CHAPTER 8

PROPERTIES OF RELATIONS

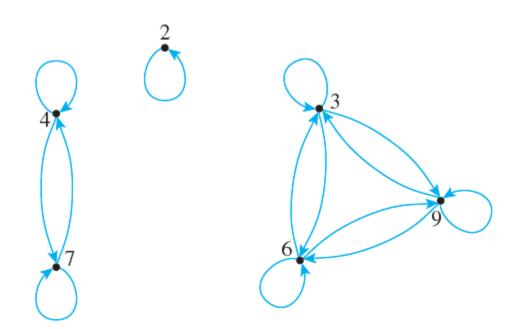
8.2

Reflexivity, Symmetry, and Transitivity

Definition

Let *R* be a relation on a set *A*.

- 1. *R* is **reflexive** if, and only if, for every $x \in A$, $x \in A$, $x \in A$.
- 2. *R* is **symmetric** if, and only if, for every $x, y \in A$, if x R y then y R x.
- 3. *R* is **transitive** if, and only if, for every $x, y, z \in A$, if x R y and y R z then x R z.

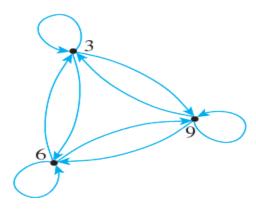


The equivalence of the expressions x R y and $(x, y) \in R$ for every x and y in A, the reflexive, symmetric, and transitive properties can also be written as follows:

- 1. *R* is reflexive \Leftrightarrow for every x in A, $(x, x) \in R$.
- 2. R is symmetric \Leftrightarrow for every x and y in A, **if** $(x, y) \in R$ then $(y, x) \in R$.
- 3. *R* is transitive \Leftrightarrow for every x, y, and z in A, **if** $(x, y) \in R$ and $(y, z) \in R$ then $(x, z) \in R$.







This means that to prove a relation has one of the properties, you use either the method of exhaustion or the method of generalizing from the generic particular.

Recall that the negation of a universal statement is existential. Hence if *R* is a relation on a set *A*, then

1. R is **not reflexive** \Leftrightarrow there is an element x in A such that $x \not R x$ [that is, such that $(x, x) \notin R$].

- 2. R is **not symmetric** \Leftrightarrow there are elements x and y in A such that x R y but $y \not R x$ [that is, such that $(x, y) \in R$ but $(y, x) \notin R$].
- 3. R is **not transitive** \Leftrightarrow there are elements x, y, and z in A such that x R y and y R z but $x \not R z$ [that is, such that $(x, y) \in R$ and $(y, z) \in R$ but $(x, z) \notin R$].

It follows that you can show that a relation does *not* have one of the properties by finding a counterexample.

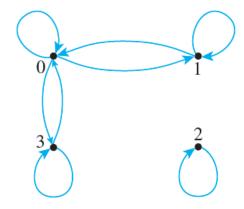
Example 8.2.1 – Properties of Relations on Finite Sets

Let $A = \{0, 1, 2, 3\}$ and define relations R, S, and T on A as follows:

$$R = \{(0, 0), (0, 1), (0, 3), (1, 0), (1, 1), (2, 2), (3, 0), (3, 3)\},\$$
 $S = \{(0, 0), (0, 2), (0, 3), (2, 3)\},\$
 $T = \{(0, 1), (2, 3)\}.$

- a. Is *R* reflexive? symmetric? transitive?
- b. Is S reflexive? symmetric? transitive?
- c. Is *T* reflexive? symmetric? transitive?

a. The directed graph of *R* has the appearance shown below.



R is reflexive: There is a loop at each point of the directed graph. This means that each element of *A* is related to itself, so *R* is reflexive.

R is symmetric: In each case where there is an arrow going from one point of the graph to a second, there is an arrow going from the second point back to the first.

This means that whenever one element of *A* is related by *R* to a second, then the second is related to the first. Hence *R* is symmetric.

R is not transitive: There is an arrow going from 1 to 0 and an arrow going from 0 to 3, but there is no arrow going from 1 to 3.

This means that there are elements of A—0, 1, and 3—such that 1 R 0 and 0 R 3 but 1 R 3. Hence R is not transitive.

b. The directed graph of S has the appearance shown below.

S is not reflexive: There is no loop at 1, for example. Thus $(1, 1) \notin S$, and so S is not reflexive.

S is not symmetric: There is an arrow from 0 to 2 but not from 2 to 0. Hence $(0, 2) \in S$ but $(2, 0) \notin S$, and so S is not symmetric.

S is transitive: There are three cases for which there is an arrow going from one point of the graph to a second and from the second point to a third.

c. The directed graph of *T* has the appearance shown below.



T is not reflexive: There is no loop at 0, for example. Thus $(0, 0) \notin T$, so *T* is not reflexive.

T is not symmetric: There is an arrow from 0 to 1 but not from 1 to 0. Thus $(0, 1) \in T$ but $(1, 0) \notin T$, and so *T* is not symmetric.

T is transitive: The transitivity condition is vacuously true for *T*. To see this, observe that the transitivity condition says that

For every $x, y, z \in A$, if $(x, y) \in T$ and $(y, z) \in T$ then $(x, z) \in T$.

The only way for this to be false would be for there to exist elements of *A* that make the hypothesis true and the conclusion false.

Properties of Relations on Infinite Sets

Example 8.2.3 – Properties of "Less Than"

Define a relation R on \mathbf{R} as follows: For all real numbers x and y,

$$x R y \Leftrightarrow x < y.$$

a. Is R reflexive? b. Is R symmetric? c. Is R transitive?

a. R is not reflexive: R is reflexive if, and only if, $\forall x \in \mathbf{R}$, $x \in \mathbf{R}$ $x \in \mathbf{R}$. By definition of R, this means that $\forall x \in \mathbf{R}$, x < x. But this is false: $\exists x \in \mathbf{R}$ such that $x \not < x$.

As a counterexample, let x = 0 and note that 0 < 0. Hence R is not reflexive.

b. R is not symmetric: R is symmetric if, and only if, $\forall x$, $y \in \mathbb{R}$, if x R y then y R x. By definition of R, this means that $\forall x, y \in \mathbb{R}$, if x < y then y < x.

But this is false: $\exists x, y \in \mathbf{R}$ such that x < y and $y \not < x$. As a counterexample, let x = 0 and y = 1 and note that 0 < 1 but $1 \not < 0$. Hence R is not symmetric.

c. R is transitive: R is transitive if, and only if, $\forall x, y, z \in R$, if x R y and y R z then x R z. By definition of R, this means that $\forall x, y, z \in R$, if x < y and y < z, then x < z. But this statement is true by the transitive law of order for real numbers. Hence R is transitive.

Example 8.2.4 – Properties of Congruence Modulo 3

Define a relation T on Z (the set of all integers) as follows: For all integers m and n,

$$m T n \Leftrightarrow 3 | (m-n).$$

This relation is called **congruence modulo 3**.

a. Is *T* reflexive? b. Is *T* symmetric? c. Is *T* transitive?

a. *T is reflexive*: To show that *T* is reflexive, it is necessary to show that For every $m \in \mathbb{Z}$, m T m.

By definition of *T*, this means that

For every
$$m \in \mathbb{Z}$$
, $3 \mid (m-m)$,

which is true because m - m = 0 and 3|0 (since $0 = 3 \cdot 0$). Hence T is reflexive.

b. *T is symmetric*: To show that *T* is symmetric, it is necessary to show that

For every $m, n \in \mathbb{Z}$, if m T n then n T m.

By definition of T this means that

For every
$$m, n \in \mathbb{Z}$$
, if $3 | (m-n)$ then $3 | (n-m)$.

Is this true? Suppose m and n are particular but arbitrarily chosen integers such that $3 \mid (m-n)$.

Must it follow that 3|(n-m)? [In other words, can we find an integer so that $n-m=3 \cdot (that integer)$?] By definition of "divides," since

$$3 | (m-n),$$

then m - n = 3k for some integer k.

The crucial observation is that n - m = -(m - n). Hence, you can multiply both sides of this equation by -1 to obtain

$$-(m-n)=-3k,$$

which is equivalent to

$$n-m=3(-k).$$

[Thus, we have found an integer, -k, so that $n - m = 3 \cdot (that integer)$.]

Since -k is an integer, this equation shows that $3 \mid (n-m)$.

It follows that *T* is symmetric.

c. *T is transitive*: To show that *T* is transitive, it is necessary to show that

For every $m, n, p \in \mathbb{Z}$, if m T n and n T p then m T p.

By definition of *T* this means that

For every $m, n \in \mathbb{Z}$, if 3 | (m-n) and 3 | (n-p) then 3 | (m-p).

Is this true? Suppose m, n, and p are particular but arbitrarily chosen integers such that $3 \mid (m-n)$ and $3 \mid (n-p)$.

Must it follow that 3|(m-p)? [In other words, can we find an integer so that $m-p=3 \cdot (that integer)$?] By definition of "divides," since

$$3 | (m-n)$$
 and $3 | (n-p)$,

then

m - n = 3r for some integer r,

and

n - p = 3s for some integer s.

The crucial observation is that (m - n) + (n - p) = m - p. Add these two equations together to obtain

$$(m-n) + (n-p) = 3r + 3s,$$

which is equivalent to

$$m - p = 3(r + s).$$

[Thus, we have found an integer so that $m - p = 3 \cdot (that integer)$.] Since r and s are integers, r + s is an integer. So, this equation shows that

$$3 | (m-p)$$
.

It follows that T is transitive.

The Transitive Closure of a Relation

The Transitive Closure of a Relation

The relation obtained by adding the least number of ordered pairs to ensure transitivity is called the *transitive closure* of the relation. More precisely, the transitive closure of a relation is the smallest transitive relation that contains the relation.

Definition

Let A be a set and R a relation on A. The **transitive closure** of R is the relation R^t on A that satisfies the following three properties:

- 1. R^{t} is transitive.
- 2. $R \subseteq R^t$.
- 3. If S is any other transitive relation that contains R, then $R^t \subseteq S$.

Example 8.2.5 – Transitive Closure of a Relation

Let $A = \{0, 1, 2, 3\}$ and consider the relation R defined on A as follows:

$$R = \{(0, 1), (1, 2), (2, 3)\}.$$

Find the transitive closure of *R*.

Every ordered pair in R is in R^t , so

$$\{(0, 1), (1, 2), (2, 3)\} \subseteq R^t$$
.

Thus, the directed graph of *R* contains the arrows shown below.



Since there are arrows going from 0 to 1 and from 1 to 2, R^t must have an arrow going from 0 to 2.

Hence $(0, 2) \in R^t$. Then $(0, 2) \in R^t$ and $(2, 3) \in R^t$, so since R^t is transitive, $(0, 3) \in R^t$. Also, since $(1, 2) \in R^t$ and $(2, 3) \in R^t$, then $(1, 3) \in R^t$. Thus R^t contains at least the following ordered pairs:

$$\{(0, 1), (0, 2), (0, 3), (1, 2), (1, 3), (2, 3)\}.$$

But this relation *is* transitive; hence it equals R^t . The directed graph of R^t is shown below.

