1 Clebsch-Gordan matrix circuit

$$S_{1} \longrightarrow H \longrightarrow \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} & 0 \\ 0 & 0 & 0 & 1 \\ 0 & \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} & 0 \end{bmatrix} \begin{bmatrix} |00\rangle \\ |01\rangle \\ |10\rangle \\ |11\rangle \end{bmatrix} = \begin{bmatrix} |J=1, M=1\rangle \\ |J=1, M=0\rangle \\ |J=1, M=-1\rangle \\ |J=0, M=0\rangle \end{bmatrix} = (Intermsof spins) \begin{bmatrix} |00\rangle \\ \frac{1}{\sqrt{2}}(|01\rangle + |10\rangle) \\ |11\rangle \\ \frac{1}{\sqrt{2}}(|01\rangle - |10\rangle) \end{bmatrix}$$
(1)

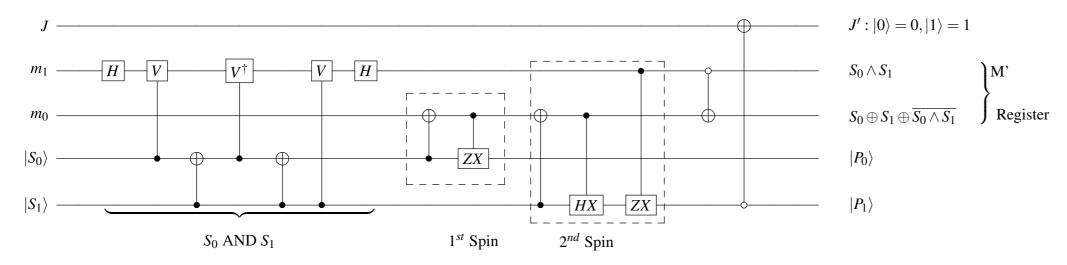
Circuit for Clebsch-Gordan transform.

$$S_{1} \longrightarrow H \longrightarrow \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} & 0 \\ 0 & \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} |00\rangle \\ |01\rangle \\ |10\rangle \\ |11\rangle \end{bmatrix} = \begin{bmatrix} |J=1, M=1\rangle \\ |J=1, M=0\rangle \\ |J=0, M=0\rangle \\ |J=1, M=-1\rangle \end{bmatrix} = (Interms of spins) \begin{bmatrix} |00\rangle \\ \frac{1}{\sqrt{2}}(|01\rangle + |10\rangle) \\ \frac{1}{\sqrt{2}}(|01\rangle - |10\rangle) \\ |11\rangle \end{bmatrix}$$

$$(2)$$

Adding a C-NOT corresponds to swapping the last two rows.

2 Spatially Multiplexed registered 2 qubit Schur transform- 12 Two-qubit gates

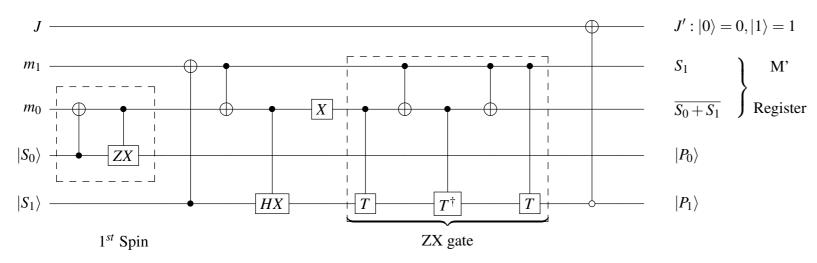


Adding another CNOT allows you to encode the M' register using the same values for M'=0.

| | Spin values | | Circuit output | | M value |
|---|-------------|-------|----------------|-------|---------|
| | S_1 | S_0 | m_1 | m_0 | M |
| Ì | 0 | 0 | 0 | 1 | M=+1 |
| | 0 | 1 | 0 | 0 | M=0 |
| ĺ | 1 | 0 | 0 | 0 | M=0 |
| | 1 | 1 | 1 | 0 | M=-1 |

Figure 1: Table giving M register decoding

3 Spatially Multiplexed Minimal gate explicit J & M recording- 11 Two-qubit gates



| Spi | n values | | Circuit output | | M value |
|-----|----------|-------|----------------|------------------------|---------|
| | S_1 | S_0 | S_1 | $\overline{S_0 + S_1}$ | M |
| | 0 | 0 | 0 | 1 | M=+1 |
| | 0 | 1 | 0 | 0 | M=0 |
| | 1 | 0 | 1 | 0 | M=0 |
| | 1 | 1 | 1 | 1 | M=-1 |

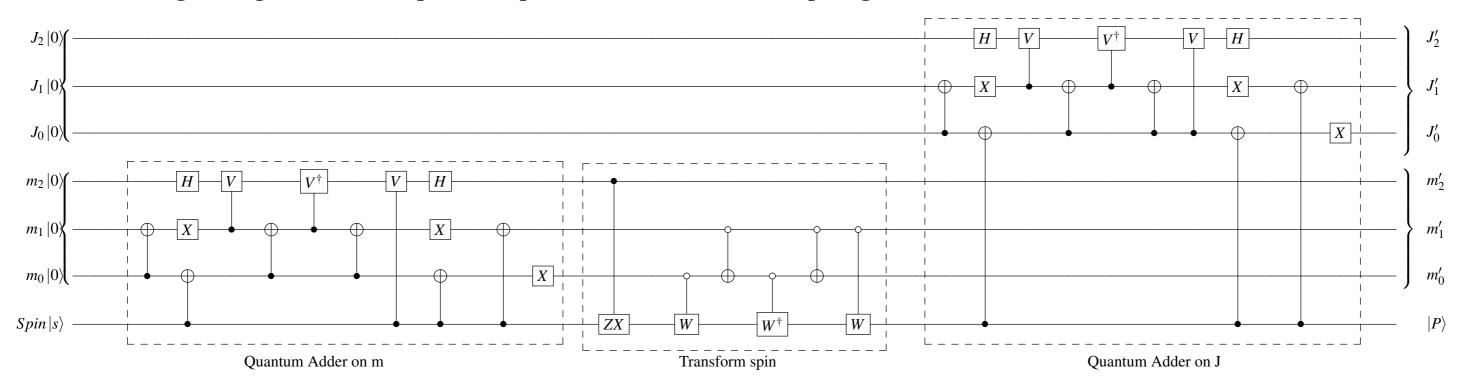
Figure 2: Table giving M register decoding

Circuit uses 11 two-qubit gates but only stores the final output values of J' & M'.

The HX gate triggers if M=0 meaning $S_0 \neq S_1$ which is implemented using an XOR between $S_0 \& S_1$. The other gate (T) is triggered when M=-1 meaning $S_0=S_1=1$ which is done using an AND (Toffoli) gate between,

 $S_0 = S_1$ AND $S_1 = 1$ which is decomposed into 5 two-qubit gates. T^2 is the ZX gate, meaning $T^2|S\rangle = XZ|S\rangle$.

Reduced general gate circuit for up to the 2 qubit Schur transform - 12 2-qubit gates



| J_2 | J_1 | J_0 | J | m_2 | m_1 | m_0 | M |
|-------|-------|-------|----------------|-------|-------|-------|----------------|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | $\frac{1}{2}$ | 0 | 0 | 1 | $\frac{1}{2}$ |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | $\frac{3}{2}$ | 0 | 1 | 1 | $\frac{3}{2}$ |
| 1 | 0 | 0 | -2 | 1 | 0 | 0 | -2 |
| 1 | 0 | 1 | $-\frac{3}{2}$ | 1 | 0 | 1 | $-\frac{3}{2}$ |
| 1 | 1 | 0 | -1 | 1 | 1 | 0 | -1 |
| 1 | 1 | 1 | $-\frac{1}{2}$ | 1 | 1 | 1 | $-\frac{1}{2}$ |

Figure 3: Tables giving binary Two's compleregisters

Circuit uses the encoding for $|S\rangle : |0\rangle \mapsto Spin =$ $+\frac{1}{2}$, $|1\rangle \mapsto Spin = -\frac{1}{2}$ and the same for $|P\rangle$.

Where V is the phase gate, $V = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix}$, $V^{\dagger}V = I$ and

 $V^2 = Z$. V is used here to expand the double controlled Toffoli gate into single control gates in the quantum adder subroutine.

The W gate, $W^2 = HX$ with $W^{\dagger} = I$, is used to expand the HX gate into single control gates in the spin transform region.

The circuit checks that if $(m_1 \text{ XNOR } m_0) \text{ AND } (m_0)$ ment encoding to spin values of the M and J XOR S_0) and will then change m_2 . Then m_1 is updated using $m_1 = m_0$ XOR S_0 . m_0 is always incremented by 1, if $|S\rangle = |0\rangle$ increment only m_0 by 1 corresponding to

adding $\frac{1}{2}$ to the M register. $|S\rangle = |1\rangle$ corresponds to subtracting $\frac{1}{2}$ from the M register by adding the string 111 bitwise to M.

For the most positive values of M the Identity is performed on the spin corresponding to the strings $M = 001(J = \frac{1}{2}, M' = \frac{1}{2})$ for the first spin and M = 010(J = 1, M' = 1) for the second coupled in spins.

The most negative values of M performs $XZ|S\rangle$ corresponding to the strings $M=111(J=\frac{1}{2},M'=-\frac{1}{2})$ for the first spin and M = 110(J = 1, M' = -1) for the second spin.

If M = 000(J = 0, M' = 0) do $XH |S\rangle$.

5 Streaming Scheme

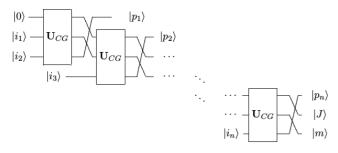


FIG. 2: Quantum circuit for the Schur transformation $\mathbf{U}_{\mathrm{Sch}}$, transforming between $|i_1i_2\cdots i_n\rangle$ and $|J,m,p\rangle$.

Figure 4: Streaming structure. Taken from Bacon, Chaung, Harrow (2004) Arxiv /0407082v4

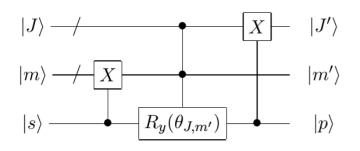


Figure 5: U_{CG} block, Qadder, controlled rotation, Qadder. Taken from Bacon, Chaung, Harrow (2004) arXiv /0407082v4

The Rotation matrix, $R_y(\theta_{J,m'})$ is given by,

$$R_{y}(\theta_{J,m'}) = \begin{bmatrix} \cos(\theta_{J,m'}) & -\sin(\theta_{J,m'}) \\ \sin(\theta_{J,m'}) & \cos(\theta_{J,m'}) \end{bmatrix}$$
(3)

=

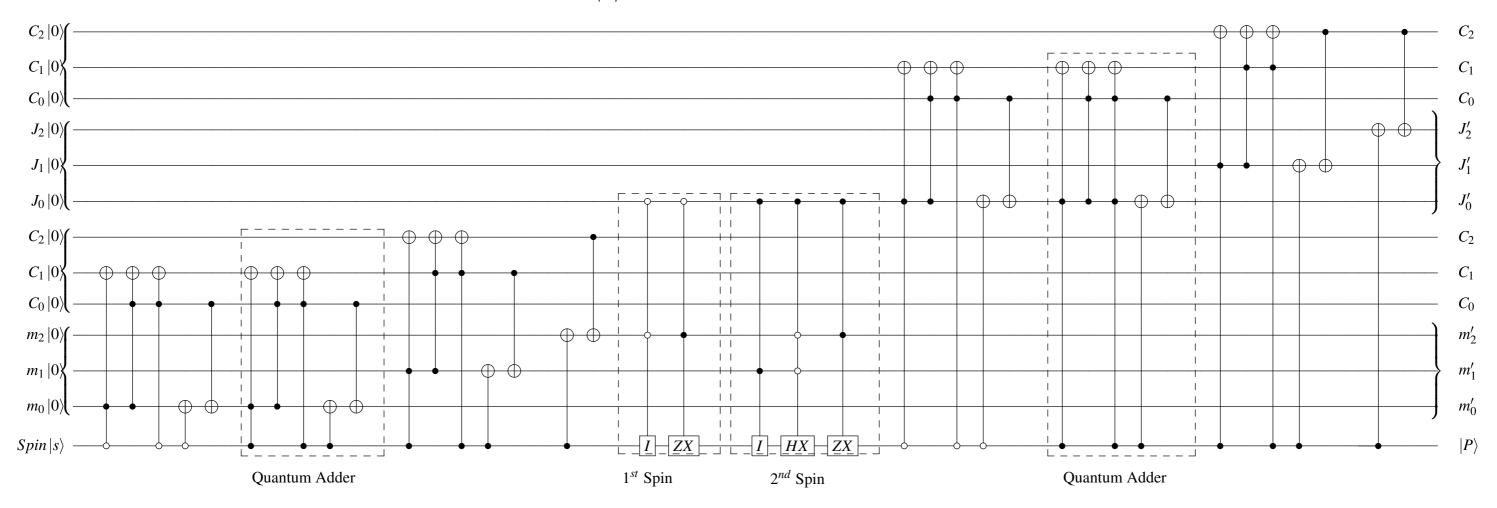
$$\frac{1}{\sqrt{2J+1}} \begin{bmatrix} \sqrt{J+\frac{1}{2}+m'} & -\sqrt{J+\frac{1}{2}-m'} \\ \sqrt{J+\frac{1}{2}-m'} & \sqrt{J+\frac{1}{2}+m'} \end{bmatrix}$$
(4)

Where primed variables means after the angular momentum addition so J is the total J that the spin is coupling to, the system will have J' total angular momentum after the coupling. m is the z component of the system before and m' is the total z component after the coupling.

The temporally multiplexed streaming scheme is shown in figures 4 & 5.

4

6 General circuit for the Quantum Schur transform ($|S\rangle$)



| J_2 | J_1 | J_0 | J | m_2 | m_1 | m_0 | M |
|-------|-------|-------|----------------|-------|-------|-------|----------------|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | $\frac{1}{2}$ | 0 | 0 | 1 | $\frac{1}{2}$ |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | $\frac{3}{2}$ | 0 | 1 | 1 | $\frac{3}{2}$ |
| 1 | 0 | 0 | -2 | 1 | 0 | 0 | -2 |
| 1 | 0 | 1 | $-\frac{3}{2}$ | 1 | 0 | 1 | $-\frac{3}{2}$ |
| 1 | 1 | 0 | -1 | 1 | 1 | 0 | -1 |
| 1 | 1 | 1 | $-\frac{1}{2}$ | 1 | 1 | 1 | $-\frac{1}{2}$ |

Figure 6: Tables giving binary Two's complement encoding to spin values of the M and J registers

Circuit uses the encoding for $|S\rangle:|0\rangle\mapsto Spin=+\frac{1}{2},|1\rangle\mapsto Spin=-\frac{1}{2}$ and the same for $|P\rangle$.

The circuit adds the value of the spin to be added, $|S\rangle$, to the M register to calculate the M' register value. This is done by implementing the quantum reversible equivalent to the digital full adder.

The case where $|S\rangle = |0\rangle$ means the spin is $+\frac{1}{2}$ so to add $\frac{1}{2}$ to M one is added to the m_0 bit. The very first Quantum Adder (QAdd) uses Toffoli gates controlled on $|0\rangle$ on $|s\rangle$ (denoted by the white control circle) with the current m_0 value and C_0 (an ancilla carry) so that in the case $m_0 = 1$ and we try and add 1 to it, m_0 goes to 0 and m_1 is increased using the carry as 001 + 1 = 010. The rest of the QAdd stages then just check the carry of the previous qubit to complete to $M + \frac{1}{2}$ addition as $|S\rangle = |0\rangle$

does not trigger any of the rest of the control gates.

The case where $|S\rangle = |1\rangle$ means the spin is $-\frac{1}{2}$ we do M $-\frac{1}{2}$ which is done by adding the binary string for $-\frac{1}{2}$ which is the all 1's string, 111. This time the very first Quantum Adder does not trigger and $|s\rangle$ is then added to all of the bits of M using C-NOT gates with carries to check for overflow.

The Unitary is then performed on $|S\rangle$ depending on the values of the newly calculated M' and J registers. The Identity is shown in the circuit for completeness on all the J and M' values. The J register is then updated to J' by adding the value of $|P\rangle$ to J using the QAdd sequence of gates.

To add the second qubit in the values of J' and M' are passed in as the initial register values. It is easy to extend this to many qubits being streamed in one at a time by carefully conditioning the controls on the unitaries, I think in the general case you need at most N controls for coupling up to N qubits in one at a time. The circuit written here has redundancy in the Identity and ZX gates appearing twice.

7 Clebsch-Gordan coefficients for 3 qubits

| $J=\frac{3}{2}$ | $S = \frac{3}{2}$ |
|--|--------------------|
| 000 | $M=\frac{2}{3}$ |
| $\sqrt{\frac{1}{3}}(001+010+100)$ | $M=\frac{1}{2}$ |
| $\sqrt{\frac{1}{3}}(110+011+101)$ | $M=-\frac{1}{2}$ |
| 111 | $M = -\frac{3}{2}$ |
| $J = \frac{1}{2}, P = 0$ | $S=\frac{1}{2}$ |
| $\sqrt{\frac{2}{3}}(001) - \sqrt{\frac{1}{6}}(010 + 100)$ | $M=\frac{1}{2}$ |
| $-\sqrt{\frac{2}{3}}(110) + \sqrt{\frac{1}{6}}(011 + 101)$ | $M=-rac{1}{2}$ |
| $J = \frac{1}{2}, P = 1$ | $S = \frac{1}{2}$ |
| $\frac{1}{\sqrt{2}}(010-100)$ | $M=\frac{1}{2}$ |
| $\frac{\sqrt{12}}{\sqrt{2}}(011-101)$ | $M=-\frac{1}{2}$ |

Figure 7: J & M values for 3 qubits using encoding 0=spin up, 1=spin down

matrix for transform is:

have to swap 011 & 100 to get block form.

$$J = \frac{3}{2} \qquad S = \frac{3}{2}$$

$$000 \qquad M = \frac{3}{2}$$

$$\sqrt{\frac{1}{3}}(001 + 010 + 100) \qquad M = \frac{1}{2}$$

$$\sqrt{\frac{1}{3}}(110 + 011 + 101) \qquad M = -\frac{1}{2}$$

$$111 \qquad M = -\frac{3}{2}$$

$$J = \frac{1}{2}, P = 0 \qquad S = \frac{1}{2}$$

$$\frac{1}{\sqrt{3}}(e^{2\pi i/3}001 + e^{4\pi i/3}010 + 100) \qquad M = \frac{1}{2}$$

$$\frac{1}{\sqrt{3}}(e^{2\pi i/3}011 + e^{4\pi i/3}101 + 110) \qquad M = -\frac{1}{2}$$

$$\frac{1}{\sqrt{3}}(e^{2\pi i/3}001 + e^{4\pi i/3}010 + 100) \qquad M = \frac{1}{2}$$

$$\frac{1}{\sqrt{3}}(e^{2\pi i/3}001 + e^{4\pi i/3}010 + 100) \qquad M = \frac{1}{2}$$

$$\frac{1}{\sqrt{3}}(e^{4\pi i/3}011 + e^{2\pi i/3}101 + 110) \qquad M = -\frac{1}{2}$$

Figure 8: Schur transform with Phase encoding?

$$\frac{1}{\sqrt{3}} \begin{bmatrix}
\sqrt{3} & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 1 & 0 & 1 & 1 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & \sqrt{3} \\
0 & e^{2\pi i/3} & e^{4\pi i/3} & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & e^{2\pi i/3} & e^{2\pi i/3} & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & e^{4\pi i/3} & e^{2\pi i/3} & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & e^{4\pi i/3} & 0 & e^{2\pi i/3} & 1 & 0 \\
0 & 0 & 0 & e^{4\pi i/3} & 0 & e^{2\pi i/3} & 1 & 0
\end{bmatrix}$$

$$\begin{vmatrix}
J = 3/2, M = 3/2 \\
J = 3/2, M = -1/2 \\
J = 3/2, M = -1/2 \\
J = 3/2, M = -1/2 \\
J = 1/2, M = -1/2, P = 0 \\
J = 1/2, M = -1/2, P = 0 \\
J = 1/2, M = -1/2, P = 1 \\
J = 1/2, M = -1/2, P = 1 \\
J = 1/2, M = -1/2, P = 1 \\
J = 1/2, M = -1/2, P = 1$$

8 Circuit for 3 qubit transform

see github for Fortran code, using Givens rotations gives 19 cc-unitary gates which is about 100 c-unitaries.

9 4 Qubit CG coefficients