

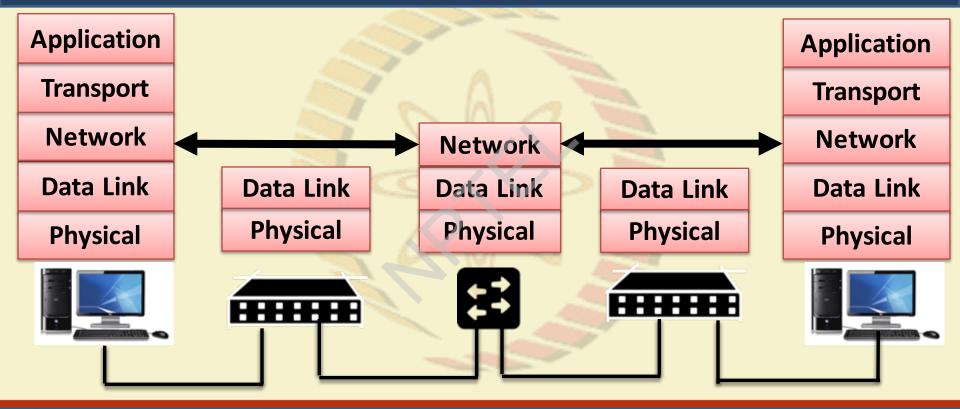


COMPUTER NETWORKS AND INTERNET PROTOCOLS

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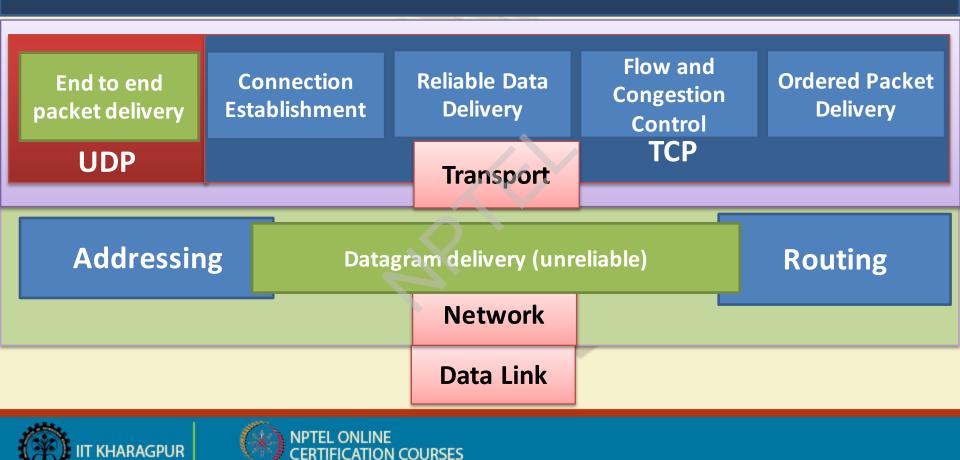
Network Layer I - Introduction

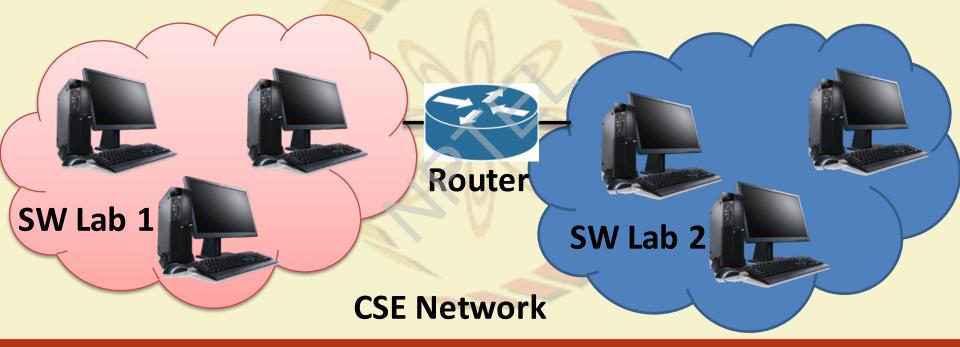






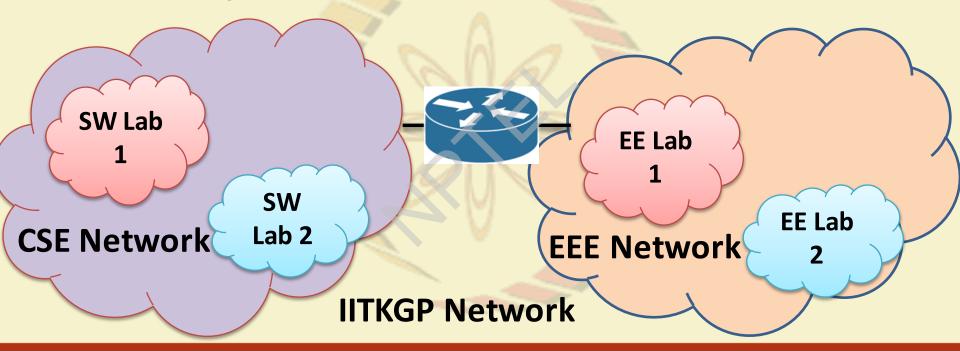
Network (Internet) Layer Services





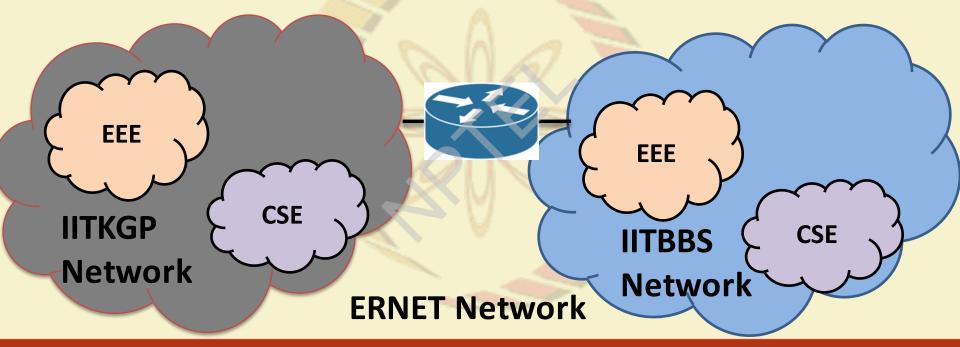






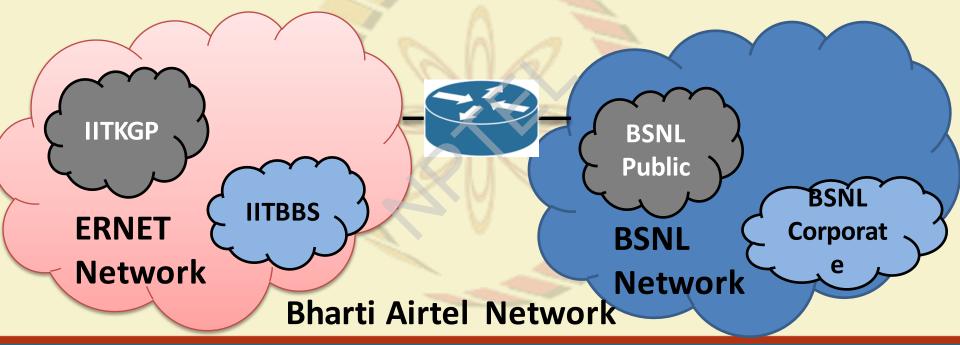












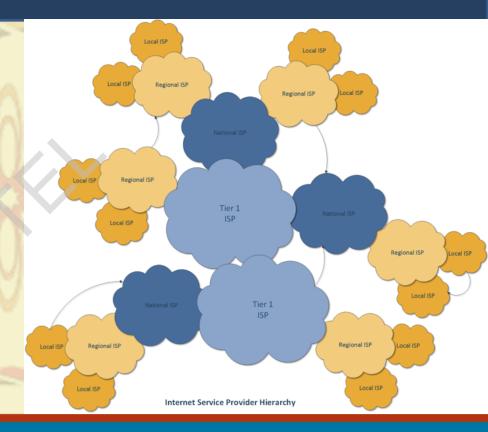




Internet Architecture

Autonomous Systems (AS) – A set of LANs for an administrative domain, identified by a unique AS number, and the routing policies are controlled by a single administrator.

Local Area Network (LAN) – A set of devices with a common layer 3 gateway

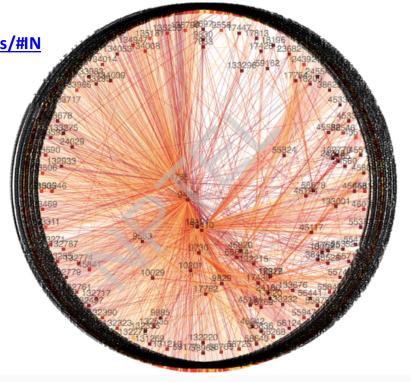






Autonomous System Graph for India

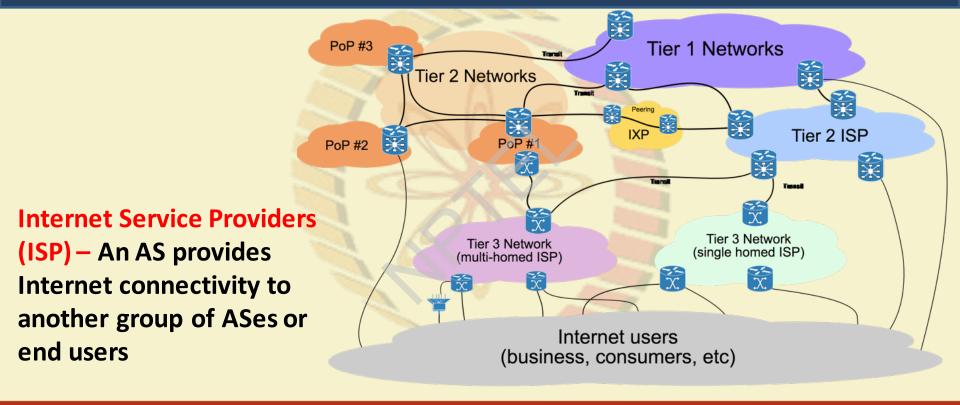
Source: https://labs.apnic.net/vizas/#IN







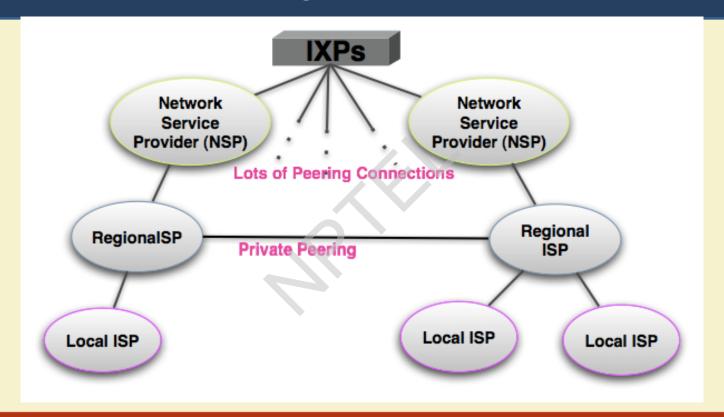
Internet Architecture







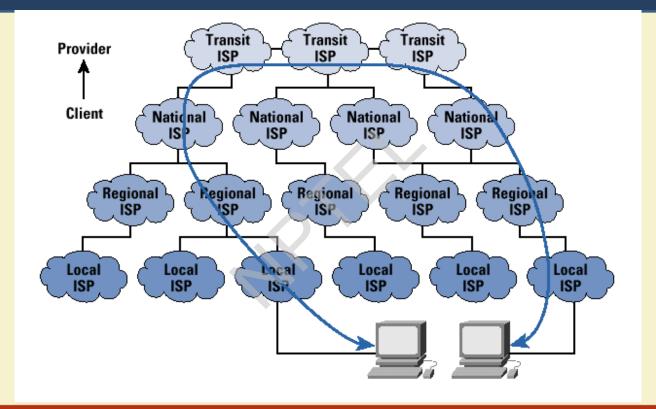
Peering between ISPs







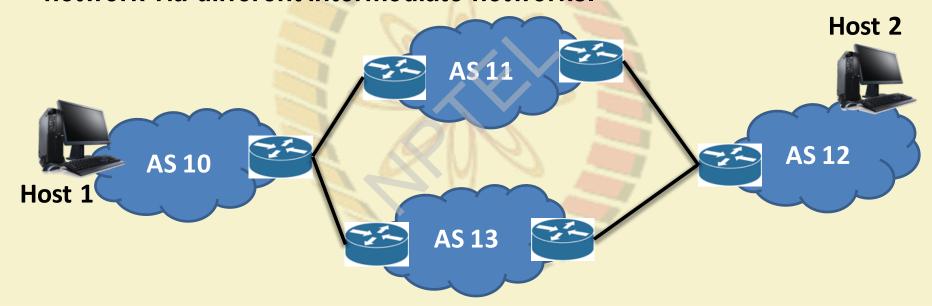
Communication between Two Nodes over ISPs





IP Addressing – Basic Principles

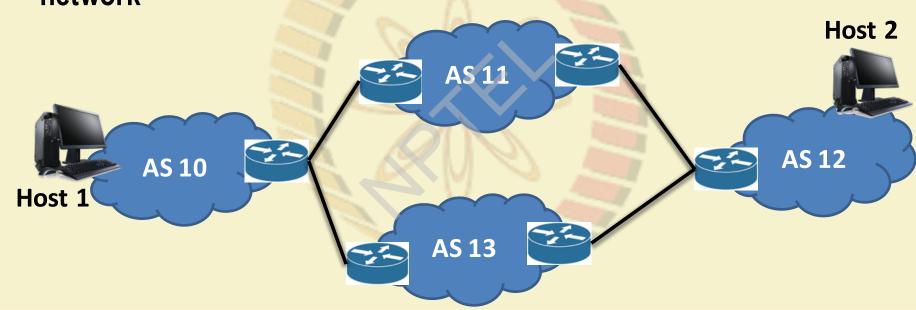
 We need to forward data packets from one network to another network via different intermediate networks.





IP Addressing – Basic Principles

The address should identify a network as well as a host inside a network















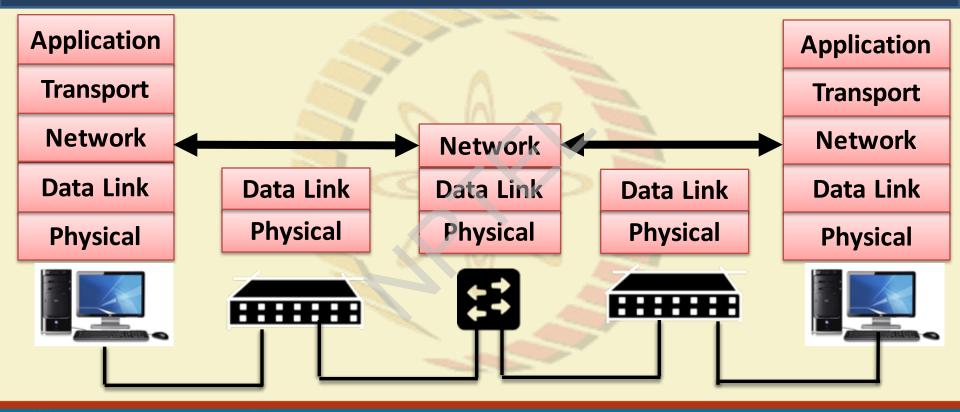


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Network Layer II - IPv4 Addressing (Classful Addresses)

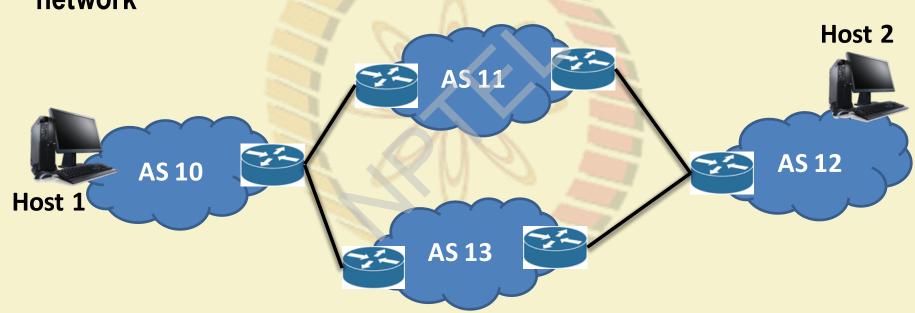






IP Addressing – Basic Principles

The address should identify a network as well as a host inside a network





IP Addressing

Network address

Host address

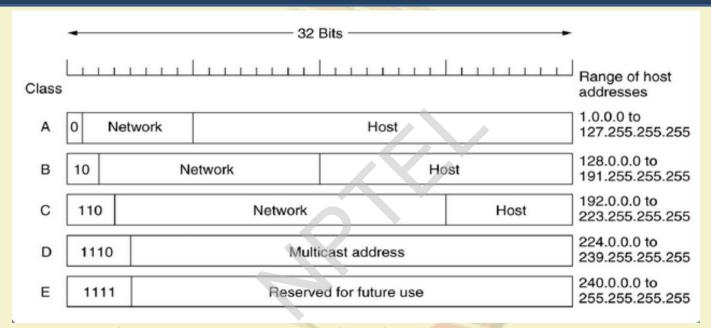
 Divide the address space (32 bit in IPv4) among network address and host address

 The old age – Classful addressing: Fixed number of bits for network address and host address





Classful Addressing



- How to identify a class use the first few bits
 - 0 Class A, 10 Class B, 110 Class C, 1110 Class D, 1111 Class E





Network Address and Broadcast Address

- Network address identify a network
 - All 0's in the host address part

 - Ex-2 (Class B): 10111101.11101001.00000000.00000000 (189.233.0.0)
- Broadcast address send the data to all the hosts of a network
 - All 1's in the host address part

 - Ex-2 (Class B): 10111101.11101001.111111111.11111111 (189.233.255.255)
- How many valid hosts can be there in a Class A, in a Class B and in a Class C IP address?





Subnetting and Supernetting – Classless Inter-domain Routing (CIDR)

- You have 255 hosts in a network. Which IPv4 address class will you use Class C or Class B?
 - Class C not possible
 - Class B huge address space is lost (using only 255 addresses out of possible 2¹⁶-2 addresses)



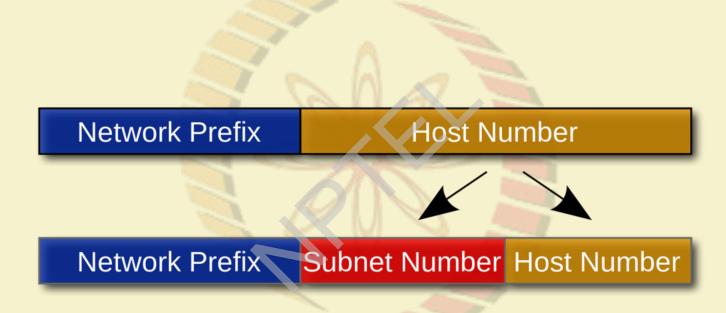
Subnetting and Supernetting – Classless Inter-domain Routing (CIDR)

- Split a large network or combine multiple small networks for efficient use of address space
 - Subnetting divide a large network into multiple small networks
 - Supernetting combine multiple small networks into a single large network

Subnet mask – denote the number of bits in the network address field



Divide a Network into Subnets



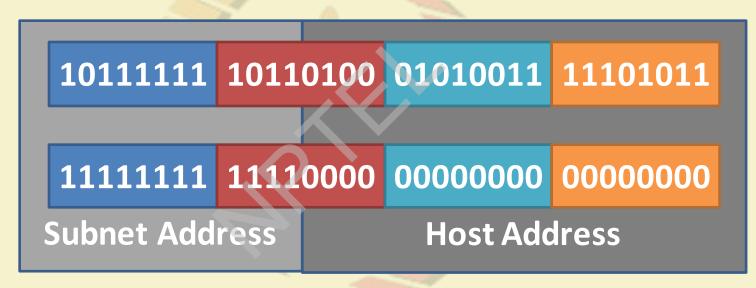




CIDR – Addressing Format

IP Address

Netmask







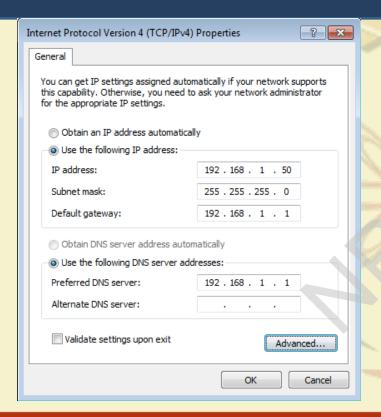
CIDR – Addressing Format

- We write the IP address as 191.180.83.235/12 in CIDR notation
 - The first 12 bits are the network address and rest (32-12)=20 bits are for host address

The subnet mask is 255.240.0.0



CIDR - Manual IP Setting in the OS

















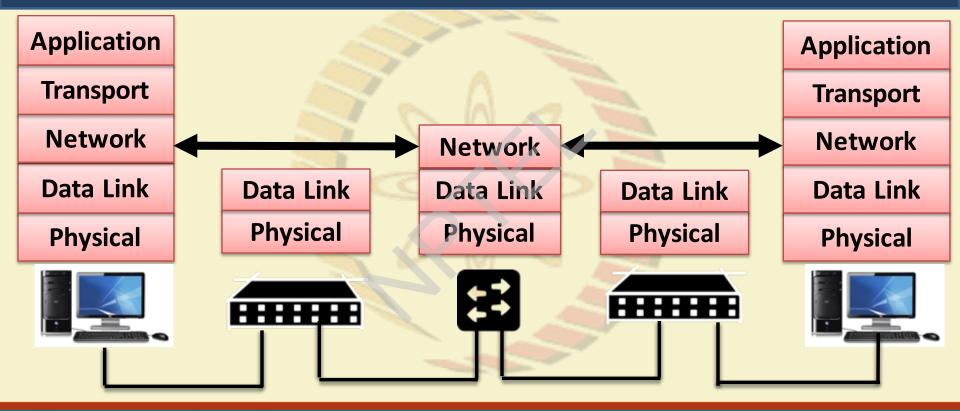


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Network Layer III - IPv4 Addressing (CIDR)



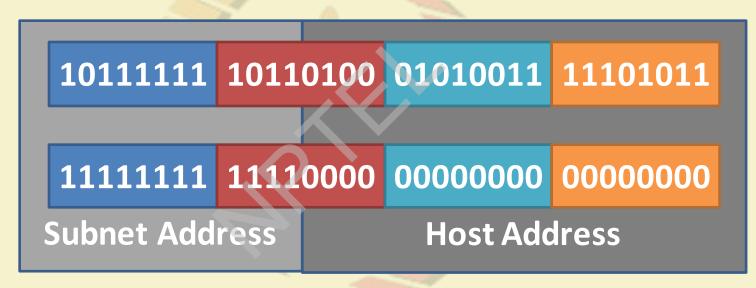




CIDR – Addressing Format

IP Address

Netmask







Divide a Network into Subnets

- Let the IP address of a network is 203.110.0.0/16
- We want to divide this network into three subnets
- We need 3 bits for subnets why not 2 bits?
 - Subnet 1 100, Subnet 2– 101, Subnet 3 110
- Rest 13 bits are used for addressing the hosts of those subnets.
- The subnets are 203.110.128.0/19, 203.110.160.0/19, 203.110.192.0/19



All Zero and All One Subnets

We normally avoid "all zero" and "all one" subnets.



192.168.0.0/16

192.168.0XXXXXXXXXX 192

192.168.0.0/17

The network address for the subnet and the original network is identical – Subnet Zero



192.168.128.0/17

192.168.1XXXXXXXXX

Broadcast address for this subnet is 192.168.255.255, broadcast address for the original network is also 192.168.255.255 – All-One Subnet





CIDR Example



CSE - 2000 Hosts





VGSOM - 500 Hosts



EE - 500 Hosts





203.110.0.0/19





CIDR Example



CSE – 2000 Hosts

11 bit hosts





VGSOM - 500 Hosts

9 bit hosts



EE – 500 Hosts

9 bit hosts





203.110.0.0/19





CIDR Example

- Address space 203.110.0.0/19
 - 13 bits are available to serve all the hosts of IITKGP network
 - We need to divide these address space among 3 subnets
- CSE 11 bits, VGSOM 9 bits, EE 9 bits for host address
- We have 2 bits left for identifying three subnets Is this possible?
 - Avoid "all zero" and "all one" subnets
- Let us apply CIDR Combine VGSOM and EE Networks together





CSE - 2000 Hosts

11 bit hosts





VGSOM - 500 Hosts 9 bit

hosts

10 bit hosts





9 bit EE - 500 Hosts

hosts





203.110.0.0/19





CSE – 11 bits, VGSOM+EE – 10 bits

CSE network address 203.110.00010XXX.XXXXXXXXX (203.110.16.0/21)







CSE – 2000 Hosts

11 bit hosts

203.110.16.0/21



hosts



9 bit VGSOM - 500 Hosts hosts



DEPARTMENT OF ELECTRICAL ENGINEERING

9 bit EE - 500 Hosts

hosts



203.110.8.0/21



203.110.0.0/19





VGSOM – 9 bits, EE – 9 bits

EE network address 203.110.0000101X.XXXXXXXXX (203.110.10.0/23)







CSE – 2000 Hosts

11 bit hosts



hosts

203.110.16.0/21



VGSOM - 500 Hosts 9 bit

hosts

203.110.12.0/23



9 bit EE - 500 Hosts

hosts



203.110.10.0/23



10 bit



203.110.0.0/19













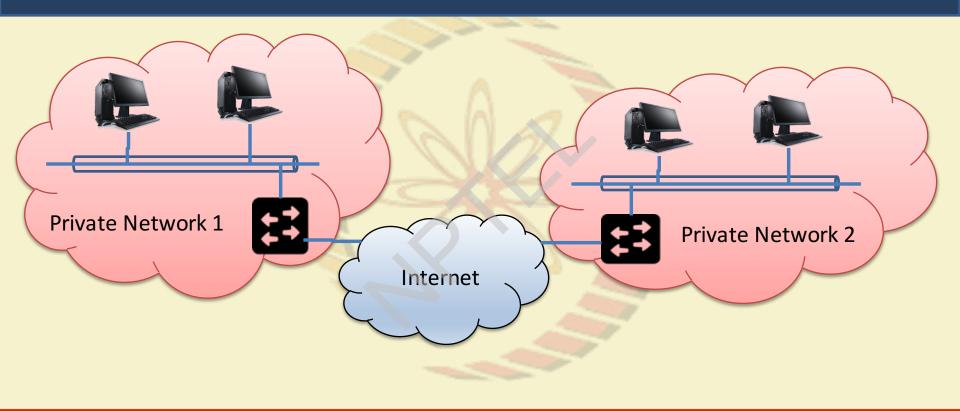


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Network Address Translation (NAT)







Issues with IPv4 Addressing

 The address space is limited - number of devices (networking equipment) are increasing exponentially.

 A large number of addresses are wasted or remain unutilized (Class D or Class E).

Solution: Make the address reusable, leveraging on the fact that not all
users or all devices will connect to the Internet at the same time.

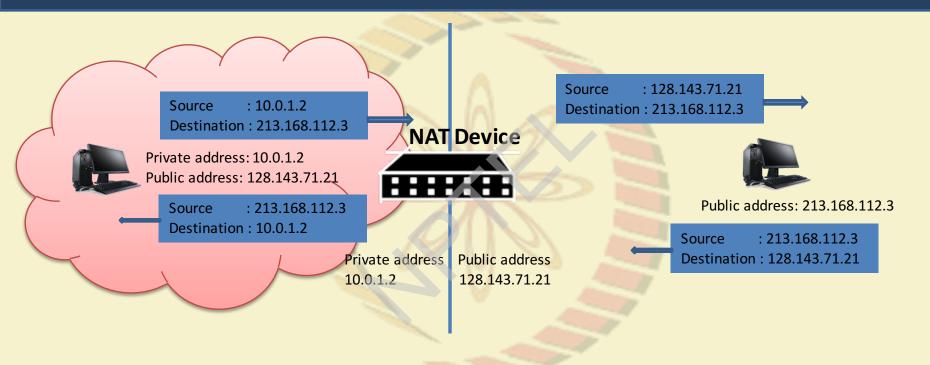


Network Address Translation (NAT)

- Divide addresses into reusable (private) and non-reusable (public) blocks
- Translate internal (private) addresses to external (public) addresses
- Hide internal machines from external devices
- Allow Internet access to large number of users via few public addresses
 - IPv4 private address
 - 10.0.0.0-10.255.255.255
 - 172.16.0.0-172.32.255.255
 - 192.168.0.0-192.168.255.255



Basic Operation of NAT







Working Principles of NAT

- Organizations manages internal private network
- NAT boxes manages a pool of public IP address
- For outgoing connections, NAT boxes selects one of the IP address from its pool, and forward packet from that IP

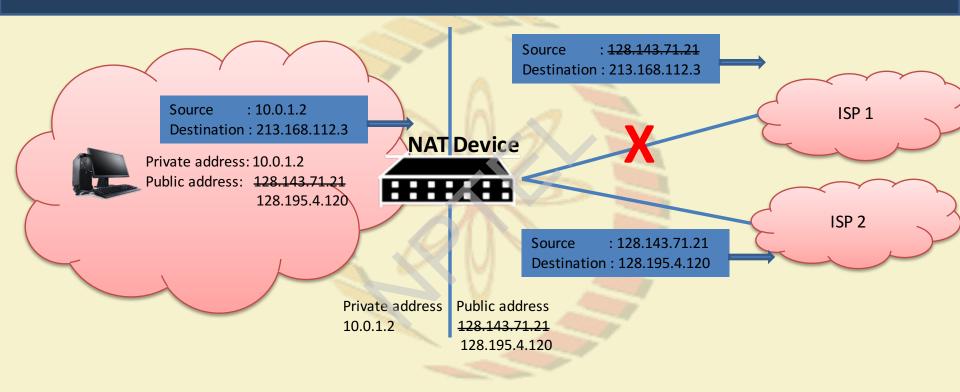


Migration between ISPs

- An organization can connect to multiple ISPs for better reliability
- NAT allow easy interchange between ISPs by changing IP addresses in NAT boxes
 - Without NAT, every internal system address need to be changed to reflect the network IP of the ISP
- NAT box can be configured to use alternative ISPs in case of a failure



Migration between Network Service Provider





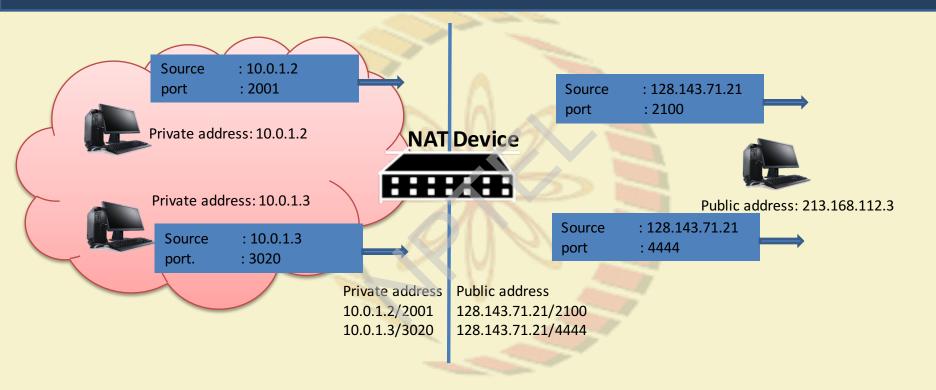


IP Masquerading

- Single public IP address is mapped to multiple hosts
- NAT box modify port address and replace private IP address to public IP address
 - Keep mapping in a table to forward incoming packet to proper internal host



IP Masquerading



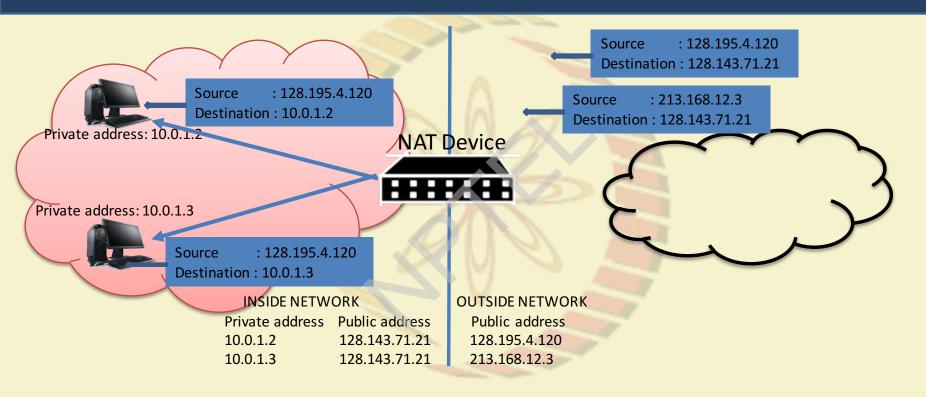


Load Balancing of Servers

- Balance the load of multiple identical server accessible from a single IP address
- NAT box translate different incoming connection to different internal IP address to balance load between server
- Internal systems are configured with private addresses



Load Balancing of Servers















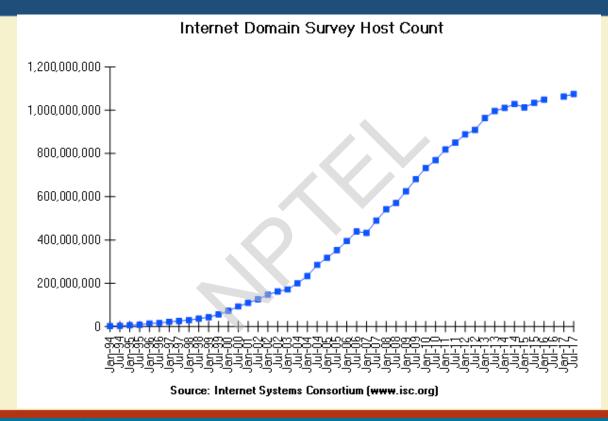


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Internet Protocol Version 6 (IPv6)



Why do We Need a New IP Structure?

- Address space is not sufficient even with CIDR.
- QoS is vaguely defined need real time service support for modernday applications
- Mobile applications are unmanageable
- There is no direct security support in IPv4

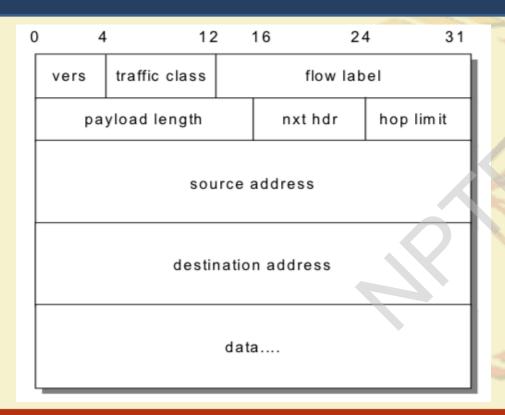
IPv6 Features

- Larger address space
- Globally unique and hierarchical addressing
- Optimized routing table using prefixes rather than address classes
- Auto-configuration of network interfaces
- Support for encapsulation
- Service class support to manage QoS classes
- Built-in authentication and encryption
- Compatibility with IPv4





IPv6 Header Format



128 bit source address and destination address





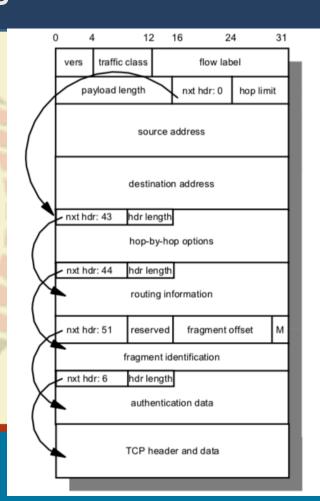
Extension Headers

 Additional information are transmitted through the extension headers.

The base header points to the extension headers







IPv6 Addressing

• 128 bit addresses - represented in 8 hexadecimal numbers FE80:0000:0000:0000:0001:0800:23E:F5DB

Leading zeros can be omitted - FE80:0:0:0:1:800:23E7:F5DB

 A group of zeros can be replaced by a double colon -FE80::1:800:23E7:F5DB (Can be used only once)





Address Space Allocation based on Prefix

Allocation	Prefix (bin)	Start of address range (hex)	Mask length (bits)	Fraction of address space
Reserved	0000 0000	0:: /8	8	1/256
Reserved for NSAP	0000 001	200:: /7	7	1/128
Reserved for IPX	0000 010	400:: /7	7	1/128
Aggregatable global unicast addresses	001	2000:: /3	3	1/8





Address Space Allocation based on Prefix

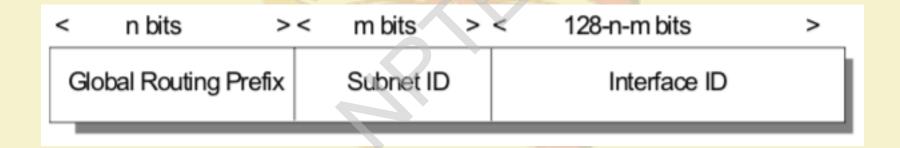
Allocation	Prefix (bin)	Start of address range (hex)	Mask length (bits)	Fraction of address space
Link-local unicast	1111 1110 10	FE80::/10	10	1/1024
Site-local unicast	1111 1110 11	FEC0:: /10	10	1/1024
Multicast	1111 1111	FF00::/8	8	1/256
Total allocation				15%





Global Unicast Address Format

 Global routing prefix: A value assigned to a site for a cluster of subnets/links. The global routing prefix is designed to be structured hierarchically







ICMPv6 - Neighbor Discovery

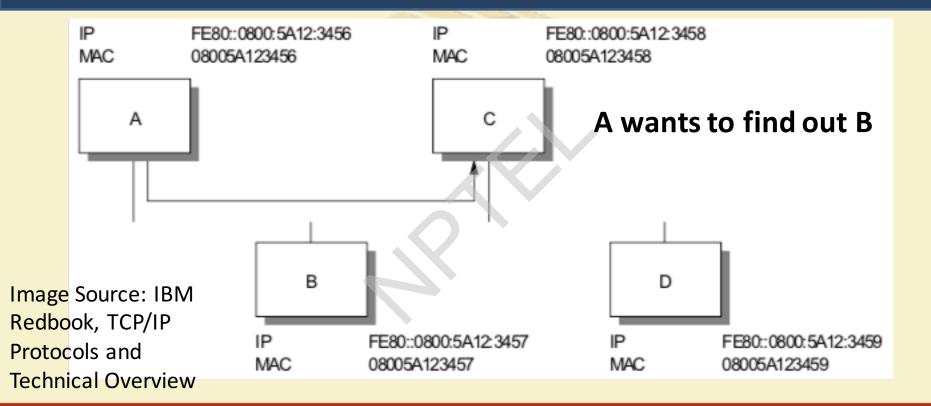
Enables a node to identify other hosts and routers on its links

 The node needs to know of at least one router so that it knows where to forward packets if a target node is not on its local link

ARP request/response in IPv4



Neighbor Discovery Example

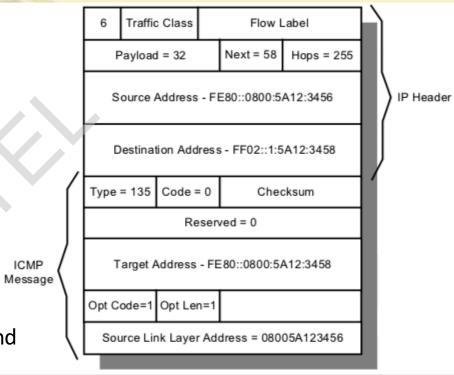






Neighbor Discovery - Neighbor Solicitation

- The destination address is the address of the solicitated node
- An improvement over ARP broadcast

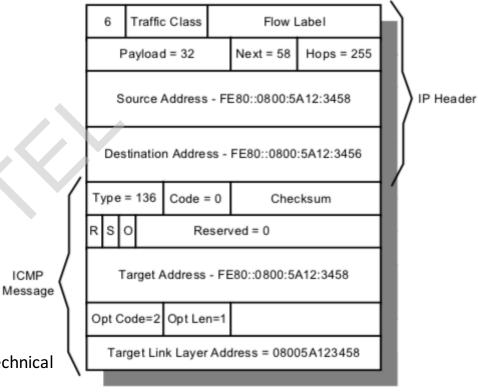






Neighbor Discovery - Neighbor Advertisement

- Response of the neighbor solicitation
- Three flags
 - R: Sender of the advertisement is a router
 - S: Advertisement is a response to a solicitation
 - O: Override, must update a cached information







IPv6 Mobility Support

- A mobile node uses a temporary address when it is away from the home location
 - Use IPv6 Destination Optional header to store its home address
- A mobile station can list the all routing header for the packets to follow a particular path for establishing connection with a service provider network
- Packets sent to a mobile node can be tunneled by IPv6 routing headers
- Do not require foreign agents like IPv4 neighbor discovery and address autoconfiguration can be used to connect a node with any network



Migrating from IPv4 to IPv6

Dual stack IP implementations

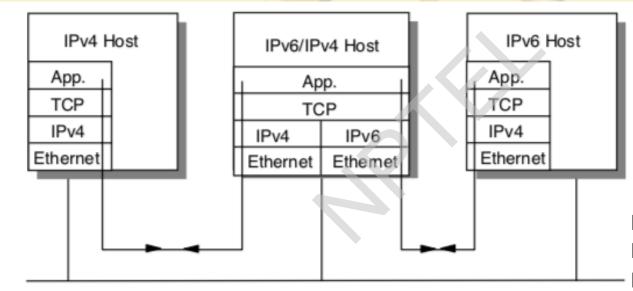


Image Source: IBM
Redbook, TCP/IP
Protocols and Technical

Overview





Migrating from IPv4 to IPv6

Tunneling: Tunnel IPv4 headers through IPv6 headers and vice-versa

- Header Translation: Translate a IPv4 header to a IPv6 header and viceversa
 - Address must be translated as well
 - Take low order 32 bits for IPv6 to IPv4
 - Append ::FFFF/96 prefix for IPv4 to IPv6



Address Translation

- IPv4 Address: 202.141.80.20
- IPv6 Address: CA8D:5014::FFFF

- IPv6 Address: FE80:2381:0000:0000:0001:0800:23E:F5DB
- IPv4 Address: 254.128.35.129





Interesting Reads

- RFC 2460 Internet Protocol, Version 6 (IPv6) (December 1998)
- RFC 4291 IP Version 6 Addressing Architecture (February 2006)
- RFC 3587 IPv6 Global Unicast Address Format (August 2003)
- IANA Assignment Documentation: INTERNET PROTOCOL VERSION 6
 MULTICAST ADDRESSES, June 2006
 http://www.iana.org/assignments/ipv6-multicast-addresses
- 6NET http://www.6net.org









