

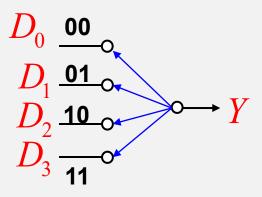
# 第四章 组合逻辑电路设计(三)

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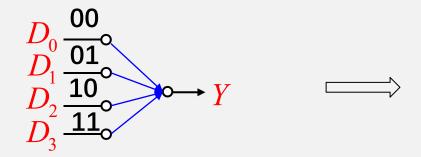
#### 1.多路选择器的基本功能

从一组输入数据中,选择出某一个数据,完成这种功能的逻辑电路称为数据选择器(或称为多路选择开关)



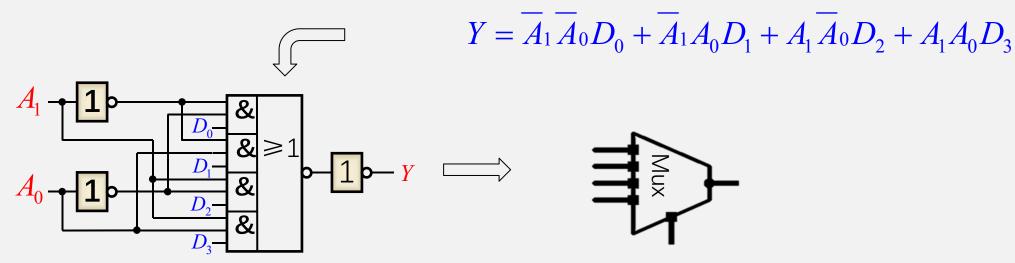


### 2. 4路数据选择器的设计 (MUX)



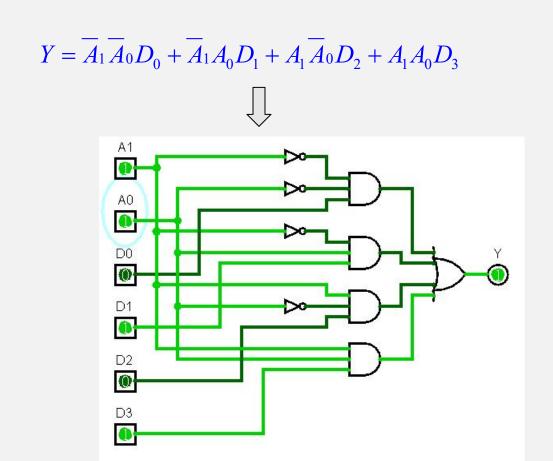
D	<b>A</b> <sub>1</sub>	$A_0$	Υ
D <sub>0</sub>	0	0	D <sub>0</sub>
D <sub>1</sub>	0	1	D <sub>1</sub>
D <sub>2</sub>	1	0	D <sub>2</sub>
$D_3$	1	1	D <sub>3</sub>





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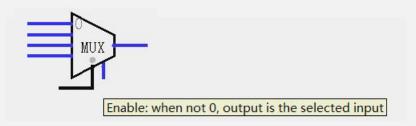


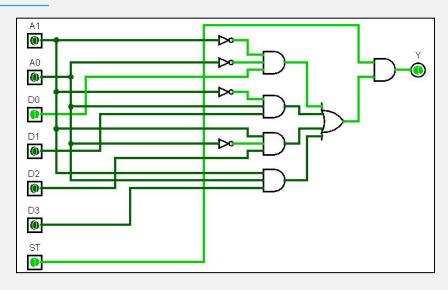


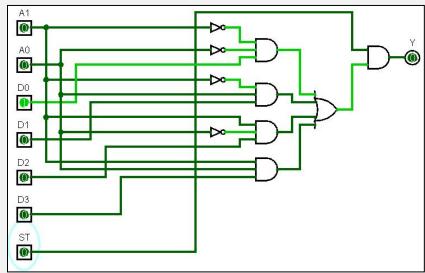


$\overline{ST}_1$	<b>A</b> <sub>1</sub>	$A_0$	<b>Y</b> <sub>1</sub>
1	X	X	0
0	0	0	$D_0$
0	0	1	D <sub>1</sub>
0	1	0	$D_2$
0	1	1	$D_3$

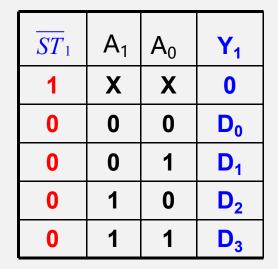
$$Y = \overline{\overline{ST_1}} (\overline{A_1} \overline{A_0} D_0 + \overline{A_1} A_0 D_1 + A_1 \overline{A_0} D_2 + A_1 A_0 D_3)$$



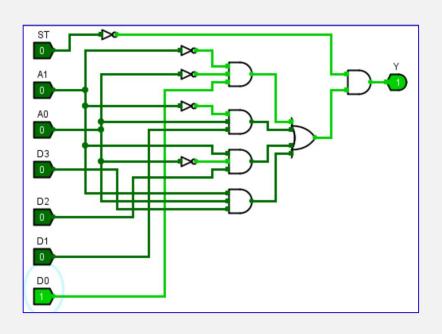




使能信号为正







使能信号为负

5 '



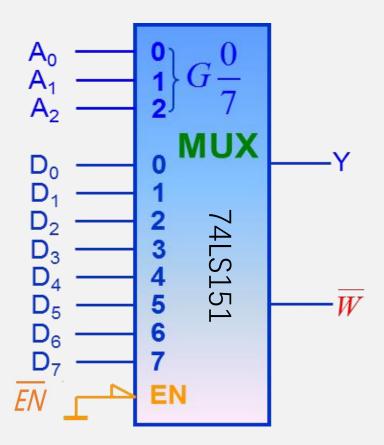
### 3.带使能和可扩展功能的8路数据选择器设计

$\overline{EN}$	$A_2$	1	0		
1	X	X	X	0	1
0	0	0	0	$D_0$	$\overline{\mathrm{D}}_{\mathrm{0}}$
0	0	0	1	$D_1$	$\overline{\mathrm{D}}_{\!1}$
0	0	1	0	$D_2$	$\overline{ extsf{D}}_{\!2}$
0	0	1	1	$D_3$	$\overline{D}_{\!3}$
0	1	0	0	$D_4$	$\overline{\mathrm{D}}_{\!4}$
0	1	1	1	$D_5$	$\overline{\mathrm{D}}_{\!5}$
0	1	0	0	$D_6$	$\overline{\mathrm{D}_{\mathrm{6}}}$
0	1	1	1	$D_7$	$\overline{\mathrm{D}_7}$

EN:选通端,低有效。

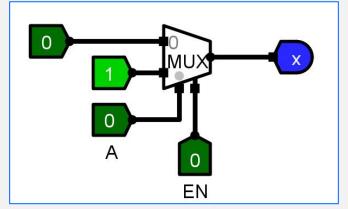
Y,W: 互补输出端。

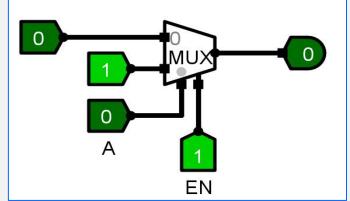
可参照4路选择器写出Y逻辑 表达式

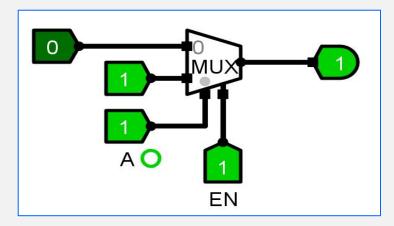


$$Y = (\overline{A}_{2}\overline{A}_{1}\overline{A}_{0} D_{0} + \overline{A}_{2}\overline{A}_{1}A_{0} D_{1} + \overline{A}_{2}A_{1}\overline{A}_{0} D_{2} + \overline{A}_{2}A_{1}A_{0} D_{3} + A_{2}\overline{A}_{1}\overline{A}_{0} D_{4} + A_{2}\overline{A}_{1}A_{0} D_{5} + A_{2}A_{1}\overline{A}_{0} D_{6} + A_{2}A_{1}A_{0} D_{7})\overline{EN}$$



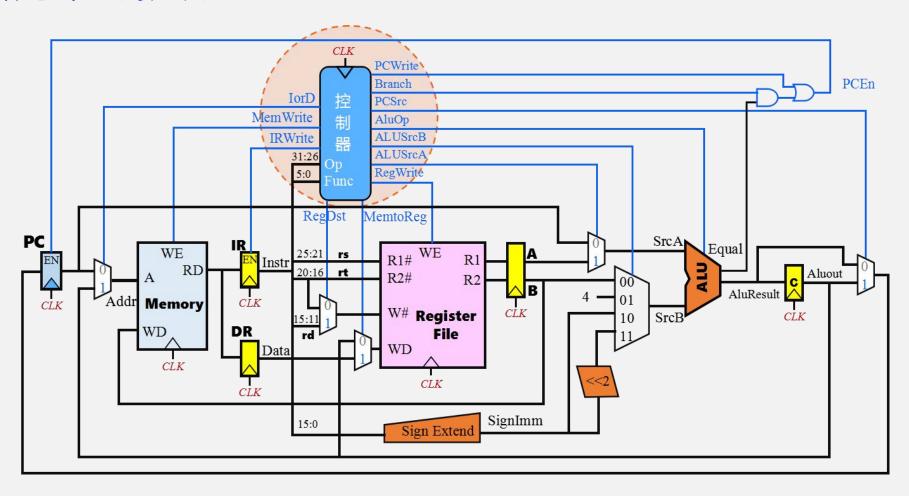








#### 4.数据选择器的应用



同一目标有多个数据来源时,在其入口处需使用多路选择器

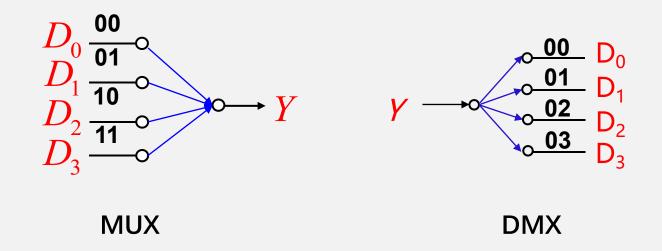
9

### 4.8 多路分配器(解复用器 Demultiplexer)



#### 1.多路分配器的基本功能

将1个输入数据,根据需要传送到m个输出端的任何一个输出端的电路, 称为数据分配器、多路分配器或解复用器, 其逻辑功能正好与多路选择器相反。

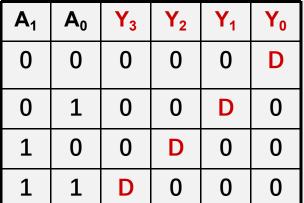


### 4.8 多路分配器(解复用器 Demultiplexer)



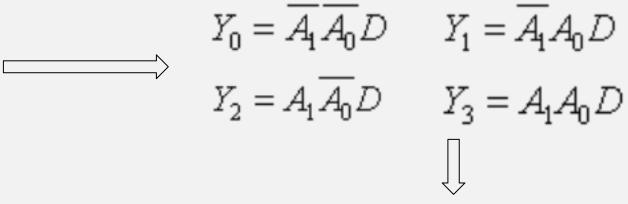
### 2.多路分配器的设计

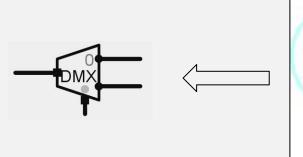
<b>A</b> <sub>1</sub>	$A_0$	Y <sub>3</sub>	Y <sub>2</sub>	<b>Y</b> <sub>1</sub>	Y <sub>0</sub>
0	0	0	0	0	D
0	1	0	0	D	0
1	0	0	D	0	0
1	1	D	0	0	0

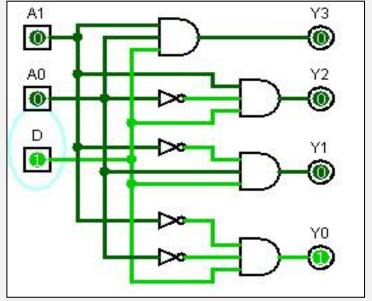


A <sub>1</sub>	$A_0$	Υ
0	0	$D_0$
0	1	D <sub>1</sub>
1	0	$D_2$
1	1	$D_3$

多路选择器真值表

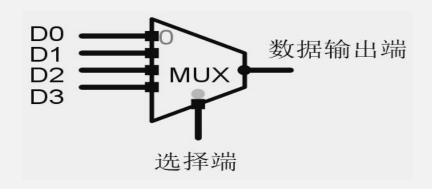




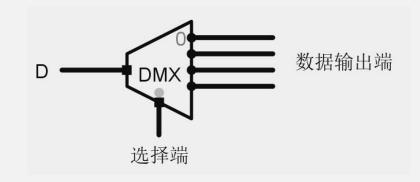


## 4.9 多路选择器、多路分配器、译码器比较





$$Y = \overline{A}_{1} \overline{A}_{0} D_{0} + \overline{A}_{1} A_{0} D_{1} + A_{1} \overline{A}_{0} D_{2} + A_{1} A_{0} D_{3}$$



$$Y_0 = \overline{A_1} \overline{A_0} D$$
  $Y_1 = \overline{A_1} A_0 D$   $Y_2 = A_1 \overline{A_0} D$   $Y_3 = A_1 A_0 D$ 

$$Y_0 = \overline{A_1} \overline{A_0}$$
  $Y_1 = \overline{A_1} A_0$   $Y_2 = A_1 \overline{A_0}$   $Y_3 = A_1 A_0$ 



#### 1.利用变量译码器实现组合逻辑函数

<b>A</b> <sub>1</sub>	$A_0$	<b>Y3</b>	Y2	<b>Y1</b>	Y0
0	0	0	0	0	1
0	1	0	0	1	0
1	0	0	1	0	0
1	1	1	0	0	0

$$Y_3 = A_1 A_0$$
  $Y_2 = A_1 \overline{A}_0$   
 $Y_1 = \overline{A}_1 A_0$   $Y_0 = \overline{A}_1 \overline{A}_0$ 

一个n变量输入的变量译码器, 其输出包含了n个输入变量的全部最小项。用n变量译码器加输出门就能实现任何形式的输入变量不大于n 的组合逻辑函数。



例1用译码器实现一组多输出函数

$$F_{1} = A\overline{B} + \overline{B}C + AC$$

$$F_{2} = \overline{AB} + B\overline{C} + ABC$$

$$F_{3} = \overline{AC} + BC + A\overline{C}$$

解: 三输入变量的多输出函数,用3-8译码器实现

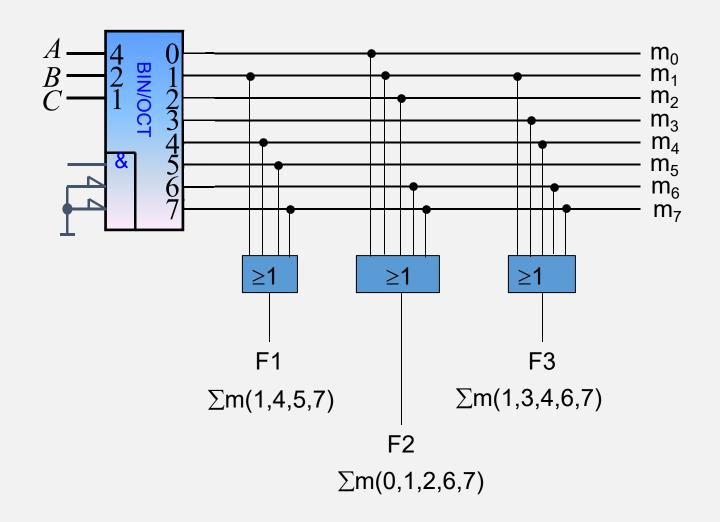
将多输出函数写成最小项之和形式,再配合适当的逻辑门即可。

$$F_1 = A\overline{B} + \overline{B}C + AC = \sum m(1,4,5,7)$$

$$F_2 = \overline{AB} + B\overline{C} + ABC = \sum m(0,1,2,6,7)$$

$$F_3 = \overline{AC} + BC + A\overline{C} = \sum m(1,3,4,6,7)$$









若译码器是以反变量形式输出,即输出的是 mi,则:

$$F_{1} = A\overline{B} + \overline{B}C + AC = m_{1} + m_{4} + m_{5} + m_{7}$$

$$= \overline{m_{1} + m_{4} + m_{5} + m_{7}} = \overline{m_{1}} \cdot \overline{m_{4}} \cdot \overline{m_{5}} \cdot \overline{m_{7}}$$

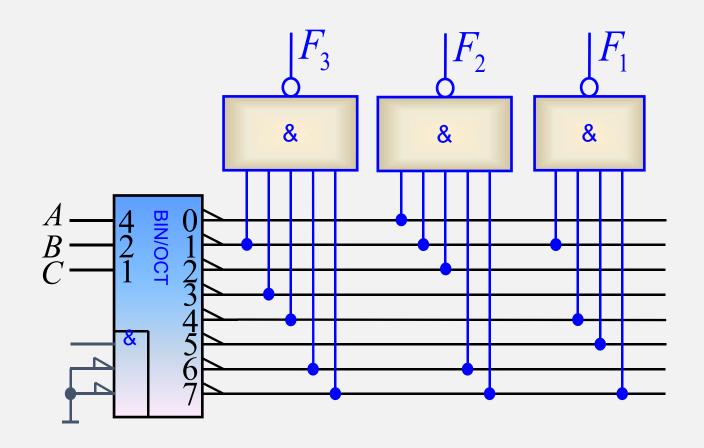
$$= \overline{\overline{Y_{1}} \cdot \overline{Y_{4}} \cdot \overline{Y_{5}} \cdot \overline{Y_{7}}}$$

$$F_{2} = \overline{AB} + B\overline{C} + ABC = \sum m(0,1,2,6,7) = \overline{m_{0}} \cdot \overline{m_{1}} \cdot \overline{m_{2}} \cdot \overline{m_{6}} \cdot \overline{m_{7}}$$

$$= \overline{\overline{Y_{0}} \cdot \overline{Y_{1}} \cdot \overline{Y_{2}} \cdot \overline{Y_{6}} \cdot \overline{Y_{7}}}$$

$$F_{3} = \overline{AC} + BC + A\overline{C} = \sum m(1,3,4,6,7) = \overline{\overline{Y_{1}} \cdot \overline{Y_{3}} \cdot \overline{Y_{4}} \cdot \overline{Y_{6}} \cdot \overline{Y_{7}}}$$





17`



例2: 用2-4译码器和适当的逻辑门实现逻辑函数

$$F_{1} = A\overline{B} + \overline{B}C + AC$$

$$F_{2} = \overline{AB} + B\overline{C} + ABC$$

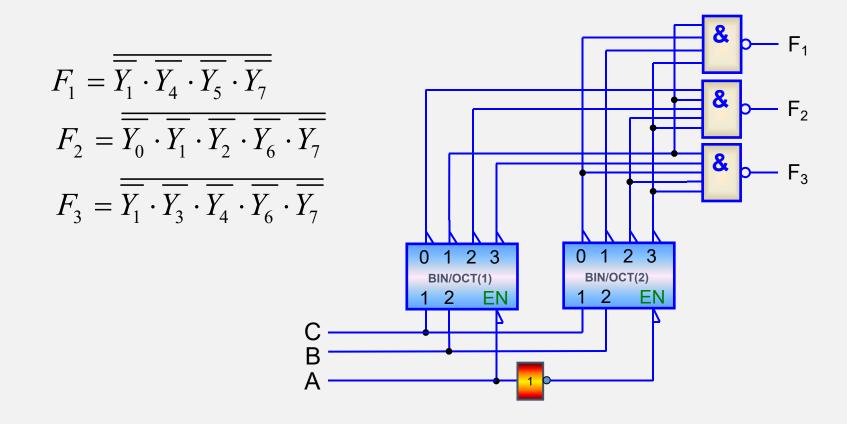
$$F_{3} = \overline{AC} + BC + A\overline{C}$$

$$F_{1} = A\overline{B} + \overline{B}C + AC = \overline{m_{1}} \cdot \overline{m_{4}} \cdot \overline{m_{5}} \cdot \overline{m_{7}} = \overline{Y_{1}} \cdot \overline{Y_{4}} \cdot \overline{Y_{5}} \cdot \overline{Y_{7}}$$

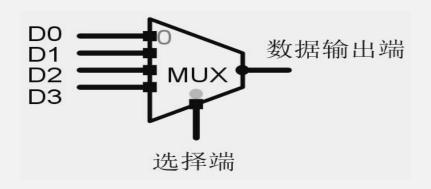
$$F_{2} = \overline{A}B + B\overline{C} + ABC = \overline{m_{0}} \cdot \overline{m_{1}} \cdot \overline{m_{2}} \cdot \overline{m_{6}} \cdot \overline{m_{7}} = \overline{Y_{0}} \cdot \overline{Y_{1}} \cdot \overline{Y_{2}} \cdot \overline{Y_{6}} \cdot \overline{Y_{7}}$$

$$F_{3} = \overline{A}C + BC + A\overline{C} = \overline{m_{1}} \cdot \overline{m_{3}} \cdot \overline{m_{4}} \cdot \overline{m_{6}} \cdot \overline{m_{7}} = \overline{Y_{1}} \cdot \overline{Y_{3}} \cdot \overline{Y_{4}} \cdot \overline{Y_{6}} \cdot \overline{Y_{7}}$$









$$Y = \overline{A}_{1} \overline{A}_{0} D_{0} + \overline{A}_{1} A_{0} D_{1} + A_{1} \overline{A}_{0} D_{2} + A_{1} A_{0} D_{3}$$

$$F_1 = A\overline{B} + \overline{B}C + AC = m_1 + m_4 + m_5 + m_7$$

$$F_2 = \overline{AB} + B\overline{C} + ABC = \sum m(0,1,2,6,7)$$

$$F_3 = \overline{AC} + BC + A\overline{C} = \sum m(1,3,4,6,7)$$



### 本节内容完

计算机组成原理