

# Discrete Mathematics: Homework 5

Name      ID: Number

2020.4.8

1. Let  $p$  be an odd prime and let  $\mathbb{Z}_p^* = \{[1]_p, [2]_p, \dots, [p-1]_p\}$ .

(a) Show that  $([a]_p)^2 = [1]_p$  iff  $[a]_p \in \{[1]_p, [p-1]_p\}$ .

(b) Show that  $[1]_p \cdot [2]_p \cdot \dots \cdot [p-1]_p = [-1]_p$  and thus conclude that  $(p-1)! \equiv -1 \pmod{p}$ .

(a) 证明.  $[1]_p = 1 + np, n \in \mathbb{N}^*$ .  $[p-1]_p = p-1 + np = (n+1)p-1$ .  $[a]_p = a + np$

prove **IF**:

when  $a + np = 1 + np$ ,  $([a]_p)^2 = (1 + np)^2 = 1 + (2n + n^2p)p = [1]_p + p$

when  $a + np = (n+1)p-1$ ,  $([a]_p)^2 = (n^2-1)p+1 = [1]_p + p$

prove **ONLY IF**:

$[a]_p^2 = a^2 + 2anp + n^2p^2 = a^2 + (2an + n^2p)p = 1 + mp = [1]_p$  where  $m, n \in \mathbb{N}^*$

so,  $a^2 = [1]_p, a = [\pm 1]_p$

□

(b) 证明. Because  $p$  is an odd prime number,

$\mathbb{Z}_p^* = \{1, 2, \dots, p-1\}$

So  $\forall a \in \mathbb{N}^*, 0 < a < p, \exists a^{-1} \in [1, p-1], s.t. aa^{-1} \equiv 1 \pmod{p}$

$\forall a \in (1, p-1)$ , we can find a unique  $a^{-1} \in (1, p-1)$

By the multiplication of  $\mathbb{Z}_p$ ,  $[2]_p \cdot \dots \cdot [p-2]_p = [1]_p$

And obviously,  $[1]_p \cdot [p-1]_p = [-1]_p$

So,  $[1]_p \cdot [2]_p \cdot \dots \cdot [p-1]_p = [(p-1)!]_p = [-1]_p$

Which equals to  $(p-1)! \equiv -1 \pmod{p}$

□

2. Let  $x, y, z$  be integers such that  $x^2 + y^2 \equiv 3z^2 \pmod{4}$ . Show that  $x, y, z$  must be all even. Based on this result, show that the equation  $x^2 + y^2 = 3z^2$  has no other integer solutions except  $(x, y, z) = (0, 0, 0)$ .

证明.  $x^2 + y^2 \equiv 3z^2 \pmod{4}$ , we have  $x^2 + y^2 = 3z^2 + 4n, n \in \mathbb{Z}$

so,  $\frac{x^2 + y^2 - 3z^2}{4}$  is an integer.  $4 \mid x^2 + y^2 - 3z^2$

if  $x, y$  are an odd and an even,  $z$  is an even.

Suppose  $x = 2k + 1, y = 2q, z = 2r$ , we have  $(2k + 1)^2 + (2q)^2 - 3(2r)^2 \pmod{4} = 1$

if  $x, y$  are an odd and an even,  $z$  is an odd.

Suppose  $x = 2k + 1, y = 2q, z = 2r + 1$ , we have  $(2k + 1)^2 + (2q)^2 - 3(2r + 1)^2 \pmod{4} = 2$

if  $x, y$  are odds,  $z$  is an even.

Suppose  $x = 2k + 1, y = 2q + 1, z = 2r$ , we have  $(2k + 1)^2 + (2q)^2 - 3(2r)^2 \pmod{4} = 2$

if  $x, y$  are odds,  $z$  is an odd.

Suppose  $x = 2k + 1, y = 2q + 1, z = 2r + 1$ , we have  $(2k + 1)^2 + (2q)^2 - 3(2r)^2 \pmod{4} = 3$

where  $k, q, r \in \mathbb{Z}$

So,  $x, y, z$  must be all even.

let  $x^2 + y^2 \equiv 3z^2 \pmod{4}$

let  $x = 2k, y = 2q, z = 2r$ , so we have  $x^2 + y^2 = 4(k^2 + q^2) = 12(r^2) = 3z^2$

and we have  $k^2 + q^2 = 3r^2$

and  $|k| < |x|, |q| < |y|, |r| < |z|$  or  $k = x = 0, q = y = 0, r = z = 0$

Assume the smallest integer positive solution is  $x_0, y_0, z_0$ , but there exists  $k_0 = \frac{1}{2}x_0, q_0 = \frac{1}{2}y_0, r_0 = \frac{1}{2}z_0$ .

So nonzero integer solution doesn't exist.

$(x, y, z) = (0, 0, 0)$

□

3. Let  $a_1, a_2, a_3, a_4$  be arbitrary integers. Find ALL integer solutions of the following equation system.

$$\begin{cases} x \equiv a_1 \pmod{11}; \\ x \equiv a_2 \pmod{13}; \\ x \equiv a_3 \pmod{17}; \\ x \equiv a_4 \pmod{19}; \end{cases}$$

**Solution:**

$$n = 11 \cdot 13 \cdot 17 \cdot 19 = 46189$$

$$m_1 = 4199 \quad m_2 = 3553 \quad m_3 = 2717 \quad m_4 = 2413$$

$$m_1^{-1} = 7 \quad m_2^{-1} = 10 \quad m_3^{-1} = 11 \quad m_4^{-1} = 18$$

$$c_1 = 29392 \quad c_2 = 35530 \quad c_3 = 29887 \quad c_4 = 43758$$

$$a = \sum_{i=1}^4 a_i c_i \pmod{n} = 29392a_1 + 35530a_2 + 29887a_3 + 43758a_4 \pmod{46189}$$

$$x \in [29392a_1 + 35530a_2 + 29887a_3 + 43758a_4]_{46189}$$

4. A composite integer  $N$  that satisfies the congruence  $b^{N-1} \equiv 1 \pmod{N}$  for all positive integers  $b$  with  $\gcd(b, N) = 1$  is called a Carmichael number. Suppose that  $N = p_1 p_2 p_3$  is an integer, where  $p_1, p_2, p_3$  are primes such that  $(p_i - 1) | (N - 1)$  for  $i = 1, 2, 3$ . Show that  $N$  is a Carmichael number.

证明. By Fermat's Little Theorem, we have,

$$b^{p_1-1} \equiv 1 \pmod{p_1}$$

$$b^{p_2-1} \equiv 1 \pmod{p_2}$$

$$b^{p_3-1} \equiv 1 \pmod{p_3}$$

$$N - 1 = \mathbb{Z}(p_i - 1), \quad b^{N-1} = b^{\mathbb{Z}(p_i-1)}$$

$$\forall x \equiv 1 \pmod{y}, \text{ we have } x^n \equiv 1 \pmod{y} \text{ Because } x = 1 + \mathbb{Z}y \text{ and } x^n = 1 + (2y + \mathbb{Z}y^2)\mathbb{Z}$$

$$\forall x \equiv 1 \pmod{y}, \text{ we have } x \equiv 1 \pmod{ky} \text{ Because } x = 1 + \mathbb{Z}y \text{ and } x = 1 + \mathbb{Z}y + \mathbb{Z}y$$

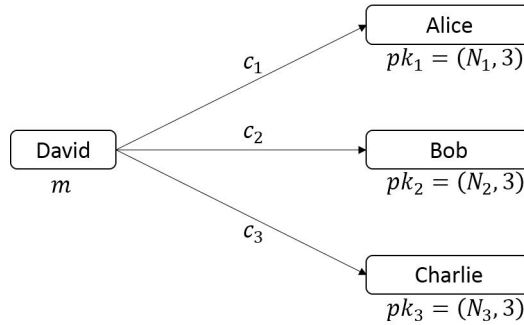
$$\text{So, } b^{N-1} \equiv b^{3(N-1)} \equiv b^{p_1-1} b^{p_2-1} b^{p_3-1} \pmod{p_1 p_2 p_3}$$

$$\text{Which has the same reminder 1 as } b^{p_1-1} \equiv 1 \pmod{p_1}$$

$$\text{So, } b^{N-1} \equiv 1 \pmod{N}$$

So,  $N$  is a Carmichael number. □

5. See the following figure. The RSA public keys of Alice, Bob and Charlie are  $pk_1 = (N_1, 3)$ ,  $pk_2 = (N_2, 3)$  and  $pk_3 = (N_3, 3)$ , respectively. David wants to send a private message  $m$  to Alice, Bob and Charlie, where  $m$  is an integer and  $0 < m < N_i$  for  $i = 1, 2, 3$ . In order to keep  $m$  secret from an eavesdropper Eve, David encrypts  $m$  as  $c_1 = m^3 \bmod N_1$ ,  $c_2 = m^3 \bmod N_2$  and  $c_3 = m^3 \bmod N_3$ ; and then sends  $c_1$  to Alice,  $c_2$  to Bob and  $c_3$  to Charlie.



Suppose that  $N_1, N_2, N_3$  are pairwise relatively prime. Show that with the knowledge of all public keys and all ciphertexts, Eve can decide the value of  $m$ .

证明. By RSA Theorem, we have

$$\begin{cases} x^3 \equiv c_1 \pmod{N_1} \\ x^3 \equiv c_2 \pmod{N_2} \\ x^3 \equiv c_3 \pmod{N_3} \end{cases}$$

so,  $n = N_1 N_2 N_3$

$$m_1 = N_2 N_3 \quad m_2 = N_1 N_3 \quad m_3 = N_1 N_2$$

$$c_1 = m_1 m_1^{-1} \quad c_2 = m_2 m_2^{-1} \quad c_3 = m_3 m_3^{-1}$$

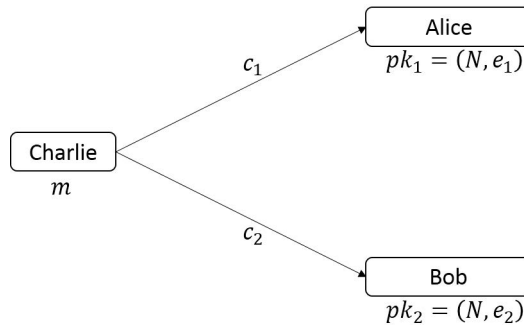
$$a = \sum_{i=1}^k a_i c_i \pmod{n}$$

$$x^3 \in [\sum_{i=1}^3 a_i c_i]_n$$

so we can try to solve every  $x$  as a message.

□

6. See the following figure. Alice and Bob trust each other very much. They set their RSA public keys as  $pk_1 = (N, e_1)$  and  $pk_2 = (N, e_2)$ , respectively. Charlie wants to send a private message  $m$  to Alice and Bob, where  $0 \leq m < N$  is an integer and  $\gcd(m, N) = 1$ . In order to keep  $m$  secret from an eavesdropper Eve, Charlie encrypts  $m$  as  $c_1 = m^{e_1} \bmod N$  and  $c_2 = m^{e_2} \bmod N$ ; and then sends  $c_1$  to Alice and  $c_2$  to Bob.



Suppose that  $\gcd(e_1, e_2) = 1$ . Show that with the knowledge of all public keys and all ciphertexts, Eve can decide the value of  $m$ .

证明. By RSA Theorem, we have

$$\begin{cases} x^{e_1} \equiv c_1 \pmod{N} \\ x^{e_2} \equiv c_2 \pmod{N} \end{cases}$$

□

Because  $\gcd(e_1, e_2) = 1$ , we have  $x_1 e_1 + y_1 e_2 = \gcd(e_1, e_2) = 1$  by EXGCD.

so,  $x^{x_1 e_1 + y_1 e_2} \equiv c_1^{x_1} + c_2^{y_1} \equiv x^1 \equiv \pmod{N}$ , and  $N$  is a prime. so, Eve can solve  $m$  by solving  $m \equiv c_1^{x_1} + c_2^{y_1} \pmod{N}$