report

Implementation

a. How do you implement the matrix multiplication (both version)?

as shown,

standard version:

blocked version:

b. How do you verify the correctness of your program?

the standard matrix multiplication program is 100% correct (has been verified by online judge), so i take it as standard program.

then i wrote a test data generator, put the test data into standard program, generating standard answers.

use diff command to find the difference between standard answers and other version's answer.

if diff detect difference, use another program to calculate the error.

c. Why the blocked version is correct? Answer by explaining the mathematics or program's calculation steps.

block matrix multiplication says:

$$A = egin{bmatrix} A_{1,1} & A_{1,2} & \dots & A_{1,k} \ A_{2,1} & \dots & & \ \dots & \dots & & \ A_{n,1} & \dots & & \ \end{bmatrix}$$

$$B = egin{bmatrix} B_{1,1} & B_{1,2} & \dots & B_{1,m} \ B_{2,1} & \dots & & \ & \dots & & \ B_{k,1} & \dots & & \ \end{pmatrix}$$

the multiplication AB can be performed blockwise.

$$C_{q,r} = \sum_{i=1}^{k'} A_{q,i} B_{i,r}$$

for example:

$$A = egin{bmatrix} 1 & 1 & 2 & 2 \ 1 & 1 & 2 & 2 \ 3 & 3 & 4 & 4 \ 3 & 3 & 4 & 4 \end{bmatrix}$$

blocked

$$A = \begin{bmatrix} A_{1,1} & A_{1,2} \\ A_{2,1} & A_{2,2} \end{bmatrix} \text{, where } A_{1,1} = \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix} \text{, } A_{1,2} = \begin{bmatrix} 2 & 2 \\ 2 & 2 \end{bmatrix} \text{, } A_{2,1} = \begin{bmatrix} 3 & 3 \\ 3 & 3 \end{bmatrix} \text{, } A_{2,2} = \begin{bmatrix} 4 & 4 \\ 4 & 4 \end{bmatrix}$$

$$A^2 = egin{bmatrix} A_{1,1} & A_{1,2} \ A_{2,1} & A_{2,2} \end{bmatrix} egin{bmatrix} A_{1,1} & A_{1,2} \ A_{2,1} & A_{2,2} \end{bmatrix} = egin{bmatrix} A_{1,1}A_{1,1} + A_{1,2}A_{2,1} & A_{1,1}A_{1,2} + A_{1,2}A_{2,2} \ A_{2,1}A_{1,2} + A_{2,2}A_{2,2} \end{bmatrix} = egin{bmatrix} C_{1,1} & C_{1,2} \ C_{2,1} & C_{2,2} \end{bmatrix} = egin{bmatrix} 1 \ 1 \ 3 \ 3 \end{bmatrix}$$

the blocked version operates like the mathematics.

Experiment

environment:

- hardware
 - CPU: AMD Ryzen 7 PRO 4750U with Radeon Graphics @ 1.70 GHz
 - RAM: DDR4 3200 MHz
 - Cache:
 - L1: 8 * 32 KB, 8-way
 - L2: 8 * 512 KB, 8-way

- L3: 2 * 4 MB, 16-way
- software
 - compiler version: g++ (Ubuntu 9.4.0-1ubuntu1~20.04.1) 9.4.0
 - compile command:
 - g++ -Wall -OX -c matrix.cpp
 - g++ -Wall -OX -DBLOCK=XX -c blocked-matrix.cpp
 - g++ -Wall -OX -DBLOCK=XX -c blocked-cache-matrix.cpp
 - g++ -msse -msse2 -msse3 -msse4 -msse4a -msse4.1 -msse4.2
 -Wall -OX -c vector-matrix.cpp
 - OS: Linux emilia 5.13.0-40-generic #45~20.04.1-Ubuntu SMP Mon Apr 4 09:38:31 UTC 2022 ×86_64 ×86_64 ×86_64 GNU/Linux

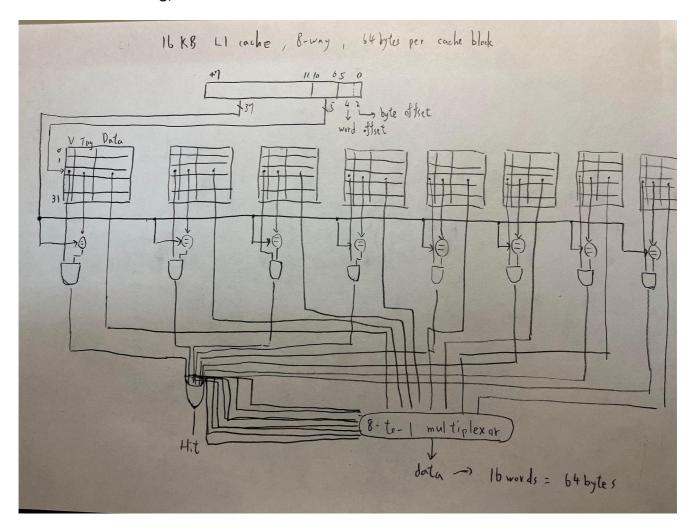
a.

my L1 cache is 32 KB, so the cache below is 16 KB.

the index counts is $16 \times 2^{10}/(8 \times 64) = 2^5$, so 5 bits is for index.

the cache line is 64 byte, namely 16 words, so the width of word offset is 4-bit, and the byte offset is 2-bit.

the left bit is for tag, which is 48 - 5 - 4 - 2 = 37-bit.



the test data have different sizes, whose n=k=m=192,384,576,768,960, produced by the geneator.

partial code of geneator:

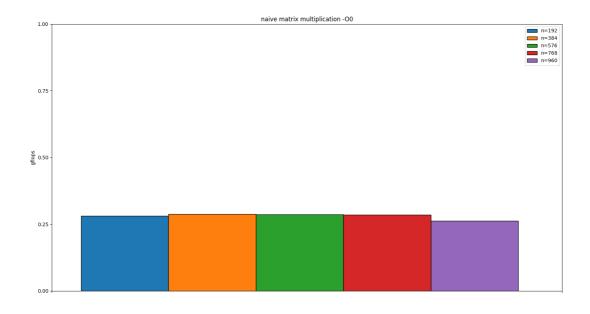
```
std::random_device rd;
std::default_random_engine eng(rd());
std::uniform_real_distribution<float> unif(-1000, 1000);

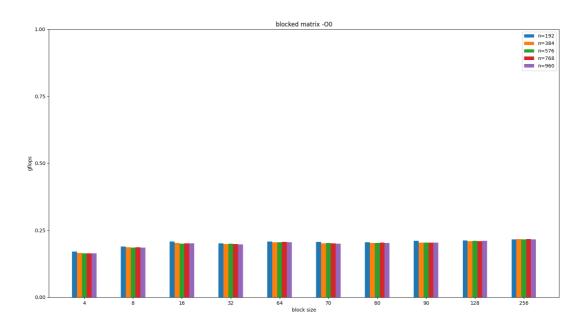
n = k = m = 192 * CNT;
std::cout << n << ' ' << k << ' ' << m << '\n';
for (int i = 0; i < n; ++i)
{
    for (int j = 0; j < k; ++j)
        std::cout << unif(eng) << ' ';
    std::cout << '\n';
}
for (int i = 0; i < k; ++i)
{
    for (int j = 0; j < m; ++j)
        std::cout << unif(eng) << ' ';
    std::cout << unif(eng) << ' ';
    std::cout << '\n';
}</pre>
```

C.

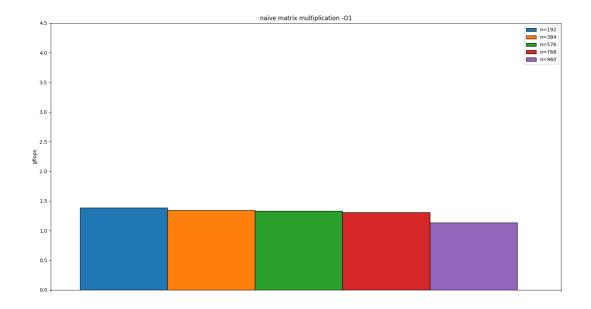
matrix multiplication programs will run 10 times on every test data (n=192,384,576,768,960), every optimization flags (O0,O1,O2,O3,Ofast), every block size (4,8,16,32,64,70,80,90,128,256), and record the execution time of MMUL(). later, a python program calculates average execution time, maximum/minimum execution time, standard deviation of execution time, and average GFLOPS.

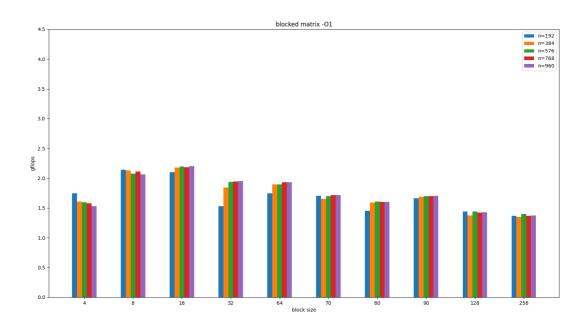
GFLOPS of standard version and blocked version in -00:



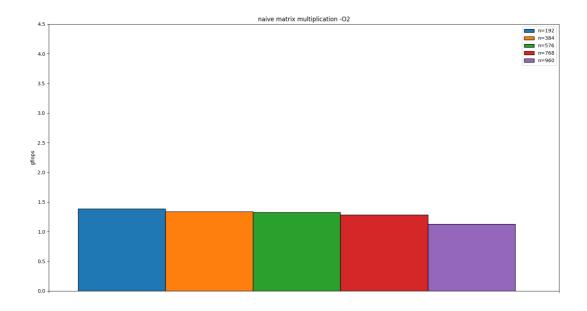


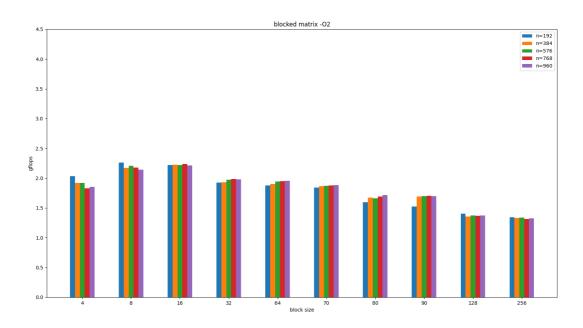
GFLOPS of standard version and blocked version in -O1:



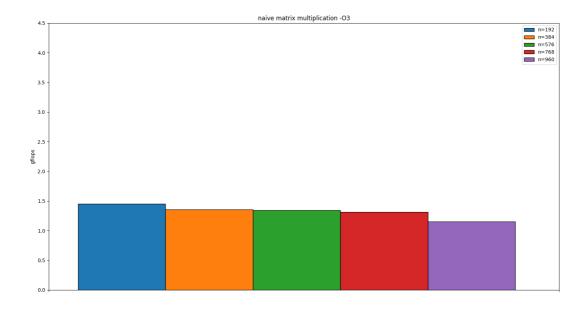


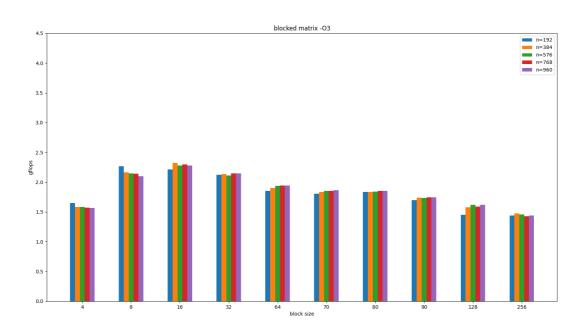
GFLOPS of standard version and blocked version in -O2:



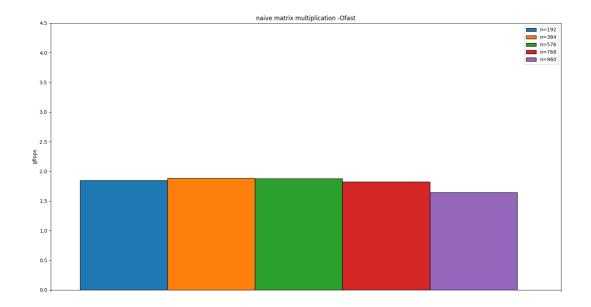


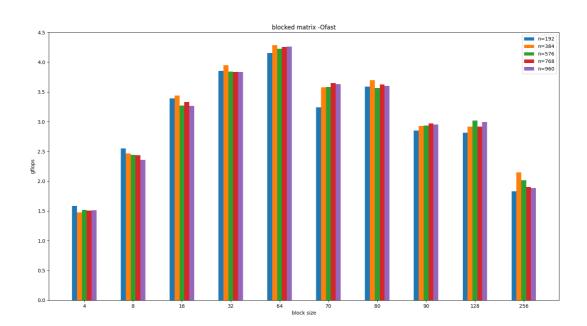
GFLOPS of standard version and blocked version in -O3:





GFLOPS of standard version and blocked version in -Ofast:





the number of float-point operations on a matrix multiplication of AB=C is $n\times k\times m$, where A is a $n\times k$ matrix and B is a $k\times m$ matrix.

so FLOPS for a matrix multiplication is
$$\frac{n \cdot k \cdot m}{time(s)} = \frac{n \cdot k \cdot m}{time \times 10^9 (ns)} \times 10^9 = \frac{n \cdot k \cdot m}{time \times 10^9 (ns)}$$
 GFLOPS

d.

i use chrono> in std c++ 11 to measure the executing time of matrix multiplication function.

```
auto begin = std::chrono::high_resolution_clock::now();
MMUL();
auto end = std::chrono::high_resolution_clock::now();
```

```
auto T = std::chrono::duration_cast<std::chrono::nanoseconds>(end -
begin);
```

i think it's an appropriate way to measure "a function" executing time.

but if i want to measure "a program" execution time, the command, time, provided by linux is the better way.

moreover, the linux kernel tool, perf, is an even better way to measure a process executing time and collect some CPU info in executing program, e.g. cache miss, cache reference, instruction counts.

e.

the difference between standard and blocked version are, the counts of **cache references** and **cache misses**.

since the blocked version divides the big matrix into pieces, the cache misses counts is lower than the standard one.

for example, this is the CPU info of standard version in n=960:

and the CPU info of blocked version in n=960 and block size is 128:

```
4.732 +- 0.284 seconds time elapsed ( +- 6.01% )
```

same version but block size is 64:

can be seen that the **cache misses in blocked version is less than standard version**_.

let's look into the impact of different data size:

• n = 192

standard version:

blocked version, block size = 128:

```
Performance counter stats for './blocked-matrixB12800' (10 runs):

98,791 cache-misses # 3.204 % of all cache
```

• n = 384

standard version:

blocked version, block size = 128:

standard version:

blocked version, block size = 128:

• n = 768

standard version:

```
2.139 +- 0.175 seconds time elapsed ( +- 8.20% )
```

blocked version, block size = 128:

• n = 960

standard version:

blocked version, block size = 128:

can be seen the cache misses/references in blocked version is much less than standard one, and, though these two counts increase when test data size increases, the amount of increasing in blocked version is much less than standard one.

obviously, the blocked version exploit more on the cache space.

so i expect the execution time of blocked version is shorter than standard one.

however, the results are not as expected.

in the O0 flag, the performance of blocked version is worse than the standard one, because the instruction counts is much more than standard one.

the turth is, without the compiler optimization, blocked version takes no advantage.

a for-loop generally need three steps:

- 1. check the stop condition
- 2. do something in the for-loop section
- 3. change the counting variable

the additional three for-loop has increase the count of checking the stop condition, and not to mention the burden added by the std::min() function.

hence, the count of instructions is much more than standard one, consequently, the GFLOPS performace is worse than the standard one.

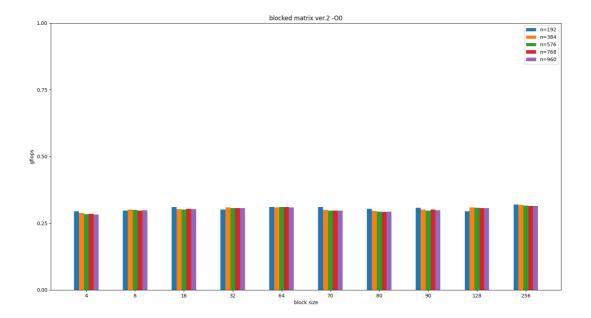
there's a simple way to imporve the performance by modifying the code:

```
#define min(a,b) ((a)>(b)?(b):(a))
inline void MMUL() noexcept
    int x, y, z, xx, yy, zz, maxx, maxy, maxz;
    for (x = 0; x < n; x += BLOCK)
        for (y = 0; y < m; y += BLOCK)
            for (z = 0; z < k; z += BLOCK)
            {
                maxx = min(n, x + BLOCK);
                maxy = min(m, y + BLOCK);
                maxz = min(k, z + BLOCK);
                for (xx = x; xx < maxx; ++xx)
                    for (yy = y; yy < maxy; ++yy)
                        for (zz = z; zz < maxz; ++zz)</pre>
                             C[xx][yy] += A[xx][zz] * B[zz][yy];
            }
}
```

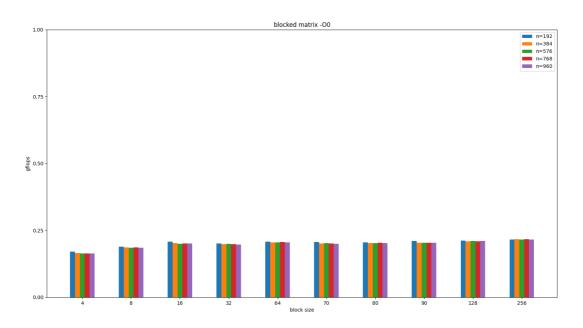
a brief example in block size = 128:

can be seen the instructions is reduced.

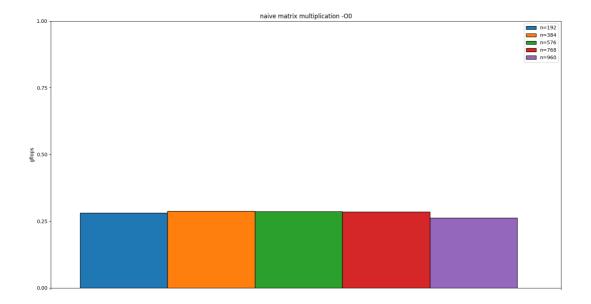
here's the GFLOPS performance in -O0 of modified version:



comparing to the unmodifying one:



and standard one:



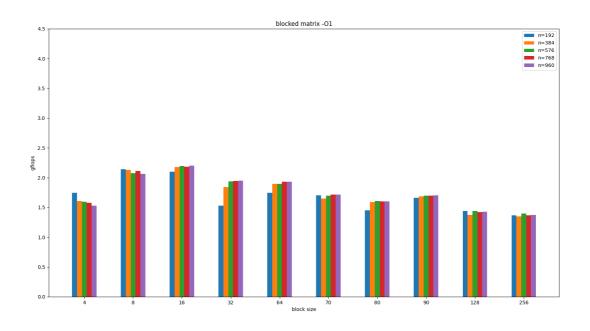
can see the performace is better than both.

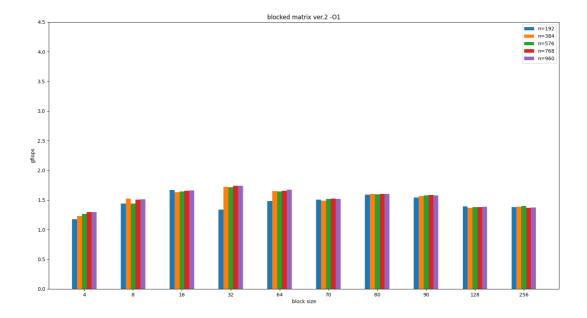
there's a more convenient way to optimize without modifying the program, by using gcc optimization flag, O1, O2, O3, Ofast, the performance can be imporved a lot.

one more thing to mention, there's a weird phenomenon that when using O1 optimization on the modified blocked version, the performance is worse than the unmodified one. but in the later optimization (O2,O3,Ofast) the performance is as good as the other one.

this phenomenon is shown below:

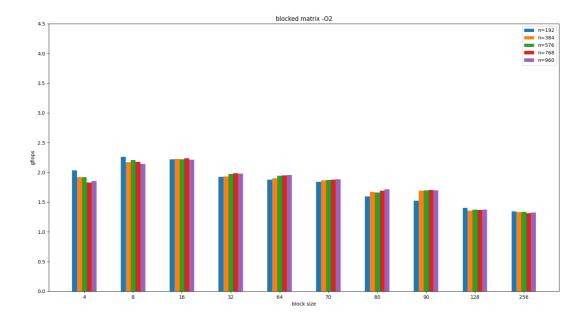
GFLOPS performace of blocked version and modified blocked version in O1:

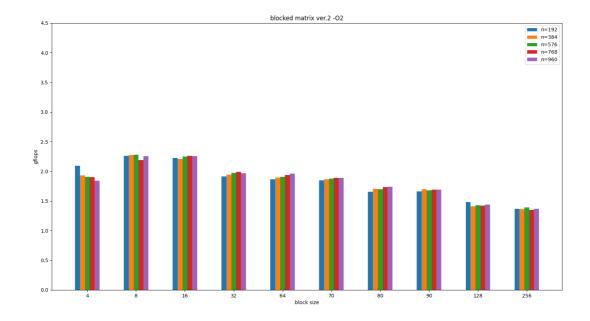




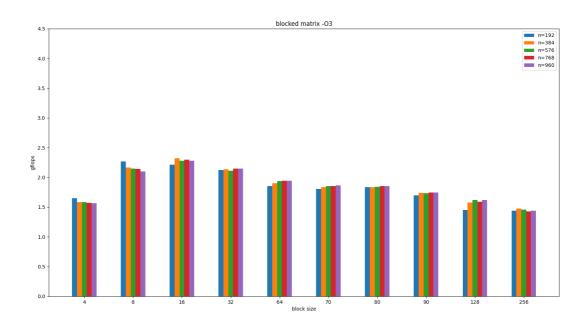
can be seen the performace is worse.

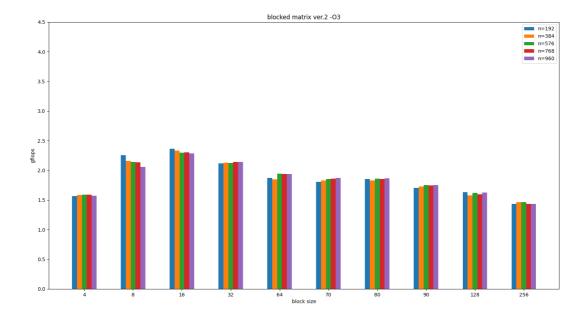
GFLOPS performace of blocked version and modified blocked version in O2:



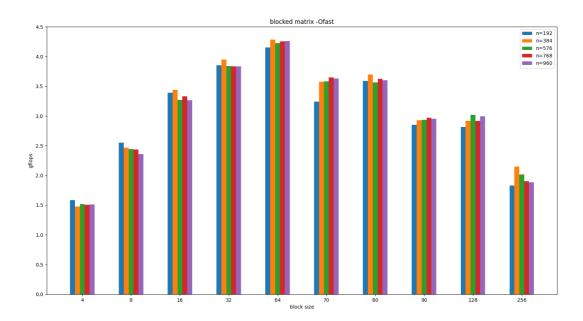


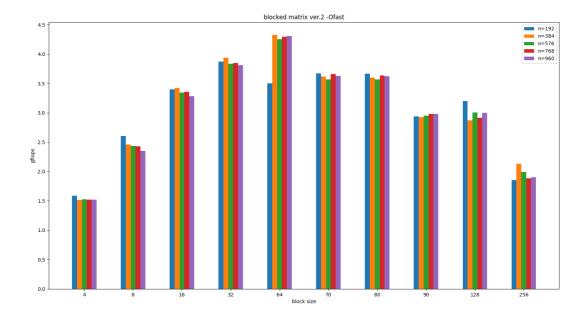
GFLOPS performace of blocked version and modified blocked version in O3:





GFLOPS performace of blocked version and modified blocked version in Ofast:





i think the reason why is that the compiler optimization in O1 has some conflicts with the manual optimization, so that the compiler can't 100% optimize the program. but in O2 and others, the optimization is aggressive enough to optimize the whole program.

so the conclusion is that, since the additional instructions in three-extra for-loop, blocked version takes advantage **only when we do some extra optimization**.

f.

00: no optimization. it's default flag if no optimization level is specified.

o1: optimize minimally. the compiler tries to reduce code size and execution time, but performing no optimizations that take a great deal of compilation time.

o2: optimize more. GCC performs nearly all supported optimizations that do not involve a space-speed tradeoff. As compared to -O1, this option increases both compilation time and the performance of the generated code.

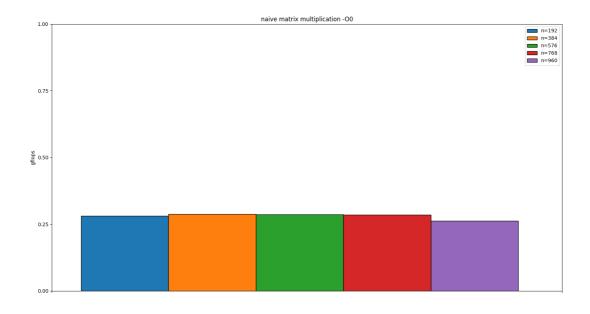
03: optimize even more.

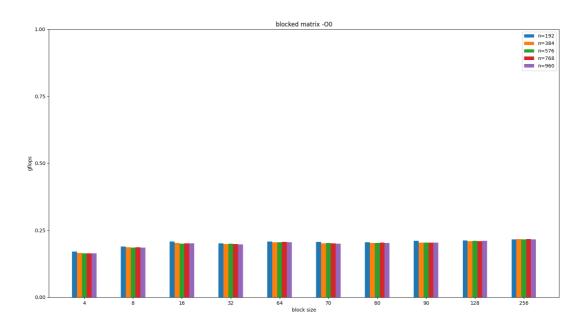
Ofast: optimize very aggressively that disregard strict standards compliance.

basicly, optimize more means faster execution time.

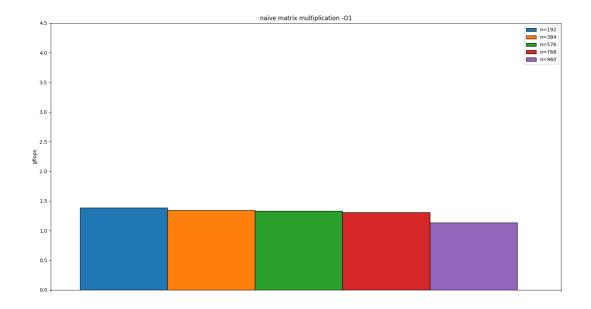
but sometimes optimization can cause unexpected result. e.g. Ofast will change the floating number operation order, make the result different.

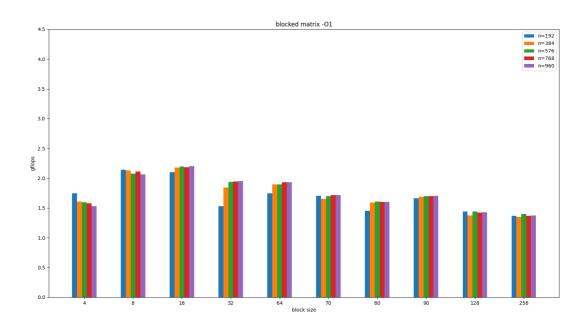
GFLOPS of standard version and blocked version in -00:



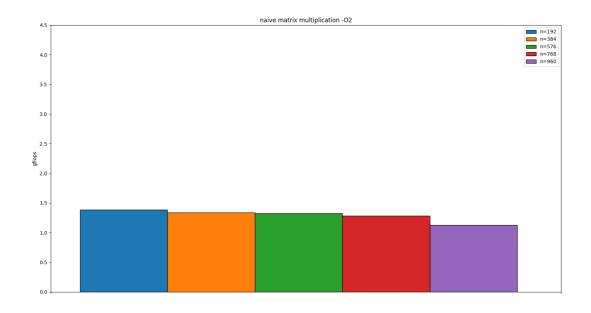


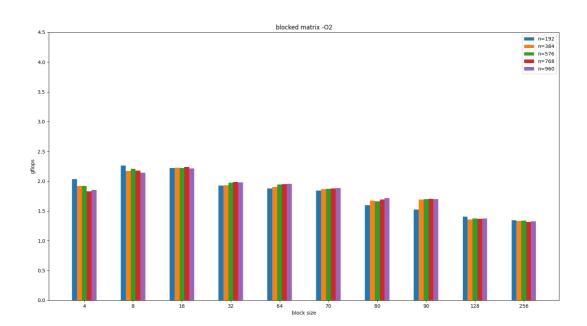
GFLOPS of standard version and blocked version in -O1:



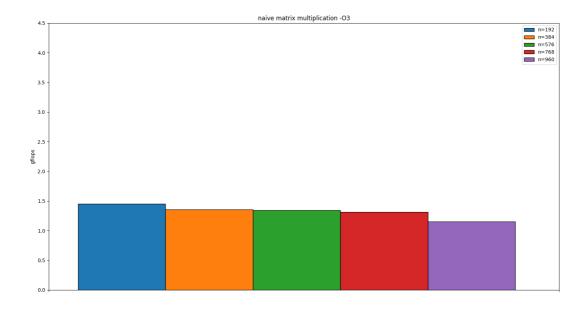


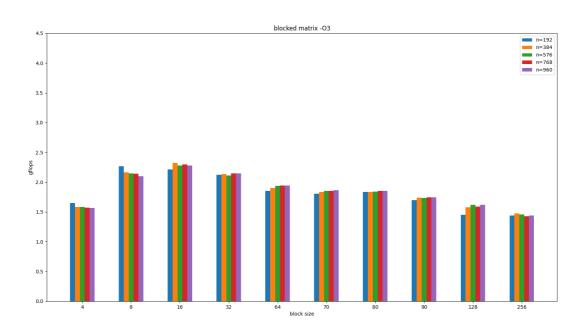
GFLOPS of standard version and blocked version in -O2:



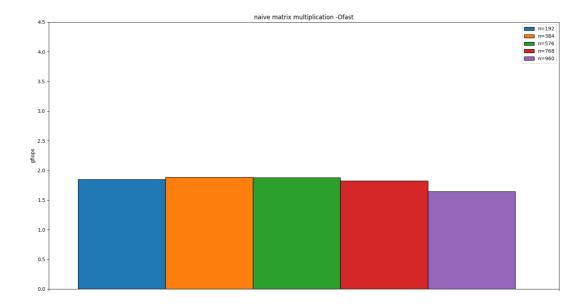


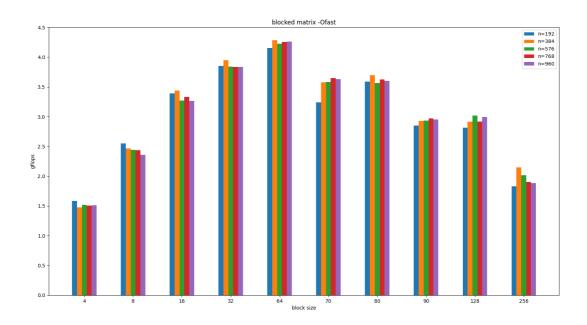
GFLOPS of standard version and blocked version in -O3:





GFLOPS of standard version and blocked version in -Ofast:





advanced questions

a.

this problem has two solutions:

1. change for-loop order:
the second and third, the fifth and sixth loop have been exchanged.
so in every iteration, the array B reads the next column instead of next row.

```
inline void MMUL() noexcept
{
   for (int x = 0; x < n; x += BLOCK)
      for (int z = 0; z < k; z += BLOCK)</pre>
```

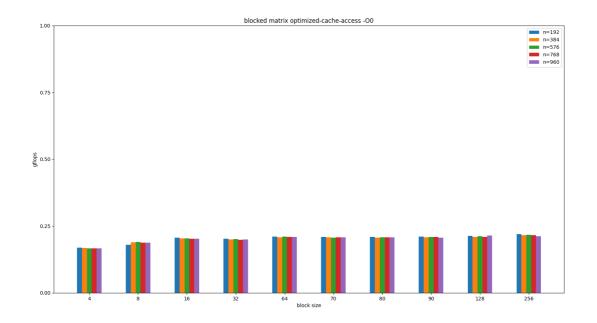
2. transpose matrix B:

when read the data into array B, transposing it.

when performing matrix multiplication, transposing it again.

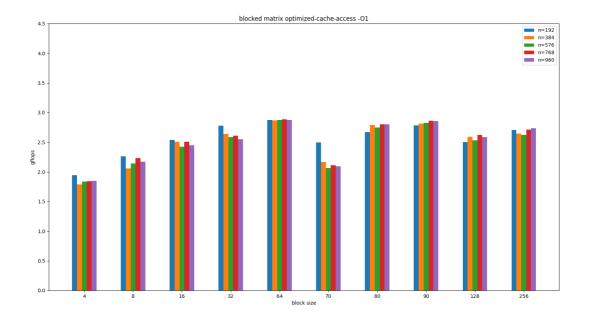
when performing matrix multiplication, transposing it again. so that in every iteration the array B reads the next column instead of next row.

GFLOPS perfomance of optimized-cache blocked matrix in O0:



it doesn't go beyond the standard version but same as blocked one, since the instruction counts is much more than standard one.

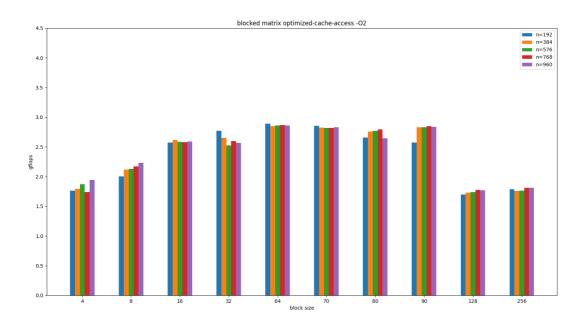
GFLOPS perfomance of optimized-cache blocked matrix in O1:



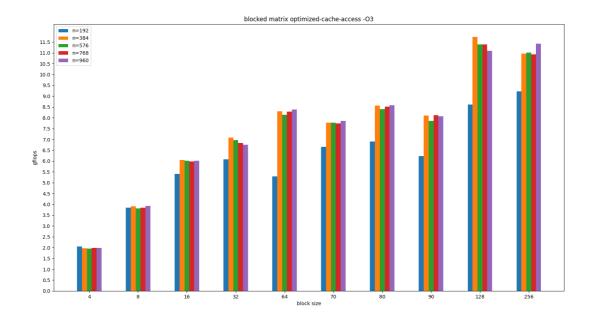
```
1.1624 +- 0.0892 seconds time elapsed ( +- 7.68% )
```

can see the performance is far beyond the blocked one and standard one.

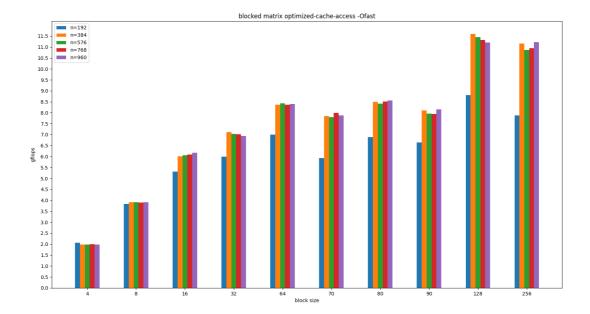
GFLOPS perfomance of optimized-cache blocked matrix in O2:



GFLOPS perfomance of optimized-cache blocked matrix in O3:



GFLOPS perfomance of optimized-cache blocked matrix in Ofast:



b.

1. ver. 1:

same as blocked version, whose block size is 4, but implement in SSE.

```
#include <xmmintrin.h>
#include <pmmintrin.h>
#include <tmmintrin.h>
#include <nmmintrin.h>
#include <nmmintrin.h>
#include <immintrin.h>

constexpr int MAXSIZE = 960;

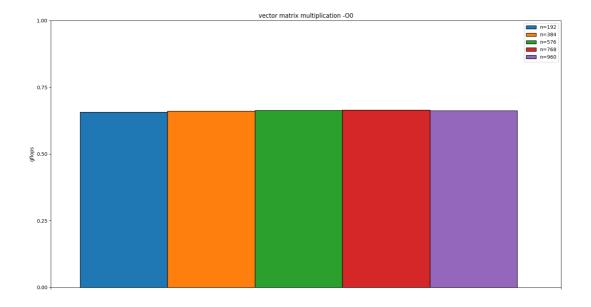
// SSE requires the data shall be aligned with 16
float A[MAXSIZE][MAXSIZE] __attribute__((aligned(16)));
```

```
float B[MAXSIZE] [MAXSIZE] __attribute__((aligned(16)));
float C[MAXSIZE] [MAXSIZE] __attribute__((aligned(16)));
int n, k, m;
inline void MMUL() noexcept
        for (int x = 0; x < n; x += 4)
                 for (int y = 0; y < m; y += 4)
                          _{m128} r0 = _{mm}load_{ps}(&C[x+0][y]);
                          __m128 r1 = _mm_load_ps(&C[x+1][y]);
                          _{m128} r2 = _{mm}load_{ps}(&C[x+2][y]);
                          _{m128} r3 = _{mm}load_{ps}(&C[x+3][y]);
                          for (int z = 0; z < k; z += 4)
                                   // load A[x:x+3][k:k+3] and B[k:k+3]
[y:y+3]
                                   _{m128 I0} = _{mm}load_{ps}(&A[x+0][z]);
                                   _{m128} I1 = _{mm}load_ps(&A[x+1][z]);
                                   _{m128} I2 = _{mm}load_ps(&A[x+2][z]);
                                   _{m128} I3 = _{mm}load_{ps}(&A[x+3][z]);
                                   _{m128} I4 = _{mm\_set\_ps}(B[z+3][y+0], B[z+2]
[y+0], B[z+1][y+0], B[z+0][y+0];
                                   _{m128} I5 = _{mm\_set\_ps}(B[z+3][y+1], B[z+2]
[y+1], B[z+1][y+1], B[z+0][y+1]);
                                   _{m128} I6 = _{mm_set_ps}(B[z+3][y+2], B[z+2]
[y+2], B[z+1][y+2], B[z+0][y+2]);
                                   _{m128} I7 = _{mm\_set\_ps}(B[z+3][y+3], B[z+2]
[y+3], B[z+1][y+3], B[z+0][y+3];
                                   // calculate
                                   _{m128} T0 = _{mm_{ul_ps}(I0, I4)};
                                   _{m128} T1 = _{mm_{ul_ps}(I0, I5)};
                                   _{m128} T2 = _{mm_{mul_ps}(I0, I6)};
                                   _{m128} T3 = _{mm_{ul_ps}(I0, I7)};
                                   _{m128} T4 = _{mm_{ul_ps}(I1, I4)};
                                   __m128 T5 = _mm_mul_ps(I1, I5);
                                   _{m128} T6 = _{mm_{ul_ps}(I1, I6)};
                                   _{m128} T7 = _{mm_{ul_ps}(I1, I7)};
                                   _{m128} T8 = _{mm_{mul_ps}(I2, I4)};
                                   _{m128} T9 = _{mm_{ul_ps}(I2, I5)};
                                   _{m128} T10 = _{mm_{mul_ps}(I2, I6)};
                                   _{m128} T11 = _{mm_{nul_ps}(I2, I7)};
                                   __m128 T12 = _mm_mul_ps(I3, I4);
                                   _{m128} T13 = _{mm_{mul_ps}(I3, I5)};
                                   __m128 T14 = _mm_mul_ps(I3, I6);
```

```
__m128 T15 = _mm_mul_ps(I3, I7);
// transpose
__m128 T16 = _mm_unpacklo_ps(T0, T1);
_{m128} T17 = _{mm_unpacklo_ps(T2, T3)};
__m128 T18 = _mm_unpackhi_ps(T0, T1);
_{m128} T19 = _{mm}_unpackhi_ps(T2, T3);
_{\text{m128 T20}} = _{\text{mm}} = _{\text{unpacklo}} (T16, T17);
__m128 T21 = _mm_unpackhi_ps(T16, T17);
__m128 T22 = _mm_unpacklo_ps(T18, T19);
__m128 T23 = _mm_unpackhi_ps(T18, T19);
T20 = _{mm\_add\_ps}(T20, T21);
T20 = _{mm\_add\_ps}(T20, T22);
T20 = _{mm\_add\_ps}(T20, T23);
T20 = _{mm\_shuffle\_ps}(T20, T20, 0xD8);
r0 = _{mm}add_{ps}(T20, r0);
T16 = _{mm\_unpacklo\_ps}(T4, T5);
T17 = _{mm\_unpacklo\_ps}(T6, T7);
T18 = _mm_unpackhi_ps(T4, T5);
T19 = _{mm\_unpackhi\_ps}(T6, T7);
T20 = _{mm\_unpacklo\_ps}(T16, T17);
T21 = _{mm}unpackhi_{ps}(T16, T17);
T22 = _{mm\_unpacklo\_ps}(T18, T19);
T23 = _{mm}unpackhi_ps(T18, T19);
T20 = _{mm\_add\_ps}(T20, T21);
T20 = _{mm\_add\_ps}(T20, T22);
T20 = \underline{mm\_add\_ps}(T20, T23);
T20 = _{mm\_shuffle\_ps}(T20, T20, 0xD8);
r1 = _{mm\_add\_ps}(T20, r1);
T16 = _{mm\_unpacklo\_ps}(T8, T9);
T17 = _{mm\_unpacklo\_ps}(T10, T11);
T18 = _mm_unpackhi_ps(T8, T9);
T19 = _mm_unpackhi_ps(T10, T11);
T20 = _{mm\_unpacklo\_ps}(T16, T17);
T21 = _{mm}unpackhi_{ps}(T16, T17);
T22 = _{mm\_unpacklo\_ps}(T18, T19);
T23 = _{mm\_unpackhi\_ps}(T18, T19);
T20 = _{mm\_add\_ps}(T20, T21);
T20 = _{mm\_add\_ps}(T20, T22);
T20 = _{mm\_add\_ps}(T20, T23);
T20 = _{mm\_shuffle\_ps}(T20, T20, 0xD8);
```

```
r2 = _{mm\_add\_ps}(T20, r2);
                                   T16 = _{mm\_unpacklo\_ps}(T12, T13);
                                   T17 = _{mm\_unpacklo\_ps}(T14, T15);
                                   T18 = _{mm\_unpackhi\_ps}(T12, T13);
                                   T19 = _{mm\_unpackhi\_ps}(T14, T15);
                                   T20 = _{mm\_unpacklo\_ps}(T16, T17);
                                   T21 = _{mm}unpackhi_ps(T16, T17);
                                   T22 = _{mm\_unpacklo\_ps}(T18, T19);
                                   T23 = _mm_unpackhi_ps(T18, T19);
                                   T20 = _{mm\_add\_ps}(T20, T21);
                                   T20 = _{mm\_add\_ps}(T20, T22);
                                   T20 = _{mm\_add\_ps}(T20, T23);
                                   T20 = _{mm\_shuffle\_ps}(T20, T20, 0xD8);
                                    r3 = _{mm}add_{ps}(T20, r3);
                          }
                           _mm_store_ps(&C[x+0][y], r0);
                           _mm_store_ps(&C[x+1][y], r1);
                          _mm_store_ps(&C[x+2][y], r2);
                           _mm_store_ps(&C[x+3][y], r3);
}
```

GFLOPS of vector matrix in O0:



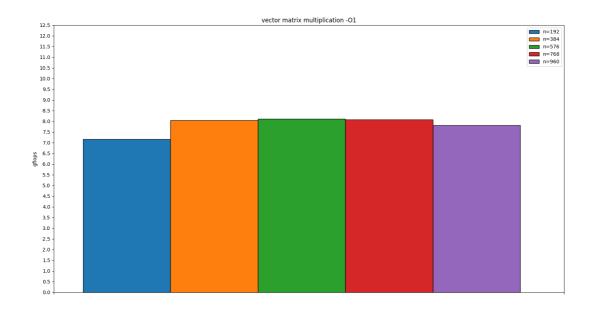
```
Performance counter stats for './vector-matrix00' (10 runs):

5,136,434 cache-misses # 4.227 % of all cache
refs (+- 24.77%)
```

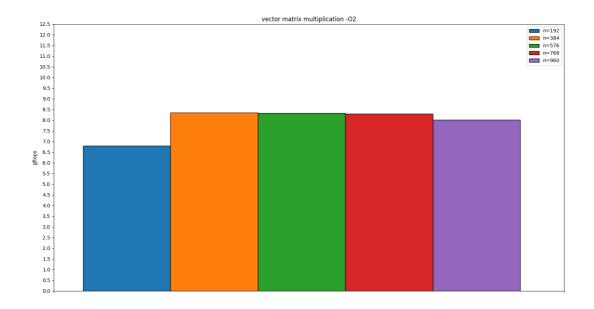
```
121,505,626 cache-references
( +- 3.66% )
    15,415,859,793 instructions # 2.07 insn per cycle
( +- 0.00% )
    7,445,628,148 cycles
( +- 1.21% )

2.140 +- 0.197 seconds time elapsed ( +- 9.19% )
```

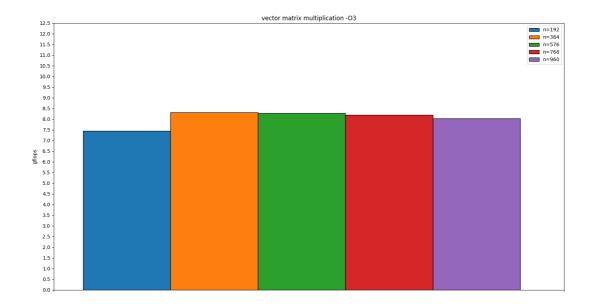
GFLOPS of vector matrix in O1:



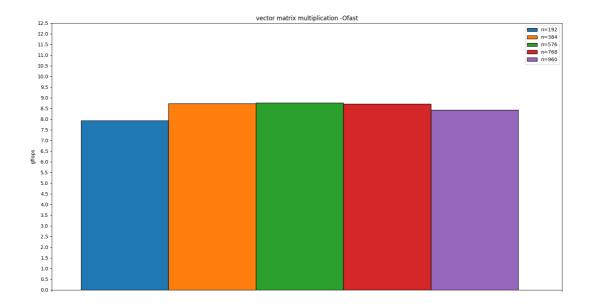
GFLOPS of vector matrix in O2:



GFLOPS of vector matrix in O3:



GFLOPS of vector matrix in Ofast:



there're lots of code for reordering the data. hence the performance is as well as expected.

2. ver. 2:

perform like this.

$$A = egin{bmatrix} a_{1,1} & a_{1,2} & \dots \ a_{2,1} & a_{2,2} & \dots \ \end{bmatrix}$$
 , $B = egin{bmatrix} b_{1,1} & b_{1,2} & \dots \ b_{2,1} & b_{2,2} & \dots \ \end{bmatrix}$

$$AB = egin{bmatrix} a_{1,1} \ [b_{1,1} & b_{1,2} & \ldots] + a_{1,2} \ [b_{2,1} & b_{2,2} & \ldots] + \ldots \ a_{2,1} \ [b_{1,1} & b_{1,2} & \ldots] + a_{2,2} \ [b_{2,1} & b_{2,2} & \ldots] + \ldots \ & \ldots \end{bmatrix}$$

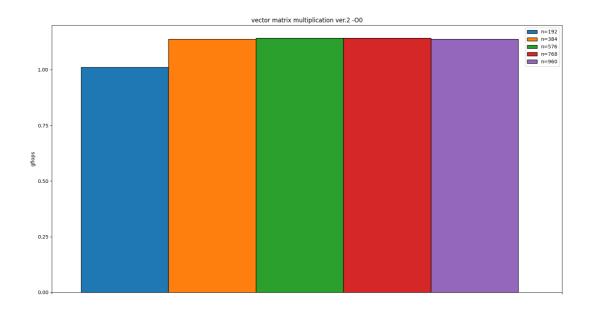
```
#include <xmmintrin.h>
#include <pmmintrin.h>
```

```
#include <tmmintrin.h>
#include <smmintrin.h>
#include <nmmintrin.h>
#include <immintrin.h>
constexpr int MAXSIZE = 960;
float A[MAXSIZE][MAXSIZE] __attribute__((aligned(16)));
float B[MAXSIZE] [MAXSIZE] __attribute__((aligned(16)));
float C[MAXSIZE] [MAXSIZE] __attribute__((aligned(16)));
int n, k, m;
inline void MMUL() noexcept
    for (int x = 0; x < n; x += 4)
                 for (int y = 0; y < m; y += 4)
                          __m128 xmm0, xmm1, xmm2, xmm3, xmm4, xmm5, xmm6,
xmm7, xmm8;
                          __m128 xmm9 = _mm_setzero_ps();
                          __m128 xmm10 = _mm_setzero_ps();
                          __m128 xmm11 = _mm_setzero_ps();
                          __m128 xmm12 = _mm_setzero_ps();
                          for (int z = 0; z < k; z += 4)
                                   xmm0 = _mm_load_ps(&B[z+0][y]);
                                   xmm1 = _mm_load_ps(&B[z+1][y]);
                                   xmm2 = _mm_load_ps(&B[z+2][y]);
                                   xmm3 = _mm_load_ps(&B[z+3][y]);
                                   // load A[x][z:z+3]
                                   xmm4 = \underline{mm}load_ps(&A[x+0][z]);
                                   xmm5 = _mm_shuffle_ps(xmm4, xmm4, 0x00);
                                   xmm6 = _mm_shuffle_ps(xmm4, xmm4, 0x55);
                                   xmm7 = _mm_shuffle_ps(xmm4, xmm4, 0xAA);
                                   xmm8 = _mm_shuffle_ps(xmm4, xmm4, 0xFF);
                                   xmm5 = \underline{mm}\underline{mul}\underline{ps}(xmm5, xmm0);
                                   xmm6 = \underline{mm}\underline{mul}\underline{ps}(xmm6, xmm1);
                                   xmm5 = _mm_add_ps(xmm5, xmm6);
                                   xmm7 = \underline{mm\_mul\_ps}(xmm7, xmm2);
                                   xmm8 = _mm_mul_ps(xmm8, xmm3);
                                   xmm7 = _mm_add_ps(xmm7, xmm8);
                                   xmm5 = _mm_add_ps(xmm5, xmm7);
                                   xmm9 = _mm_add_ps(xmm9, xmm5);
                                   // load A[x+1][z:z+3]
```

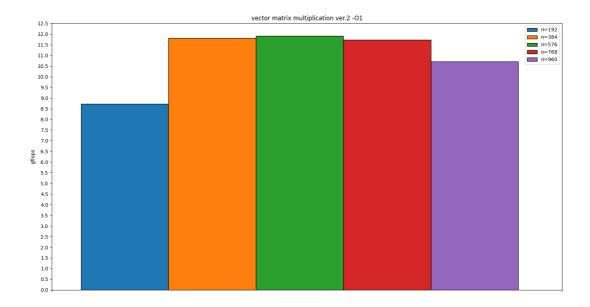
```
xmm4 = \underline{mm}load_ps(&A[x+1][z]);
xmm5 = _mm_shuffle_ps(xmm4, xmm4, 0x00);
xmm6 = _mm_shuffle_ps(xmm4, xmm4, 0x55);
xmm7 = _mm_shuffle_ps(xmm4, xmm4, 0xAA);
xmm8 = _mm_shuffle_ps(xmm4, xmm4, 0xFF);
xmm5 = \underline{mm}\underline{mul}\underline{ps}(xmm5, xmm0);
xmm6 = \underline{mm}\underline{mul}\underline{ps}(xmm6, xmm1);
xmm5 = _mm_add_ps(xmm5, xmm6);
xmm7 = \underline{mm}\underline{mul}\underline{ps}(xmm7, xmm2);
xmm8 = _mm_mul_ps(xmm8, xmm3);
xmm7 = _mm_add_ps(xmm7, xmm8);
xmm5 = _mm_add_ps(xmm5, xmm7);
xmm10 = _mm_add_ps(xmm10, xmm5);
// load A[x+2][z:z+3]
xmm4 = \underline{mm}load_ps(&A[x+2][z]);
xmm5 = _mm_shuffle_ps(xmm4, xmm4, 0x00);
xmm6 = _mm_shuffle_ps(xmm4, xmm4, 0x55);
xmm7 = _mm_shuffle_ps(xmm4, xmm4, 0xAA);
xmm8 = _mm_shuffle_ps(xmm4, xmm4, 0xFF);
xmm5 = \underline{mm\_mul\_ps}(xmm5, xmm0);
xmm6 = \underline{mm}\underline{mul}\underline{ps}(xmm6, xmm1);
xmm5 = _mm_add_ps(xmm5, xmm6);
xmm7 = \underline{mm}\underline{mul}\underline{ps}(xmm7, xmm2);
xmm8 = _mm_mul_ps(xmm8, xmm3);
xmm7 = _mm_add_ps(xmm7, xmm8);
xmm5 = _mm_add_ps(xmm5, xmm7);
xmm11 = _mm_add_ps(xmm11, xmm5);
// load A[x+3][z:z+3]
xmm4 = \underline{mm}load_ps(&A[x+3][z]);
xmm5 = _mm_shuffle_ps(xmm4, xmm4, 0x00);
xmm6 = _mm_shuffle_ps(xmm4, xmm4, 0x55);
xmm7 = _mm_shuffle_ps(xmm4, xmm4, 0xAA);
xmm8 = _mm_shuffle_ps(xmm4, xmm4, 0xFF);
xmm5 = \underline{mm}\underline{mul}\underline{ps}(xmm5, xmm0);
xmm6 = \underline{mm}\underline{mul}\underline{ps}(xmm6, xmm1);
xmm5 = _mm_add_ps(xmm5, xmm6);
xmm7 = \underline{mm\_mul\_ps}(xmm7, xmm2);
xmm8 = _mm_mul_ps(xmm8, xmm3);
xmm7 = _mm_add_ps(xmm7, xmm8);
xmm5 = _mm_add_ps(xmm5, xmm7);
xmm12 = _mm_add_ps(xmm12, xmm5);
```

```
}
_mm_store_ps(&C[x+0][y], xmm9);
_mm_store_ps(&C[x+1][y], xmm10);
_mm_store_ps(&C[x+2][y], xmm11);
_mm_store_ps(&C[x+3][y], xmm12);
}
```

GFLOPS of vector matrix ver.2 in O0:



GFLOPS of vector matrix ver.2 in O1:



```
Performance counter stats for './vector-matrix01V2' (10 runs):

4,735,092 cache-misses # 3.961 % of all cache

refs (+- 34.17%)

119,535,729 cache-references

(+- 4.75%)

9,645,654,696 instructions # 1.97 insn per cycle

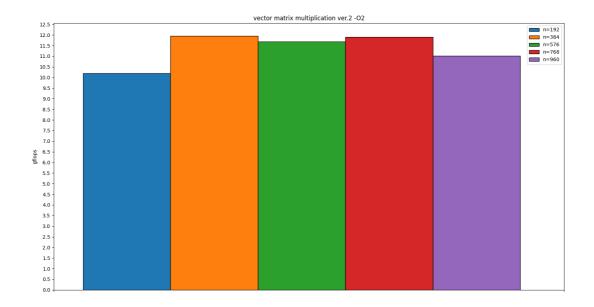
(+- 0.00%)

4,896,032,179 cycles

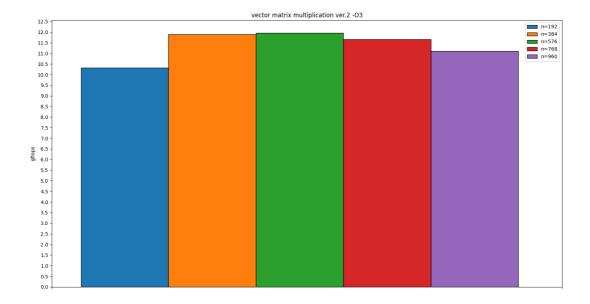
(+- 2.38%)

1.541 +- 0.147 seconds time elapsed (+- 9.53%)
```

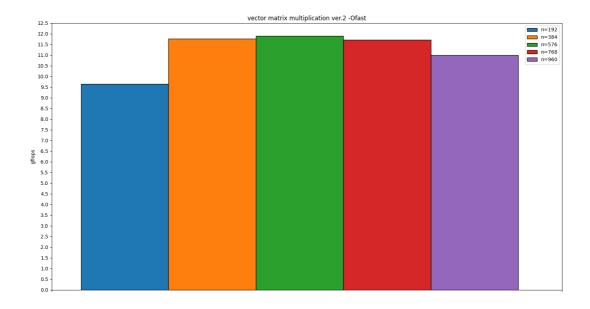
GFLOPS of vector matrix ver.2 in O2:



GFLOPS of vector matrix ver.2 in O3:



GFLOPS of vector matrix ver.2 in Ofast:



comparing with version 1, this version doesn't need the data reordering (less instructions), more the cache miss counts but less the cache reference counts, so the performance is much better.

source code

https://github.com/OEmiliatanO/CSE_computer_organzation_final_project

reference

2017q1 matrix

Matrix Multiplication using SIMD

How to vectorize with gcc?

How many GCC optimization levels are there?

Options That Control Optimization

Introduction to SSE Programming

Intel® Intrinsics Guide