

1. Simplify the logical expression $q \vee \neg(\neg(p \wedge r) \rightarrow (p \vee \neg q))$ using logical equivalence rules.

Answer:

$$\begin{aligned}
 q \vee \neg(\neg(p \wedge r) \rightarrow (p \vee \neg q)) &\equiv q \vee \neg((p \wedge r) \vee (p \vee \neg q)) && \text{(Implication Equivalence)} \\
 &\equiv q \vee ((\neg p \vee \neg r) \wedge (\neg p \wedge q)) && \text{(De Morgan's Law)} \\
 &\equiv q \vee [(\neg p \wedge \neg p \wedge q) \vee (\neg r \wedge \neg p \wedge q)] && \text{(Distributive Law)} \\
 &\equiv q \vee [(\neg p \wedge q) \vee (\neg r \wedge \neg p \wedge q)] && \text{(Idempotent Law)} \\
 &\equiv q \vee (\neg p \wedge q) && \text{(Absorption Law)} \\
 &\equiv q && \text{(Absorption Law)}
 \end{aligned}$$

2. Use the logical equivalence properties below to verify the logical equivalence.

$$(p \rightarrow \neg q) \wedge (p \rightarrow \neg r) \equiv \neg(p \wedge (q \vee r))$$

Answer:

$$\begin{aligned}
 (p \rightarrow \neg q) \wedge (p \rightarrow \neg r) &\equiv (\neg p \vee \neg q) \wedge (\neg p \vee \neg r) && \text{(Implication Equivalence)} \\
 &\equiv \neg p \vee (\neg q \wedge \neg r) && \text{(Distributive Law)} \\
 &\equiv \neg p \vee \neg(q \vee r) && \text{(De Morgan's Law)} \\
 &\equiv \neg(p \wedge (q \vee r)) && \text{(De Morgan's Law)}
 \end{aligned}$$

3. Using truth tables, determine whether $(p \wedge q) \rightarrow r \equiv (p \rightarrow r) \wedge (q \rightarrow r)$

Answer:

p	q	r	$(p \wedge q)$	$(p \wedge q) \rightarrow r$	$(p \rightarrow r)$	$(q \rightarrow r)$	$(p \rightarrow r) \wedge (q \rightarrow r)$
T	T	T	T	T	T	T	T
T	T	F	T	F	F	F	F
T	F	T	F	T	T	T	T
T	F	F	F	T	F	T	F
F	T	T	F	T	T	T	T
F	T	F	F	T	T	F	F
F	F	T	F	T	T	T	T
F	F	F	F	T	T	T	T

The two propositions are not logically equivalent.

4. (A) Express the following using predicates, quantifiers, logical connectives, and mathematical operators if necessary.
- (a) Every positive integer is the sum of the squares of four integers. (The universe of discourse contains all integers)
 - (b) Every user has access to exactly one mailbox. (Assume that the domain consists of all users and all mailboxes)
- (B) Let $G(x, y)$ mean that child x has played video-game y , where the domain for x consists of all the children in your school and the domain for y consists of all video-games. Express these statements as an English sentence.
- (a) $\exists a \forall b (a \neq (child_1) \wedge (G(child_1, b) \rightarrow G(a, b)))$
 - (b) $\exists x \exists y \forall z ((x \neq y) \wedge (G(x, z) \leftrightarrow G(y, z)))$

Answer:

- (A) (a) $\forall x \exists a \exists b \exists c \exists d ((x > 0) \rightarrow x = a^2 + b^2 + c^2 + d^2)$
- (b) Let $A(x, y)$ mean that user x has access to mailbox y .
 $\forall x \exists y (A(x, y) \wedge \forall z (z \neq y \rightarrow \neg A(x, z)))$
- (B) (a) There is a child other than $child_1$ who has played all the video-games that $child_1$ has played. (This does not mean that this child has only played the video-games played by $child_1$. The child could have also played other games.)
- (b) There exist two different children who have played exactly the same video-games.
5. Prove or disprove that the following compound proposition is a contingency.

$$((p \vee r) \vee ((q \wedge p) \vee (q \wedge r))) \wedge \bar{r} \wedge \bar{p}$$

Answer:

$$\begin{aligned} & ((p \vee r) \vee ((q \wedge p) \vee (q \wedge r))) \wedge \bar{r} \wedge \bar{p} \\ \equiv & ((p \vee r) \vee (q \wedge (p \vee r))) \wedge \bar{r} \wedge \bar{p} \text{ (Distributive Law)} \\ \equiv & (p \vee r) \wedge \bar{r} \wedge \bar{p} \text{ (Absorption Law)} \\ \equiv & (p \vee r) \wedge \overline{(p \vee r)} \text{ (DeMorgan's Law)} \\ \equiv & F \text{ (Complement Law)} \end{aligned}$$