity: ??

Space Complexity: $O(1)$