

Welcome to 01418333

Formal Language and Automata Theory

Why Study Automata?

The slides are created by Jeffrey D. Ullman

<http://infolab.stanford.edu/~ullman/ialc/spr10/spr10.html#LECTURE>

How Could That Be?

- Regular expressions are used in many systems.

- E.g., UNIX `a.*b`.

✗.txt

- E.g., DTD's describe XML tags with a RE format like `person (name, addr, child*)`.

- Finite automata model protocols, electronic circuits.

- Theory is used in *model-checking*.

CFG

How? – (2)

grammar notation

if (_)
if a > b

- Context-free grammars are used to describe the syntax of essentially every programming language.
 - Not to forget their important role in describing natural languages.
- And DTD's taken as a whole, are really CFG's.

How? – (3)

- When developing solutions to real problems, we often confront the limitations of what software can do.
 - ^{ไม่มีคำตอบ} *Undecidable* things – no program whatever can do it.
 - ^{ไม่มีคำตอบที่เร็ว} *Intractable* things – there are programs, but no fast programs.

Course Outline

□ Regular Languages and their descriptors:

- FA □ Finite automata, nondeterministic finite automata, regular expressions.
- Algorithms to decide questions about regular languages, e.g., is it empty?
- Closure properties of regular languages.

Course Outline – (2)

CFL

□ Context-free languages and their descriptors:

- Context-free ^{CFG} grammars, pushdown automata.
- Decision and closure properties.

Course Outline – (3)

- Recursive and recursively enumerable languages.
 - Turing machines, decidability of problems.
 - The limit of what can be computed.
- Intractable problems.
 - Problems that (appear to) require exponential time.
 - NP-completeness and beyond.

Text

- Hopcroft, Motwani, Ullman, *Automata Theory, Languages, and Computation* 3rd Edition.
- Course covers essentially the entire book.

Finite Automata

Motivation

An Example

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Informal Explanation

- Finite automata are finite collections of states with transition rules that take you from one state to another.
- Original application was sequential switching circuits, where the “state” was the settings of internal bits.
- Today, several kinds of software can be modeled by FA.

Representing FA

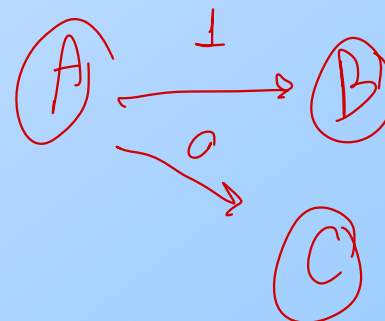
□ Simplest representation is often a graph.

□ Nodes = states.

□ ^{เส้นเชื่อม} Arcs indicate state transitions.

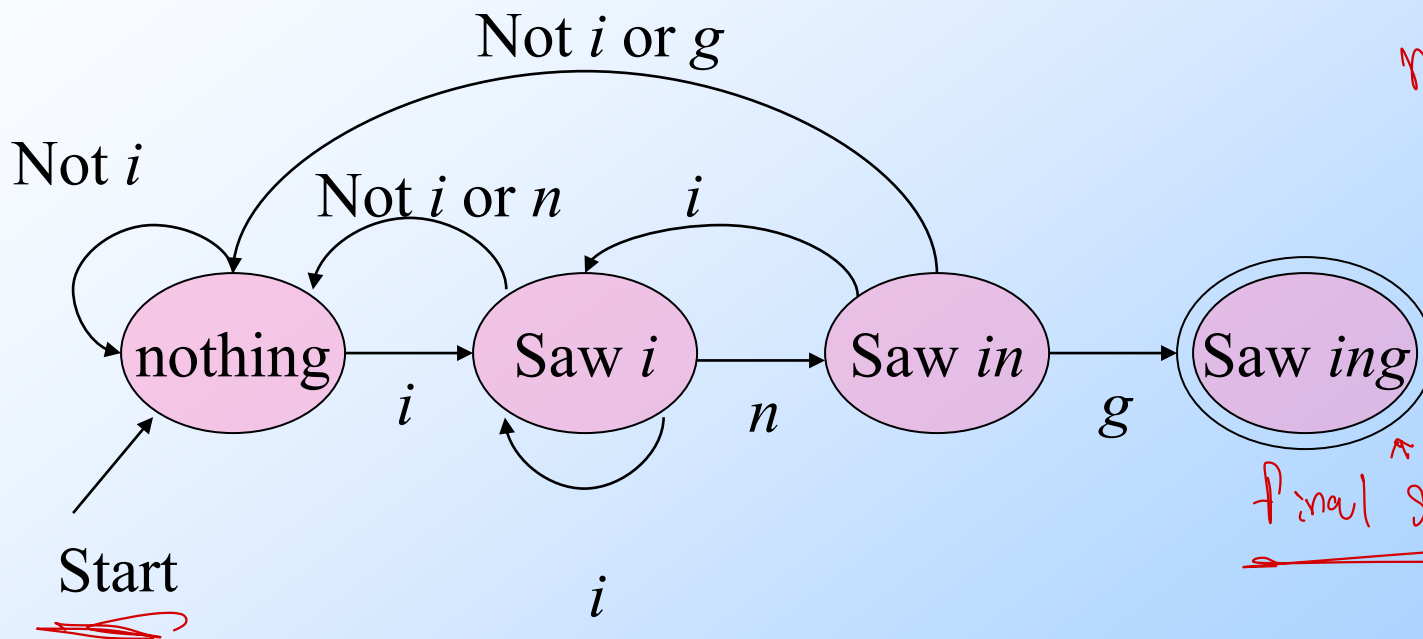
□ Labels on arcs tell what causes the transition. ^{input}

$A \rightarrow B$



Example: Recognizing Strings Ending in "ing"

Accept and Final state can be multiple but start must have only one



being ✓
accept ✓
reject ✗

AN
non state nae 126 12 822. 7001100100.

play → reject
ingby → reject

Automata to Code

- In C/C++, make a piece of code for each state. This code:
 1. Reads the next input.
 2. Decides on the next state.
 3. Jumps to the beginning of the code for that state.

Example: Automata to Code

```
2: /* i seen */  
   c = getNextInput();  
   if (c == 'n') goto 3;  
   else if (c == 'i') goto 2;  
   else goto 1;  
3: /* "in" seen */  
   . . .
```

Automata to Code – Thoughts

- How would you do this in Java, which has no goto?
- You don't really write code like this.
- Rather, a code generator takes a "regular expression" describing the pattern(s) you are looking for.
 - Example: `.*ing` works in grep.

Example: Protocol for Sending Data

