# Design and Implement Binary Fuzzing based on LibFuzzer

Chun-Ying Huang, Wei-Chieh Chao, Si-Chen Lin, Yi-Hsien Chen Department of Computer Science, College of Computer Science, National Chiao Tung University

Abstract—Libfuzzer is a coverage-guided fuzzing engine. But not like AFL, which can do the fuzzing without source code, Libfuzzer needs source code of the target to compile and do the fuzzing. This paper presents a way to use Libfuzzer to fuzz the binary without source code. Our implementation link the IO of binary with Libfuzzer, and use qemu to collect runtime information.

#### I. Introduction

Fuzzing is a technique to automatically test software by providing random data as input. When fuzzing, we usually build a fuzzer to generate and send random inputs to the program. After program exit, it will send execution information back to the fuzzer. Then, the fuzzer will generate new input based on the execution information.

When doing fuzz testing, we focus on crashes. That is, we try to generate input data that make program crashed. Therefore, fuzzer analyzes execution information such as path, coverage, and exceptions to decide what kind of input data is next time.

American Fuzzy Lop (AFL) is one of the most famous fuzzing tools. There are two modes in AFL. One is source code mode, and the other is binary (QEMU) mode. Source code mode uses aff-gcc to compile the source code, in the meanwhile, it inserts lots of checkpoints to send execution information back. On the other hand, binary (QEMU) mode uses QEMU to run the binary, and it analyzes execution information from QEMU directly. Because of the implementation, its no doubt that source code mode would find crashes faster than binary (QEMU) mode.

LibFuzzer is a library for coverage-guided fuzz testing. It only support source code mode. Different from AFL, LibFuzzer considers target source code as library, and it send inputs through call function. It compile its source code with target source code together.

In some case, LibFuzzer found crashes much faster than AFL did. Heartbleed (CVE-2014-0160) was a critical security bug in the OpenSSL cryptography. It took couples of hours to find crash using AFL. However, it only took ten seconds to find crash using LibFuzzer. Thats one of the reasons that we wanted to build binary mode for LibFuzzer.

Generally, we built a bridge for communication between LibFuzzer and binary. We used QEMU, as same as AFL, to generate runtime information. In detail, we executed binary through QEMU and analyzed runtime information from QEMU after exited. LibFuzzer need lots of information from binary, such as coverage, caller, and cmp. We focused on coverage here. Coverage is how many percent of codes have been executed. First, we used radare2 to generate basic blocks infromation. Moreover, we did static analysis before executing binary. Second, we executed binary through QEMU and analyzed runtime information dynamically. Finally, we sent coverage infromation back to LibFuzzer.

II. RELATED WORK

Sample citations [1], [2], [3], [4].

#### III. METHODOLOGY

## A. Link IO together

Normally, when using Libfuzzer, we put the source code need to be test inside LLVMFuzzerTestOneInput as describe in the introduction section. But when we want to run a binary, we don't have source code. Hence, what we do is actually putting execv("our-binary", ..., ...) inside the LLVMFuzzerTestOneInput to invoke our binary. And before execute execv we need to use dup2 to link

- 1. fuzzer test input  $\rightarrow$  binary stdin.
- 2. binary stdout  $\rightarrow$  /dev/null.
- 3. binary stderr  $\rightarrow$  /dev/null

The simplified code is something like below

```
extern "C" int LLVMFuzzerTestOneInput(const
    uint8_t *Data, size_t Size) {
    int P_IN[2]; pipe(P_IN);
    if(Size) write(P_IN[1], Data, Size);
    ...
    int pid = fork();
    if(pid == 0) {
        dup2(P_IN[0], STDIN_FILENO);
        dup2(dev_null, STDOUT_FILENO);
        dup2(dev_null, STDERR_FILENO);
        ...
        execv(..., ..., ...);
    }
    ...
}
```

# B. Collect Runtime Information

After we link IO together, the binary should get the input of the libfuzzer test input now. Though it is runnable, Libufuzzer don't have any runtime information like code coverage and will stop running after a few rounds. First, we need to figure out what Libfuzzer collect while running. Libfuzzer collect runtime information through **Clang SanitizerCoverage**, which provides simple code coverage instrumentation and has hook function for customization. Now we use qemu to

collect right information for those hook function implemented by LibFuzzer.

- 1) trace pc guard: LibFuzzer use two array to store the information of code coverage.
  - 1. \_\_sancov\_trace\_pc\_pcs : store the program counter of the beginning of a code block.
  - \_\_sancov\_trace\_pc\_guard\_8bit\_counters:
     store how many times the code block was hit.

```
__sancov_trace_pcs[Idx] = PC;
__sancov_trace_pc_guard_8bit_counters[Idx]++;
```

We need to fill those array for LibFuzzer to work. First, we use radare2 api **r2pipe** to do static analysis and index those code block. While doing static analysis with r2pipe, we filter out addresses we don't want, like library code. Second, we use qemu to check whether we hit a code block, and fill in those value. The qemu command we use is qemu-x86\_64 -D logfile -d in\_asm binary. in\_asm provides us the assembly code and its address being executed. ( see example below )

```
...
IN:
0x0000004000802c90: mov %r11,%rcx
0x0000004000802c93: sub %rax,%rcx
0x0000004000802c96: cmp $0xb,%rcx
0x0000004000802c9a: ja 0x4000802dd8
```

After parsing qemu output and fill in those value, LibFuzzer is now runnable. But we still got plenty of hook function we need to implement to make it works exactly like it should be.

## IV. EVALUATION

In this section, we will compare the performance of original LibFuzzer, our LibFuzzer and AFL. We use fuzzer-test-suite, a set of tests for fuzzing engines provided by Google, as the test cases. In order to support our LibFuzzer, we modify the target.cc in some test cases. Take c-ares-CVE-2016-5180 for example, we modified the extern "C" int LLVMFuzzerTestOneInput(const uint8\_t \*Data, size\_t Size) function declaration to int

\*Data, size\_t Size) function declaration to int main(). And use read function to read input data, the size is the length of input data. Check out more detail of patching fuzzer-test-suite in wiki page of our repository. The following specification is our machine information and all experiments were performed on it.

```
virtual private server on Google Cloud Platform
n1-standard-1 (one vCPU, 3.75 GB RAM)
```

# A. Compared with original LibFuzzer

We choose some CVE test cases in fuzzer-test-suite, including CVE-2016-5180, CVE-2015-8317 and HeartBleed (CVE-2014-0160), to evaluate the performance of our LibFuzzer. These are the output of running test cases by our LibFuzzer.

```
1. CVE-2016-5180
```

```
# ()
       READ units: 1
#2
       INITED cov: 4 ft: 20 corp: 1/1b
    exec/s: 2 rss: 25Mb
#3
       NEW
              cov: 4 ft: 39 corp: 2/3b
     exec/s: 3 rss: 25Mb L: 2/2 MS: 1
    CopyPart-
       pulse cov: 4 ft: 59 corp: 2/3b
#4
     exec/s: 2 rss: 25Mb
            cov: 4 ft: 59 corp: 3/5b
#4
       NEW
     exec/s: 2 rss: 25Mb L: 2/2 MS: 2
    CopyPart-ShuffleBytes-
Error status : 139
Signal: Segmentation fault
```

2. CVE-2015-8317

#0 READ units: 1
Error status : 139
Signal : Segmentation fault

3. HeartBleed (CVE-2014-0160)

#0 READ units: 1
WARNING: block size exceeding max
 block size at 0x004959e0
[+] Try changing it with e anal.bb.
 maxsize

WARNING : block size exceeding max block size at 0x004d7920

[+] Try changing it with e anal.bb.
maxsize

WARNING: block size exceeding max block size at 0x0043e9d0

[+] Try changing it with e anal.bb.
maxsize

WARNING: block size exceeding max block size at 0x0043f160

[+] Try changing it with e anal.bb.
maxsize

WARNING: block size exceeding max block size at 0x004962f0 [+] Try changing it with e anal.bb.

maxsize
WARNING: block size exceeding max

block size at 0x00496e70

[+] Try changing it with e anal.bb.
maxsize

WARNING: block size exceeding max block size at 0x00457460

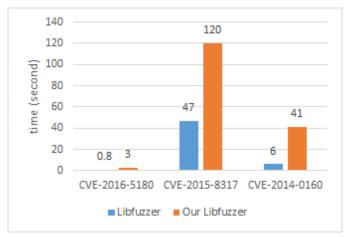
[+] Try changing it with e anal.bb.
 maxsize

WARNING : block size exceeding max block size at 0x004f2ace

[+] Try changing it with e anal.bb.
maxsize

Error status : 134 Signal : Aborted

Compared with original LibFuzzer and plot the chart.



As we can see, because the implementaion of our LibFuzzer, the IO speed and analysis take a lot of time. And the preformance is slower than original LibFuzzer. But compared the time of original LibFuzzer and our LibFuzzer took, we still can find the bugs efficiently in these test cases.

# B. Compared with AFL

### V. CONCLUSION

#### REFERENCES

- [1] K. Serebryany, "libFuzzer a library for coverage-guided fuzz testing," LLVM project, 2015, https://llvm.org/docs/LibFuzzer.html.
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- [4] K. Serebryany, "OSS-Fuzz Google's continuous fuzzing service for open source software," 2017, USENIX Security. [Online]. Available: https://www.usenix.org/conference/usenixsecurity17/technicalsessions/presentation/serebryany