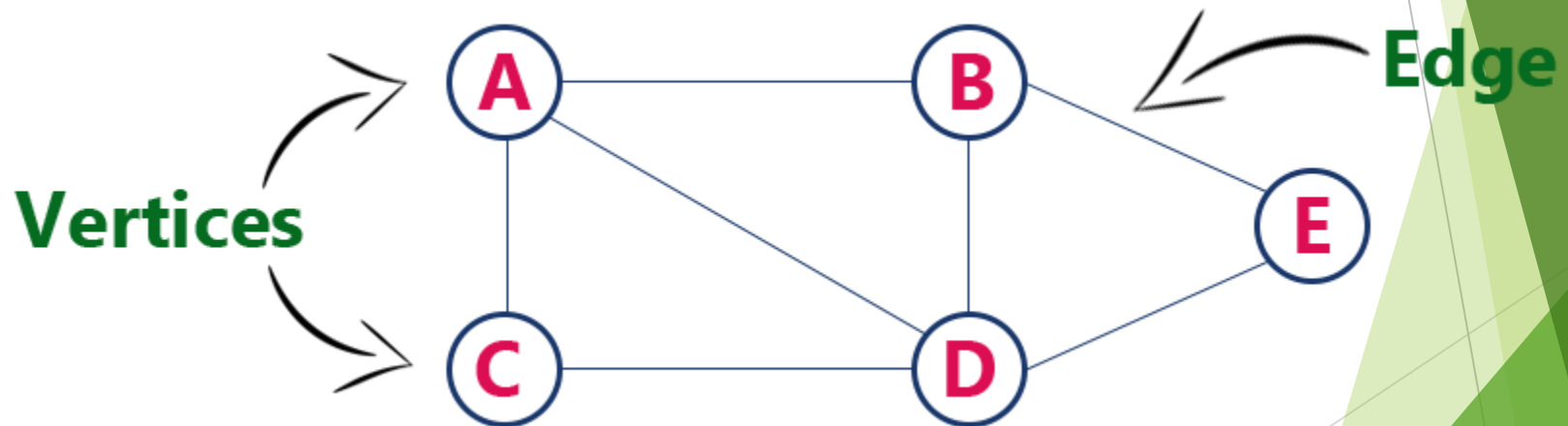


01418231

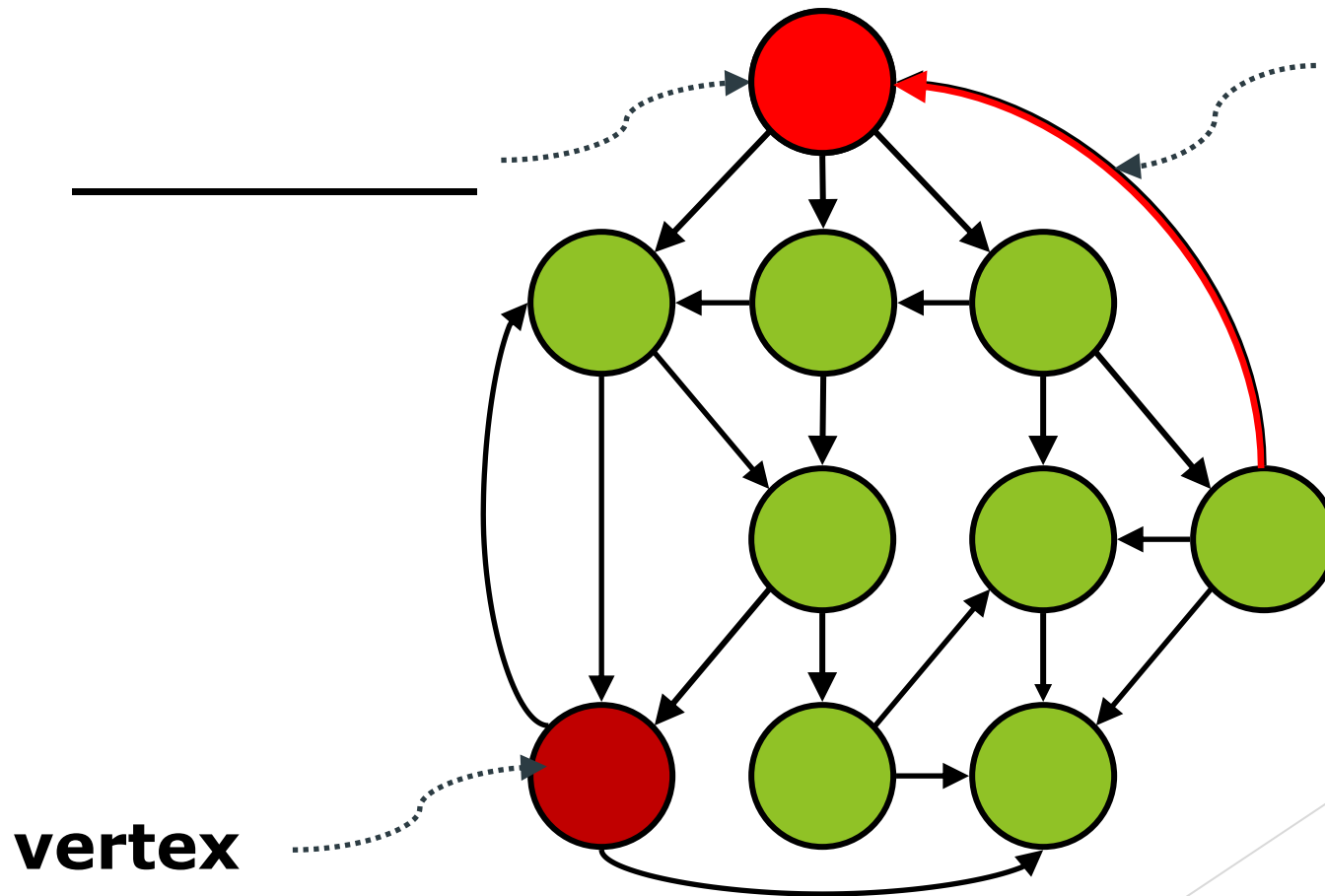
Data Structure

Graph



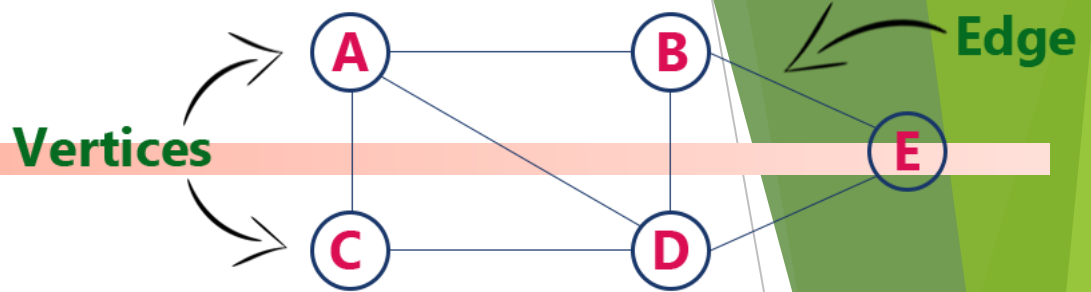
Graph and Traversal

Graph



vertex

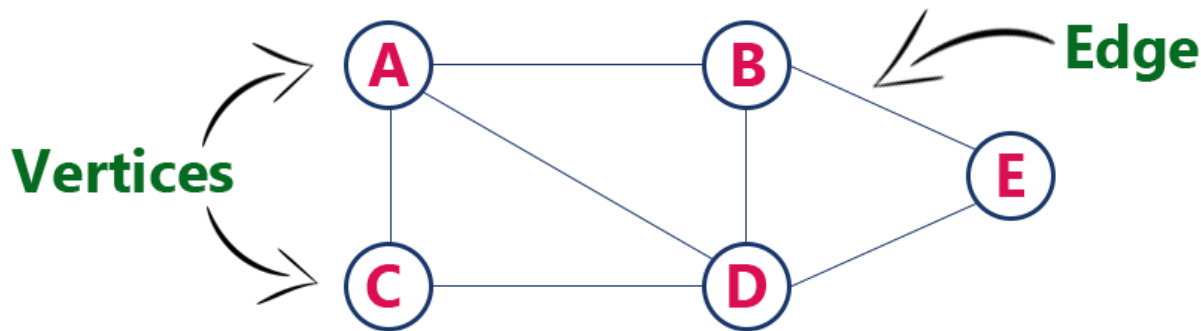
Graphs



- ▶ $G = (V, E)$
- ▶ V is the vertex set
 - ▶ Vertices are also called nodes and points.
 - ▶ A, B, C, D & E are known as vertices.
- ▶ E is edge or arc that connects two between vertices.
 - ▶ (A,B) is the link between vertices A and B
- ▶ Edges are three types.
 1. _____ - An undirected edge is a bidirectional edge. If there is a undirected edge between vertices A and B then edge (A , B) is equal to edge (B , A).
 2. _____ - A directed edge is a unidirectional edge. If there is a directed edge between vertices A and B then edge (A , B) is not equal to edge (B , A).
 3. _____ - A weighted edge is an edge with cost on it.

Graphs

- ▶ This graph G can be defined as $G = (V, E)$
- ▶ Where $V = \{A, B, C, D, E\}$
- ▶ and $E = \{(A, B), (A, C), (A, D), (B, D), (C, D), (D, E)\}$



Graphs

- ▶ Undirected edge has no orientation (u,v) .

u v

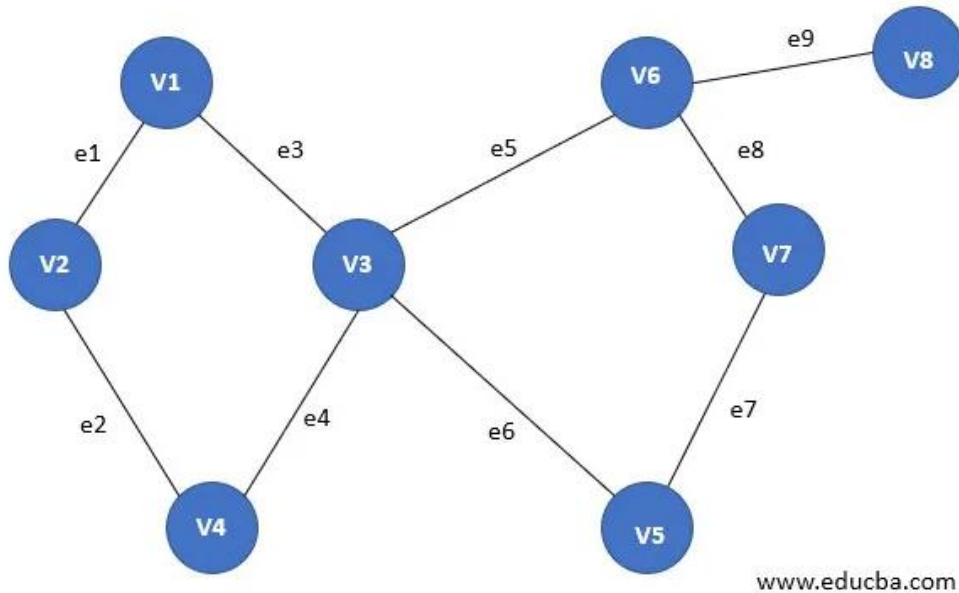
- ▶ Directed edge has an orientation (u,v)

u v

- ▶ Weighted Edge has an orientation (u,v) or no orientation

u v

Graphs



_____ =
The number of vertices in the graph.

_____ =
The number of edges in the graph.

_____ =
Number of edges incident to the vertex.

Types of Graphs in Data Structures

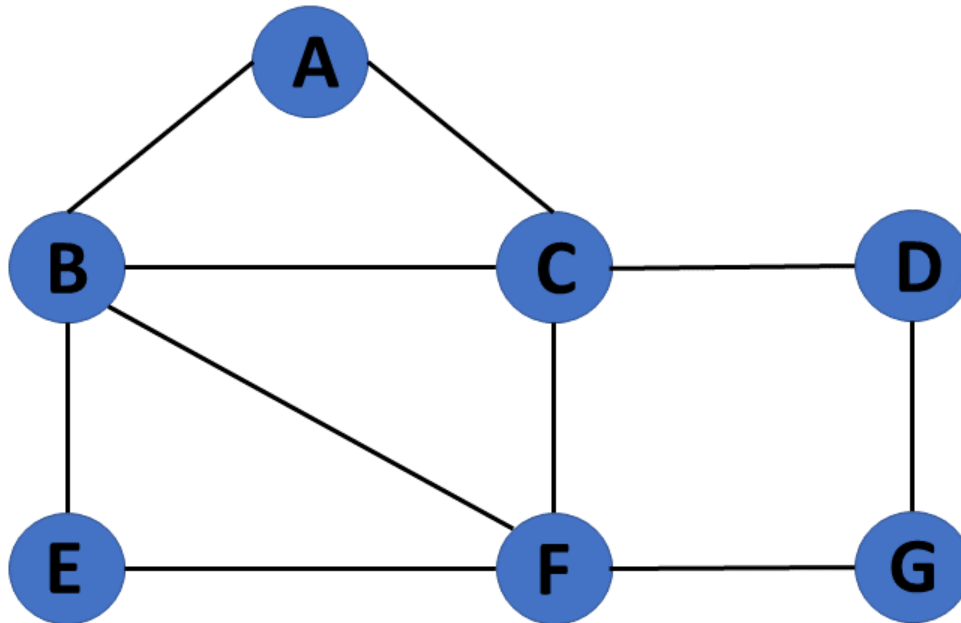
Types of Graphs in Data Structures

1. Finite Graph
2. Infinite Graph
3. Trivial Graph
4. Simple Graph
5. Multi Graph
6. Null Graph
7. Complete Graph
8. Pseudo Graph
9. Regular Graph
10. Weighted Graph
11. Directed Graph
12. Undirected Graph
13. Connected Graph
14. Disconnected Graph
15. Cyclic Graph
16. Acyclic Graph
17. Directed Acyclic Graph
18. Subgraph

<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

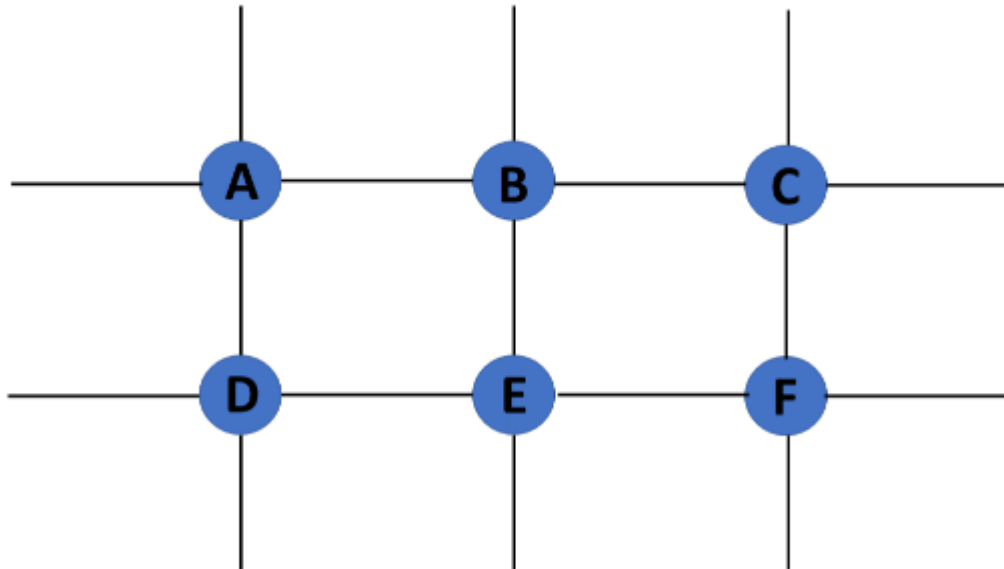
- The graph $G=(V, E)$ is called a finite graph if the number of vertices and edges in the graph is **limited** in number



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

- ▶ The graph $G=(V, E)$ is called a infinite graph if the number of vertices and edges in the graph is **interminable**.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

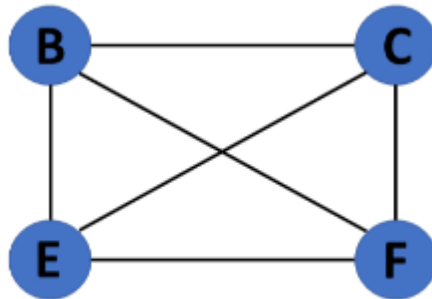
- ▶ A graph $G = (V, E)$ is trivial if it contains only a **single vertex and no edges**.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

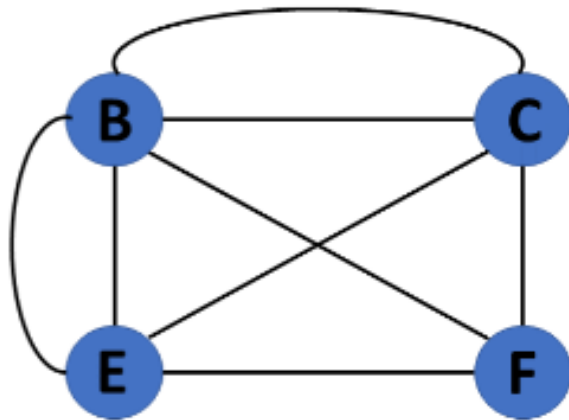
- ▶ If each pair of nodes or vertices in a graph $G=(V, E)$ has only one edge, it is a simple graph.
- ▶ As a result, there is just one edge linking two vertices, depicting one-to-one interactions between two elements.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

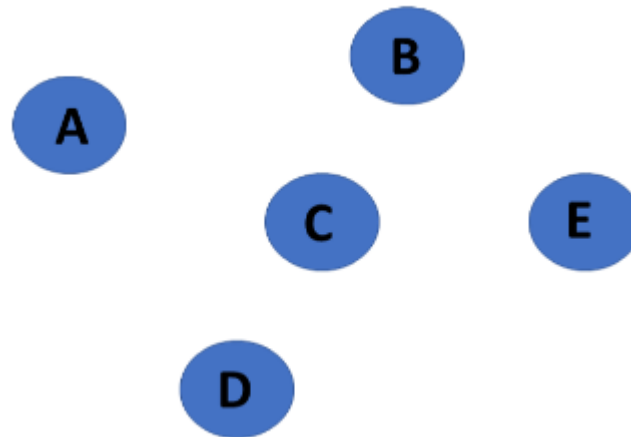
- ▶ If there are numerous edges between a pair of vertices in a graph $G = (V, E)$, the graph is referred to as a multigraph.
- ▶ There are **no self-loops** in a Multigraph.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

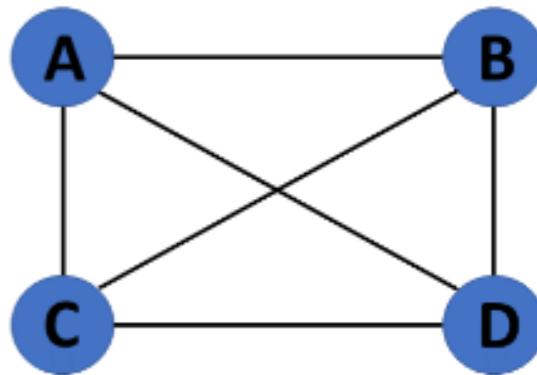
- It's a reworked version of a trivial graph. If several vertices but **no edges connect them**, a graph $G = (V, E)$ is a null graph.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

- ▶ If a graph $G = (V, E)$ is also a simple graph, it is complete.
- ▶ Using the edges, **with n number of vertices must be connected.**
- ▶ It's also known as a full graph because each **vertex's degree must be n-1.**



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

Have all possible edges.

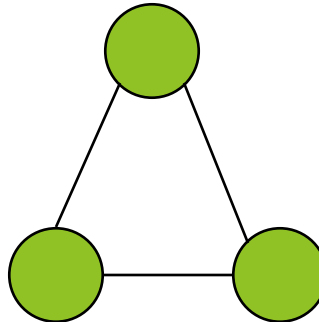
- Undirected Graph , Directed Graph



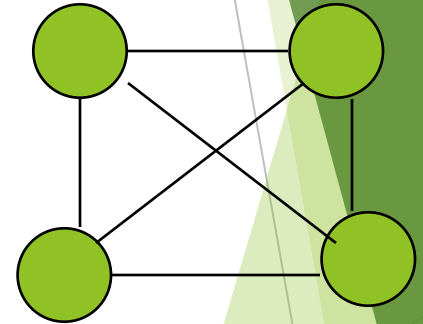
$n = 1$



$n = 2$

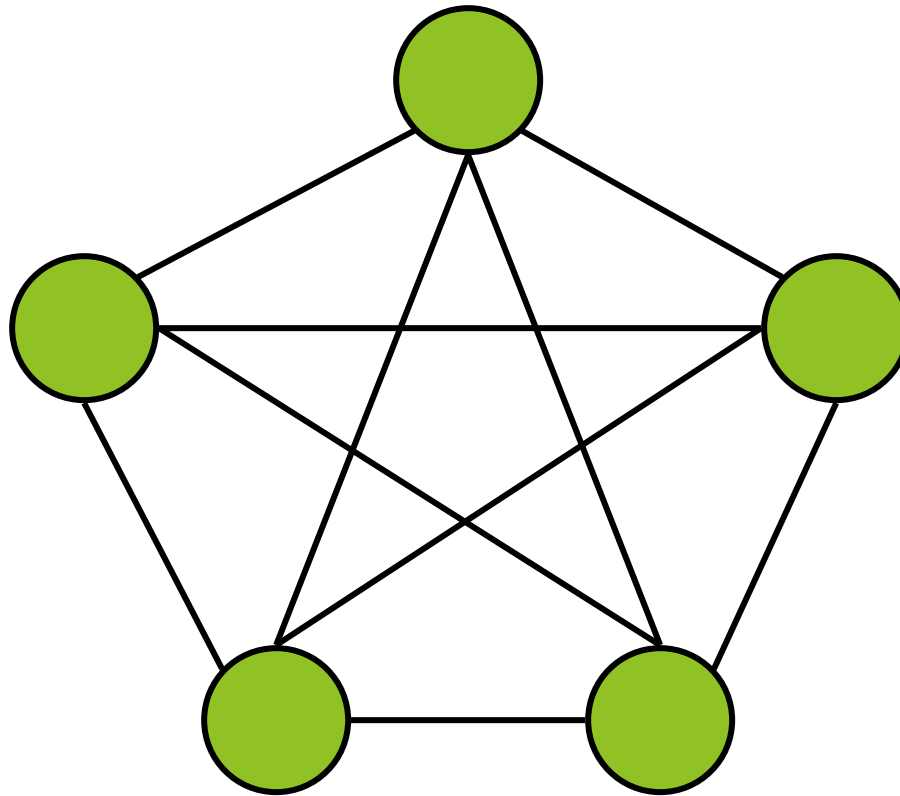


$n = 3$



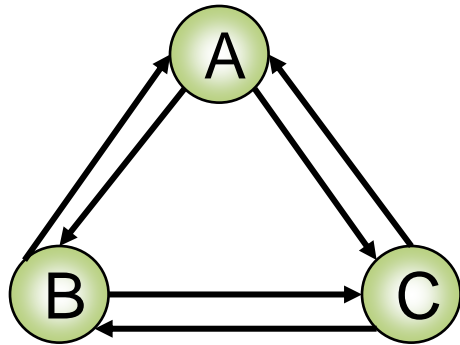
$n = 4$

Graph (Clique)



Graph (Clique)

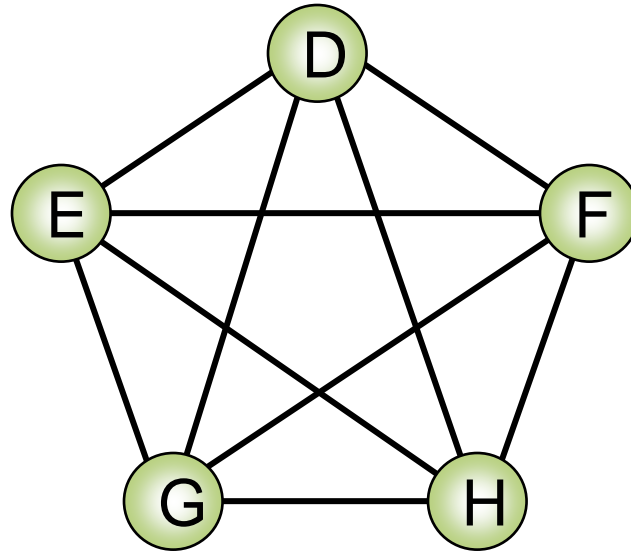
► Number of edges of graph



Directed Graph

$$\text{Number of Edges} = N*(N-1)$$

$$3*(3-1) = 6$$



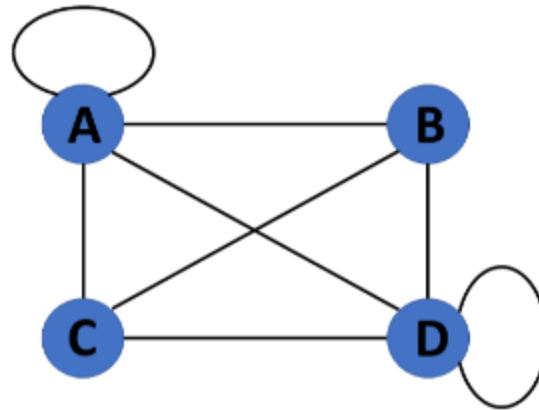
Undirected Graph

$$\text{Number of Edges} = [N*(N-1)] / 2$$

$$5*(5-1) / 2 = 10$$

Graph

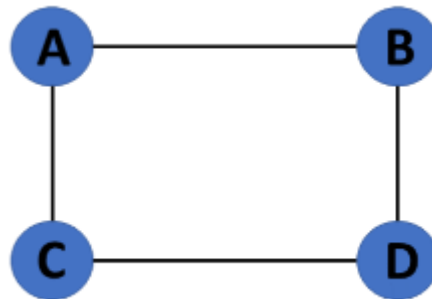
- If a graph $G = (V, E)$ contains a **self-loop** besides other edges, it is a pseudograph.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

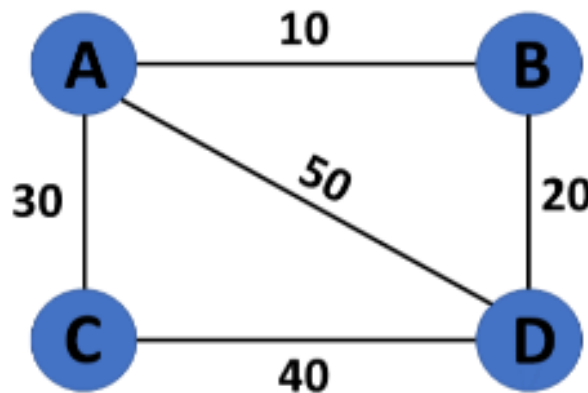
- ▶ If a graph $G = (V, E)$ is a simple graph with the **same degree at each vertex**, it is a regular graph.
- ▶ As a result, every whole graph is a regular graph.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

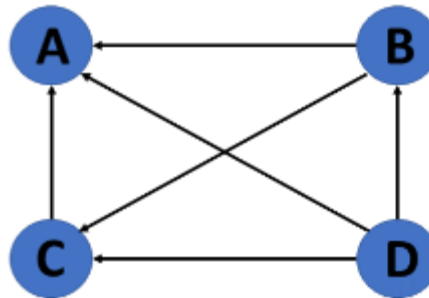
- ▶ A graph $G = (V, E)$ is called a **labeled or weighted** graph because each edge has a value or weight representing the cost of traversing that edge.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

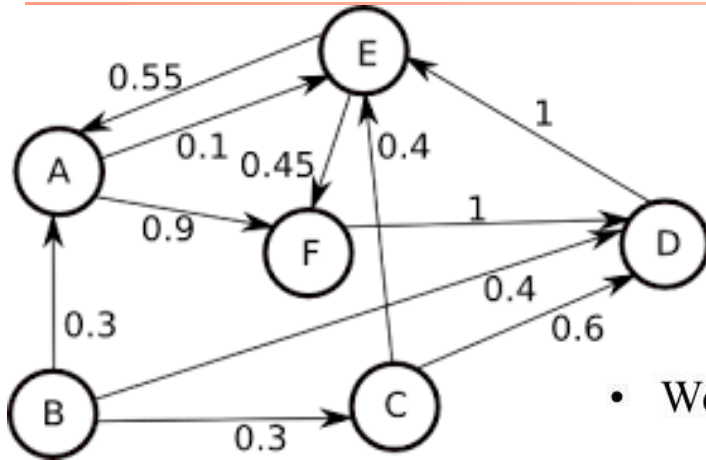
Graph

- ▶ A directed graph also referred to as a **digraph**, is a set of nodes connected by edges, each with a **direction**.



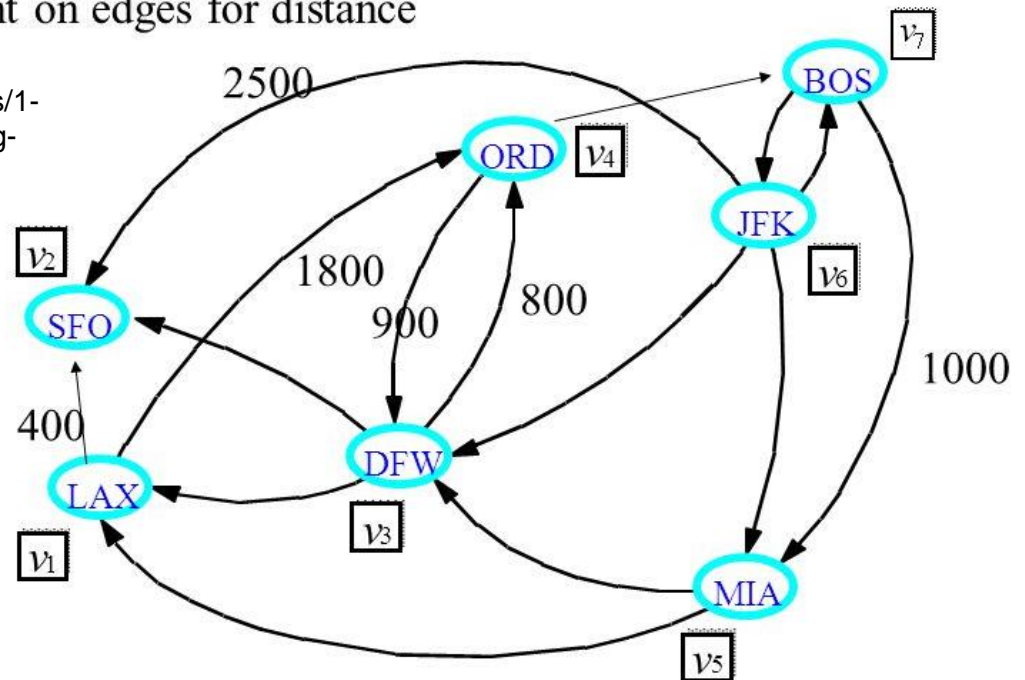
<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph with weighted edges



- Weight on edges for distance

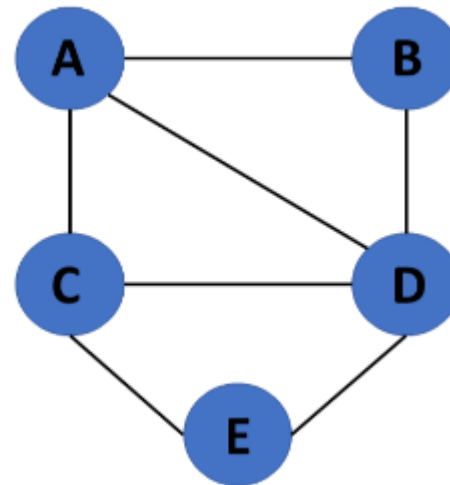
<https://www.chegg.com/homework-help/questions-and-answers/1-consider-directed-graph-weighted-edges-represent-graph-using-adjacency-list-adjacency-ma-q20254763>



<https://slideplayer.com/slide/4953310/>

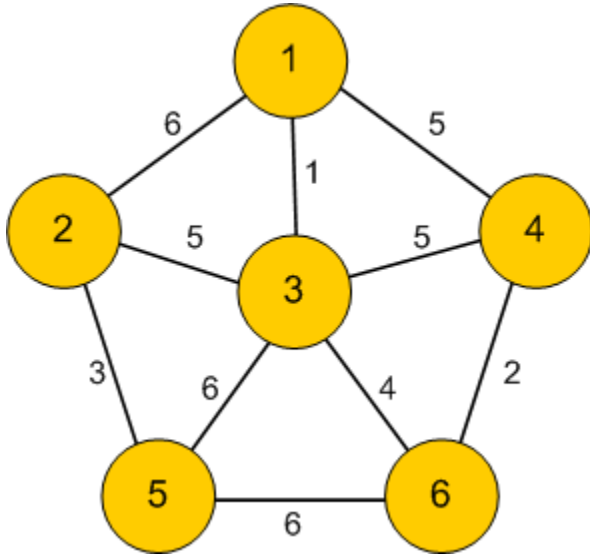
Graph

- ▶ An undirected graph comprises a set of nodes and links connecting them.
- ▶ The order of the two connected vertices is irrelevant and has **no direction**.
- ▶ You can form an undirected graph with a **finite number of vertices and edges**.

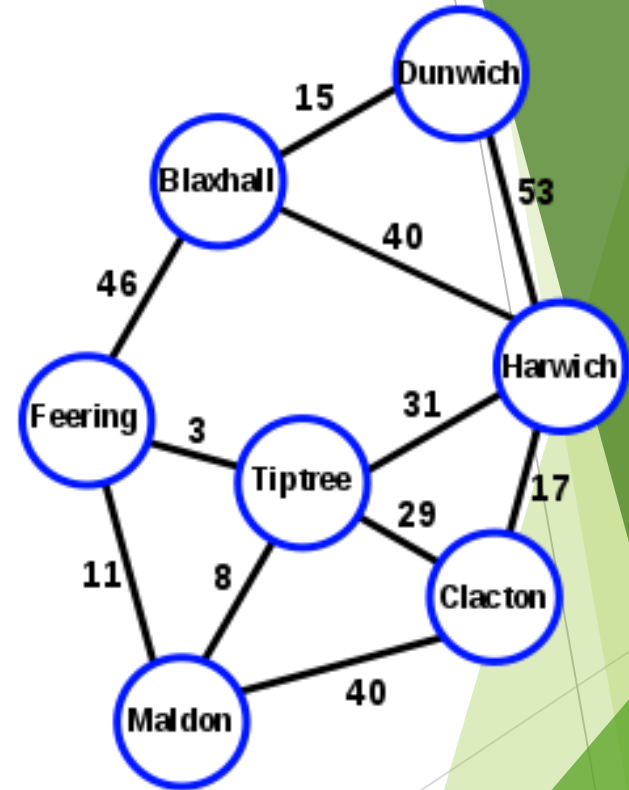


<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph with weighted edges



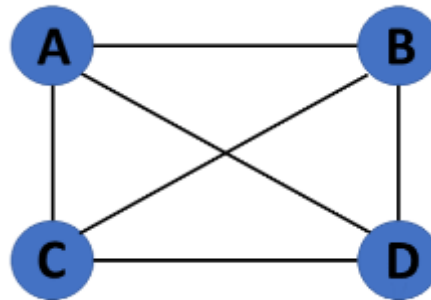
<https://forum.qt.io/topic/77736/qt-simple-library-for-drawing-undirected-weighted-graph>



https://en.wikipedia.org/wiki/Widest_path_problem

Graph

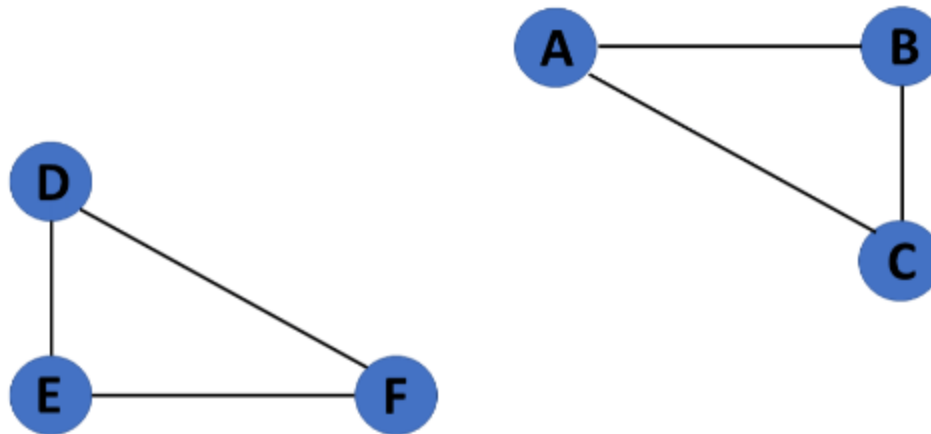
- If there is a path between one vertex of a graph data structure and any other vertex, the graph is connected.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

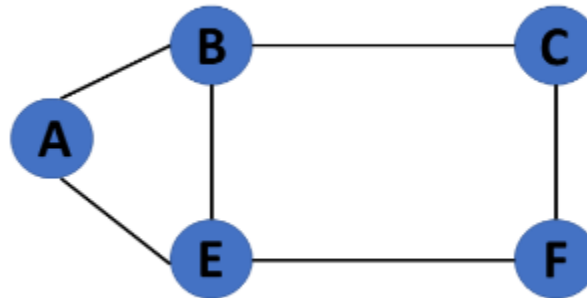
- ▶ When there is **no edge linking** the vertices, you refer to the null graph as a disconnected graph.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

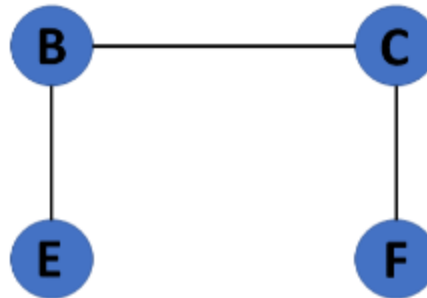
- If a graph contains at least one graph cycle, it is considered to be **cyclic**.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

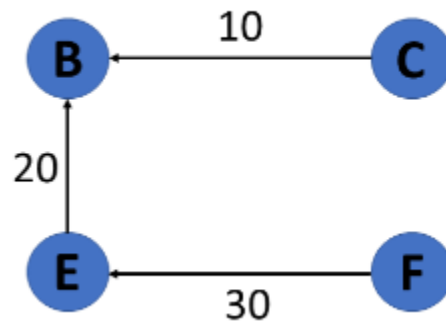
- ▶ When there are no cycles in a graph, it is called an **acyclic** graph.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

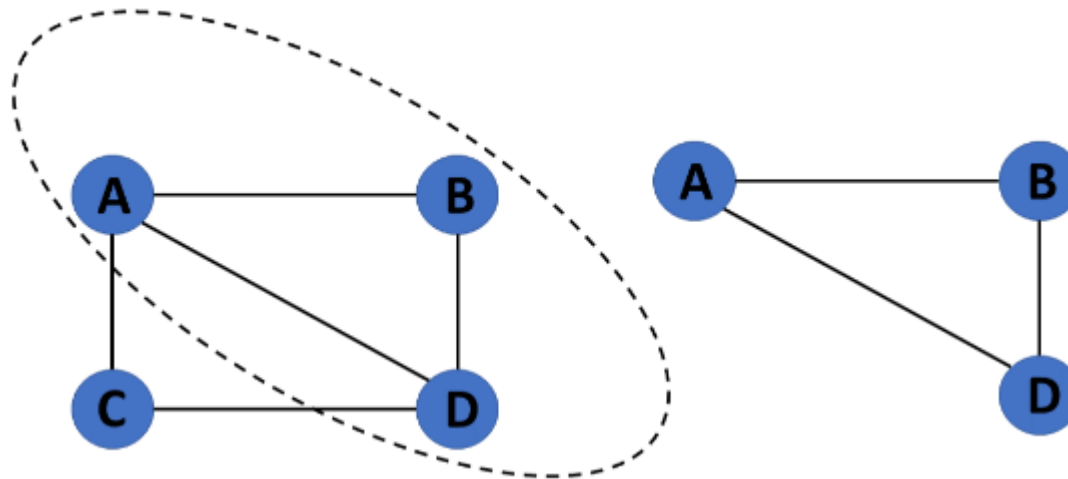
- It's also known as a directed acyclic graph (**DAG**), and it's a graph with directed edges but no cycle. It represents the edges using an ordered pair of vertices since it directs the vertices and stores some data.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

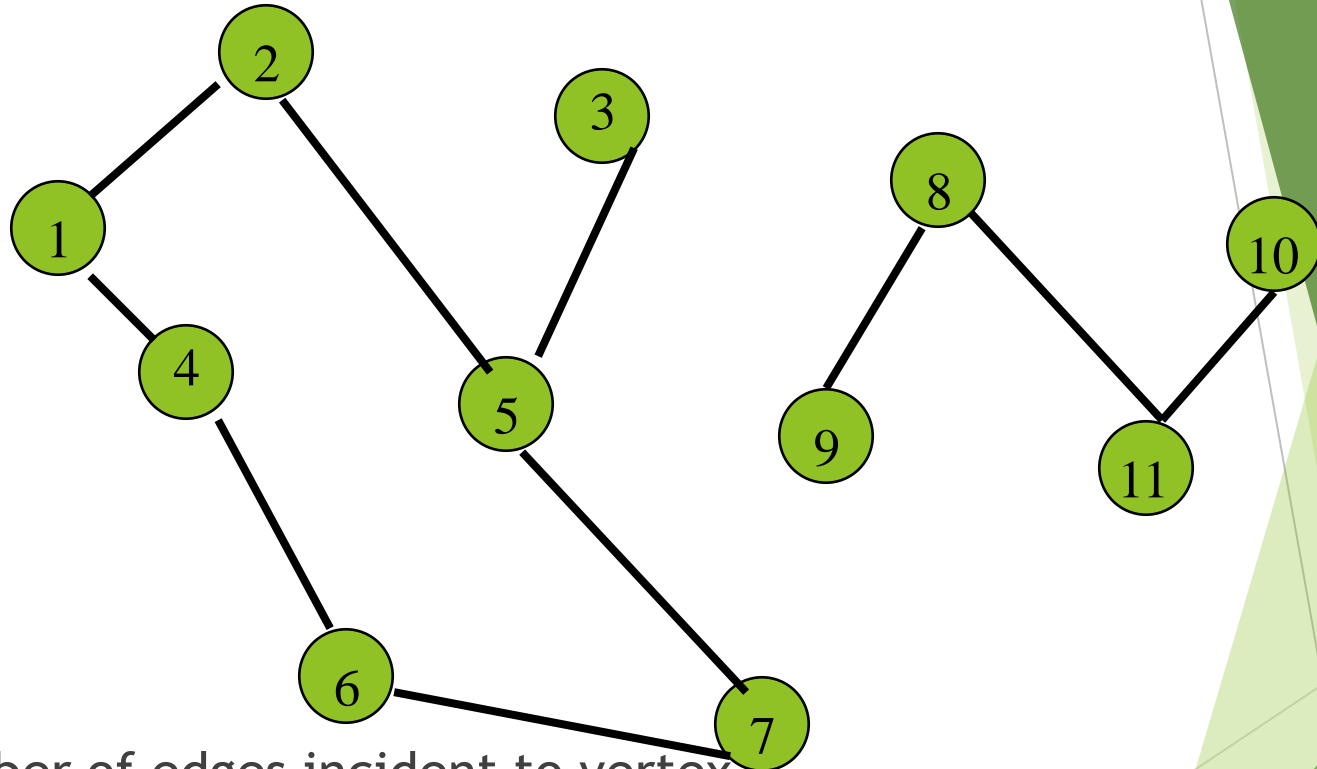
- ▶ The vertices and edges of a graph that are **subsets** of another graph are known as a subgraph



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Graph

Sum of degrees = $2e$ (e is number of edges)



- ▶ Number of edges incident to vertex.
- ▶ Undirected graph
- ▶ $\text{degree}(2) = 2, \text{degree}(5) = 3, \text{degree}(3) = 1$

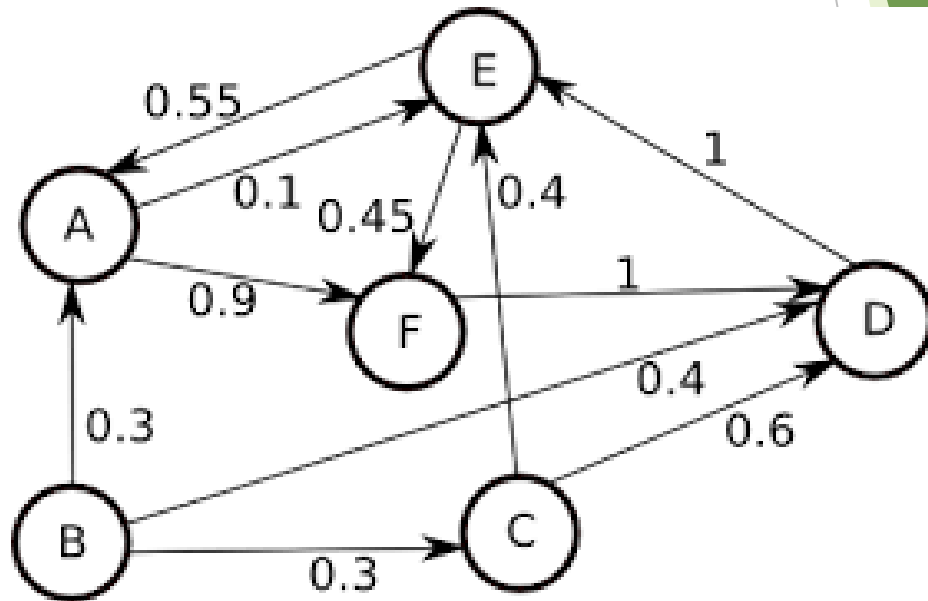
Vertex Degree

- ▶ Directed graph

- ▶ Sum of Degree (degree)= _____

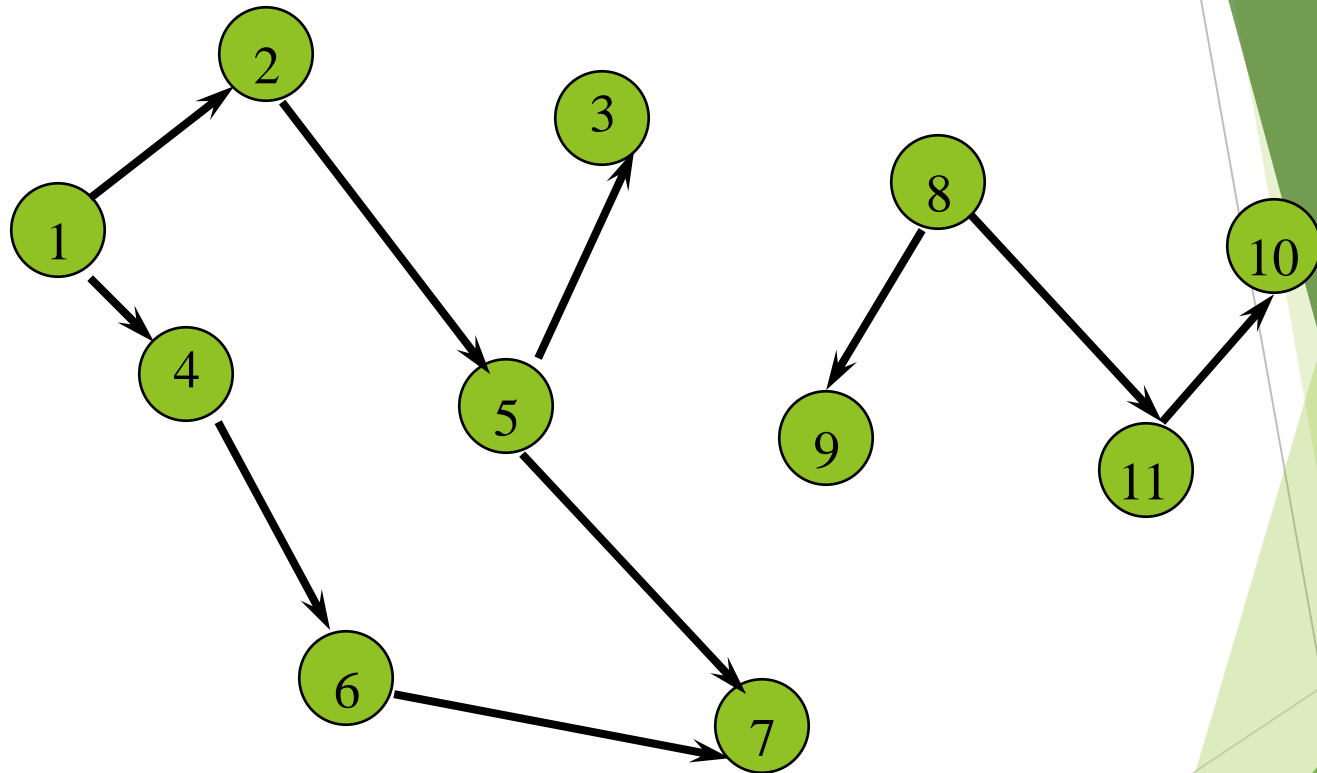
- ▶ In-Degree

- ▶ Out-Degree



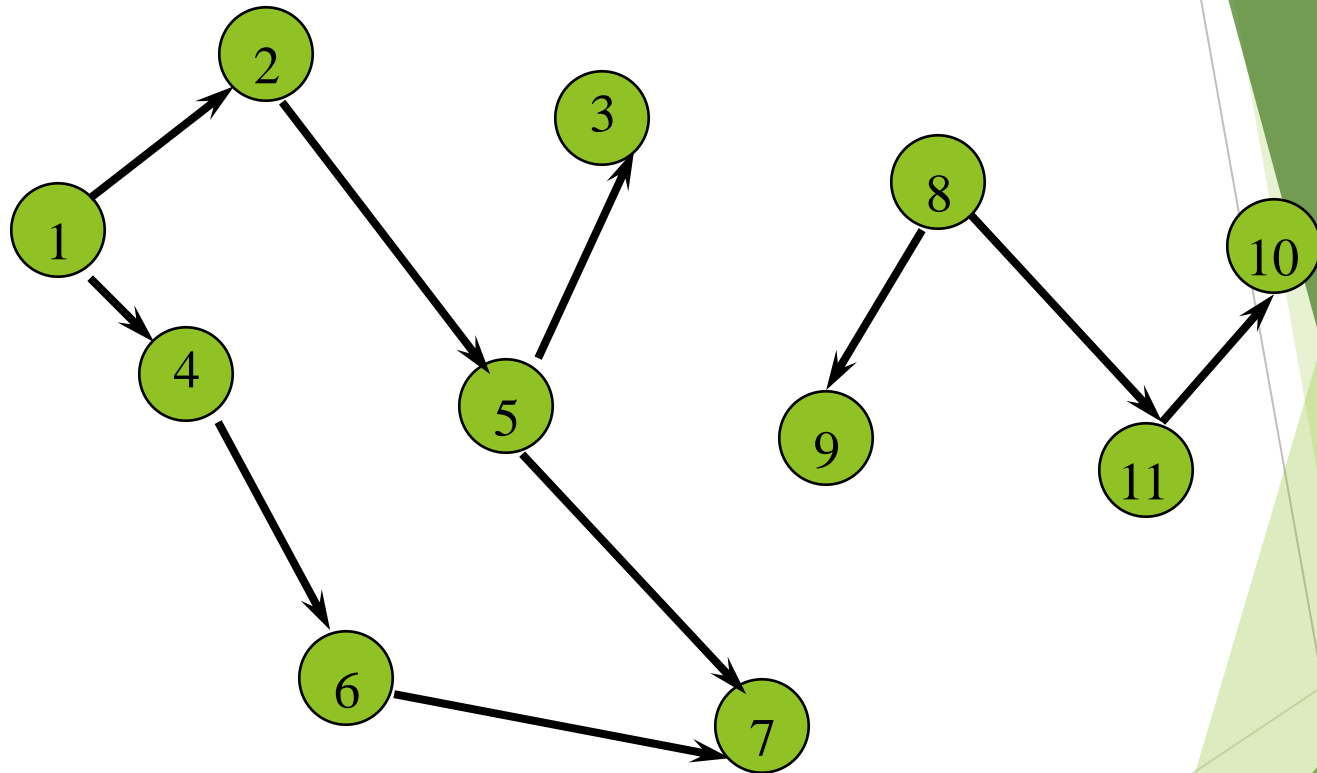
<https://www.chegg.com/homework-help/questions-and-answers/1-consider-directed-graph-weighted-edges-represent-graph-using-adjacency-list-adjacency-ma-q20254763>

In-Degree



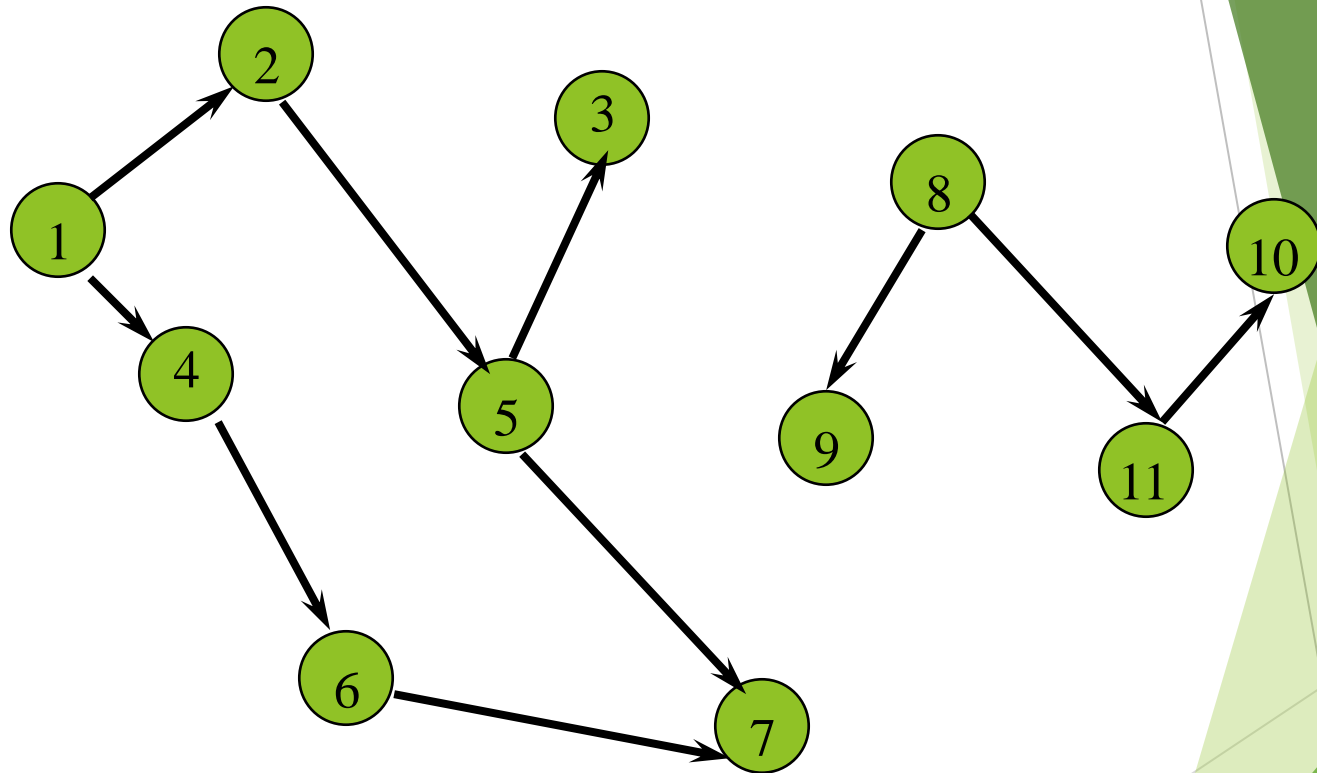
- ▶ in-degree is number of incoming edges
- ▶ indegree _____

Out-Degree



- ▶ out-degree is number of outbound edges
- ▶ $\text{outdegree}(2) = 1$, $\text{outdegree}(8) = 2$

Degree

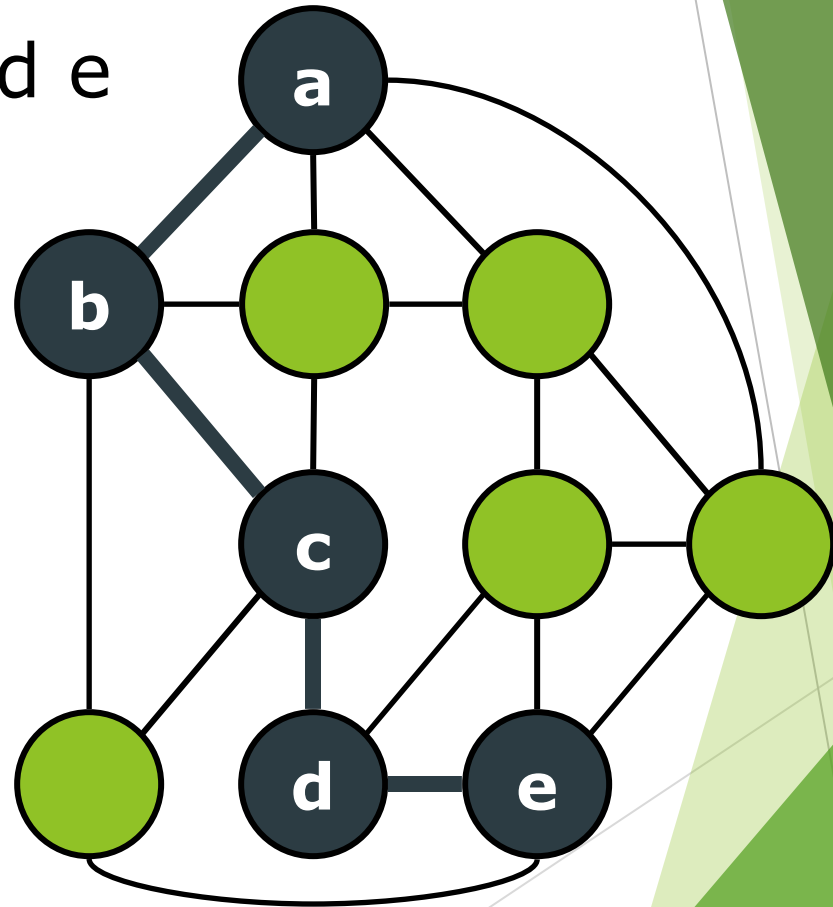


► Degree (2) = _____

► Degree (5) = _____

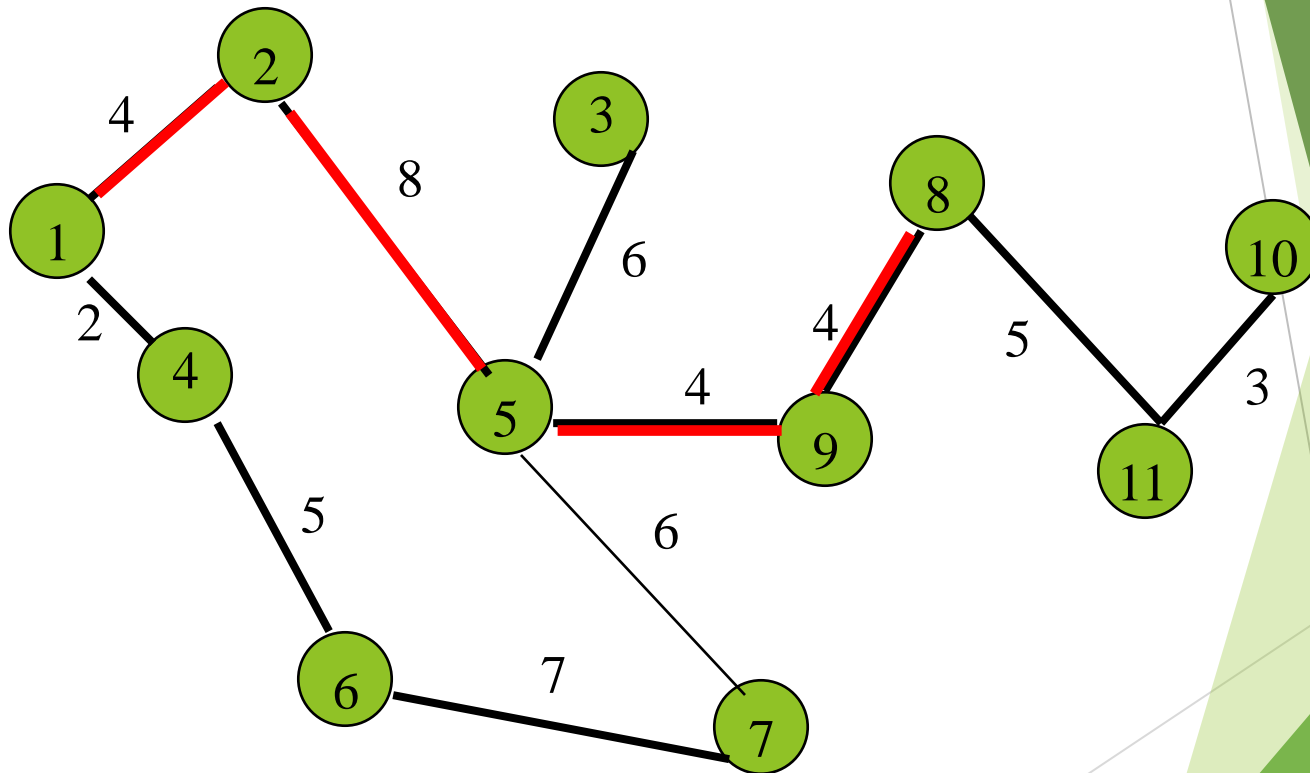
Path

- Path between a and e
- Length =



Path Finding

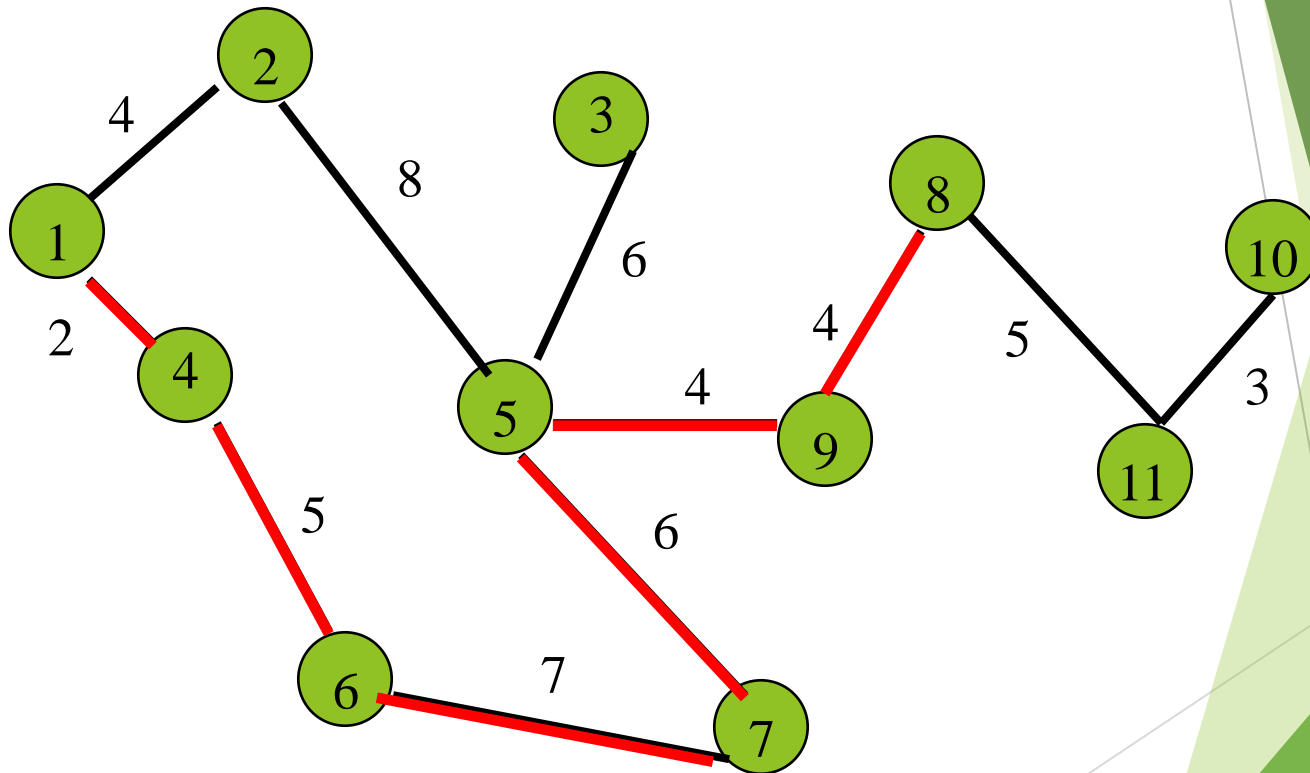
- Path between 1 and 8.



Path length

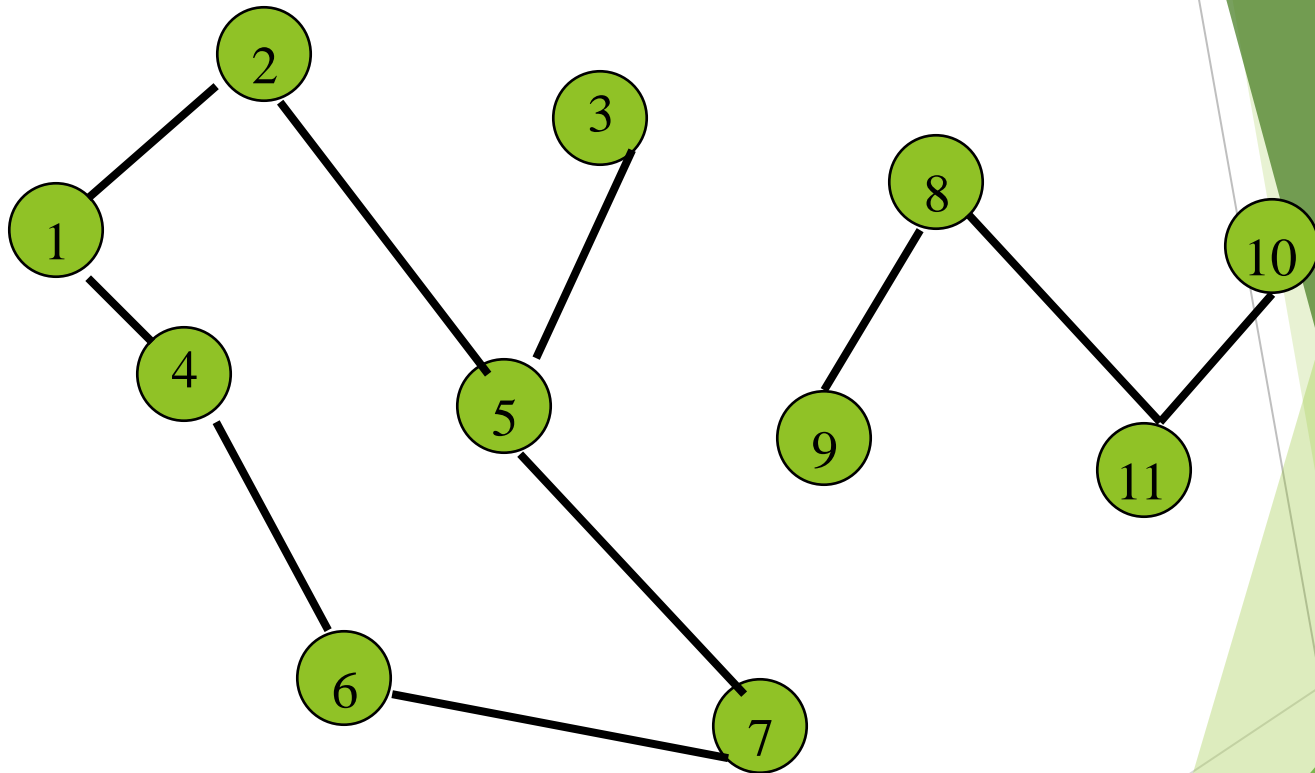
Path Finding

- ▶ Path between 1 and 8.



Path length is _____

Path Finding

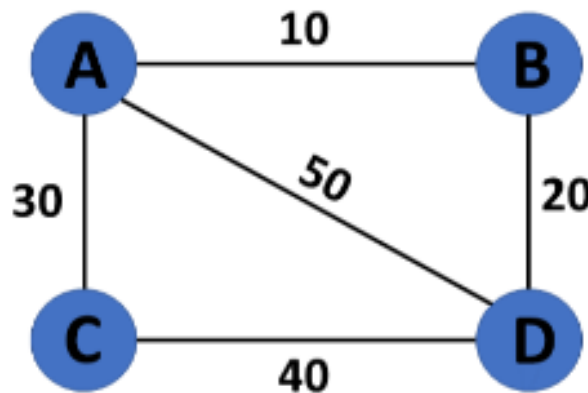


No path between 2 and 9.

Weighted Graph

Weighted Graph

- ▶ A graph $G = (V, E)$ is called a **labeled or weighted** graph because each edge has a value or weight representing the cost of traversing that edge.



<https://www.simplilearn.com/tutorials/data-structure-tutorial/graphs-in-data-structure>

Formally

A weighted graph $G =$ _____

where

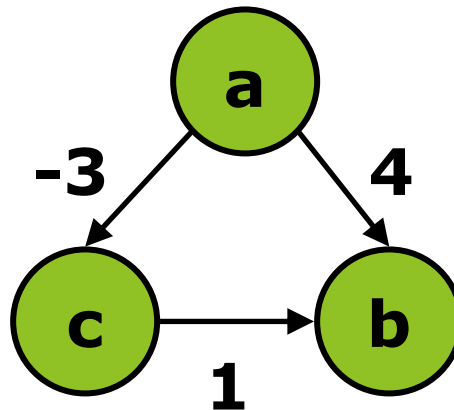
- ▶ V is the set of vertices
- ▶ E is the set of edges
- ▶ W is the weight function

Example

$V = \{ a, b, c \}$

$E = \{ (a,b), \text{_____} \}$

$W = \{ ((a,b), 4), \text{_____} \}$



Adjacent Vertices

- ▶ $\text{adj}(v)$ = set of vertices adjacent to v

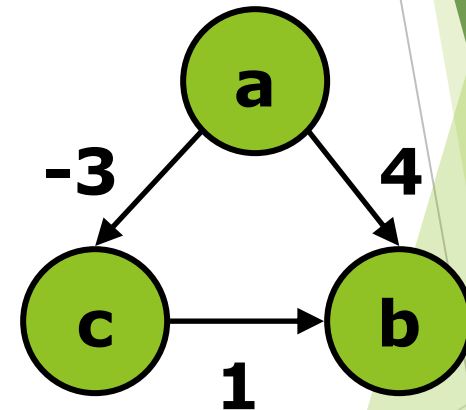
$\text{adj}(a) = \{\underline{\hspace{2cm}}\}$

$\text{adj}(b) = \{\underline{\hspace{2cm}}\}$

$\text{adj}(c) = \{\underline{\hspace{2cm}}\}$

- ▶ $\sum_v |\text{adj}(v)| = |E|$

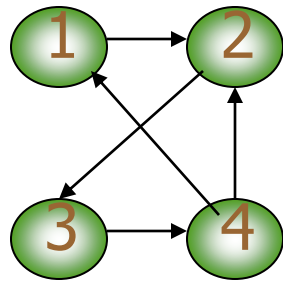
- ▶ $\text{adj}(v)$: Neighbours of v



Graph Representation

1. Adjacency Matrix
2. Incidence Matrix
3. Adjacency Lists
 1. Array Adjacency Lists
 2. Linked Adjacency Lists

Adjacency Matrix

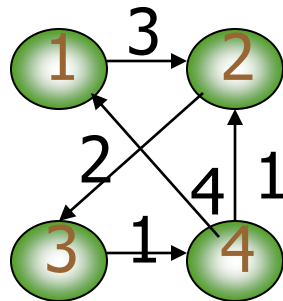


graph

	1	2	3	4
1		T		
2			T	
3				T
4	T	T		

OR

	1	2	3	4
1		1		
2			1	
3				1
4	1	1		



graph

	1	2	3	4
1		3		
2			2	
3				1
4	4	1		

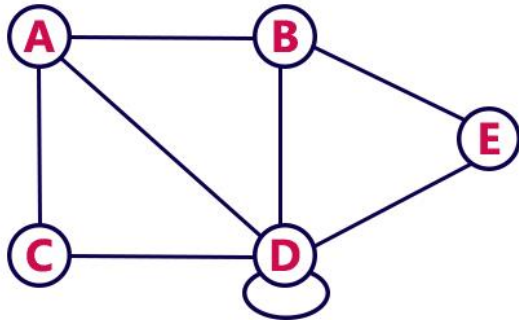
Ideal view

01418231: Data Structure

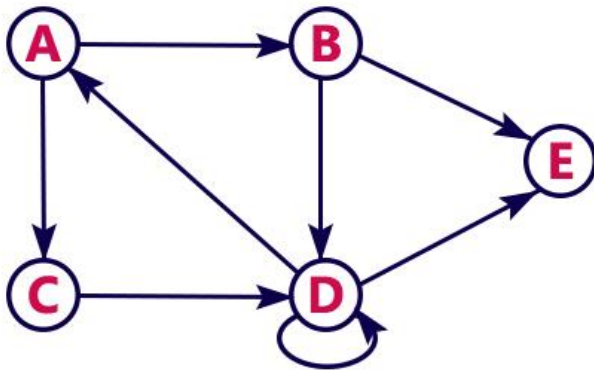
Computerized view

Powered By: Jirawan Charoensuk

Adjacency Matrix



	A	B	C	D	E
A	0	1	1	1	0
B	1	0	0	1	1
C	1	0	0	1	0
D	1	1	1	1	1
E	0	1	0	1	0

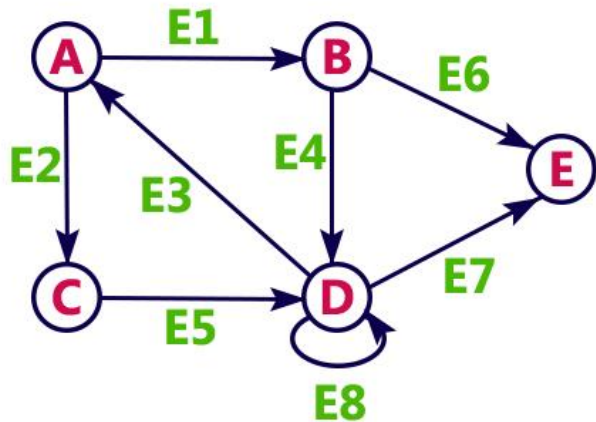


	A	B	C	D	E
A	0	1	1	0	0
B	0	0	0	1	1
C	0	0	0	1	0
D	1	0	0	1	1
E	0	0	0	0	0

Incidence Matrix

- ▶ graph can be represented using a matrix of size total number of vertices by total number of edges.
- ▶ That means if a graph with 4 vertices and 6 edges can be represented using a matrix of 4X6 class.
- ▶ In this matrix, rows represents vertices and columns represents edges.
 - ▶ This matrix is filled with either 0 or 1 or -1.
 - ▶ 0 represents row edge is not connected to column vertex,
 - ▶ 1 represents row edge is connected as outgoing edge to column vertex
 - ▶ -1 represents row edge is connected as incoming edge to column vertex.

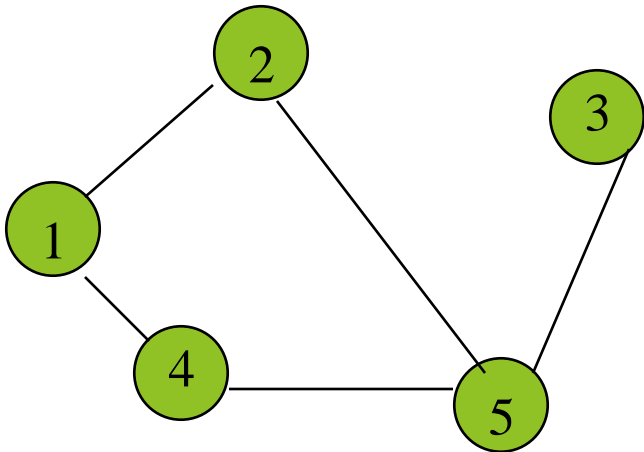
Incidence Matrix



	E1	E2	E3	E4	E5	E6	E7	E8
A	1	1	-1	0	0	0	0	0
B	-1	0	0	1	0	1	0	0
C	0	-1	0	0	1	0	0	0
D	0	0	1	-1	-1	0	1	1
E	0	0	0	0	0	-1	-1	0

Array Adjacency Lists

- ▶ Adjacency list for vertex i is a linear list of vertices adjacent from vertex i .
- ▶ An array of n adjacency lists.



aList[1] = _____

aList[2] = _____

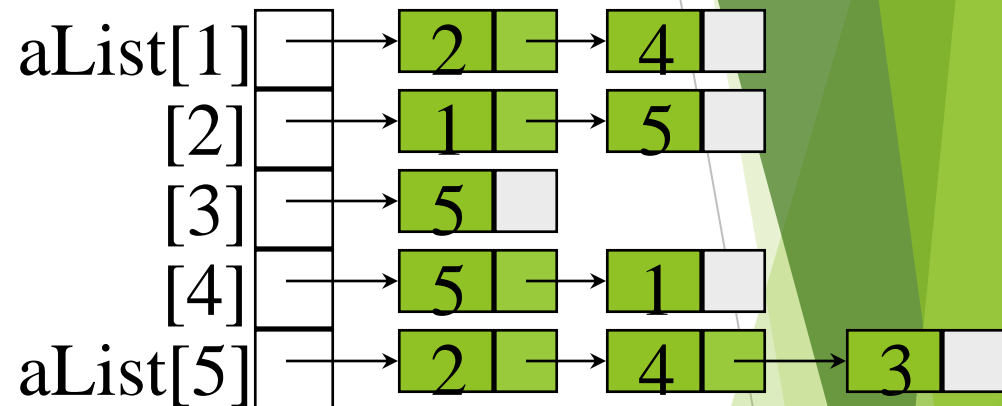
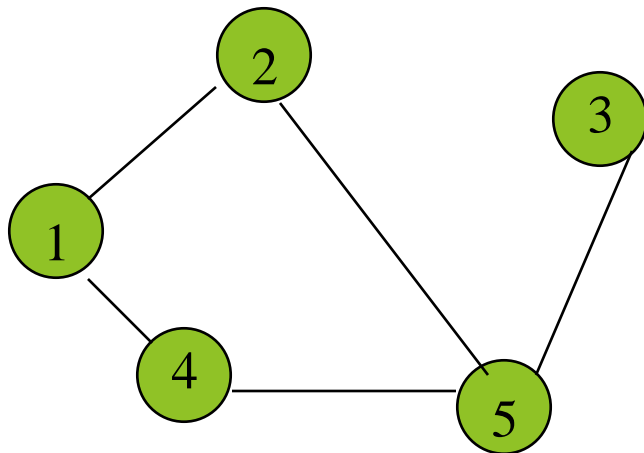
aList[3] = _____

aList[4] = _____

aList[5] = _____

Linked Adjacency Lists

- Each adjacency list is a chain.



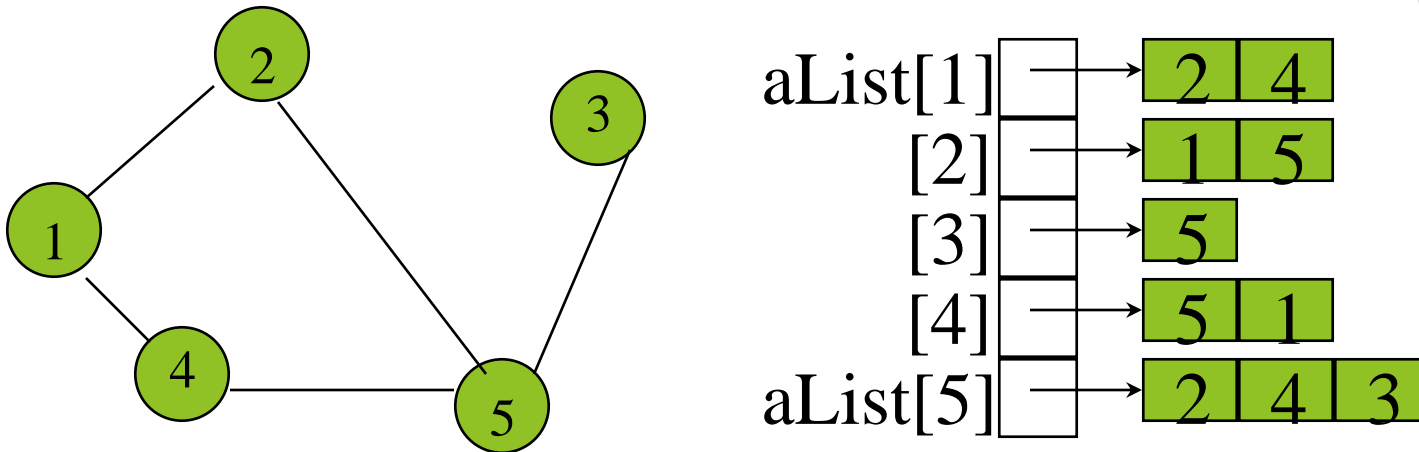
Array Length = n

of chain nodes = $2e$ (undirected graph)

of chain nodes = e (digraph)

Array Adjacency Lists

- Each adjacency list is an array list.



Array Length = n

of list elements = $2e$ (undirected graph)

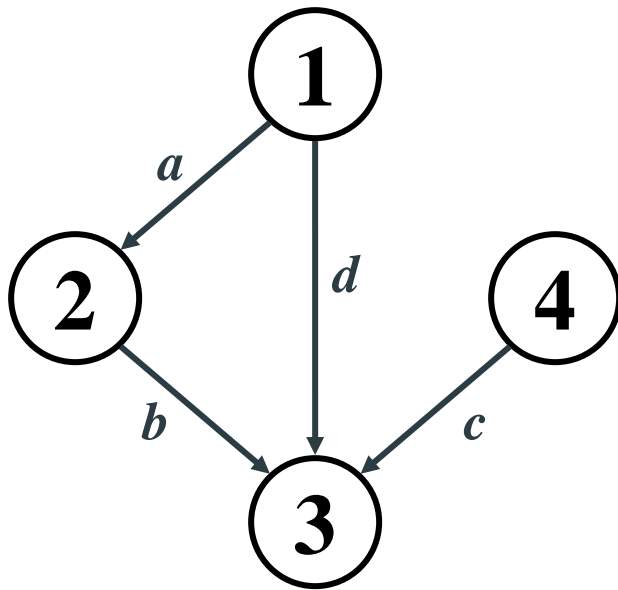
of list elements = e (digraph)

The background features abstract, overlapping green geometric shapes, primarily triangles and polygons, in various shades of green, creating a modern, layered effect. The word "Question?" is centered in a green, sans-serif font.

Question?

Graphs: Adjacency Matrix

► Example:

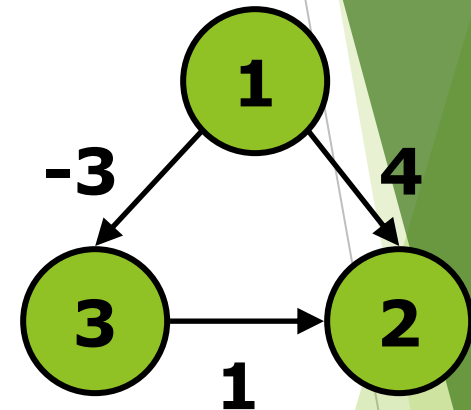


A	1	2	3	4
1				
2				
3			??	
4				

Adjacency Matrix (weight)

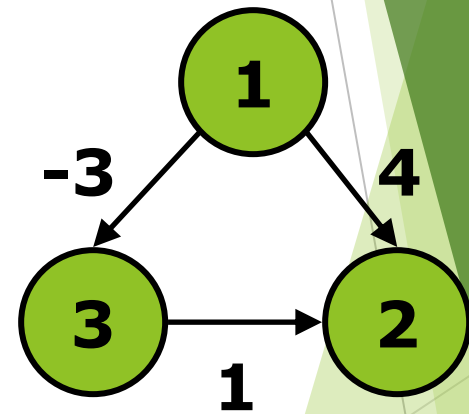
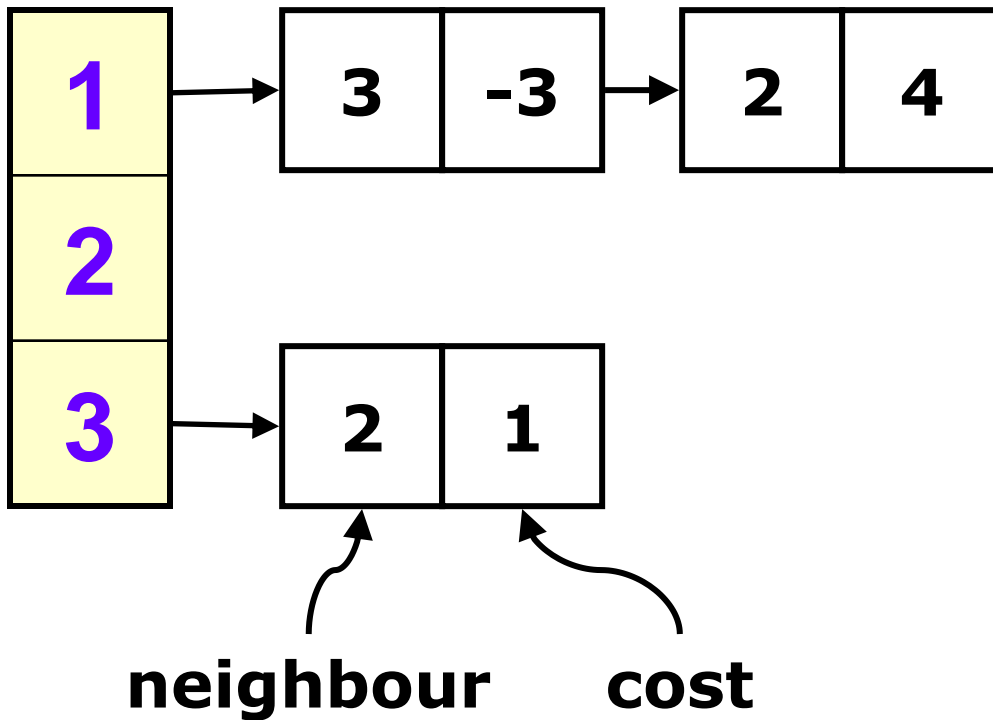
```
double vertex[][];
```

	1	2	3
1			
2			
3			



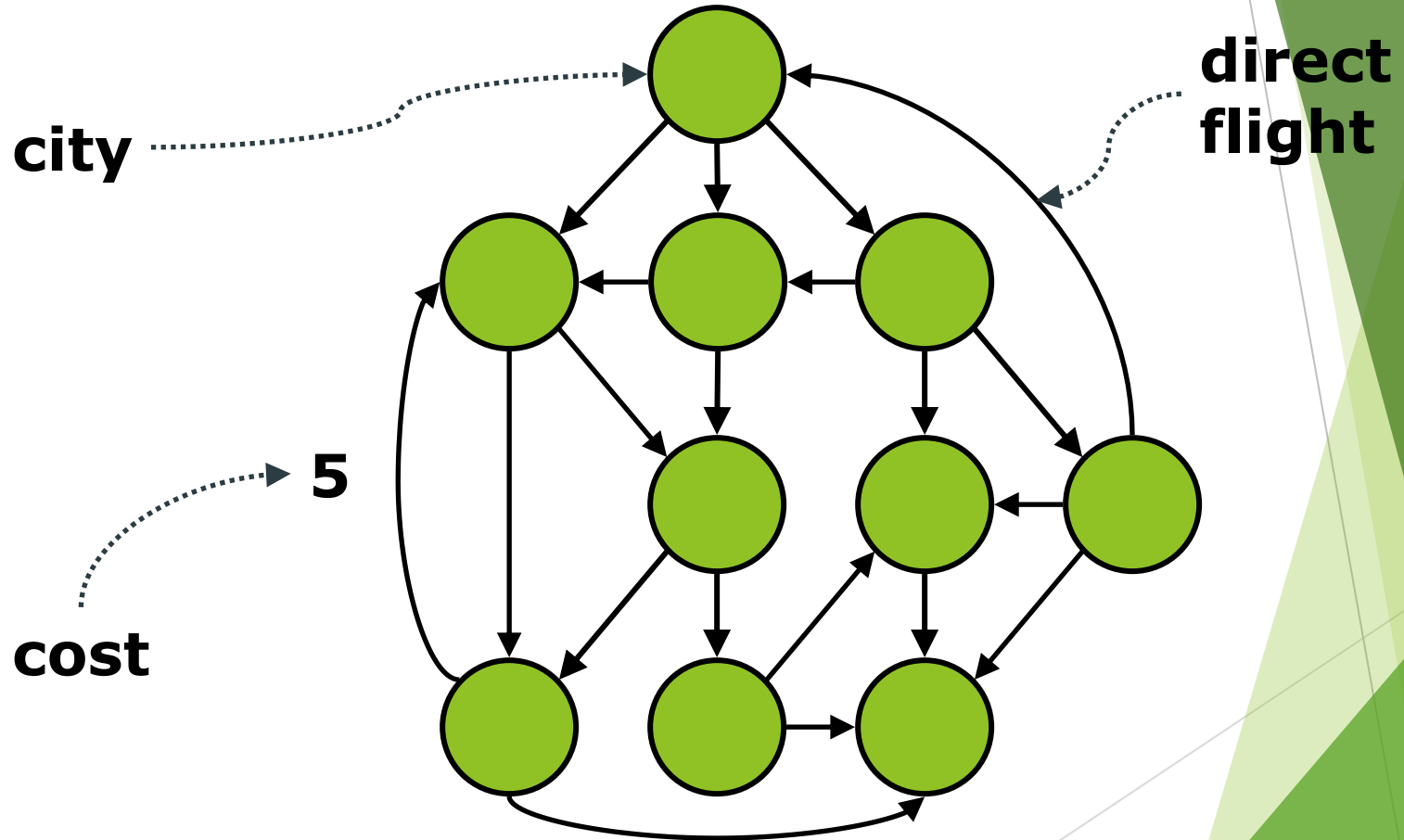
Adjacency List (weight)

EdgeList vertex[];



Applications

Travel Planning



Question

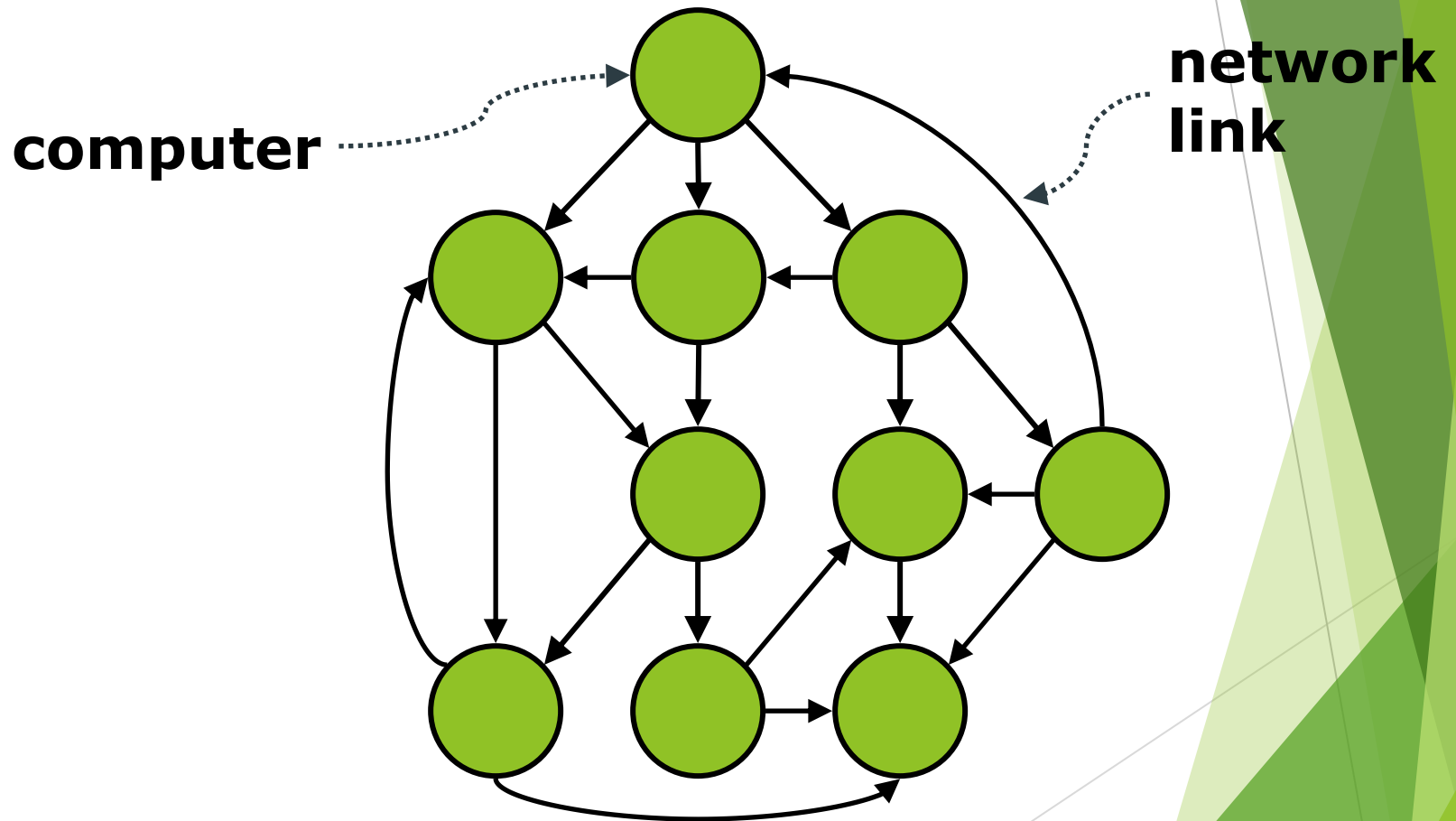
- ▶ What is the shortest way to travel between A and B?

“SHORTEST PATH PROBLEM”

- ▶ How to minimize the cost of visiting n cities such that we visit each city exactly once, and finishing at the city where we start from?

“TRAVELING SALESMAN PROBLEM”

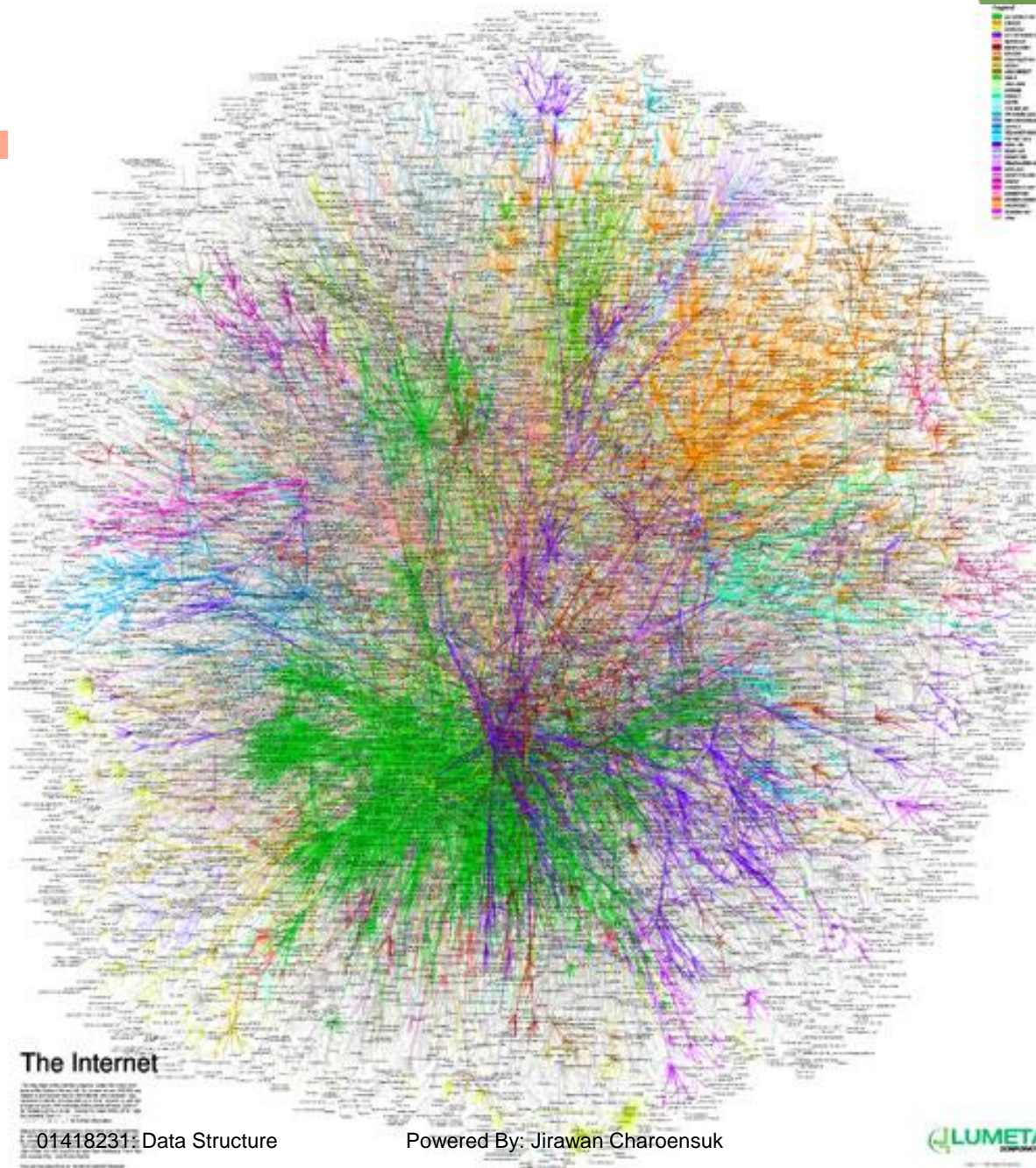
Internet



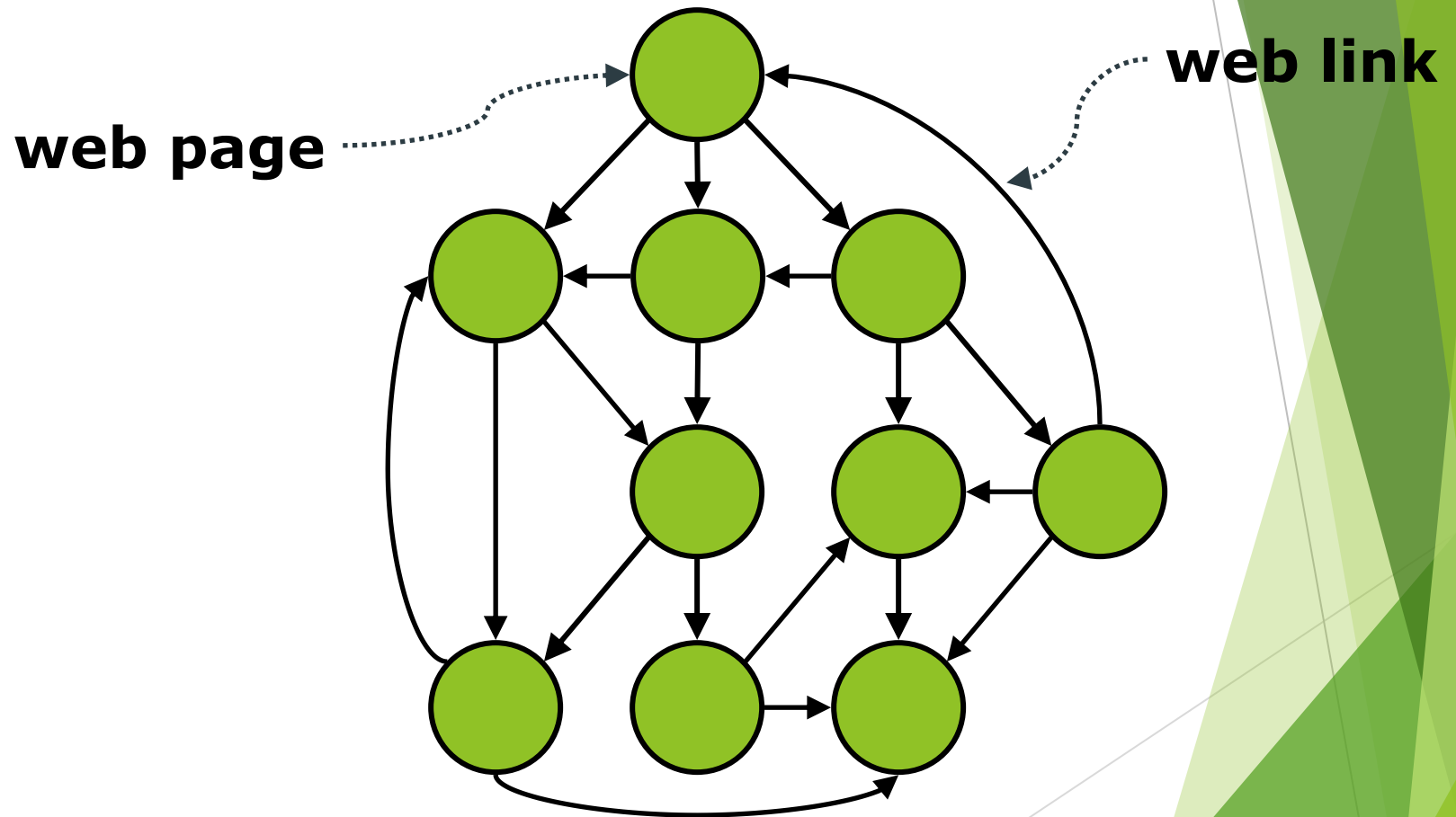
Question

- ▶ What is the shortest route to send a packet from A to B?

“SHORTEST PATH PROBLEM”



The Web

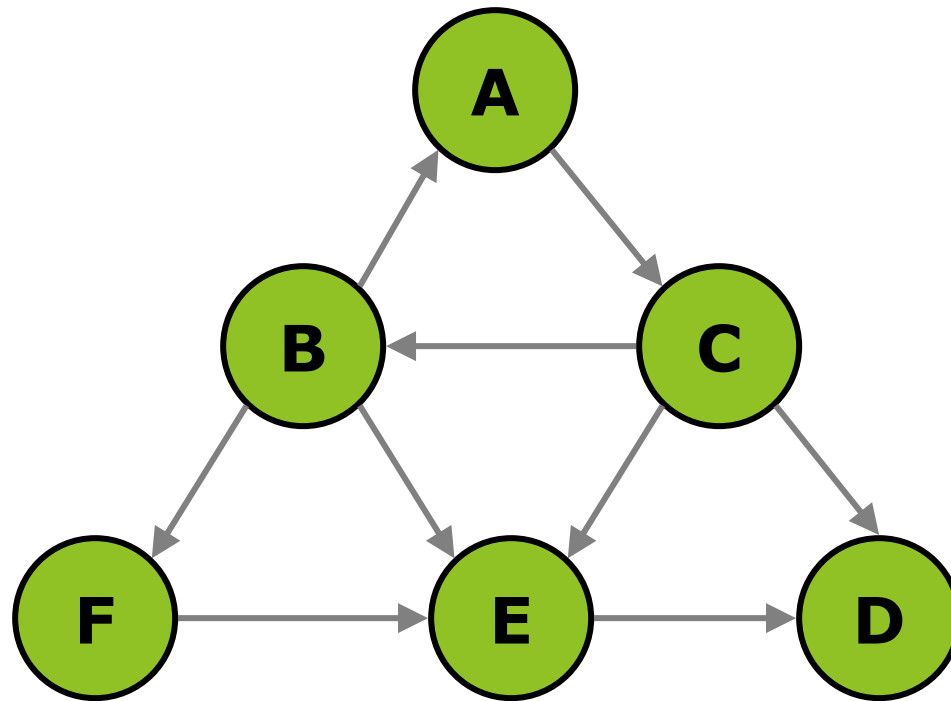


Question

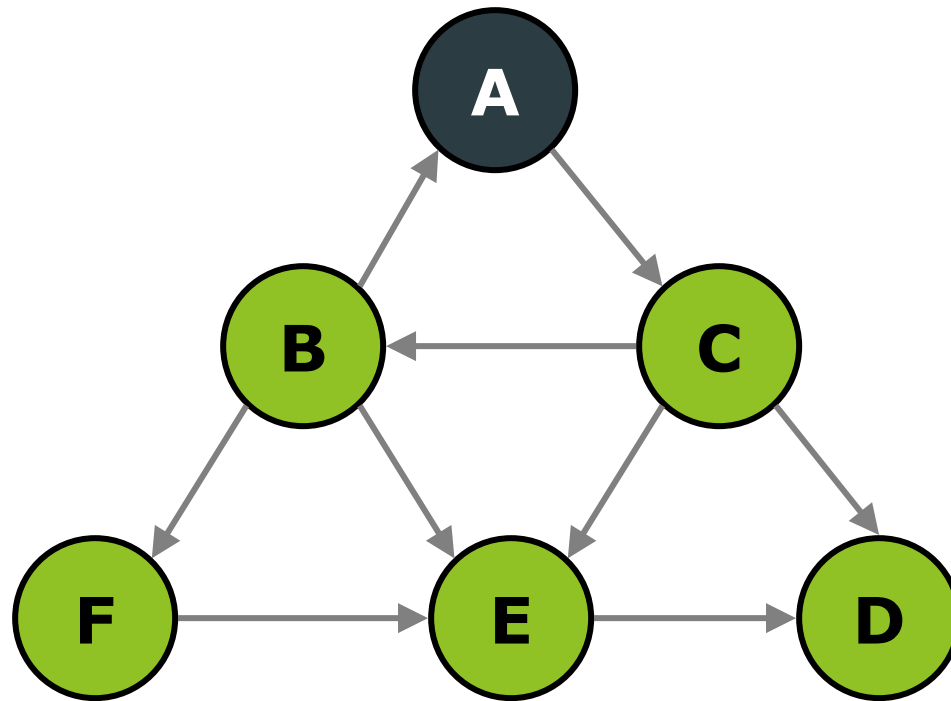
- ▶ Which web pages are important?
- ▶ Which set of web pages are likely to be of the same topic?

Breadth-First Search

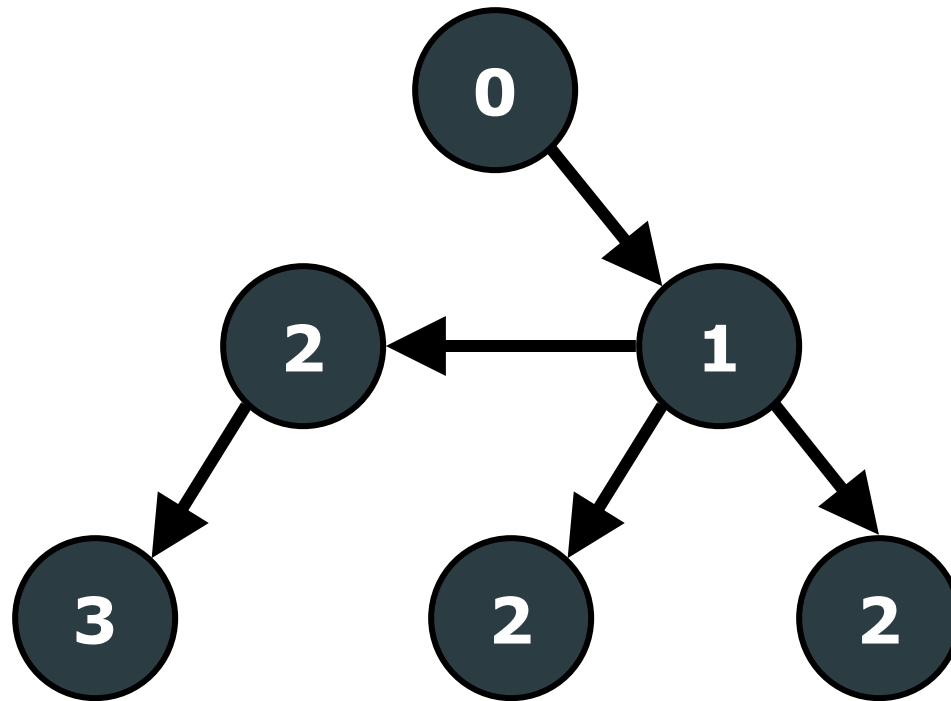
Breadth-First Search



Breadth-First Search



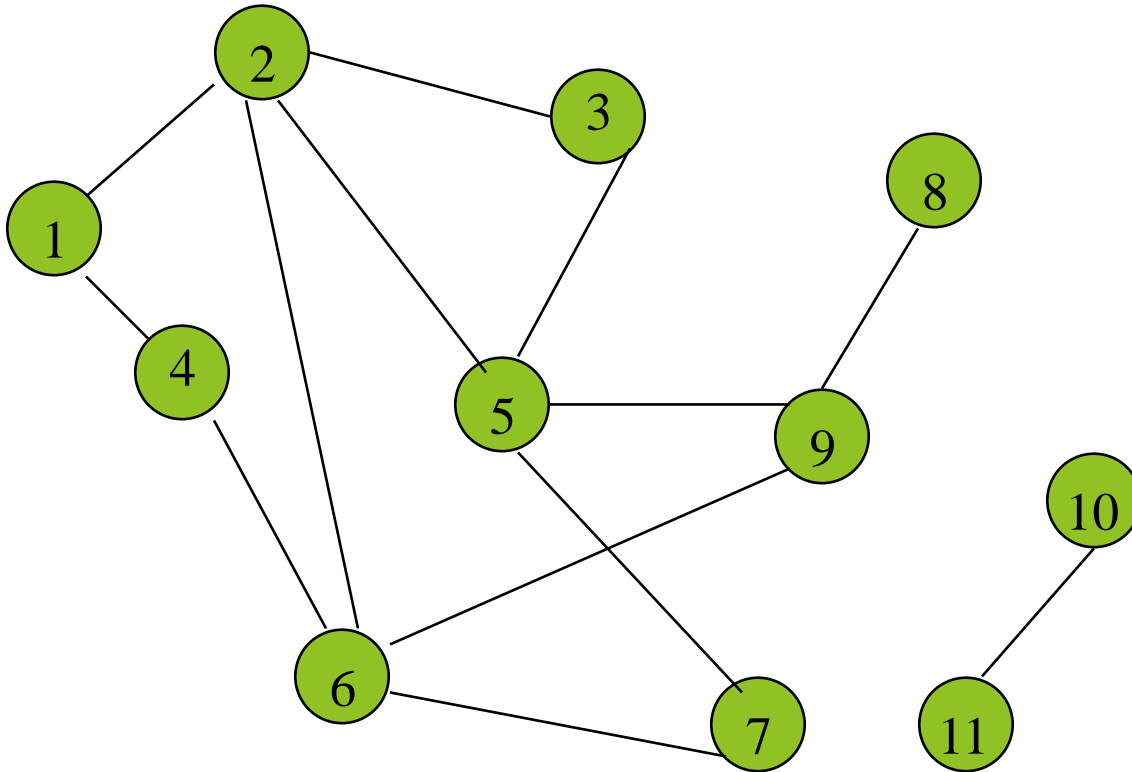
Breadth-First Search



Breadth-First Search

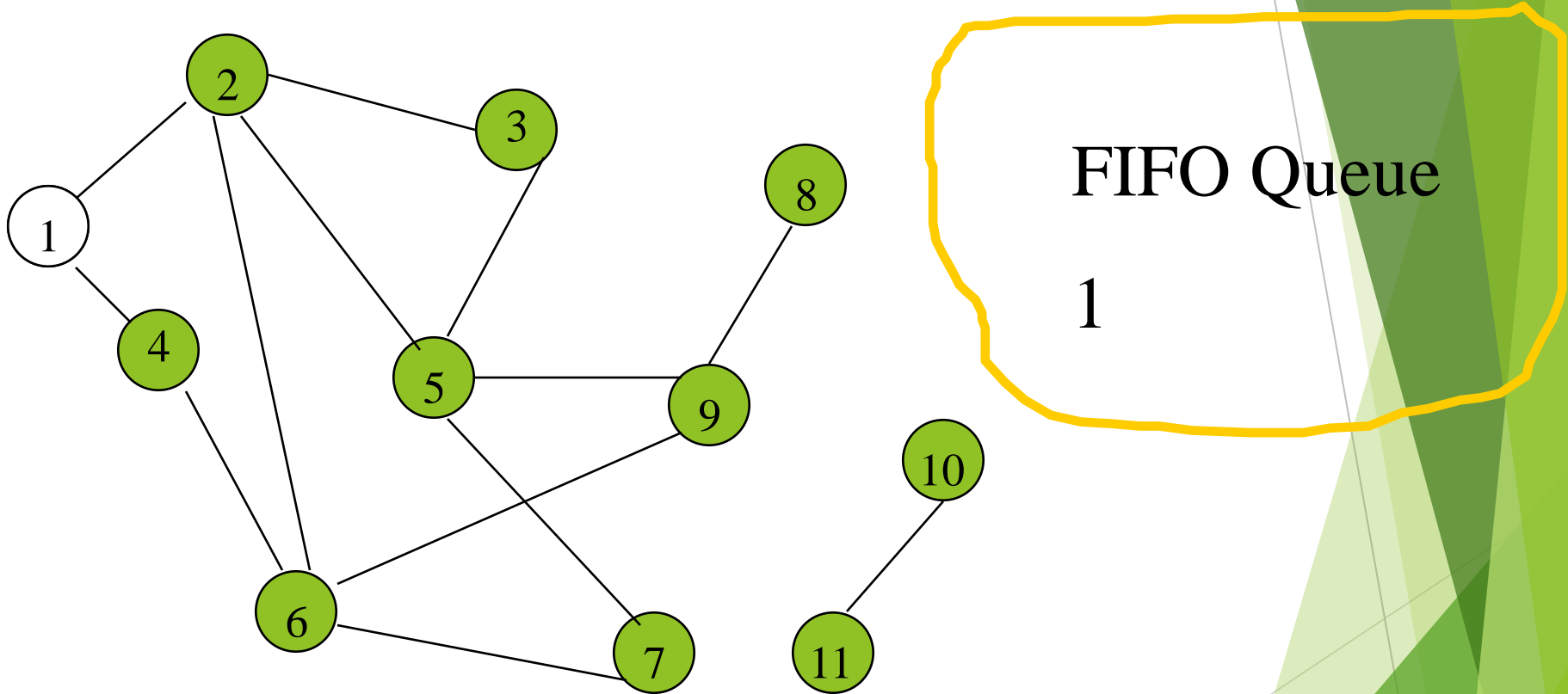
Example 2

Breadth-First Search



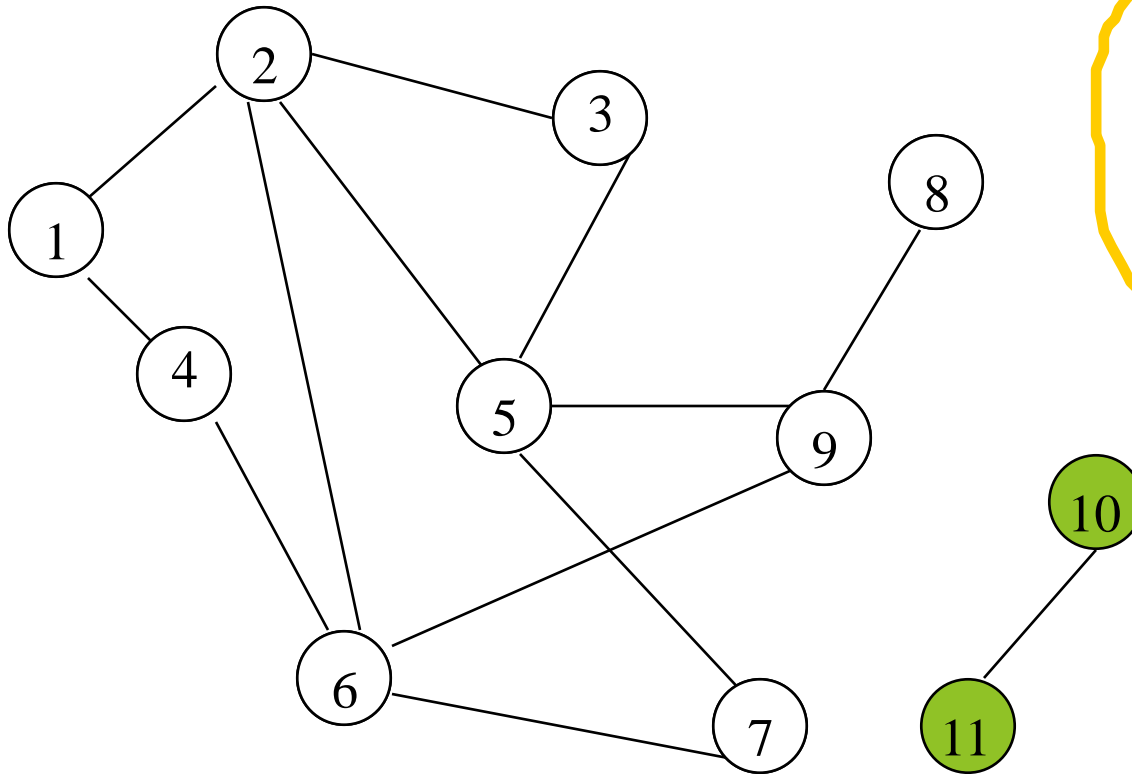
Start search at vertex 1.

Breadth-First Search



Mark/label start vertex and put in a FIFO queue.

Breadth-First Search

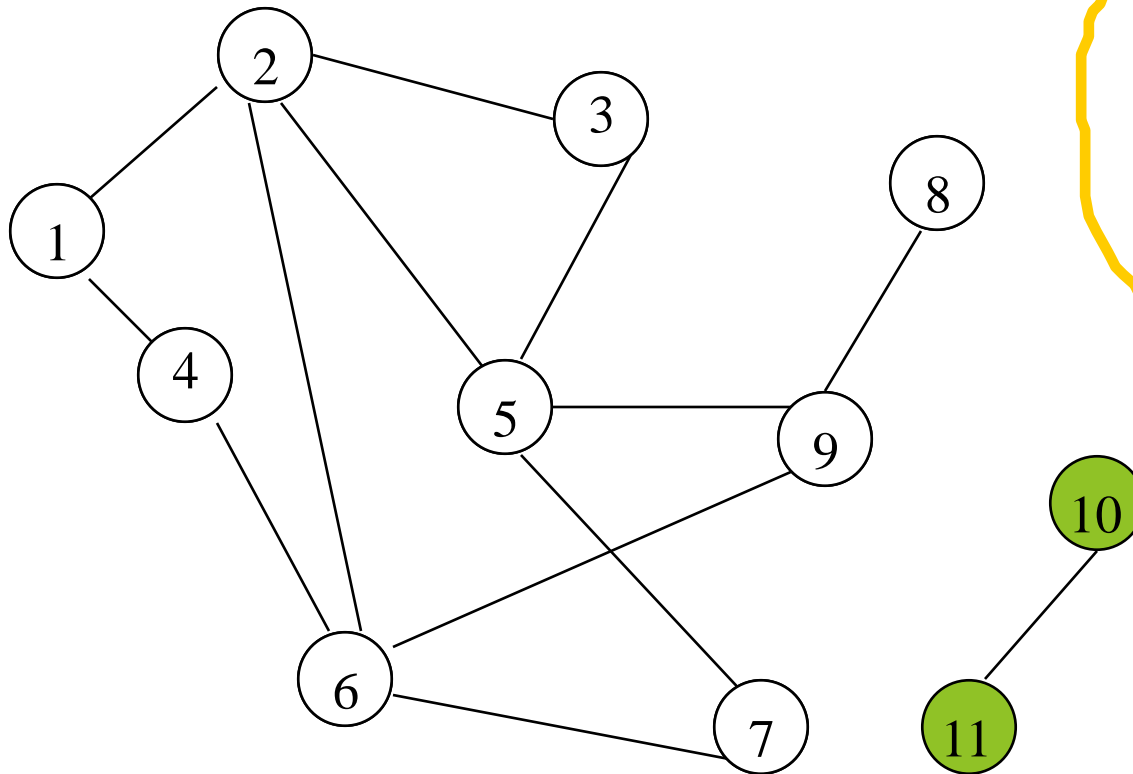


FIFO Queue

8

Remove 8 from Q; visit adjacent unvisited vertices;
put in Q.

Breadth-First Search



FIFO Queue

Queue is empty. Search terminates.

Breadth-First Search

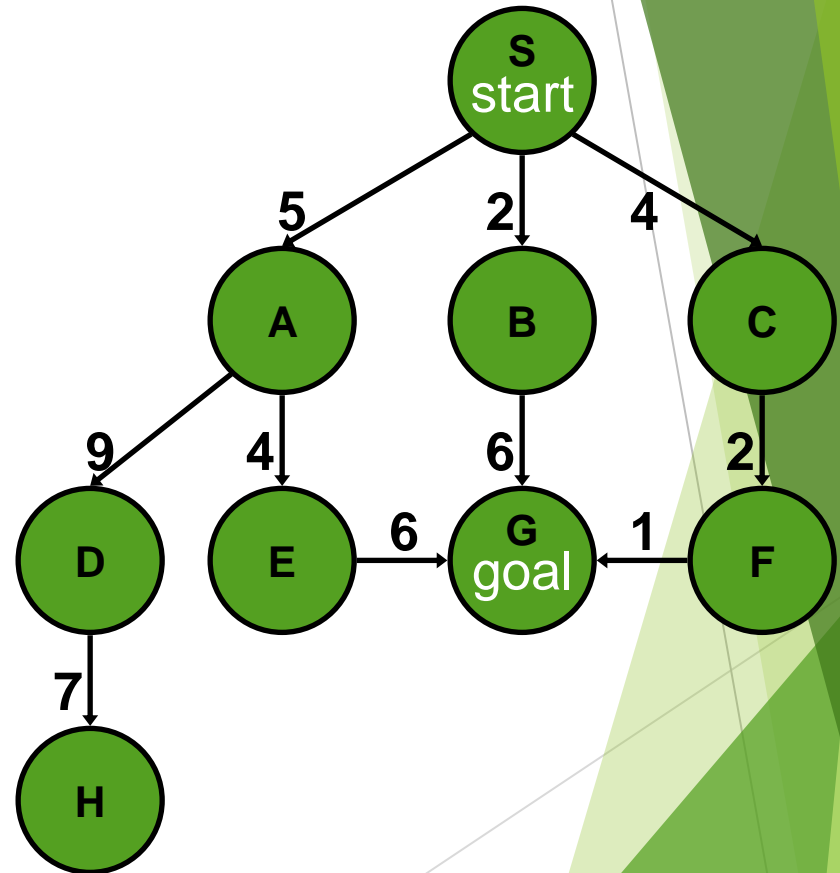
Example 3

Breadth-First Search (BFS)

generalSearch(problem, queue)

of nodes tested: 0, expanded: 0

current	nodes list
	{S}

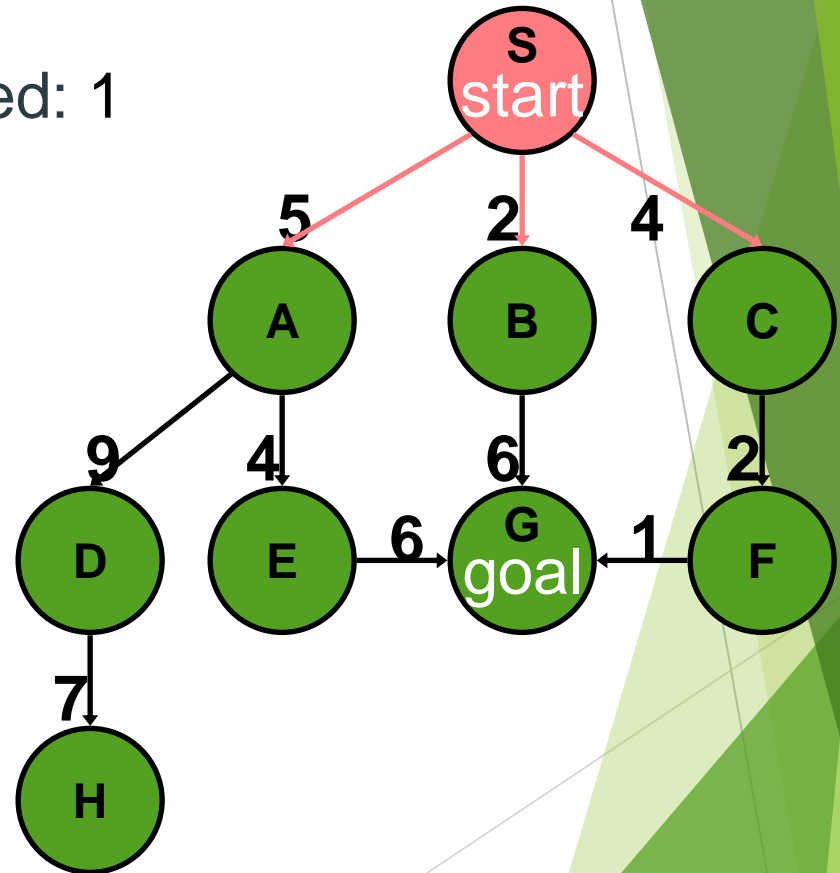


Breadth-First Search (BFS)

`generalSearch(problem, queue)`

of nodes tested: 1, expanded: 1

current	nodes list
	{S}
S not goal	{A,B,C}

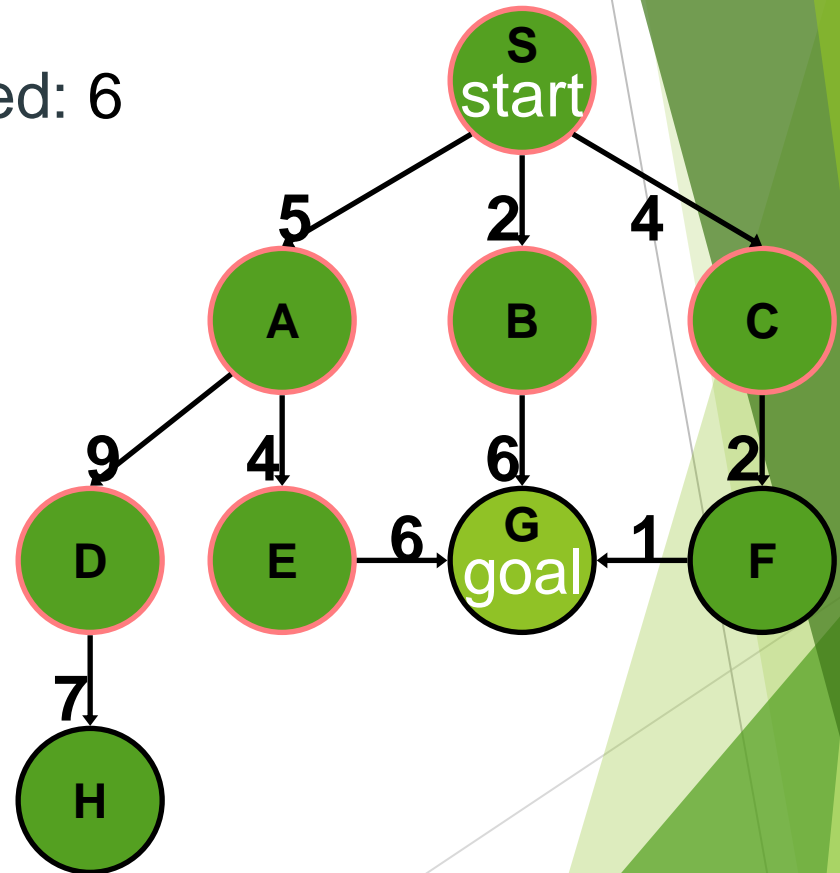


Breadth-First Search (BFS)

`generalSearch(problem, queue)`

of nodes tested: 7, expanded: 6

current	nodes list
	{S}
S	{A,B,C}
A	{B,C,D,E}
B	{C,D,E,G}
C	{D,E,G,F}
D	{E,G,F,H}
E	{G,F,H}
G goal	{F,H} no expand

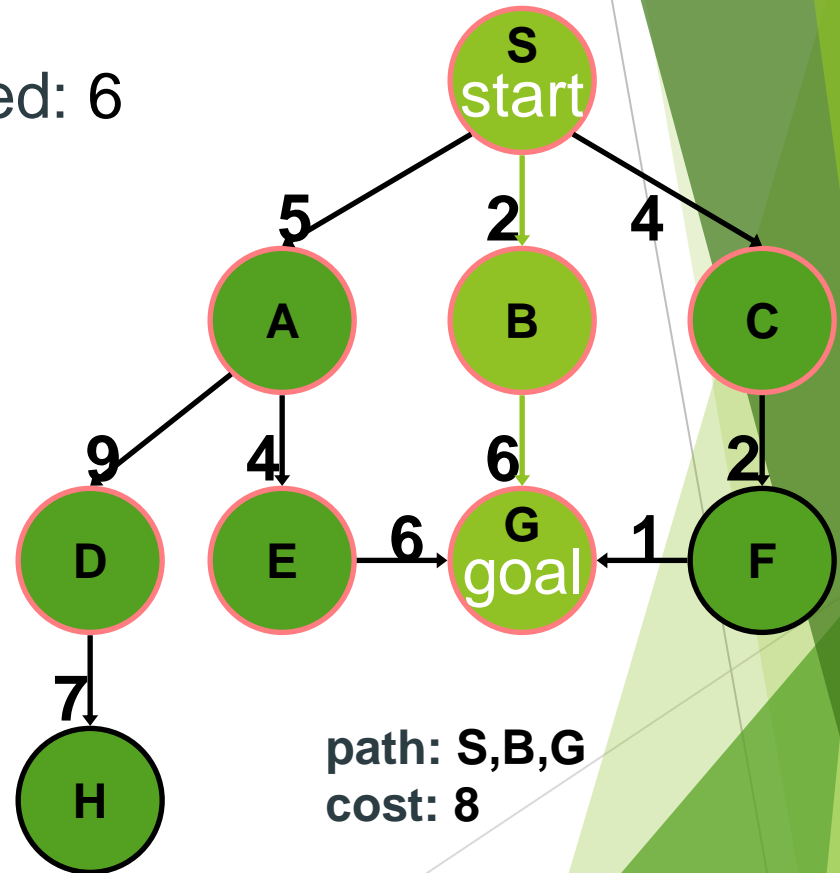


Breadth-First Search (BFS)

`generalSearch(problem, queue)`

of nodes tested: 7, expanded: 6

current	nodes list
	{S}
S	{A,B,C}
A	{B,C,D,E}
B	{C,D,E,G}
C	{D,E,G,F}
D	{E,G,F,H}
E	{G,F,H,G}
G	{F,H,G}



Breadth-First Search

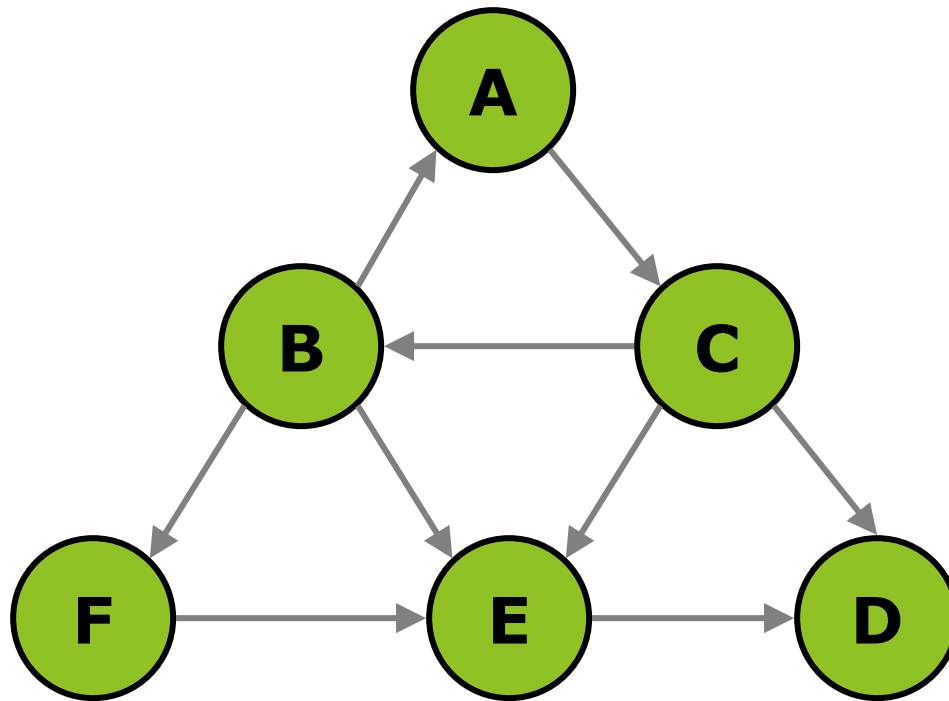
- ▶ “Explore” a graph, turning it into a tree
 - ▶ One vertex at a time
 - ▶ Expand frontier of explored vertices across the *breadth* of the frontier
- ▶ Builds a tree over the graph
 - ▶ Pick a *source vertex* to be the root
 - ▶ Find (“discover”) its children, then their children, etc.

Breadth-First Search

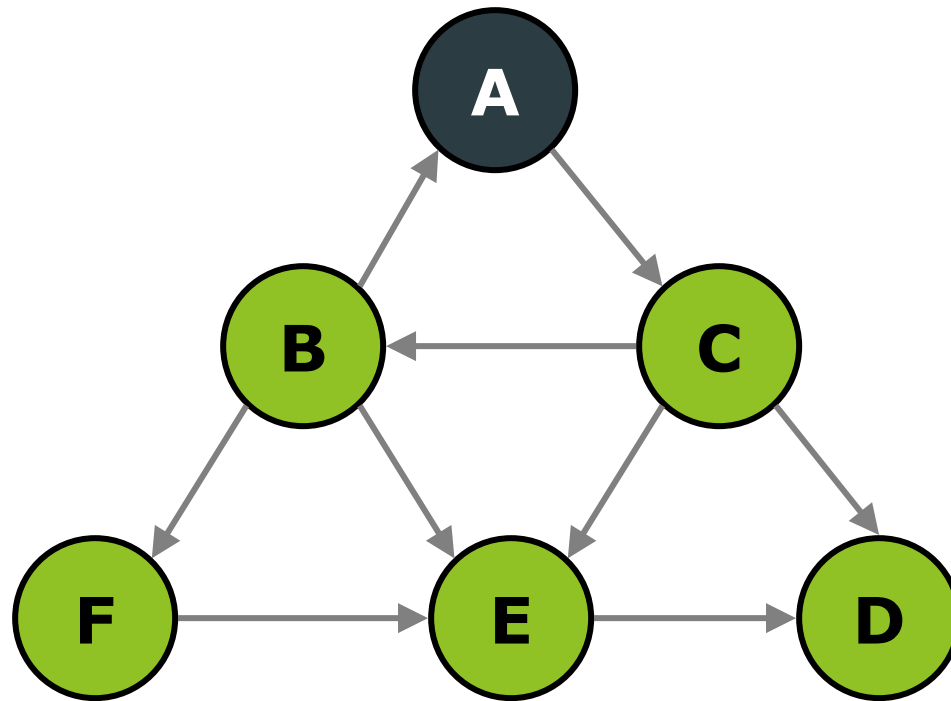
- ▶ Visit start vertex and put into a FIFO queue.
- ▶ Repeatedly remove a vertex from the queue, visit its unvisited adjacent vertices, put newly visited vertices into the queue.

Depth-First Search

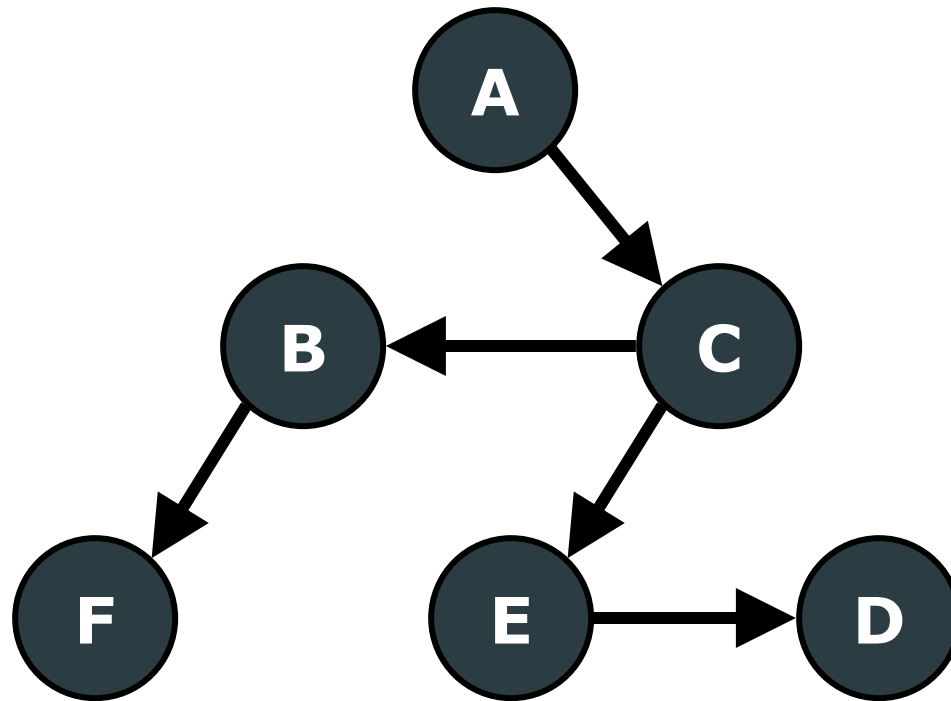
Depth-First Search



Depth-First Search



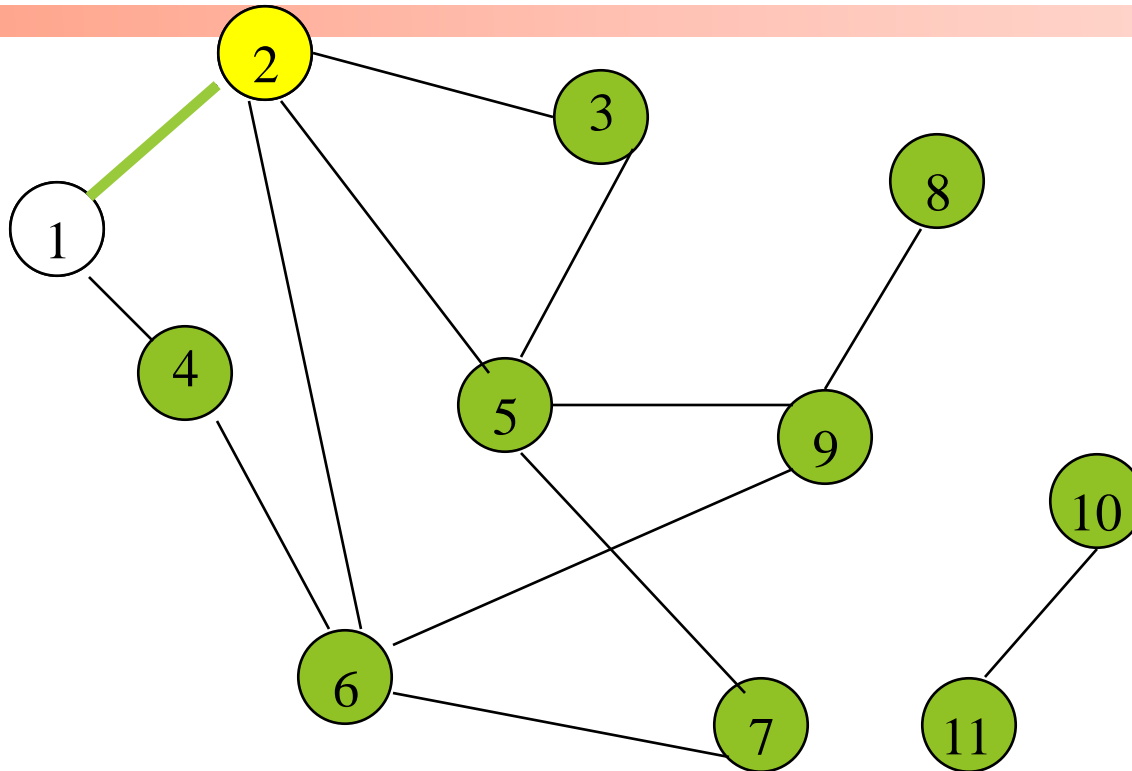
Depth-First Search



Depth-First Search

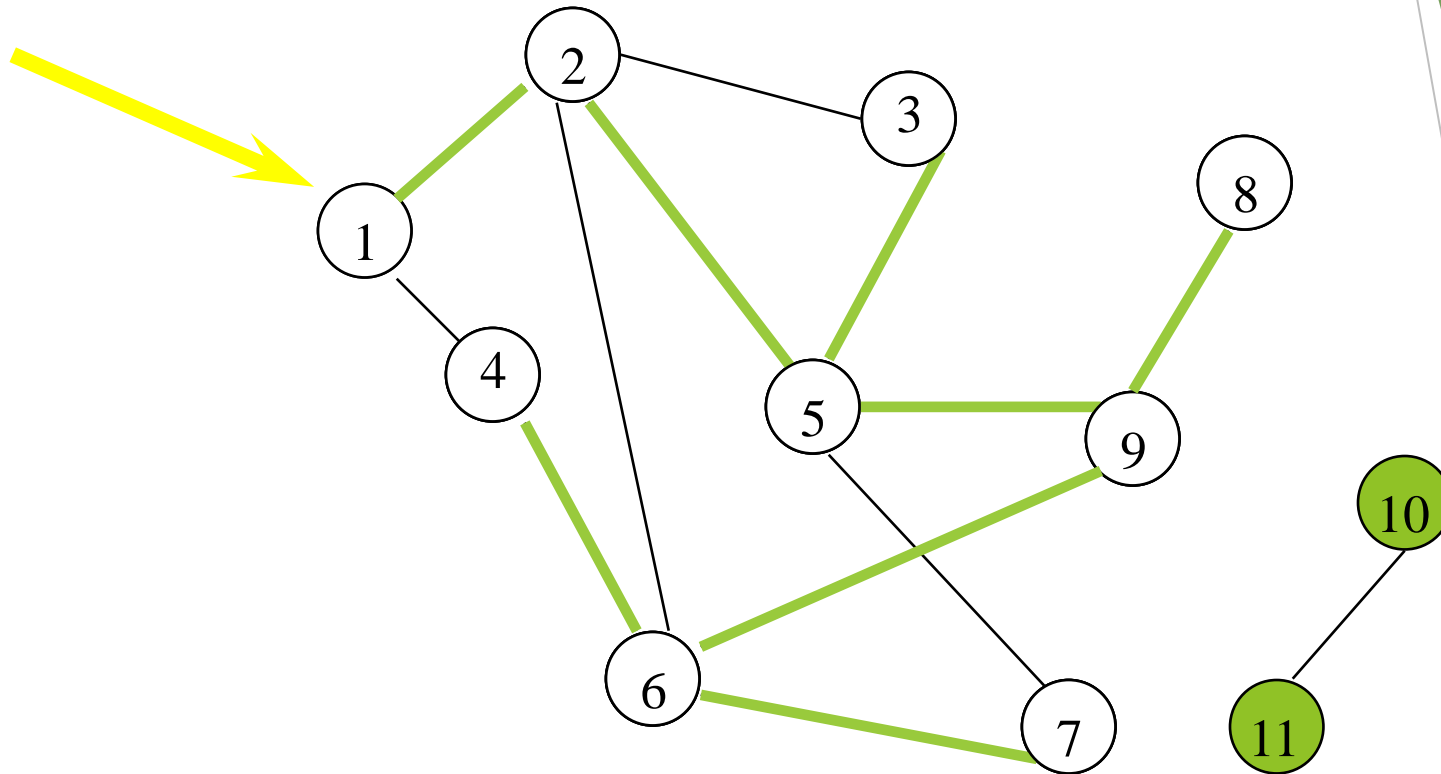
Example 2

Depth-First Search



Start search at vertex 1.
Label vertex 1 and do a depth first
search from either 2 or 4.

Depth-First Search



Return to invoking method.

Depth-First Search

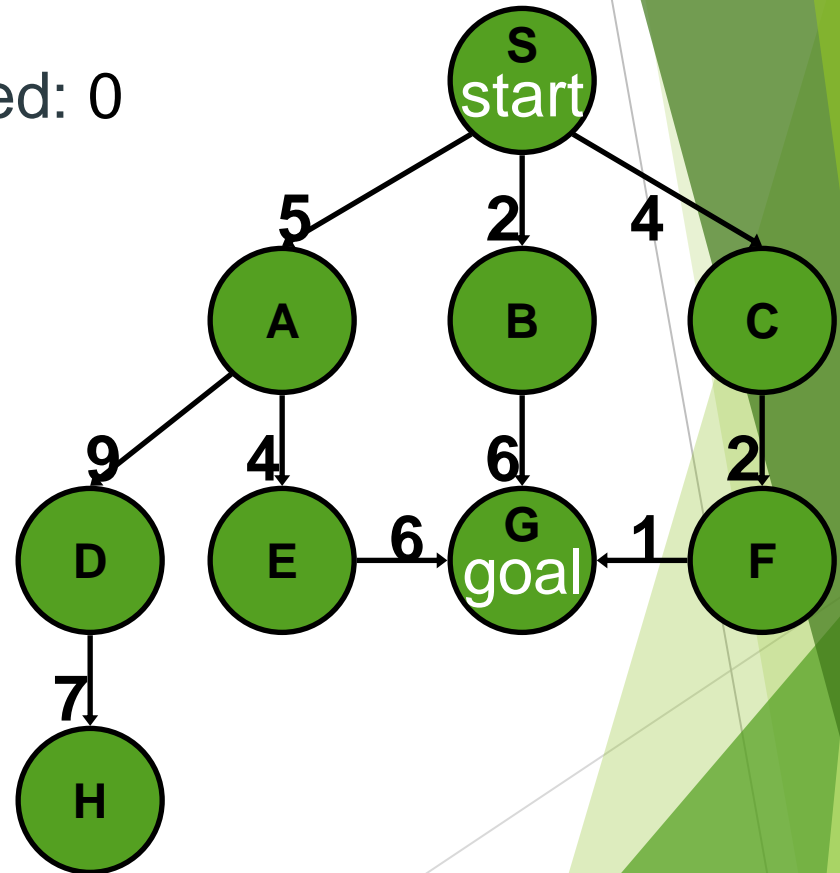
Example 3

Depth-First Search (DFS)

generalSearch(problem, stack)

of nodes tested: 0, expanded: 0

current	nodes list
	{S}

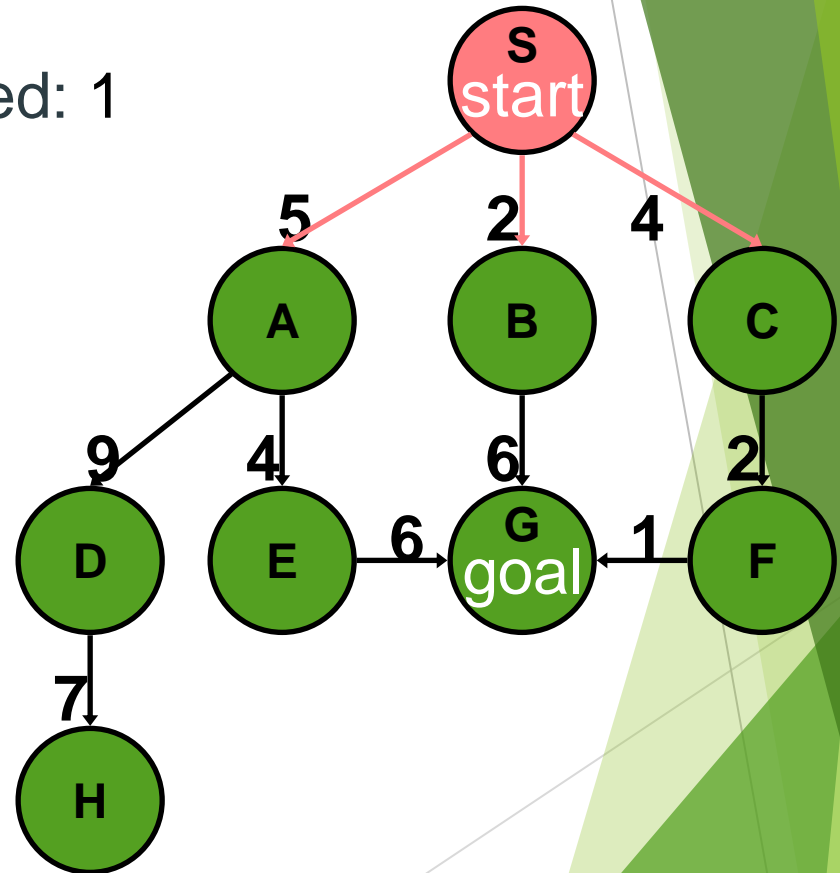


Depth-First Search (DFS)

`generalSearch(problem, stack)`

of nodes tested: 1, expanded: 1

current	nodes list
	{S}
S not goal	{A,B,C}

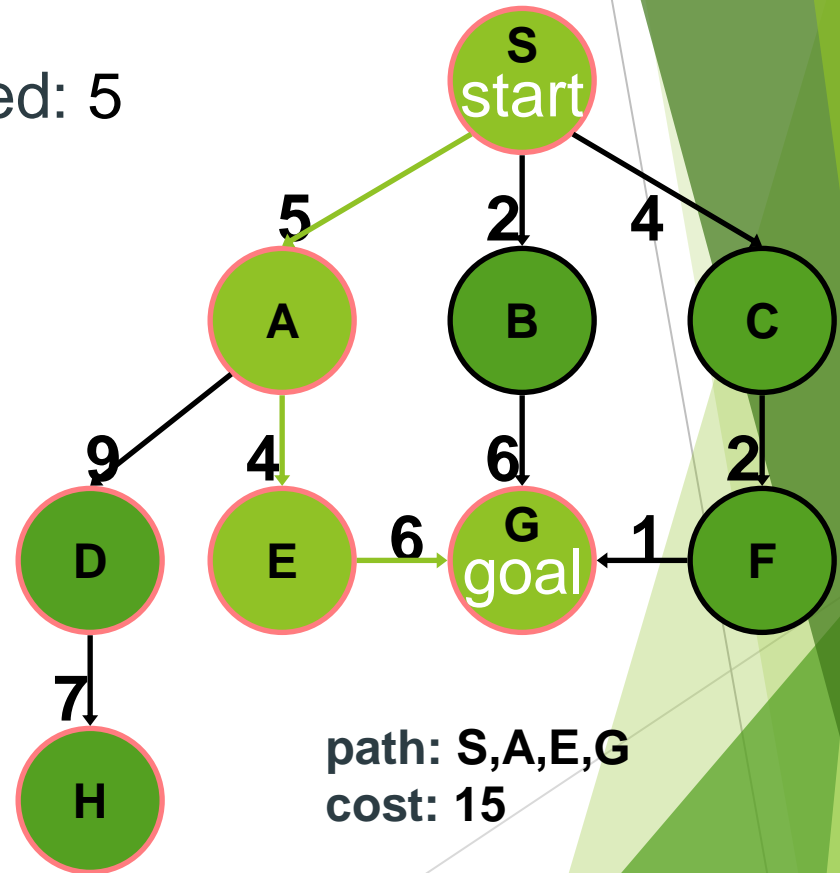


Depth-First Search (DFS)

generalSearch(problem, stack)

of nodes tested: 6, expanded: 5

current	nodes list
	{S}
S	{A,B,C}
A	{D,E,B,C}
D	{H,E,B,C}
H	{E,B,C}
E	{G,B,C}
G	{B,C}



Summary

Breadth-First Search

- ▶ Visit start vertex and put into a FIFO queue.
- ▶ Repeatedly remove a vertex from the queue, visit its unvisited adjacent vertices, put newly visited vertices into the queue.

Depth-First Search

- ▶ A depth-first search (DFS) in an undirected graph G is like wandering in a labyrinth with a string and following one path to the end
- ▶ We then backtrack by rolling up our string until we get back to a previously visited vertex v .
- ▶ v becomes our current vertex and we repeat the previous steps

Question

