# EECS 545: Machine Learning Lecture 14.Learning Theory

# Honglak Lee and Michał Dereziński 02/23/2022





#### Outline

- PAC theory
- Generalization bounds for finite hypothesis classes
  - Zero training error case
  - Non-zero training error case

VC Dimension

#### Overview

- Probably approximately correct learning (PAC learning)
  - framework for mathematical analysis of machine learning
  - proposed in 1984 by Leslie Valiant (Turing Award winner 2010)
- In this framework, the learner receives samples and must select a generalization function (**hypothesis**) from a certain class of possible functions.
- Goal: with high probability (the "probably" part), the selected function will have low generalization error (the "approximately correct" part).
- PAC theory introduces computational complexity theory.
  - E.g., How many samples do you need to guarantee "probably approximately correctness"
  - − E.g., With N( $\epsilon$ ,  $\delta$ ) samples, empirical error is within  $\epsilon$  of the generalization error with high probability ( $\geq 1-\delta$ ).

# Generalization bounds for zero training error learning

# Finite hypothesis classes

- Consider the case where our hypothesis class is finite.
- Our learning algorithm selects a hypothesis h
  from H
  - based on a sample D of n i.i.d. examples  $\{x^{(i)}, y^{(i)}\}$  from P(x; y) denoted as  $D \sim P^n$

#### Generalization error

 We will denote the 0/1 training data error (also called loss) of a hypothesis h as:

$$L_D(h) = \frac{1}{n} \sum_{i} \mathbf{1} \{ h(x^{(i)}) \neq y^{(i)} \}$$

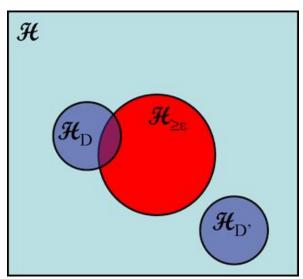
and we'll denote the generalization (true) error as

$$L_P(h) = \mathbf{E}_{(\mathbf{x}, y) \sim P}[\mathbf{1}(h(\mathbf{x}) \neq y)]$$

 The true error of the hypothesis we select based on low training error is most likely larger than the training error.

- Suppose our algorithm finds a hypothesis consistent with the data,  $L_D(h)=0$ .
- We are interested in the probability (over all possible training samples D) that h could have  $L_P(h) \ge \epsilon$ .
- Let's denote the set of hypothesis consistent with D as  $H_D \subset H$  and the set of hypotheses that have true error greater than  $\epsilon$  as  $H_{>\epsilon} \subset H$ .
- Note that  $H_D$  is a random subset of H, while  $H_{\geq \epsilon}$  is fixed.

- The problem for the learner occurs when the two overlap, since in the worst case, a consistent hypothesis can end up being from  $H_{>\epsilon}$ .
- Can we bound the difference in terms of the complexity of our hypothesis class H?



- We're going to bound the probability that the overlap is non-empty.
- Consider a given hypothesis h that has true error greater than  $\epsilon$ , i.e.:  $h \in \mathcal{H}_{\geq \epsilon}$
- The probability that it has made zero errors on D,  $h \in \mathcal{H}_D$ , is exponentially small in n:

$$Pr_{D\sim P^n}(L_D(h)=0) \leq (1-\varepsilon)^n$$

- Now suppose that there are K hypotheses with true error greater than  $\epsilon$ , i.e.  $\mathcal{H}_{\geq \epsilon} = \{h_1, \dots, h_K\}$
- What is the probability that any of them is consistent with the data *D*?

$$Pr_{D \sim P^n}(\exists h \in \mathcal{H}_{\geq \epsilon} : h \in \mathcal{H}_D)$$
  
=  $Pr_{D \sim P^n}(h_1 \in \mathcal{H}_D \lor \dots \lor h_K \in \mathcal{H}_D)$ 

– where ∨ is the logical OR symbol.

 To bound the probability of a union of events, we use the the union bound:

**Union Bound:** 
$$P(A \cup B) \leq P(A) + P(B)$$

• Hence

$$Pr_{D\sim P^n}(\exists h\in\mathcal{H}_{\geq \epsilon}:h\in\mathcal{H}_D)\leq K(1-\epsilon)^n$$

• To make this bound useful, we will simplify it at the expense of further looseness.

- Since we don't know K in general, we will upper-bound it by |H|.
- Since  $(1-\varepsilon) \le e^{-\varepsilon}$  for  $\varepsilon \in [0,1]$ , we will write:

$$Pr_{D\sim P^n}(\exists h \in \mathcal{H}_{\geq \epsilon} : h \in \mathcal{H}_D) \leq K(1-\epsilon)^n$$
  
  $\leq |\mathcal{H}|e^{-n\epsilon}$ 

• Hence the probability that our algorithm will select a hypothesis with true error greater than  $\epsilon$  given that it selected a hypothesis with zero training error is bounded by  $|\mathcal{H}|e^{-n\epsilon}$ , which decreases exponentially with n.

# What's the bound good for? (1)

• There are two ways to use the bound. One is to set the probability of failure  $\delta$ , and the number of examples n and ask about the largest  $\epsilon$ .

$$\delta = |\mathcal{H}|e^{-n\varepsilon} \to \varepsilon = \frac{\log |\mathcal{H}| + \log \frac{1}{\delta}}{n}$$

• So with prob. 1- $\delta$  over the choice of training sample of size n, for any hypothesis h with zero training error,  $L_P(h) \leq \frac{\log |\mathcal{H}| + \log \frac{1}{\delta}}{n}$ 

# What's the bound good for? (2)

- The other way to use it is to fix  $\delta$  and  $\epsilon$  and ask how many examples n are needed to guarantee them (sample complexity).
- That is:

$$\delta = |\mathcal{H}|e^{-n\varepsilon} \to n = \frac{1}{\varepsilon}(\log|\mathcal{H}| + \log\frac{1}{\delta})$$

# PAC (Probably Approximately Correct) framework

- The PAC (Probably Approximately Correct) framework asks these questions about different hypothesis classes
- A class is **PAC-learnable** if the number of examples needed to learn a probably (with prob.  $\delta$ ), approximately (true error of at most  $\epsilon$ ) correct hypothesis is polynomial in parameters of the class
  - (e.g. depth of the tree, dimension of hyperplane, etc.)
  - as well as in  $\epsilon$  and  $1/\delta$  (log  $1/\delta$ )
  - for any distribution
- If the time complexity is also polynomial, the class is efficiently
  PAC-learnable

#### Generalizing for non-zero training error

- This is the simplest of generalization bounds but their form remains the same:
  - the difference in the true error and training error is bounded in terms of complexity of the hypothesis class.
- The limitation of the result so far:
  - it considers finite classes
  - even for finite classes, we have only considered hypotheses with zero training error (which may not exist).

# Generalization bounds for non-zero training error learning

- We will need a form of Chernoff bound for biased coins
- Suppose we have n i.i.d. examples  $z^{(i)} \in \{0, 1\}$ - where  $P(z^{(i)} = 1) = p$
- Let  $\bar{z} = \frac{1}{n} \sum_{i} z^{(i)}$  be the proportion of 1s in the sample.
- Then

$$Pr(p-\overline{z}\geq \varepsilon)\leq e^{-2n\varepsilon^2}$$

• Given a hypothesis h, the probability that the difference between its true error and training error is greater than  $\epsilon$  is bounded:

$$Pr_{D\sim P^n}(L_P(h)-L_D(h)\geq \varepsilon)\leq e^{-2n\varepsilon^2}$$

$$Pr_{D\sim P^n}(L_P(h)-L_D(h)\geq \varepsilon)\leq e^{-2n\varepsilon^2}$$

 Now we need to bound the probability of observing the difference for all hypotheses, so that we can pick the one with lowest error safely. Again, we use the union bound:

$$Pr_{D\sim P^n}(\exists h\in\mathcal{H}: L_P(h)-L_D(h)\geq \varepsilon)\leq |\mathcal{H}|e^{-2n\varepsilon^2}$$

• Setting  $\delta = |\mathcal{H}|e^{-2n\epsilon^2}$  and solving, we get: with prob.  $1-\delta$  over the choice of training sample of size n, for all hypotheses in H:

$$L_P(h) - L_D(h) \le \sqrt{\frac{\log |\mathcal{H}| + \log \frac{1}{\delta}}{2n}}$$

 In case of non-zero training error, the generalization bound decreases with the square root of 1/n, as opposed to with 1/n, which is much slower.

# Complexity of hypothesis classes

- Consider the sizes of some simple hypothesis classes. The class H of all Boolean functions on m binary attributes  $|\mathcal{H}|=2^{2^m}$ 
  - Since we can choose any output for any of the  $2^{2^m}$  possible inputs.
  - Clearly, this class is too rich, since  $\log_2 |\mathcal{H}| = 2^m$ .
- These bounds can be essentially tight,
  - i.e. there are pathological examples where we need that many samples.

# Summary so far

- To learn from a small number of examples, we need a small hypothesis class.
- We have only considered finite hypothesis class.
- Finer-grained notions of complexity: VC-dimension.

# Generalization bounds for infinite classes: VC dimension

# Generalization bounds for infinite classes: VC dimension

 We will now consider the case where our hypothesis class is **infinite**, for example, hyperplanes of dimension m.

- As before, our learning algorithm selects a hypothesis h from H
  - based on a sample D of n i.i.d. examples  $\{x^{(i)}, y^{(i)}\}$  from P(x;y), denoted as  $D \sim P^n$ .

### Bounding the complexity of hypothesis class

- We now consider the case of infinite hypothesis classes, for example, hyperplanes.
- To bound the error of a classifier from a class H using n examples, what matters is <u>not</u> the pure size of H, but its **richness**:
  - The size of the largest dataset such that the model can perfectly classify every possible labeling of that dataset

#### **VC** Dimension

 Vapnik-Chervonenkis dimension is one of the most fundamental measures of richness (or power or complexity or variance) of a hypothesis class studied in machine learning.

 VC theory provides such a general measure of complexity, and gives associated bounds on the optimism.

#### Main result

- A fundamental result by Vapnik and Chervonenkis is that the log |H| in the bounds can be replaced by a function VC(H).
- **Theorem**: With probability  $1-\delta$  over the choice of training sample of size n, for all hypotheses h in H:

$$L_P(h) - L_D(h) \le \sqrt{\frac{VC(\mathcal{H})(\log \frac{2n}{VC(\mathcal{H})} + 1) + \log \frac{4}{\delta}}{n}}$$

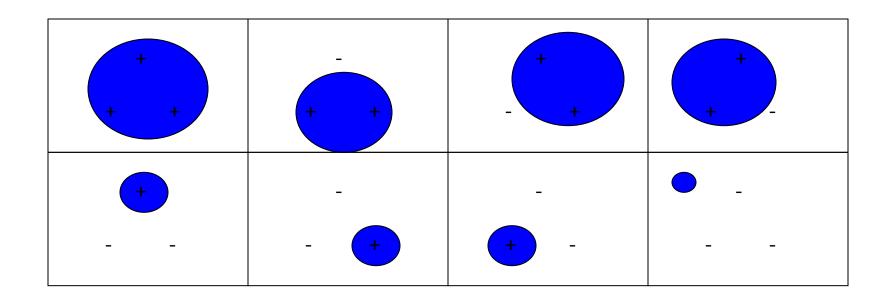
## Shattering:

- A set of instances S is <u>shattered</u> by a hypothesis class H if for every <u>dichotomy</u> (every possible labeling) of S there is a consistent hypothesis in H
  - Intuition: The hypothesis class H is rich enough so that it can handle any possible labeling of dataset S.

# **Example: Shattering**

Is this set of points shattered by the hypothesis class *H* of all circles?

# Example: Shattering

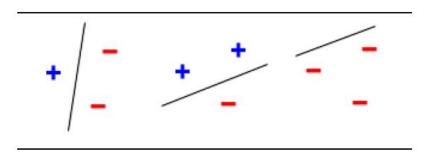


#### Is this set of points shattered by circles?

## Shattering hyperplanes

• Given a set of three points in 2 dimensions, can the set of hyperplanes shatter it?

$$\mathcal{H} = \{h(\mathbf{x}) = sign(w_1x_1 + w_2x_2 + b)\}$$



#### VC dimension

- **Definition:** The VC-dimension of a class H over the input space X is the size of the largest finite set shattered by H.
- To show the VC-dimension of a class is v, we need to show:
  - There exists a set of size v shattered by H (usually easy)
  - There does not exists a set of size v+1 shattered by H (usually harder).

#### **VC** dimension

 What is the VC dimension of the previous examples?

$$\mathcal{H} = \{h(\mathbf{x}) = sign(w_1x_1 + w_2x_2 + b)\}$$

#### VC dimension

 What is the VC dimension of the previous examples?

$$\mathcal{H} = \{h(\mathbf{x}) = sign(w_1x_1 + w_2x_2 + b)\}$$

- There exists a set of three points that can be shattered by H.
- There does not exist a set S of size 4 shattered by
  H!

# VC dimension for hyperplanes

• In general, we can show the VC dimension of hyperplanes in m dimensions is m+1.

 However, in general, the number of parameters of a classifier is not necessarily it's VC dimension.

## Summary

• Given the VC dimension of the hypothesis class VC(H), with probability  $1-\delta$  over the choice of training sample of size n, for all hypotheses h in H:

$$L_P(h) - L_D(h) \le \sqrt{\frac{VC(\mathcal{H})(\log \frac{2n}{VC(\mathcal{H})} + 1) + \log \frac{4}{\delta}}{n}}$$

• i.e., this gives a generalization error bound, as well as sample complexity.

#### Conclusion

 PAC theory provides a bound on generalization error and sample complexity

• The bound critically depends on the complexity of the hypothesis class.

 VC dimension can provide a measure for the complexity of the hypothesis class.

# Quiz

https://forms.gle/9NvcsUsLcQ3CZeDQA

