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Cholesky Decomposition - Locaks of A

is symmetric
and positive definite

Coustin Flimination (contract pivots) works to find solution under certain conditions

Gaussian Elimination (with pivots) — works well to that solutions to Ax=bIf we about have an ill-conditioned system.

The most extreme case of an ill-conditioned system is where A is singular, det(A)=0

Ax=b does not have one unique solution,

n A = 4 singular madrix in the process of

If A is a singular matrix in the process of Gaussian elémination there will be a step where all possible pivots are zera

But with each step of the algorithm the entries of the matrix are subject to round off

And we may arrive at a "solution".

But it is not in fact a solution,

The additional cost of cow interchanges is not great small we make  $n^2/2$  comparisons is relatively 2 compared,

O(n3) of the algorithm,

Roughly the cost of Gaussian dimination with partial privating is 3n3 flops.

to exemple from last Phursday

A = 
$$\begin{bmatrix} 0 & 4 & 1 \\ 1 & 1 & 3 \\ 2 & -2 & 1 \end{bmatrix}$$
 and we some  $A = \begin{bmatrix} 2 & -2 & 1 \\ 0 & 4 & 1 \\ 1 & 1 & 3 \end{bmatrix}$ 

Find P such that  $\widehat{A} = PA$  $\begin{bmatrix} 0 & 0 & 1 & 0 & 4 & 1 \\ 1 & 6 & 0 & 27 & -2 & 1 \\ 0 & 1 & 0 & 27 & -2 & 1 \end{bmatrix} = \begin{bmatrix} 2 & -2 & 1 & 1 \\ 0 & 4 & 1 & 1 \\ 1 & 1 & 3 & 1 \end{bmatrix}$ 

Definition A permutation matrix is a matrix that has exactly one I in each row one in each Column, all other entires are zero.

Exercise Show that if P is a pomutitus motifix PTP = PPT = I, Mus P is non singular tren

then  $P^TP = PP^T = I$ , Thus P is mon singular and  $\rho^{-1} = \rho^{T}$  $P = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} \qquad P^{T} = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix}$  $PP^{T} = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$ Let P be a permutation matrix that is nxn Let Q = PT PQ is = (row i of P) x (column J of Q) This is nonzero if and only if Since I one in each row pone I are In each column of Q. Since Q = PT it would be nonzero If i=j. POis = EPix Oxi Cet's assume its and Paij = 1 run rue exists a x 12x2n Pik = QKJ = 1 Such that Since Okj = 1 then Pjk=1 since Q=pT But we said Pik=1 and itj. So Phes two I's in column K. This contradicts Pheny a permutation And I=5 must be me for matrix, Pais to be non zero. Theorem Gaussian elimination with partial privoting matrix A wooduces

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Theorem Gaussian elimination with partial pivoling

applied to an axa matrix A produces
a unit lower triangular matrix L such that

[listed, an upper triangular matrix U, and
a permutation matrix P such that

A = PTLU

Exercise  $A = \begin{bmatrix} 2 & 2 & -4 \\ 1 & 1 & 5 \\ 1 & 3 & 6 \end{bmatrix}$ 

Find P, L, and U

Such Mg+  $A = P^T L U$ we already know  $L = \begin{bmatrix} 1 & 0 & 0 \\ \frac{1}{2} & 0 & 1 \end{bmatrix}$ 

 $LU = A = \begin{bmatrix} 2 & 2 - 4 \\ 1 & 3 & 6 \\ 1 & 1 & 5 \end{bmatrix}$  and  $U = \begin{bmatrix} 2 & 2 - 4 \\ 0 & 2 & 4 \\ 0 & 0 & 3 \end{bmatrix}$  $P = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix}$ 

Thus  $p^{T} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$   $A = P^{T}LU \quad \text{where} \quad P^{T} = P = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$   $A = V^{T}LU \quad \text{where} \quad P^{T} = P = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$   $A = V^{T}LU \quad \text{where} \quad P^{T} = P = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$   $A = V^{T}LU \quad \text{where} \quad P^{T} = P = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$   $A = V^{T}LU \quad \text{where} \quad P^{T} = P = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$   $A = V^{T}LU \quad \text{where} \quad P^{T} = P = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$   $A = V^{T}LU \quad \text{where} \quad P^{T} = P = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$ 

Computing A-1
We know A-1A = I

$$I_{\chi} = \chi$$

$$(A^{-1}A)\chi = \chi$$

A X = I where X is the columns of X.

A  $[X_1 X_2 X_3 X_4... X_n] = [e_1...e_n]$  some matrix X and X gives the columns of X.

AX: = e; i=1, -- 1 If we solve those n-systems we obtain A.

Complete Pivoting

Here both row and column interelarges are allowed,

allowed,

allowed,

we start by finding the largest (0:3) and move to 0...

after Zerbing out the first column we know bous on an (n-1) x(n-1) matrix and report the process

· Rouides extra protection from roundoff error

. The disadvantage is seeking out the largest entry.

step 1 You must compare nº entries step 2 You must compare (n-1)2 entries

Z (n-K) 2 2 3 /3. Comparison 9,

Since pertial privating works fairly this is the more often used algorithm,

I Diemein in Comesian Flimmoth.

Error Discussion in Gaussian Elimnation

· Propagation of randoff error in Gaussian Elimination.

Example Consider the ill-conditioned matrix  $A = \begin{bmatrix} 1000 & 9997 \\ 999 & 998 \end{bmatrix}$ 

How would we get Land U?

ROW 2 - 999 ROW |

[ 1000 999 ] 998 - 999 (999) 0 -0.001 ] 998 - 998.001 = -0.001

Example The ill-conditioned Hilbert matrices, defined by his = (15-1)

for example  $t_{4} = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 2 & 1 & 1 & 1 \\ 3 & 1 & 1 & 1 \\ 2 & 1 & 1 & 1 \\ 3 & 1 & 1 & 1$ 

See if you can find Lad

U= [1 \frac{1}{2} \frac{1}{3} \frac{1}{4}]
0 \frac{1}{2} \frac{1}{2} \frac{3}{340}

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