Write a comparison table of all time complexities of all sorting algorithms?

The efficiency of an algorithm depends on two parameters. They are:-

- 1. Time Complexity
- 2. Space complexity.
- 1. Time Complexity:- Time complexity is defined as the number of times a partiallar instruction Set is executed rather than the total time taken. It is because the total time took also depends on some external factors like the compiler used, Processor's speed etc.
- 2. Space Complexity: Space complexity is the total memory space required by the program for its execution.

Both are calculated as the function of input Size (n). One important thing here is that inspite of these parameters the efficiency of an algorithm also depends upon the nature and size of the input

Types of Time Complexity:

- 1. Best Time Complexity: Define the input for which algorithm takes less time (04) minimum time. In the best case, Calculate the lower bound of an algorithm. Example: In the linear search, when search data is present at the first location of large data then the best case occurs.
- 2. Average Time Complexity: In the overage case take all random inputs and calculate the computation time for all inputs.

And we divide it by the total number of inputs.

3. Worst Time Complexity: Defines the input for which algorithm takes (place) a long time (or) maximum time. In the worst calculate the upper bound of an algorithm.

Example: - In the linear Search, When Search data is present at the last location

Of large data then the worst case occurs.

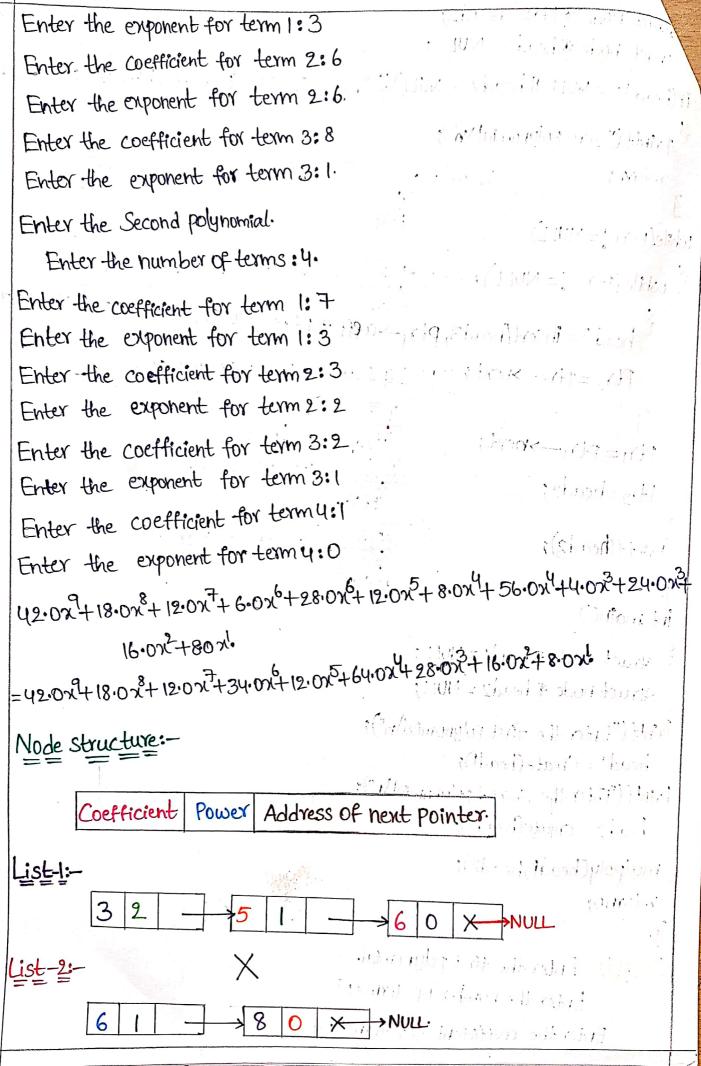
The below table Compares the time complexities and Space complexities					
SIND	Algorithm	Time Complexity			Space Complexity
,		Best	Average	Worst	(Morst).
10.7	M. 14	· I MILL	- Tanyan, ,		
31.	Selection Sort	1 (n2)	$\Theta(n^2)$	O(n2)	0(1).
y£.	the state of the said of	· / * 4/	7 1 1 1 1 1		
9.	Bubble Sort	-0 (h)	$O(n^2)$	O(n²)	0(1).
Lori	ede i nate i navolavi iz	3 .		1 10	
_3.	Insertion Sort	0 (n)	$O(n^2)$	O(n²)	O(i).
			* # W 1 *	of a mari	000.
4.	Heap Sort	a (nlogen)	O(n log(n))	O(n log(n)	000
TOTAL.		0	e (i (se fui)	Ochrogon)	<u>-O(j)</u> .
5.	Quick Sort	2 (n log(m)	O(n log(n))	(12)	0.000
	•	J	<u> </u>	U(N)	O(n).
6.	Merge Sort	-a (n logun)	O(n login)	O(n log(n)	O(n).
	The red hop to his you	राजित में मुख्ये अंती	28 Jan - 12		
	Bucket Sort	o (n+k)	O(ntk)	O(n2)	0(n).
8.	Radix Sort	12 2 4 30 3	0.		Tarista Y.
<u> </u>	Manx Sore	-0 (nk)	9(nk)	O(nk)	O(n+k).
	Charles			Firm Str	(6) 1
9.	Count Sort	o (n+k)	O(n+k)	O(MK)	O(k).
10	Chuca		A	1000	
10.	Shell Sort	a (nlogun)	O(nlogin)	O(12)	Oa).
odqisi.	Little to the little				10 bak
11.	Tim Sort	_n(n)	O(n log(n))	O(n logun)	O(n).
12.	Tree Sort	n (nlogen)	O(n log(n))	O(h2)	O(n)
13.	Cube Sort	.o.(n)	O(nlog(n))	O(nloch)	0(n).

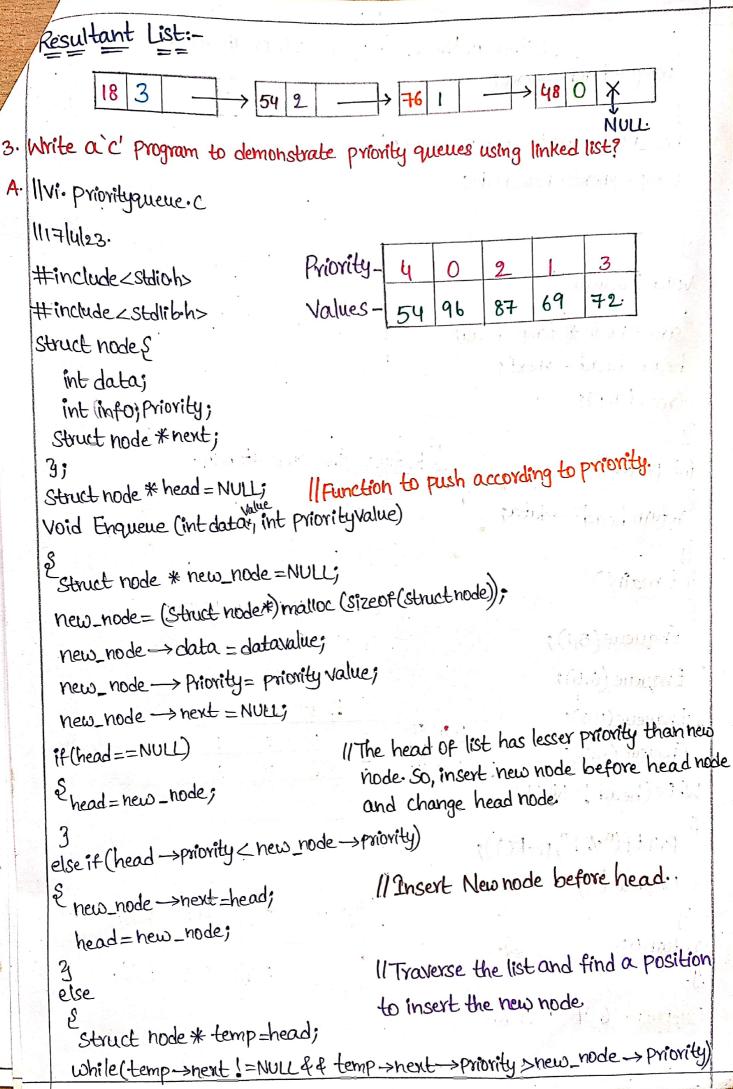
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lite a'c' program for multiplication of two polynomials using linked lists?
11 vi. mulpoly.c
                                  Example: - Poly 1 = 2x3+3x2+4x+7=0
117/4/23.
                                            Poly2=8x2+9x+3=0
                                          Poly 3=16x5+18x4+6x3+24x4+27x3+
#include_stdio.h>
                                                   1221+ 3223+3622+122+5622+
#include < Stdlib.h>
                                   1011 / 10 10 10 1 63x+21=0
                                            = 16x5+42x4+65x3+92x2+87x+21+0.
Struct nodes
        float coeff;
                                     Poly1= 23 -> 32 -> 41 -> 70 X
          int expo;
                                     Poly 2=82 -> 911 -> 30 X
        Struct node * next;
                            Poly3=1615 -> 424 -> 653 -> 622 -> 8711 -> 210 X·
  4;
 Struct node * temp;
temp=(Struct node*)malloc(sizeof(struct rode));
      new → coeff = co;
       new-expo= ex;
       new-next=NULLi
if (head == NULL || ex>head -> expo) | I If there is no node in the list.
                            11 Given exponent is greater than the first node exponent
 & new→next=head;
   head = new;
 y
else
                                   rainstiff ( of it x 1 id ) house - scoot (bit 1 x 7 is o) thinks
   Etemp=head;
   while (temp->next!=NULLP& temp->next->expo>=ex); 11/=1 grost) +
  temp=temp->nent;
  new->next=temp->next;
  temp next=new;
 return head;
Struct node * create(struct node * head)
```

Scanned with CamScanner

```
float coeff;
int expo;
Printf("Enter the number of terms of the polynomial: ");
Scanf (" 1.d", &ri);
for(1=0; icn; itt)
& Printf("Enter the co-efficient of 1.d:", iti);
 Scanf (" 1:04", & coeff);
Printf("Enter the exponent Of 16d: ", i+1);
Scanf (" 1.d", f expo);
head=insert(head, coeff, expo);
 return head;
Void print (struct node * head)
Pif(head==NULL)
  Printf("No polynomial exists");
  else
  Struct node * temp = head;
  while (temp!=NULL)
 \mathcal{E}_{\text{Printf("(%.1f } \times^{\wedge} \%d)", \text{temp} \longrightarrow \text{coeff, temp} \longrightarrow \text{expo)};}
  temp=temp->hent;
                               co- from the junt -
   if (temp! = NULL)
   & Printf(" +");
  else ş
         Printf ("\n");
                                                                        Harrid repuls
  Void mulpoly (struct node * head 1, struct node * head 2)
   & Struct node * ptr_= head;
```

```
struct node * Ptr2=head2;
 Struct hode * head 3= NULL;
if (head 1 == NULL 11 head 2 == NULL) // Check if first (or) Second Polynomial is Null.
? Printf("Zero polynomial \n");
  return;
                       head 1 > 4 3
                                 PLYI
While (Pty != NULL)
& while (Ptro != NULL) head 2->53
    S head 3 = insert(head3, Ptr, -> coeff * ptr2 -> coeff, Ptr, -
      Ptr2=ptr2->next; // Multiplication of two polynomials.
    Ptr = ptr -> next;
    Ptr2=head2;
  Print (head 2);
 20.00 + 1000 + 100.00 + 100.00 1000 10008 +0
 int main()
 & struct hode * head 1 = NULL;
   struct node * head 2 = NULL;
 frintf("Enter the first polynomial\n");
    head 1 = Create (head 1);
 Printf ("Enter the second polynomial \n");
    head 2 = create(head2);
  mulpoly(head1, head2);
  returno;
 Output: - Enter the first polynomial.
           Enter the number of terms:3
        Enter the coefficient for term 1: 4
```





```
// Either at the ends of the list (or) at required position.
  temp=temp-next;
 new_hode->next = temp -> hext;
 temp-next=new_node;
Void Dequeuel)
Struct node * temp=head;
 head=head ->next;
 free (temp);
4
                         // Return the value at head
int peck()
int main()
                              1/ Create a priority queue
  Enqueue (5,1);
                                  7->4->5->6-
                                  ביותול וב יחושולע לחיטבן
 Enqueue (610);
                                     Enqueue.
 Enqueue (4,2);
 Enqueue (7,3);
                     DataValue
While (head != NULL)
  Printf("/d", peek());
  Dequeue();
                                  Linked
                                             Push()
                                                   POPC)
                                                          Peek
                                   list
returnoj.
                                  Time
                                             (O(n)
                                                    0(1)
                                  Complexite
Output: - 67456.
```

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