

# **Fragmentation**

- Why needed?
  - The IP layer injects a packet into the datalink layer.
    - Not responsible for the reliable transport of these packets.
  - Each layer imposes some maximum size of packets, due to various reasons.
    - Called Maximum Transfer Unit (MTU).
  - Suppose a large packet travels through a network whose MTU is too small.
    - Fragmentation (and also reassembly) is required.
    - Each fragment is transmitted as a separate IP packet.
    - Fragmentation is typically done by routers.
- Fragments reassembled later: transparent or non-transparent.







Interconnection of Networks

N1

R

N3

R

N4

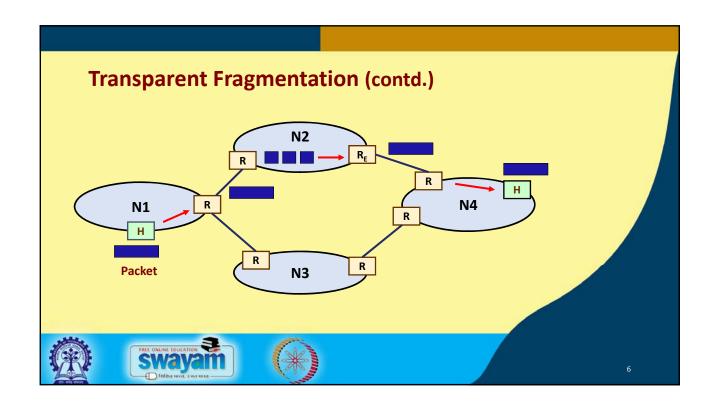
# **Transparent Fragmentation**

- Fragmentation is *transparent* to subsequent networks, through which the packet pass.
- Basic concept:
  - An oversized packet reaches a router, which breaks it up into fragments.
  - All fragments sent to the same exit router (say, R<sub>F</sub>).
  - R<sub>E</sub> reassembles the fragments before forwarding to the next network.
- Why called transparent?
  - Subsequent networks are not even aware that fragmentation had occurred.
- A packet may get fragmented several times.









# **Transparent Fragmentation (contd.)**

- Drawbacks:
  - All packets must be routed via the same exit router.
  - Exit router must know when all the pieces have been received.
    - Either a *count* field or *end-of-packet* field must be stored in each packet.
  - Lot of overhead.
    - A large packet may be fragmented and reassembled repeatedly.







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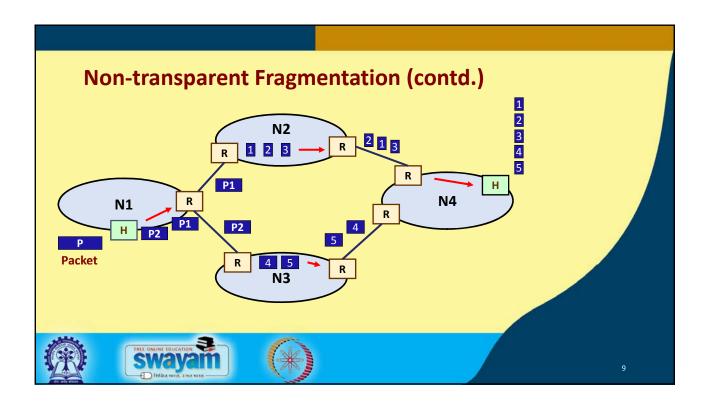
### **Non-transparent Fragmentation**

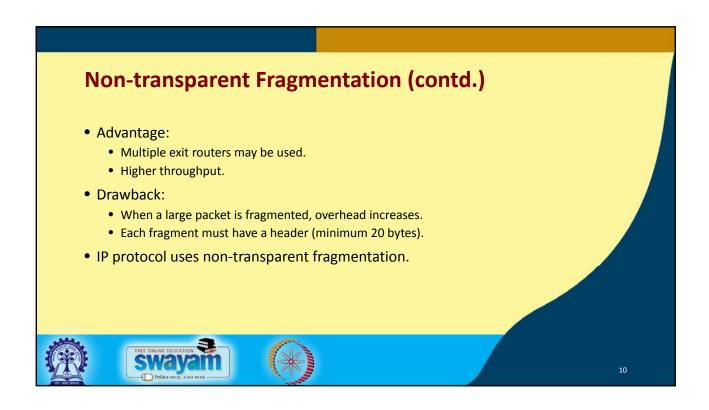
- Fragmentation is not transparent to subsequent networks.
- Basic concept:
  - Packet fragments are not reassembled at any intermediate router.
  - Each fragment is treated as an independent packet.
  - The fragments are reassembled at the final destination host.
- IP uses this philosophy.

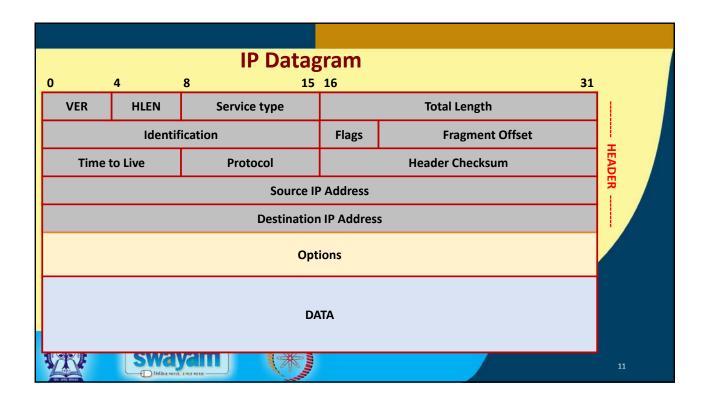


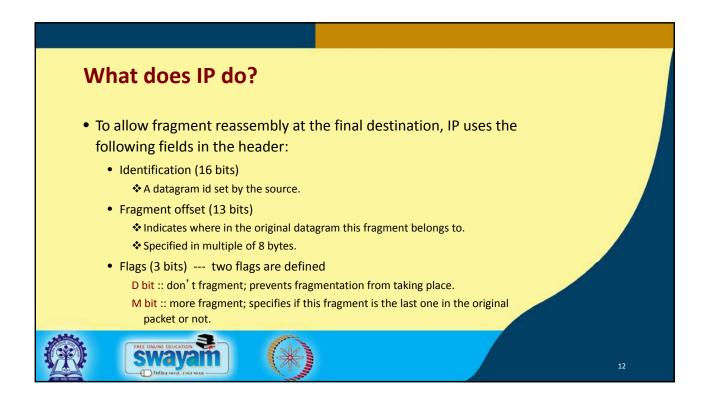


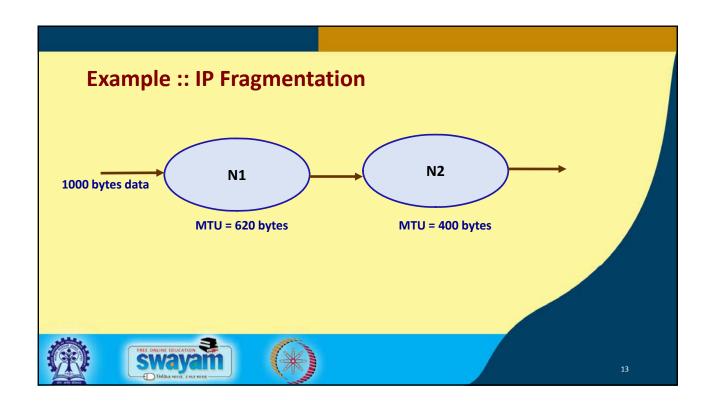


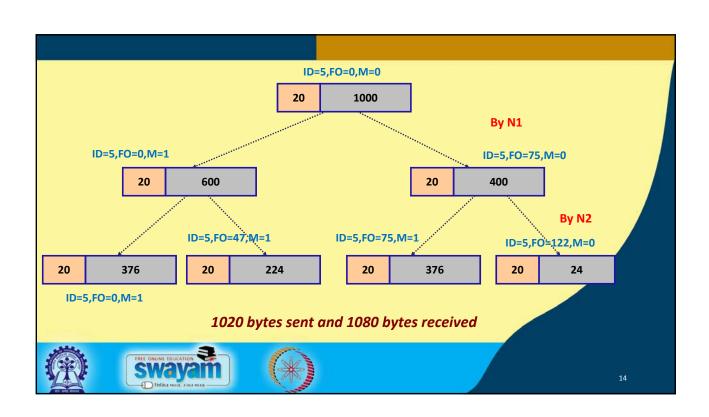






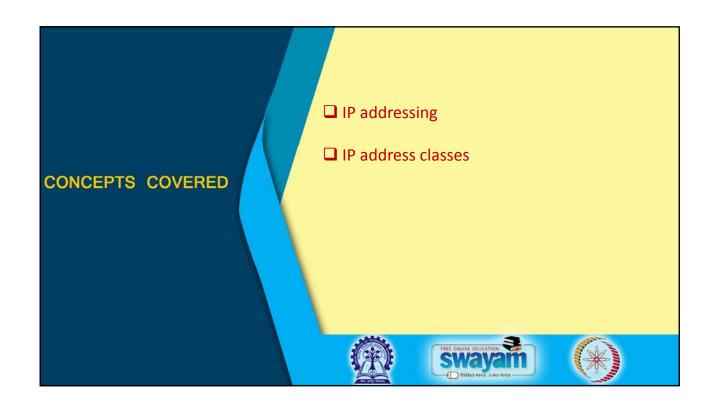












# **Basic IP Addressing**

- Each host connected to the Internet is identified by a unique IP address.
- An IP address is a 32-bit quantity.
  - Expressed as a dotted-decimal notation W.X.Y.Z, where dots are used to separate each of the four octets of the address.
  - Consists of two logical parts:
    - a) A network number
    - b) A host number
  - This partition defines the IP address classes.







### **Hierarchical Addressing**

- A computer on the Internet is addressed using a two-tuple:
  - The network number
    - ❖ Assigned and managed by central authority.
  - The host number
    - ❖ Assigned and managed by local network administrator.
- When routing a packet to the destination network, only the network number is looked at.







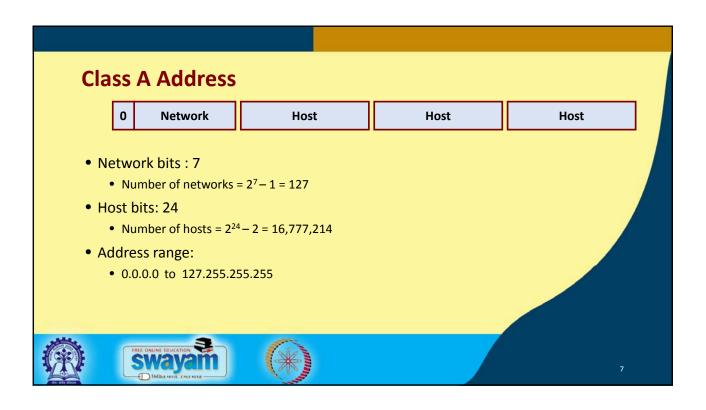
### **IP Address Classes**

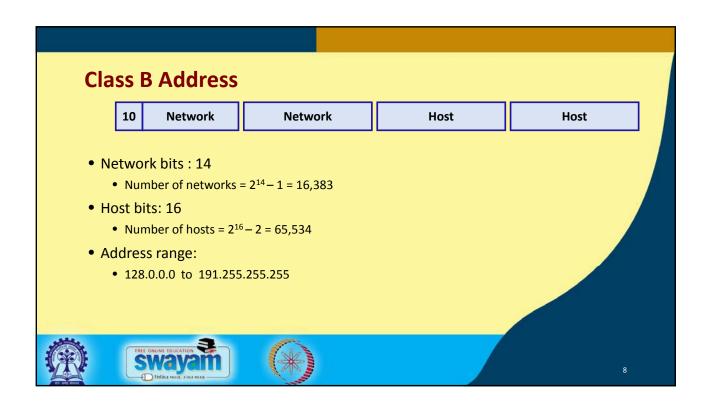
- There are five defined IP address classes.
  - Class A UNICAST
  - Class B UNICAST
  - Class C UNICAST
  - Class D MULTICAST
  - Class E RESERVED
- Identified by the first few bits in the IP address.
- There also exists some special-purpose IP addresses.
- The class-based addressing is also known as the *classful model*.

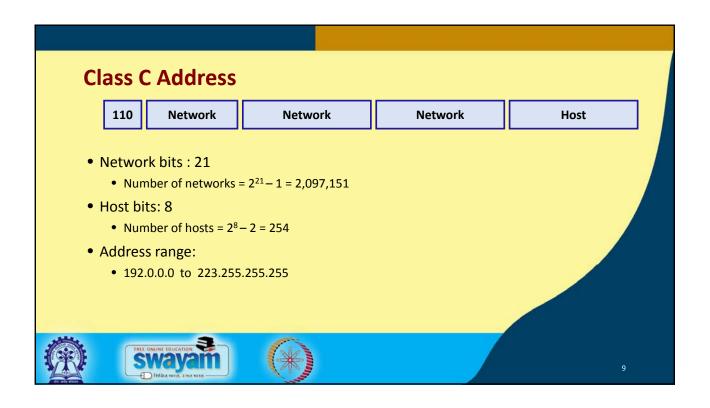




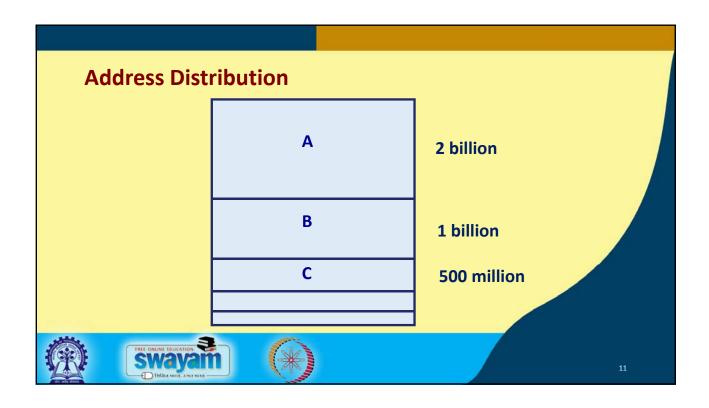


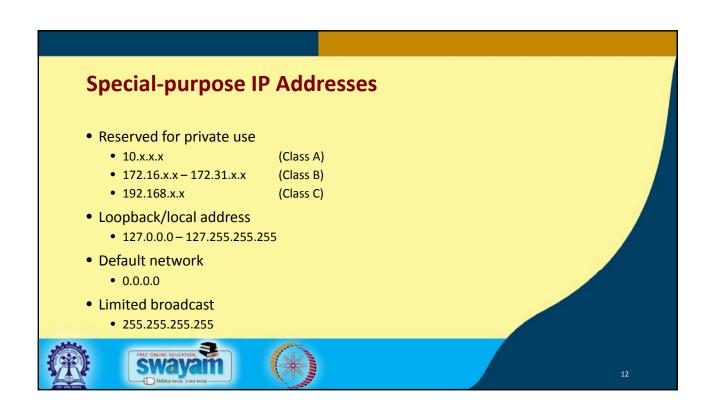












### **Some Conventions**

- Within a particular network (Class A, B or C), the first and last addresses serve special functions.
  - The first address represents the *network number*.
    - ❖ For example, 118.0.0.0
  - The last address represents the directed *broadcast address* of the network.
    - For example, 118.255.255.255

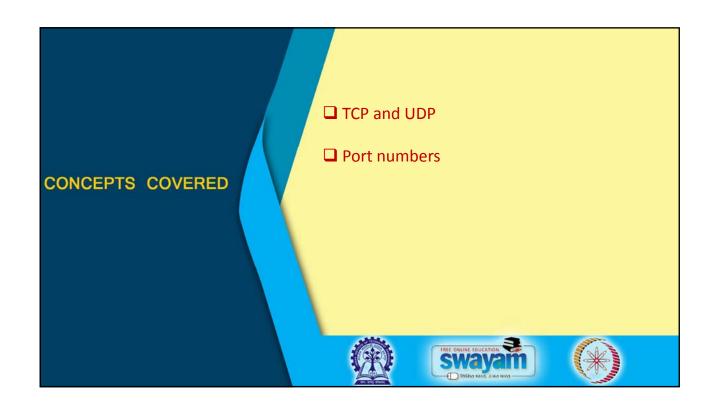












### Introduction

- In TCP/IP, the transport layer consists of two different protocols.
  - a) Transmission control protocol (TCP).
  - b) User datagram protocol (UDP).
- Basic idea:
  - User processes (applications) interact with the TCP/IP protocol suite by sending/receiving TCP or UDP data.
  - Both TCP and UDP in turn uses the IP layer for delivery of packets.







TCP and UDP

User Process

User Process

UDP

Datalink and Hardware Layer (e.g., Ethernet)

### **Role of TCP**

- Provides a connection-oriented, reliable, full-duplex, byte-stream service.
  - Underlying IP layer is unreliable and provides connectionless delivery service.
  - TCP provides end-to-end reliability using
    - Checksum
    - Positive acknowledgements
    - Timeouts
    - End-to-end flow control.
- TCP also handles
  - Establishment and termination of connections between processes.
  - Sequencing of data that might reach the destination in any arbitrary order.







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### **Role of UDP**

- UDP provides a connectionless and unreliable datagram service.
  - Very similar to IP in this respect.
  - Provides two features that are not there in IP:
    - ❖ A checksum to verify the integrity of the UDP packet.
    - ❖ Port numbers to identify the processes at the two ends.







### **Port Numbers**

- Multiple user processes on a machine may use TCP or UDP at the same time.
- There is need for a mechanism to uniquely identify the data packets associated with each process.





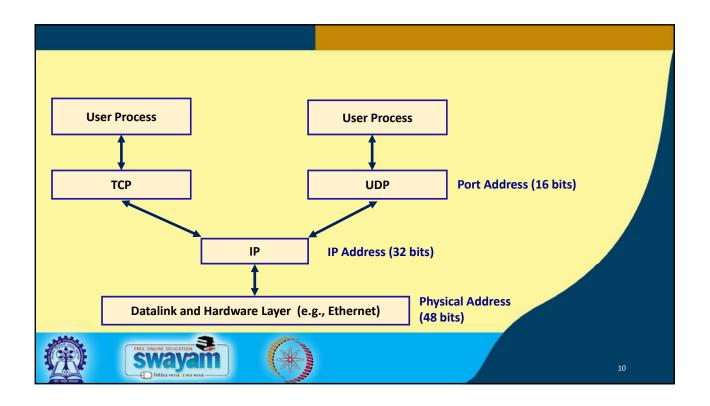


Port Numbers (contd.)

Port 10
Process 1
Process 2
Process 3

A host on the Internet

# Port Numbers (contd.) • How this is done? • Both TCP and UDP uses 16-bit integer port numbers. • Different applications are identified by different port numbers. • Port numbers are stored in the headers of TCP or UDP packets.



# Port Numbers (contd.)

- Client-server scenario
  - By knowing the 32-bit IP address of the server host, a client host can connect to the server.
  - To identify a particular process running on the server host, the client must also know the corresponding port number.
- Well-known port numbers
  - Predefined, and publicly known.
  - FTP uses port 21, SMTP uses port 25.







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### **Port Numbers (contd.)**

- Well-known port numbers are stored in a particular file on the host machine.
  - Unix:: /etc/services
  - Windows:: C:\WINDOWS\system32\drivers\etc\services
  - Each line has the format:

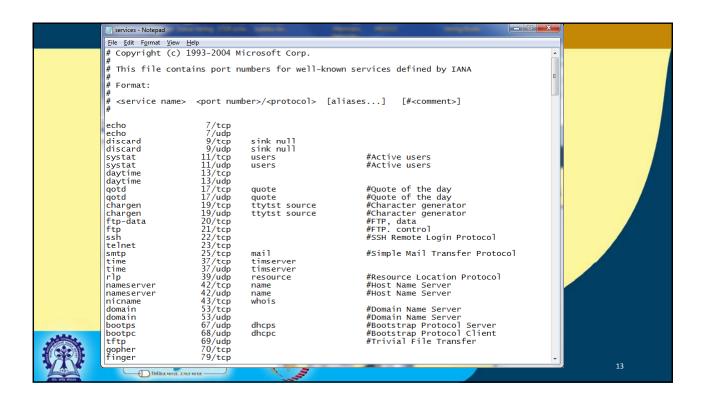
<service name> <port number>/<protocol> [aliases...] [#<comment>]

• Few lines of the file are shown next.









### **Ephemeral Port Numbers**

- A typical scenario:
  - A client process sends a message to a server process located on some host at port 1534.
  - How will the server know where to respond?
    - ❖ Client process requests an unused port number from the TCP/UDP module on its local host.
    - ❖ These are temporary port numbers, called *ephemeral port numbers*.
    - ❖ Send along with the TCP or UDP header.
- How are the port numbers assigned?
  - Port numbers from 1 to 1023 are reserved for well-known ports.
    - ❖ Has been extended to 4095.
  - Numbers beyond this and up to 65535 used as ephemeral port numbers.







### **Connection Establishment**

- A hierarchical addressing scheme is used to define a connection path between two hosts.
  - IP address
    - Identifies the communicating hosts.
  - Protocol identifier
    - ❖ Identifies the transport later protocol being used (TCP, UDP or anything else).
  - Port number
    - ❖ Identifies the communicating processes in the two hosts.







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### **Association**

- A set of five values that describe a unique process-to-process connection is called an *association*.
  - The protocol (TCP or UDP).
  - Local host IP address (32-bit value).
  - Local port number (16-bit value).
  - Remote host IP address (32-bit value).
  - Remote port number (16-bit value).
- Example of an association:

{TCP, 144.16.192.5, 1785, 144.16.202.57,21}

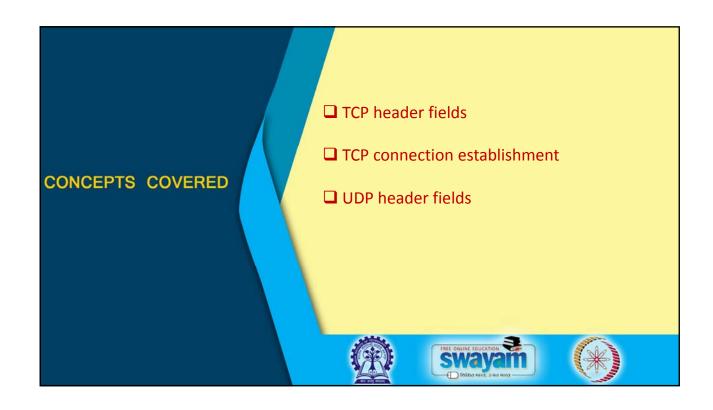


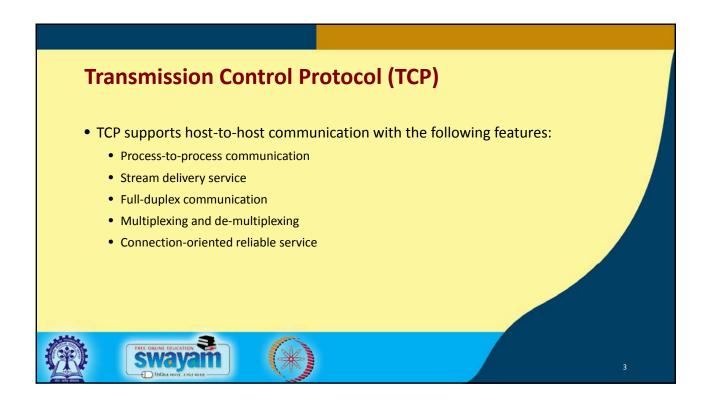


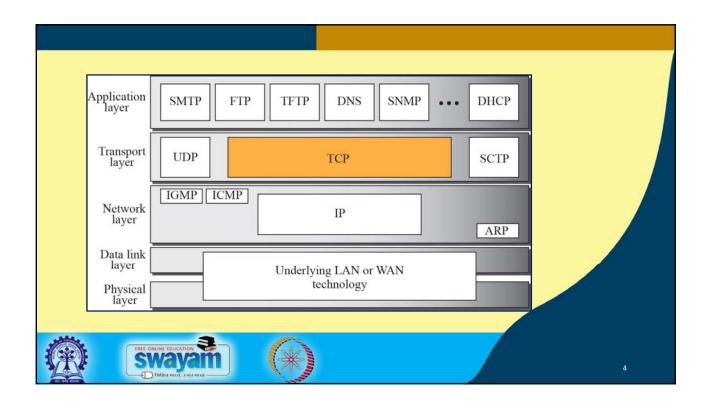


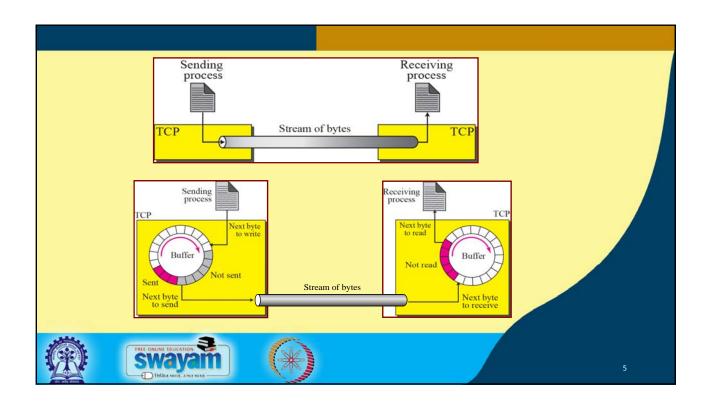


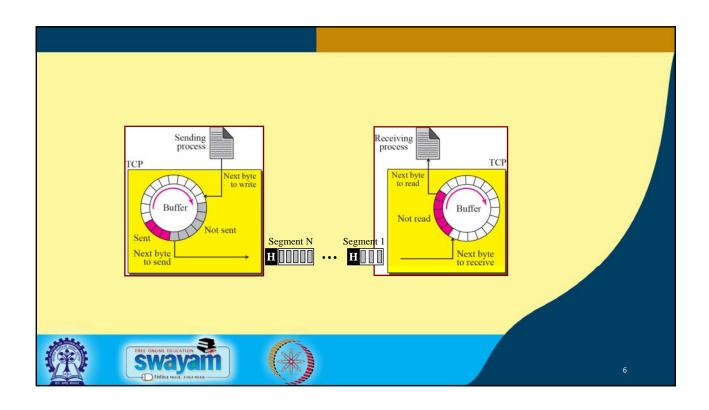


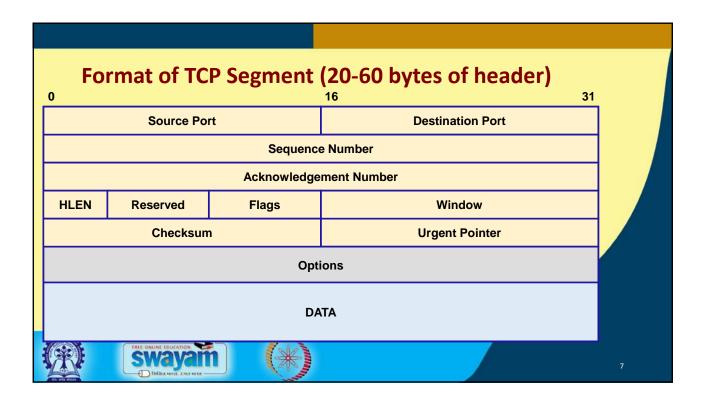












### **TCP Header Fields**

- Source port (16 bits)
  - Identifies the process at the local end.
- Destination port (16 bits)
  - Identifies the process at the remote end.
- Sequence number (32 bits)
  - Used for reliable delivery of message.
  - Each byte of message is assigned a 32-bit number that is incremented sequentially.
  - The field holds the number of the first byte in that TCP segment.







# **TCP Header Fields (contd.)**

- Acknowledgement Number (32 bits)
  - Used by remote host to acknowledge receipt of data.
  - Contains the number of the next byte expected to be received.
- HLEN (4 bits)
  - Specifies the header length in number of 32-bit words.







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### **TCP Header Fields**

- Flags (6 bits)
  - There are six flags.
    - ❖ URG is set to 1 if the urgent pointer is in use.
    - ❖ A connection request is sent by making SYN=1 and ACK=0.
    - ❖ A connection is confirmed by sending SYN=1 and ACK=1.
    - ❖ When the sender has no more data, FIN=1 is sent to release the connection.
    - $\ensuremath{ \diamondsuit }$  RST bit is used to reset a connection. It is also used to reject a connection attempt.
    - ❖ PSH bit indicates the push function. Used to indicate end of message.







# **TCP Header Fields (contd.)**

- Window (16 bits)
  - Specifies how many bytes may be sent beyond the byte acknowledged.
  - This number, called window advertisement, can increase or decrease as needed.
  - A value of zero closes the window altogether.







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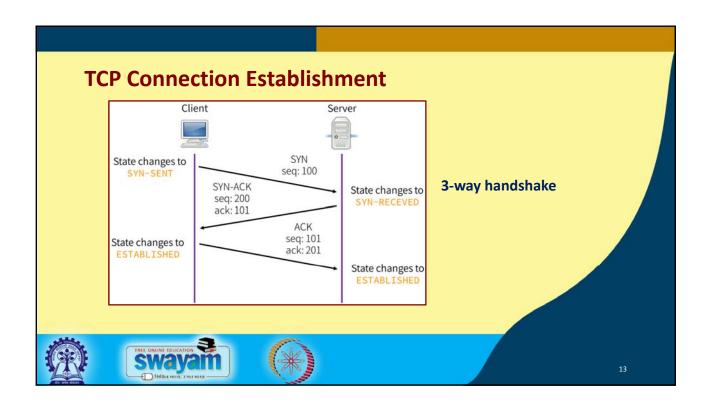
# **TCP Header Fields (contd.)**

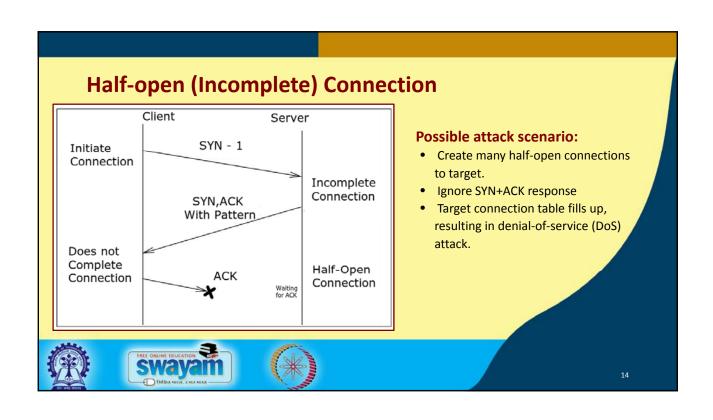
- Checksum (16 bits)
  - Applies to the entire segment and a pseudo-header.
  - The pseudo-header contains the following IP header fields:
    - ❖ Source IP address, destination IP address, protocol, segment length.
    - ❖ TCP protects itself from misdelivery by IP (delivered to wrong host).
  - Same algorithm as used in IP.

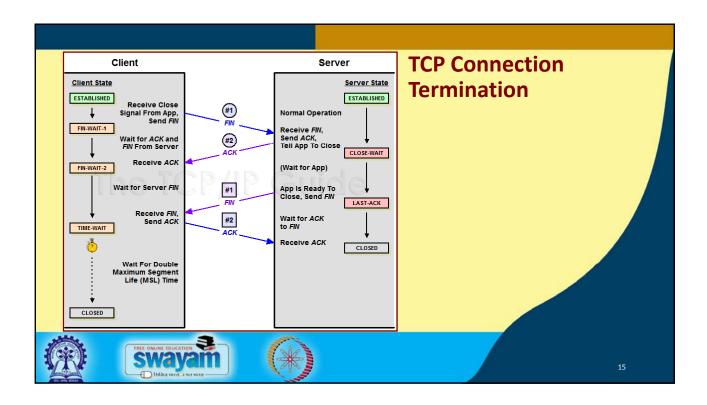


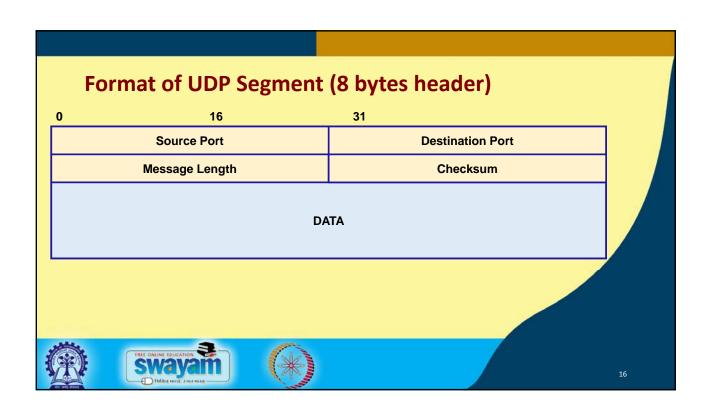


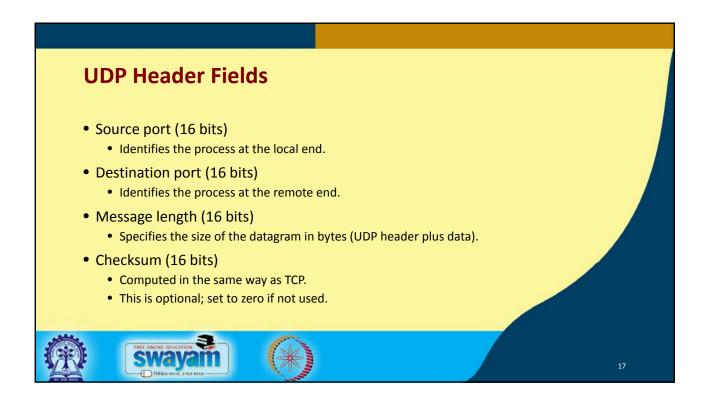






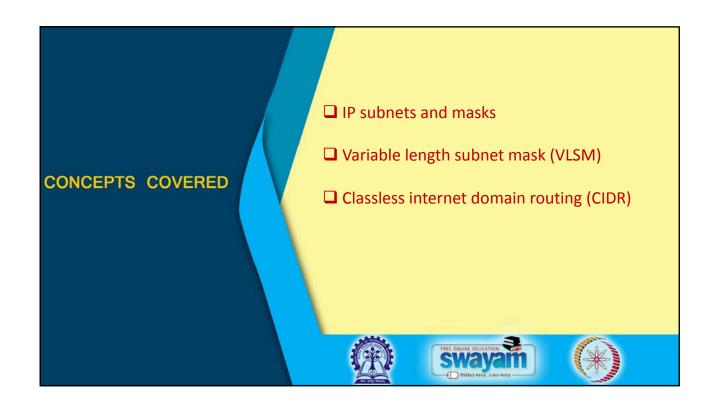












### **IP Subnet**

- A subnet is a subset of a class A, B or C network.
- IP addresses without subnets consists of a network portion, and a host portion.
  - Represents a static two-level hierarchical addressing model.
- IP subnets introduces a third level of hierarchy.
  - a) a network portion
  - b) a subnet portion
  - c) a host portion
- Allow more efficient (and structured) utilization of the addresses.
- Uses network masks.







# Natural Masks • Network mask 255.0.0.0 is applied to a class A network 10.0.0.0. • In binary, the mask is a series of contiguous 1's followed by a series of contiguous 0's. 11111111 00000000 00000000 000000000 Network portion Host portion

# **Natural Masks (contd.)**

- Provide a mechanism to split the IP address 10.0.0.20 into
  - a network portion of 10, and
  - a host portion of 20.

<u>Decimal</u> <u>Binary</u>

Network Host







# **Natural Masks (contd.)**

- Class A, B and C addresses
  - Have fixed division of network and host portions.
  - Can be expressed as masks.
    - Called *natural masks*.
- Natural Masks

Class A :: 255.0.0.0Class B :: 255.255.0.0Class C :: 255.255.255.0







# **Creating Subnets using Masks**

- Masks are very flexible.
  - Using masks, networks can be divided into smaller subnets.
  - By extending the network portion of the address into the host portion.
- Advantage:
  - We can create a large number of subnets from one network.
  - Can have less number of hosts per network.







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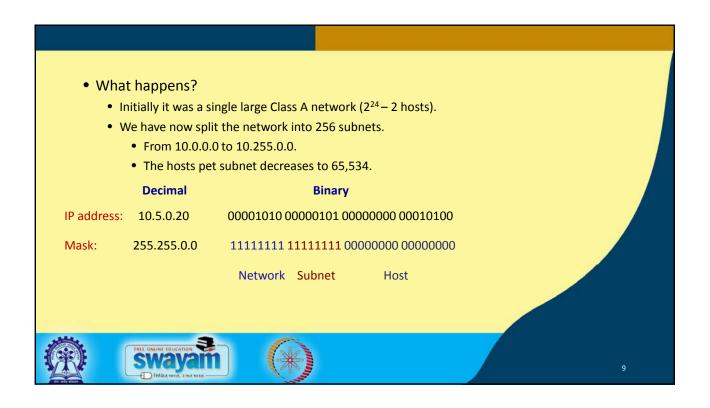
### **Example: Subnets**

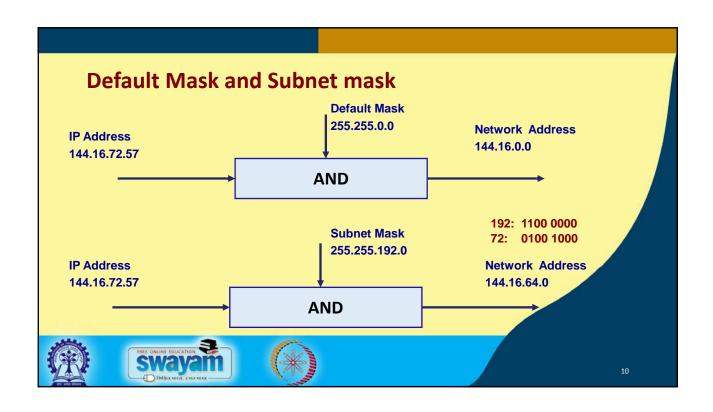
- Network mask 255.255.0.0 is applied to a class A network 10.0.0.0.
  - This divides the IP address 10.5.0.20 into
    - a network portion of 10,
    - a subnet portion of 5, and
    - a host portion of 20.
- The 255.255.0.0 mask borrows a portion of the host space, and applies it to network space.











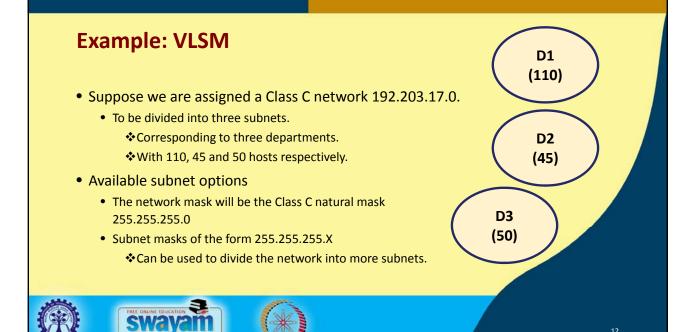
# **Variable Length Subnet Masks (VLSM)**

- Basic concept
  - The same network can be configured with different masks.
  - Can have subnets of different sizes.
  - Allows better utilization of available addresses.









# **The Subnet Options**

| Х   | X (in binary) | No. of Subnets | No. of Hosts |
|-----|---------------|----------------|--------------|
| 128 | 1000 0000     | 2              | 128          |
| 192 | 1100 0000     | 4              | 64           |
| 224 | 1110 0000     | 8              | 32           |
| 240 | 1111 0000     | 16             | 16           |
| 248 | 1111 1000     | 32             | 8            |
| 252 | 1111 1100     | 64             | 4            |

Cannot satisfy the requirements.







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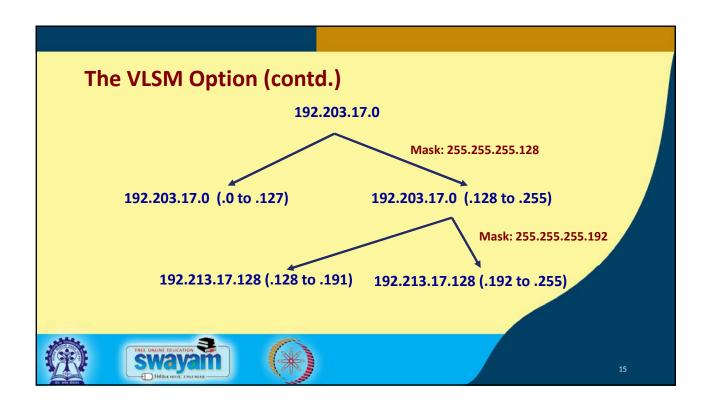
# **The VLSM Option**

- Basic concept:
  - Use the mask 255.255.255.128 to divide the network address into two subnets with 128 hosts each.
    - 192.203.17.0 (.0 to .127)
    - 192.203.17.0 (.128 to .255)
  - Next subnet the second .128 subnet using a mask of 255.255.255.192. (Creates two subnets, 64 hosts each)
    - 192.213.17.128 (.128 to .191)
    - 192.213.17.128 (.192 to .255)









### **Classless Internet Domain Routing (CIDR)**

- CIDR is a new concept to manage IP networks.
  - Classless Inter Domain Routing.
  - No concept of class A, B, C networks.
  - Reduces sizes of routing tables.
- An IP address is represented by a prefix, which is the IP address of the network.
- It is followed by a slash, followed by a number M.
  - M: number of leftmost contiguous bits to be used for the network mask.
  - Example: 144.16.192.57 / 18







### **CIDR: An Important Rule**

- The number of addresses in each block must be a power of 2.
- The beginning address in each block must be divisible by the number of addresses in the block.
  - A block that contains 16 addresses cannot have beginning address as 144.16.223.36.
  - But the address 144.16.192.64 is possible.







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### **Example: CIDR**

An organization is allotted a block with beginning address:

144.16.192.24 / 29

What is the range of the block?

 ${\sf Start\ addr:\ } 10010000\ 00011000\ 11000000\ 00011000$ 

End addr: 10010000 00011000 11000000 00011111

There are 8 addresses in the block.







