



ECOAR – Computer Architecture

Module 2 - Data

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#### Outline

- · Data types
- · Binary data representation
- Data in computer's memory
- Vector data types





## Data types

- Boolean (True/False)
- · Text characters
- Numbers
  - · integer
    - unsigned
    - signed
  - non-integer
    - Fixed point
    - Floating point
- Sounds and other unidimensional signals
- Raster images





### Binary data representation

- Computer operates solely on groups of bits (binary digits), treating them as binary numbers
  - Contemporary computers usually operate on binary words of lengths equal 8×2<sup>n</sup> bits (8, 16, 32, 64)
- Non-numeric data must be expressed using numbers









### Alphanumeric data

- Every character is represented by a number denoting its position in code table
- Most frequently used codes:
  - ASCII 128 positions including upper- and lowercase Latin letters
  - Extended ASCII 256 positions
    - First 128 positions identical to ASCII, the remaining 128 positions used for national characters and special symbols
    - problem: different codes are needed for various parts of the world
  - EBCDIC family of codes 256 symbols, used mainly by IBM
  - UNICODE
    - Initially 2<sup>16</sup>, currently up to 2<sup>21</sup> positions
    - Currently contains approx. 150000 characters
    - Covers all alphabetic characters used in the world







#### **ASCII**

- American Standard Code for Information Interchange
- Developed for teletype devices, later used in computers
- 128 positions, 95 visible and 33 invisible
- Invisible use codes 0..31 (0x00..0x1f) and 127 (0x7f):
  - White spaces and formatting codes (tabs, new line, new page)
  - Transmission and device control
  - CAN cancel/delete code 127 (0x7f)
- Visible use codes 32..126 (0x20..0x7e):
  - space (code 32)
  - Digits
  - Latin letters
  - Punctuation marks
  - Common mathematical symbols





### What a programmer should know about ASCII

- Control codes positions 0..31, including
  - CR carriage return 13
  - LF line feed 10
  - Also important: HT, FF, BSP, BEL
- Space 32 (0x20)
- digits 0..9 codes 48..57 (0x30..0x39)
  - to get digit's value, subtract `0` (0x30) from character representing it
- Letters in alphabetic order
  - Upper case 65..90 (0x41..0x5a)
  - Lower case 97..122 (0x61..0x7a)
  - Distance between lower and upper case of a letter is 32 (0x20)
- Other visible chars fit in the range 33..126
- 127 the very last special code (cancel/delete)







#### Extended codes based on ASCII

- 256 code points 8-bit representation
  - First 128 positions identical to ASCII
  - Next 128 positions used to represent characters from a selected set of alphabets, ex.
    - Slavic
    - Nordic
    - Cyrillic
    - Greek
- Many code tables of different origin used in some countries
  - ISO8859 family (Polish characters present in ISO8859-2)
  - Microsoft "code pages" CP (Poland CP852, CP1250)
  - Local usage (Polish Mazovia, Polgaz)

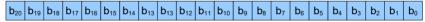






#### **UNICODE and UTF-8**

- Unicode, introduced in late 1980s, is a 21-bit character code designed to include all alphabetic characters used in the world
- Symbolic notation of Unicode character: U+<hex\_code>
- UTF-8 is the most common Unicode representation
- A character is represented by 1..4 octets character binary code



#### is represented as:

- ASCII chars  $U+00..U+7F \rightarrow 1$  octet:  $0b_6b_5b_4b_3b_2b_1b_0$
- U+80..U+7FF  $\rightarrow$  2 octets: 110b<sub>10</sub>b<sub>9</sub>b<sub>8</sub>b<sub>7</sub>b<sub>6</sub> 10b<sub>5</sub>b<sub>4</sub>b<sub>3</sub>b<sub>2</sub>b<sub>1</sub>b<sub>0</sub>
- U+800..U+FFFF  $\rightarrow$  3 octets: 1110b<sub>15</sub>b<sub>14</sub>b<sub>13</sub>b<sub>12</sub> 10b<sub>11</sub>b<sub>10</sub>b<sub>9</sub>b<sub>8</sub>b<sub>7</sub>b<sub>6</sub> 10b<sub>5</sub>b<sub>4</sub>b<sub>3</sub>b<sub>2</sub>b<sub>1</sub>b<sub>0</sub>
- U+10000...U+10FFFF → 4 octets:
   11110b<sub>20</sub>b<sub>19</sub>b<sub>18</sub> 10b<sub>17</sub>b<sub>16</sub>b<sub>15</sub>b<sub>14</sub>b<sub>13</sub>b<sub>12</sub> 10b<sub>11</sub>b<sub>10</sub>b<sub>9</sub>b<sub>8</sub>b<sub>7</sub>b<sub>6</sub> 10b<sub>5</sub>b<sub>4</sub>b<sub>3</sub>b<sub>2</sub>b<sub>1</sub>b<sub>0</sub>









### Text strings

- Text string is a sequence of characters
- The software must know where the string ends
- Two common conventions:
  - Character string end identified by a special character
    - In C and similar languages, the string ends with NUL (zero) code
    - in other environments '\$' or LF was used
  - An explicit string length specified as binary number, stored before the text
    - common in Basic language





## Sounds and images

#### · Sound:

- Voltage analog of acoustic pressure sampled with frequency dependent on the quality requirements (from 8 to 48 kHz)
- Sample values stored as integer numbers
- Raster images
  - Represented as rectangular arrays of square picture elements (pixels)
  - Single color is assigned to every pixel
  - Color is represented using three primaries values of basic lights (red, green, blue)
  - Values of primaries stored as unsigned integer numbers



#### Units

- bit (Binary digIT) abbreviation "b"- smallest information unit, may represent two-state information YES-NO, 1-0, TRUE-FALSE
- octet (abbrev. 'o') 8 bits, tetrade/nibble 4 bits
- byte abbreviation "B" smallest unit of information addressed by computer – in contemporary g.p computers – 8 bits
- word unit of information operated on by the computer
  - 1, 2, 4, 8, 16 bytes
- Processor word –unit of information on which a given processor may operate
  - Maximum word size equals register width in contemporary g.p. computers 64 or 32 bits
- Memory word unit of information that may be transferred in a single transfer cycle between processor and memory
  - Typ. 128 or 64 bits
- Longer memory word = faster data transfer







#### Data formats

- Computers operate on words groups of (usually) 2<sup>n</sup> bits
  - Common word lengths in g.p. computers 8, 16, 32, 64, 128 bits
  - Specialized computers may use different word lengths 24 bits
- Some computers are also capable of operating on single bits and bit fields
- Data stored as words in memory





## Logical (Boolean) data

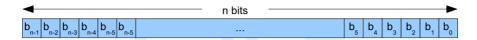
- Single bit is sufficient to store Boolean data
  - Single bit Boolean format is used in computers providing hardware handling of single bit data
- Boolean data is commonly represented using word of a length natural for a given computer
  - Bool/bool type in C/C++ is implemented as byte
  - Some other languages use the same size as used for integers typ.
     32 bits
- Representation depends on the programming environment and programming language
  - FALSE is usually represented by a word of value 0
  - TRUE may be represented by
    - Value 1 (C language result)
    - Any non-zero value (most languages source arguments)
    - Word made of all bits set to 1 (Visual Basic result)







## **Unsigned integers**



Natural binary code

$$v_N = \sum_{i=0}^{n-1} b_i \cdot 2^i$$

- BCD (Binary-Coded Decimal)
  - Used for fixed point decimal numbers
  - Decimal digits encoded in binary 4 bits (nibble/tetrade) per digit
  - Allowed nibble values 0..9
  - · Formats:
    - packed 2 digits per byte; range 0..99
    - Unpacked ("ASCII") one digit per byte, range 0..9

## Signed integers

Two's complement

Ones' complement

Sign-magnitude

- Biased
  - typically
     BIAS = 2<sup>n-1</sup>-1

$$v_{2C} = -b_{n-1} \cdot 2^{n-1} + \sum_{i=0}^{n-2} b_i \cdot 2^i$$

$$v_{1C} = -b_{n-1} \cdot (2^{n-1} - 1) + \sum_{i=0}^{n-2} b_i \cdot 2^i$$

$$v_{S-M} = (-1)^{b_{n-1}} \cdot \sum_{i=0}^{n-2} b_i \cdot 2^i$$

$$v_{B} = -BIAS + \sum_{i=0}^{n-1} b_{i} \cdot 2^{i}$$





## Important features of numeric codes

- · Representation of zero
  - Number of distinct representations
    - Two representations of 0 in ones complement and sign-magnitude
  - Fase of detection
- Symmetry of range (for signed numbers)
- Representation of sign
  - Ease of sign determination
- Change of sign
  - Ones complement bit negation
  - Two's complement bit negation and increment
  - Sign-magnitude negation of sign bit
- Implementation of arithmetic operations
  - Addition and subtraction of two's complement numbers may be performed using the hardware designed for unsigned numbers
    - Only the overflow detection is different







#### Numeric codes and their features

| Code     | Binary patterns |      |    |    | 8-bit byte value |    |    | 16-k | it wc | ord va | alue | 32-bit word value |          |      |         |
|----------|-----------------|------|----|----|------------------|----|----|------|-------|--------|------|-------------------|----------|------|---------|
|          | min             | zero | ma | ЭX | min              | ma | Χ  | m    | in    | ma     | ax.  |                   | min      | ı    | max     |
| Unsigned | 000             | 000  | 11 | 1  | 0                | 2  | 55 |      | 0     | 65     | 535  |                   | 0        | 4294 | 1967295 |
| 2C       | 100             | 000  | 01 | .1 | -128             | 1  | 27 | -32  | 768   | 32     | 767  | -21               | 47483648 | 2147 | 7483647 |
| S-M      | 111             | 000  | 01 | .1 | -127             | 1  | 27 | -32  | 767   | 32     | 767  | -21               | 47483647 | 2147 | 7483647 |
|          |                 | 100  |    |    |                  |    |    |      |       |        |      |                   |          |      |         |
| 1C       | 100             | 000  | 01 | .1 | -127             | 1  | 27 | -32  | 767   | 32     | 767  | -21               | 47483647 | 2147 | 7483647 |
|          |                 | 111  |    |    |                  |    |    |      |       |        |      |                   |          |      |         |
| Biased   | 000             | 011  | 11 | .1 | -127             | 1  | 28 | -32  | 767   | 32     | 768  | -21               | 47483647 | 2147 | 7483648 |







### Change of sign, 2's complement dec→bin conversion

- · Change of sign:
  - Ones' complement: ~x
  - Two's complement: ~x+1
  - Biased: BIAS x
  - Sign-magnitude: negate the sign bit only
  - In codes with asymmetric range (2's C, biased), sign change may cause overflow
- Decimal to 2's complement conversion
  - For non-negative numbers, 2's complement representation is the same as unsigned
    - The representable range limit is smaller 2<sup>n-1</sup> 1
  - For negative numbers:
    - Convert |x| to unsigned binary
    - Change the sign: negate all bits, then add 1





### Fixed point notation

- Obtained by shifting bit weights in integer notation
  - Equivalent to multiplication of integer value by 2<sup>-f</sup> (f no. of bits in fractional part)
- Used for 2's complement and unsigned numbers
- Commonly used formats:
  - One or two bits of integer part, rest of word used as fractional part
  - Half of the word used as integer part, the other half as fraction
- Arithmetic operations similar to integer operations
  - Scaling (shifting) required during multiplication and division
  - Do not require special instructions or hardware structures



### Fixed point notation - examples



Fixed point weights, 16 bits in integer part, 16 bits in fractional part

|             |     |     |                       |      |      |      |     |             |               |              | - 40         |
|-------------|-----|-----|-----------------------|------|------|------|-----|-------------|---------------|--------------|--------------|
| <b>7</b> 15 | 214 | 213 | <b>2</b> <sup>2</sup> | 1 21 | 1 20 | 1つ-1 | 2-2 | <b>ე</b> -3 | <b>1</b> 2-14 | <b>ე</b> -15 | <b>7</b> -16 |
| _           | _   | _   | <br>  _               | _    | _    | _    | _   |             | <br>  _       |              | _            |
|             |     |     |                       |      |      |      |     |             |               |              |              |

Fixed point weights, 1 bit of integer part, 31 bits in fractional parti

| <b>2</b> ° | 2-1 | 2-2 | 2-3 | 2-4 | 2-5 |  |  |  |  |  | <b>2</b> <sup>-26</sup> | 2-27 | 2-28 | 2-29 | 2-30 | 2-31 |
|------------|-----|-----|-----|-----|-----|--|--|--|--|--|-------------------------|------|------|------|------|------|





## Binary vs. decimal fractions - accuracy

- In a fractional part of a binary number, bit weights are 2-
  - ½, ¼, 1/8, 1/16, 1/32, etc...
- As 10 is not equal to any power of 2, decimal fractions cannot be in general precisely represented as binary fractions with finite number of digits

  - Only fractions being sums of binary fractions may be represented precisely:

• • •







### Floating point notation - basics

- Examples of decimal FP notation
  - -1.234 10<sup>5</sup> -0.1234 10<sup>6</sup> -12.34 10<sup>4</sup>
- · Elements:
  - Sign of a number
  - Significand (unsigned fixed point number)
  - Exponent (signed integer)
- System base (10) is fixed there is no need to store it
  - It is only a graphic element of the notation
- Normalized form form in which integer part of significand is expressed using single non-zero digit
  - In the above example: -1.234 10<sup>5</sup>
  - Other forms are denormalized (non normalized)
  - It is not possible to express zero in a normalized form







### Binary floating point notation – IEEE754, IEC60559

- If possible, numbers are stored in normalized form
  - Exceptions: zero and numbers of a very small magnitude
  - The only binary digit different from 0 is 1
    - Every normalized number has an integer part of significand equal to 1
    - A non-normalized number has an integer part of significand equal to 0
    - There is no need to store integer part of significand if we know whether the number is normalized or not
- Elements of notation
  - Sign single bit: 0 non-negative, 1 non-positive
  - Exponent field
    - In normalized form exponent in biased code
    - Special values: 00...00 denormalized form, 11...11 Not-a-Number
  - Trailing significand (mantissa) field contains only fractional part of significand
    - Normalized form integer part is 1
    - denormalized form integer part is 0







### IEEE 754/IEC60559 - exponent field encoding

- exponent stored in e-bit wide field as biased value
- bias =  $2^{e-1} 1$  (01..11 binary)

| exp field | expo <mark>ne</mark> nt v | alue   |         | meaning                              |
|-----------|---------------------------|--------|---------|--------------------------------------|
| 0000      | -bias+1                   | 1      | zero o  | <mark>r d</mark> enormal             |
| 0001      | exp-bias = -l             | oias+1 | norma   | lized value                          |
| 0xxx      | exp-bia                   | S      | norma   | lized value                          |
| 0111      | exp <mark>-bi</mark> as   | = 0    | normal  | l <mark>iz</mark> ed value           |
| 1xxx      | exp-bia                   | s      | norma   | l <mark>iz</mark> ed value           |
| 1110      | exp-bia                   | s      | bigges  | t <mark>r</mark> epresentable number |
| 1111      | ( <mark>no</mark> ne)     |        | Not-a-l | <mark>Nu</mark> mber                 |







#### IEEE754/IEC60559 - values

| Sign | Biased exponent | Trailing significand (mantissa) field |
|------|-----------------|---------------------------------------|
| S    | Е               | Т                                     |

| Sign bit | Exponent field            | Exp.value | <mark>M</mark> antissa field Significand valu | е |
|----------|---------------------------|-----------|---|---|
| Х        | 0000                      | -bias+1   | m 0.m   |   |
| Х        | from 0 <mark>0</mark> .01 | exp-bias  | m 1.m   |   |
|          | to 1110                   |           |   |   |
| 0        | 11 <mark>11</mark>        | -         | 0000 + infinity                               |   |
| 1        | 1111                      | -         | 0000 - infinity                               |   |
| Х        | 1111                      | -         | 0xxx Signaling NaN                            |   |
| Х        | 11 <mark>11</mark>        |           | 1xxx Quiet NaN                                |   |







#### IEEE754/IEC60559 - common formats

| Word     | Symbolic | Field v  | vidths   | Application                             |  |  |  |  |  |
|----------|----------|----------|----------|---|--|--|--|--|--|
| size [b] | notation | mantissa | exponent | Application                             |  |  |  |  |  |
| 16       | s10e5    | 10       | 5        | Computer graphics data storage          |  |  |  |  |  |
| 24       | s16e7    | 16       | 7        | Computer graphics                       |  |  |  |  |  |
| 16       | s7e8     | 7        | 8        | artificial intelligence (BF16)          |  |  |  |  |  |
| 32       | s23e8    | 23       | 8        | binary32 (IEEE single), general purpose |  |  |  |  |  |
| 64       | s52e11   | 52       | 11       | binary64 (IEEE double), general purpose |  |  |  |  |  |
| 80       | s64e15   | 64       | 15       | Extended double – x87                   |  |  |  |  |  |
| 96       | s80e15   | 80       | 15       | DSP                                     |  |  |  |  |  |
| 128      | s112e15  | 112      | 15       | binary128, high-precision arithmetics   |  |  |  |  |  |



## Floating point arithmetic

- Floating point representation is an approximate one
  - Results of arithmetic operations are also approximate
- Result may depend on order of operations
  - a + b + c may not be the same as c + b + a
  - if |a| << |b|, then a + b may be equal to b
  - Computing the sum or average of many values is not a trivial task
  - Addition and subtraction of multiple arguments should be performed in order of growing magnitude
- Equality test usually gives false result
  - Use abs(a-b) < ε instead</li>
- Precision of IEEE single (24 bits) is smaller than that of 32-bit integer or fixed point
  - Precise representation of integers in IEEE single is limited to absolute value range of 2<sup>24</sup> only





## Memory organization

- In general purpose computers 8-bit byte is the smallest addressable unit of memory
- Multibyte data occupy the appropriate number of consecutive byte-sized memory locations

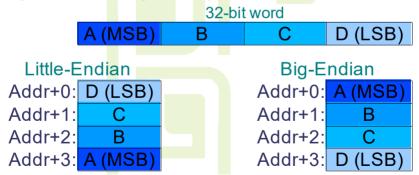






### Multibyte data addressing - byte order

- Little Endian least significant byte of a data word at the lowest address, more significant bytes follow
- Big Endian most significant byte at the lowest address









#### Little-Endian

- Byte numbering corresponds to bit weights
- Natural for computers
- Not so natural for humans
- Type casting preserves pointer value

|         | 32-bit word, value = 5 |   |  |         |        |  |  |  |  |  |
|---------|------------------------|---|--|---------|--------|--|--|--|--|--|
|         | 0                      | 0 |  | 0       | 5      |  |  |  |  |  |
|         |                        |   |  |         |        |  |  |  |  |  |
| 32-bit  | access                 |   |  | 16-bit  | access |  |  |  |  |  |
| Addr+0: | 5                      |   |  | Addr+0: | 5      |  |  |  |  |  |
| Addr+1: | 0                      |   |  | Addr+1: | 0      |  |  |  |  |  |
| Addr+2: | 0                      |   |  | 8-bit a | iccess |  |  |  |  |  |
| Addr+3: | 0                      |   |  | Addr+0: | 5      |  |  |  |  |  |







### Big-Endian

- Natural for humans
- Type cast changes the pointer value
- Fast string compare is possible

#### "abc" string in memory

Addr+0: 'a'
Addr+1: 'b'
Addr+2: 'c'
Addr+3: 0

#### 32-bit word read from Addr

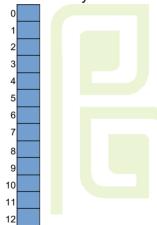
**'a'** 'b' 'c' 0



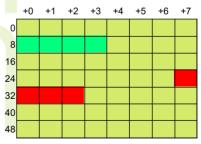


## Logical vs. physical memory organization

Logical – vector of bytes



- Physical "2-dimensional" bytes grouped in words, usually 32-, 64- or 128-bit
- Bytes within a word may be accessed simultaneously





### Data alignment

- Physically the memory is organized as a vector of words, each word being a vector of bytes
  - Any number of bytes within a single word may be accessed at the same time
- Memory word is usually at least as long as the processor's word
  - Typically 64 or 128 bits in 32-/64-bit processors (2 or 4 times longer than processor's word)
  - In older designs (before 1990) memory was sometimes narrower than processor
- The access to multibyte data contained in a single memory word is faster than in case when the memory words
  - In the latter case two physical accesses are needed
- New architectures enforce placement of every data item ensuring the fastest possible access



## Size alignment

- For the fastest access regardless of the physical memory width, each data item should be placed in memory at the address divisible by its length (rounded up to 2<sup>n</sup>)
- Size alignment is enforced by hardware in newer architectures
  - Unaligned access attempt results in execution error
- Even if not enforced in hardware, size alignment boosts the efficiency
  - Compilers enforce size alignment even if underlying hardware doesn't
- Each data type is characterized by two values:
  - Alignment the data starting address must be a multiple of alignment
    - In C language returned by \_Alignof() operator
  - Size must be a multiple of alignment
    - In C language returned by sizeof() operator
    - in modern environments, size is always 2<sup>n</sup> and it is equal to alignment







## Vectors and arrays

- The elements of a vector occupy the consecutive memory locations, starting with the first element (placed at the lowest address – the address of a vector)
  - In C-like languages, the first element has an index of 0
    - v[N] vector elements are v[0], v[1], ..., v[N-1]
- Vector alignment is the same as its element alignment
- An n-dimensional array is a vector of (n-1)-dimensional arrays







# Vectors and arrays

| char s[12]; |       | ]; cha | ar t[3][ | 4]; | int32_t v[12]; | int32_t a[3][4]; |         |  |
|-------------|-------|--------|----------|-----|----------------|------------------|---------|--|
| 0           | s[0]  | 0      | t[0][0]  | 0   | v[0]           | 0                | a[0][0] |  |
| 1           | s[1]  | 1      | t[0][1]  | 4   | v[1]           | 4                | a[0][1] |  |
| 2           | s[2]  | 2      | t[0][2]  | 8   | v[2]           | 8                | a[0][2] |  |
| 3           | s[3]  | 3      | t[0][3]  | 12  | v[3]           | 12               | a[0][3] |  |
| 4           | s[4]  | 4      | t[1][0]  | 16  | v[4]           | 16               | a[1][0] |  |
| 5           | s[5]  | 5      | t[1][1]  | 20  | v[5]           | 20               | a[1][1] |  |
| 6           | s[6]  | 6      | t[1][2]  | 24  | v[6]           | 24               | a[1][2] |  |
| 7           | s[7]  | 7      | t[1][3]  | 28  | v[7]           | 28               | a[1][3] |  |
| 8           | s[8]  | 8      | t[2][0]  | 32  | v[8]           | 32               | a[2][0] |  |
| 9           | s[9]  | 9      | t[2][1]  | 36  | v[9]           | 36               | a[2][1] |  |
| 10          | s[10] | 10     | t[2][2]  | 40  | v[10]          | 40               | a[2][2] |  |
| 11          | s[11] | 11     | t[2][3]  | 44  | v[11]          | 44               | a[2][3] |  |
| 12          |       | 12     |          | 48  |                | 48               |         |  |





#### **Structures**

- Compiler is required to preserve the order of fields
  - It cannot optimize the layout
- Each field must be aligned according to its type requirements
  - Padding (unused space) may be added between fields to achieve the proper alignment
- Structure must be aligned according to the alignment requirement of the field with the biggest alignment
- The structure is padded to the multiple of its alignment
- C language sizeof() operator returns the address distance between two structures of a given type, necessary for allocating a vector of structures.







### Order of structure fields vs. memory footprint

```
+3
struct st1 {
                                                 с1
     char c1:
                                               8
                                                                d1
     double d1:
                                              16
                                                    s1
     short int s1;
                                              24
                                                                d2
    double d2:
                                              32
                                                  c2
                                                                         i1
    char c2:
                                              40
                                                 с3
                                                            s2
    int i1:
                                                                   c4
     char c3;
                                              48
     short int s2;
                                              56
     char c4:
                                                     sizeof(struct st1) == 56
    int i2;
} e1;
struct st2 {
                                                                d1
     double d1, d2;
                                               8
                                                                d2
     int i1, i2;
     short int s1, s2;
                                              16
                                                        i1
                                                                         i2
     char c1, c2, c3, c4;
                                              24
                                                    s1
                                                            s2
                                                                   c1
                                                                       c2
                                                                           с3
} e2;
                                              32
                                                     sizeof(struct st2) == 32
```





#### Vector data formats

- · Contemporary processors may operate on long data words
- Sounds and pictures are expressed using short numbers
- Signal processing algorithms execute the same set of operations on all data samples
- It is possible to treat the long word as a vector of several shorter values
- The introduction of vector data formats enables better utilization of processor's resources







#### Vector data formats - x86 SSE

- SSE is a vector unit implemented in x86 processors from 1999
- Data word 128 bits
- Data formats supported (SSE2 and newer versions, 2000→):
  - 16 8-bit integers
  - 8 16-bit integers
  - 4 32-bit integers
  - 2 64-bit integers
  - 4 32-bit IEEE single
  - 2 64-bit IEEE double
  - 8 16-bit s10e5 floats (SSE5)







#### x86 AVX vector unit

- new vector unit of x86 processors, introduced in 2011 in "2<sup>nd</sup> generation Intel Core i" processors
- 16 256-bit registers
- FP data formats similar to those of SSE but with 256-bit vectors 4×64-bit or 8×32-bit
- AVX2 version also supports 256-bit integer vectors 32×8b, 16×16b, 8×32b, 4×64b
- New extension called AVX512 supports 512-bit vectors







## Requirements for the test

- Integer conversion between decimal and 2's complement, signmagnitude
- Structure and vector layout (drawing, sizeof operator value)
- Big-endian/little-endian data placement in memory







ECOAR – Computer Architecture

**End of module** 

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