

- Set up ^{Discord} server for communication

Gaussian Elimination

We are trying to solve

$$Ax = b$$

is through elementary operations.

set an upper triangular matrix U such that $Ux = y$ is an equivalent linear system.

A_k - k^{th} leading principal submatrix of A obtained by intersecting the first k rows and columns.

To proceed we must assume that A_k is non-singular $k = 1, 2, \dots, n$

$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} & \dots & a_{1n} \\ a_{21} & a_{22} & a_{23} & \dots & a_{2n} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & a_{n3} & \dots & a_{nn} \end{bmatrix}$$

Note
WANT $a_{21} - M_{21}a_{11} = 0$
 $a_{21} = M_{21}a_{11}$
 $\frac{a_{21}}{a_{11}} = M_{21}$

In general,
For the i^{th} row
 $a_{i1} - M_{i1}a_{11} = 0$
 $M_{i1} = \frac{a_{i1}}{a_{11}}$

These are the multiples we need to get the zeros in the first column
 $Ax = b$

$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} & \dots & a_{1n} \\ M_{21} & a_{22}^{(1)} & a_{23}^{(1)} & \dots & a_{2n}^{(1)} \\ M_{31} & a_{32}^{(1)} & a_{33}^{(1)} & \dots & a_{3n}^{(1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ M_{n1} & a_{n2}^{(1)} & a_{n3}^{(1)} & \dots & a_{nn}^{(1)} \end{bmatrix} \quad \begin{bmatrix} (n-1) \times (n-1) \\ a_{22}^{(1)} & \dots & a_{2n}^{(1)} \\ \vdots & \ddots & \vdots \\ a_{n2}^{(1)} & \dots & a_{nn}^{(1)} \end{bmatrix}$$

$$M_{i2} = \frac{a_{i2}^{(1)}}{a_{22}^{(1)}}$$

$$\begin{bmatrix} a_{11}^{(1)} & a_{12}^{(1)} & \dots & a_{1n}^{(1)} \\ M_{32} & a_{22}^{(1)} & \dots & a_{2n}^{(1)} \\ M_{42} & a_{32}^{(1)} & \dots & a_{3n}^{(1)} \\ \vdots & \vdots & \ddots & \vdots \\ M_{n2} & a_{n2}^{(1)} & \dots & a_{nn}^{(1)} \end{bmatrix}$$

$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} & \dots & a_{1n} \\ M_{21} & a_{22}^{(1)} & a_{23}^{(1)} & \dots & a_{2n}^{(1)} \\ M_{31} & M_{32} & a_{33}^{(2)} & \dots & a_{3n}^{(2)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \end{bmatrix} b$$

After $n-1$ iterations we get an upper triangular matrix

$$A = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ m_{21} & a_{22} & \dots & a_{2n} \\ m_{31} & m_{32} & a_{33} & \dots & a_{3n} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ m_{n1} & m_{n2} & \dots & \dots & a_{nn} \end{bmatrix} b$$

We get an upper triangular matrix

What happens to b ?

$$b_i = b_i + \sum$$

$$b_1 = b_1$$

$$b_2^{(1)} = b_2 - m_{21} b_1$$

$$b_3^{(2)} = b_3 - m_{32} b_2^{(1)} - m_{31} b_1$$

$$b_3^{(2)} = b_3 - m_{32} (b_2 - m_{21} b_1) - m_{31} b_1$$

$$b_3^{(2)} = b_3 - m_{32} b_2 + m_{32} m_{21} b_1 - m_{31} b_1$$

What is the flop count to get the upper triangular form?

For each row (eliminating 1st column)

2n flops and n-1 rows

$$2n(n-1) + 2(n-1)(n-2) + \dots = 2 \sum_{k=1}^n k(k-1)$$

first iteration

The sum is approximately $\frac{2}{3} n^3$ flops

If you have an augmented matrix

$$2n^2 + 2(n-1)^2 + \dots = 2 \sum_{k=1}^n k^2 \approx \frac{2}{3} n^3$$

With back substitution there are n^2 flops.

The total cost of solving $Ax = b$ using this method is

$$\frac{2}{3} n^3 \text{ flops}$$

Example

$$A = \begin{bmatrix} 2 & 1 & -1 \\ 2 & 2 & -1 \\ 4 & -1 & 6 \end{bmatrix}$$

$$\text{and } b = \begin{bmatrix} 9 \\ 9 \\ 16 \end{bmatrix}$$

We want solve $Ax = b$

$$[A | b] = \left[\begin{array}{ccc|c} 2 & 1 & -1 & 9 \\ 2 & 2 & -1 & 9 \\ 4 & -1 & 6 & 16 \end{array} \right]$$

$$m_{21} = \frac{a_{21}}{a_{11}} = \frac{2}{2} = 1$$

$$m_{31} = \frac{a_{31}}{a_{11}} = \frac{4}{2} = 2$$

$$\rightarrow \left[\begin{array}{ccc|c} 2 & 1 & -1 & 9 \\ 1 & 1 & -2 & 0 \\ 2 & -3 & 4 & -2 \end{array} \right]$$

row3 - 2(row1)

row 3 - 2(row 1) \rightarrow $\begin{bmatrix} 2 & 1 & 1 & 9 \\ 1 & 1 & -2 & 0 \\ 2 & -3 & -2 & -2 \end{bmatrix}$

$$\rightarrow \begin{bmatrix} 2 & 1 & 1 & 9 \\ 1 & 1 & -2 & 0 \\ 2 & -3 & -2 & -2 \end{bmatrix}$$

$$m_{32} = \frac{a_{32}^{(1)}}{a_{22}^{(1)}} = \frac{-3}{1} = -3$$

The upper triangular matrix $\begin{bmatrix} 2 & 1 & 1 \\ 0 & 1 & -2 \\ 0 & 0 & -2 \end{bmatrix} x = \begin{bmatrix} 9 \\ 0 \\ -2 \end{bmatrix}$

$$x = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix}$$

$$\begin{aligned} 2x_1 + x_2 + x_3 &= 9 \\ 2x_1 + 2 + 1 &= 9 \\ 2x_1 &= 6 \\ x_1 &= 3 \end{aligned}$$

$$\begin{aligned} -2x_3 &= -2 \\ x_3 &= 1 \\ x_2 - 2x_3 &= 0 \\ x_2 - 2(1) &= 0 \\ x_2 &= 2 \end{aligned}$$

$$x = \begin{bmatrix} 3 \\ 2 \\ 1 \end{bmatrix}$$

Interpretation of the Multipliers

The operations of b

$$\begin{aligned} b_i^{(1)} &= b_i - m_{i1} b_1 \\ b_i^{(2)} &= b_i^{(1)} - m_{i2} b_2^{(1)} \end{aligned}$$

\vdots

$$b_i^{(n-1)} = b_i^{(n-2)} - m_{i,n-1} b_{n-1}^{(n-2)}$$

$$i = 2, 3, \dots, n$$

$$i = 3, 4, \dots, n$$

\vdots

$$i = n$$

At the end we have

$$\begin{bmatrix} b_1 \\ b_2^{(1)} \\ b_3^{(2)} \\ \vdots \\ b_n^{(n-1)} \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{bmatrix} = y$$

If we have $Ax = b$ could be solved using
on equivalent upper-triangular system $Ux = y$.
in the original

If we have an equivalent upper-triangular system $Ux = y$ where U is the upper triangular matrix obtained in the original reduction of A .

Example

Suppose we want to solve

$$Ax = \vec{b} \quad \text{where}$$

$$A = \begin{bmatrix} 2 & 2 & -1 \\ 2 & 2 & -1 \\ 4 & -1 & 6 \end{bmatrix}$$

$$\vec{b} = \begin{bmatrix} 7 \\ 3 \\ 20 \end{bmatrix}$$

We already know $U = \begin{bmatrix} 2 & 1 & 1 \\ 0 & 1 & -2 \\ 0 & 0 & -2 \end{bmatrix}$

$$b_2^{(1)} = 3 - (1)(7) = -4$$

$$b_3^{(1)} = 20 - 2(7) = 6 \Rightarrow b_3^{(2)} = 6 - (-3)(b_2^{(1)}) = 6 + 3(-4) = -6$$

$$Ux = \begin{bmatrix} 7 \\ -4 \\ -6 \end{bmatrix}$$

$$\begin{bmatrix} 2 & 1 & 1 \\ 0 & 1 & -2 \\ 0 & 0 & -2 \end{bmatrix} x = \begin{bmatrix} 7 \\ -4 \\ -6 \end{bmatrix}$$

$$x = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix}$$

$$-2x_3 = -6$$

$$x_3 = 3$$

$$x_2 - 2(3) = -4$$

$$x_2 = 2$$

$$2x_1 + 2 + 3 = 7$$

$$2x_1 = 2$$

$$x_1 = 1$$

$$x = \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix}$$

We start out with $Ax = b$

and end up with $Ux = y$

We know

$$b_i^{(1)} = b_i - m_{i1}y_1 \quad i = 2, 3, \dots, n$$

$$b_i^{(2)} = b_i^{(1)} - m_{i2}y_2 \quad i = 3, 4, \dots, n$$

$$b_i^{(n-1)} = b_i^{(n)} - m_{i,n-1}y_{n-1} \quad i = n$$

This leads us to

$$y_i = b_i - \sum_{j=1}^{i-1} m_{ij} y_j \quad i=1, 2, \dots, n$$

Let's

set

$$L = \begin{bmatrix} 1 & 0 & \dots & 0 \\ m_{21} & 1 & \dots & 0 \\ m_{31} & m_{32} & 1 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ m_{n1} & m_{n2} & m_{n3} & \dots & 1 \end{bmatrix}$$

Unit
Lower triangular
matrix

$$Ly = b$$

In summary, we start with $Ax = b$

we get $Ux = y$
where U is upper triangular and y is the solution
of a unit lower triangular system

$$Ly = b$$

$$Ux = y$$

$$LUx = Ly$$

$$LUx = b$$

$$Ax = b$$

$$\text{Thus } A = LU$$

Example Solve the system

$$A = \begin{bmatrix} 2 & 1 & -1 \\ 2 & 2 & -1 \\ 4 & -1 & 6 \end{bmatrix}$$

$Ax = b$ where

$$\text{and } b = \begin{bmatrix} 3 \\ 0 \\ 11 \end{bmatrix}$$

$$U = \begin{bmatrix} 2 & 1 & -1 \\ 0 & 1 & -2 \\ 0 & 0 & -2 \end{bmatrix}$$

$$L = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 2 & -3 & 1 \end{bmatrix}$$

$$\left[\begin{array}{ccc|c} 1 & 0 & 0 & 3 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & -2 & 11 \end{array} \right] \left[\begin{array}{c} y_1 \\ y_2 \\ y_3 \end{array} \right] = \left[\begin{array}{c} 3 \\ 0 \\ 11 \end{array} \right]$$

$$Ly = b \quad \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 2 & -3 & 1 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix} = \begin{bmatrix} 3 \\ 0 \\ 11 \end{bmatrix}$$

Forward substitution

$$y_1 = 3$$

$$y_1 + y_2 = 0$$

$$y_2 = -3$$

$$Ux = y$$

$$\begin{bmatrix} 2 & 1 & 1 \\ 0 & 1 & -2 \\ 0 & 0 & -2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 3 \\ -3 \\ -4 \end{bmatrix}$$

$$x_3 = 2$$

$$x_2 - 2(2) = -3$$

$$x_2 = 1$$

$$2x_1 + x_2 + x_3 = 3$$

$$2x_1 + 1 + 2 = 3$$

$$x_1 = 0$$

$$x = \begin{bmatrix} 0 \\ 1 \\ 2 \end{bmatrix}$$

$$2y_1 - 3y_2 + y_3 = 11$$

$$2(3) - 3(-3) + y_3 = 11$$

$$15 + y_3 = 11$$

$$y_3 = -4$$

Theorem (LU decomposition Theorem)

Let A be an $n \times n$ matrix whose leading principal submatrices are all nonsingular. Then A can be decomposed in exactly one way into a product $A = LU$ such that L is a unit lower triangular and U is upper triangular.

Proof start of class on Thursday.

Cholesky Decomposition

If we have a positive definite matrix

that A is $n \times n$ real symmetric satisfies the property

$$x^T A x > 0$$

that it is ...
property $x^T A x > 0$

Theorem (Cholesky Decomposition Theorem) Let A be positive definite. Then A can be decomposed in exactly one way into a product

$$A = R^T R$$

such that R is upper triangular and has all main diagonal entries r_{ii} positive. R is called the Cholesky factor of A .

Example Let $A = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 2 & 2 \\ 1 & 2 & 3 \end{bmatrix}$

and $R = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \end{bmatrix}$

show $A = R^T R$

$$R^T = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 1 \end{bmatrix}$$